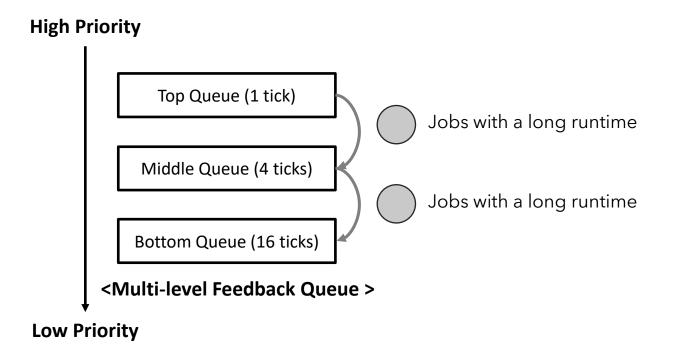
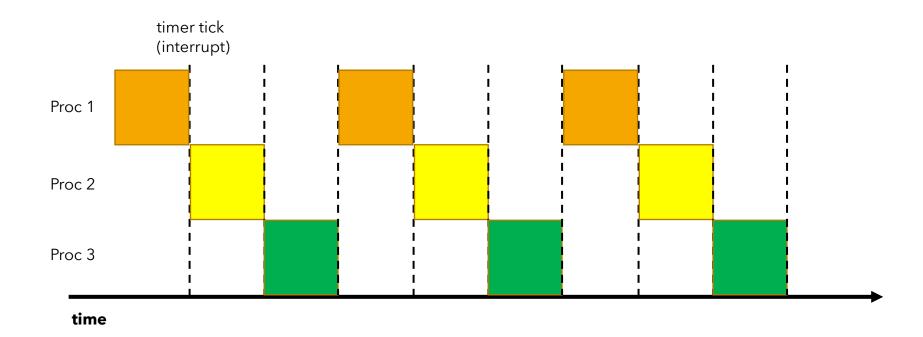


- Implement a basic MLFQ scheduler in the xv6
 - Modify the default xv6 scheduler to Multi-level Feedback Queue



- What is the default xv6 scheduler?
 - Round-robin scheduler
 - For each timer tick (~10ms), a timer interrupt occurs to incur a context switch
 - Change this round robin to the multi-level feedback!



(1 tick)

- Core function: void scheduler() in proc.c
- First for loop is looping forever Infinite loop
 - This function never returns
 - Find a new process to be scheduled, run it until it yields
- Scan ptable to find RUNNABLE process

```
10 struct {
11    struct spinlock lock;
12    struct proc proc[NPROC];
13 } ptable;
```

- Switch from scheduler to the process
 - After the process yields by timer interrupt, come back to swtch()

<default round-robin scheduler code>

```
oid scheduler(void) {
struct proc *p;
struct cpu *c = mycpu();
c \rightarrow proc = 0;
for(;;){
  sti();
  acquire(&ptable.lock):
   for(p = ptable.proc; p < &ptable.proc[NPROC]; p++){</pre>
    if(p->state != RUNNABLE)
       continue;
    c->proc = p;
     switchuvm(p);
     p->state = RUNNING;
     swtch(&(c->scheduler), p->context);
    switchkvm();
     c \rightarrow proc = 0;
  release(&ptable.lock);
```

- Objectives of this project
 - Understand how context-switches are performed in xv6 code
 - Implement a basic MLFQ scheduler
- Where to look and write code:
 - proc.c, proc.h (+ etc)

- Rules for MLFQ scheduler
 - I. There are 3 priority levels top, middle, bottom
 - 2. Whenever a timer tick interrupt occurs, a process in the top queue is scheduled
 - 1. Similarly, a process in the top queue is scheduled after previously running process yields the CPU or exits
 - 2. When a new process is created, it starts from top queue
 - 3. When scheduling, find a process in the lower queue if there are no processes in the upper queue
 - 3. Processes in the same queue are scheduled by a round-robin policy
 - 1. However, the time slice of each queue should be considered
 - 2. Time-slice: # of ticks allowed to each process according to the priority level
 - 4. Time-slice of top/middle/bottom queue is 1 tick / 4 ticks / 16 ticks
 - 5. If a process consumes its time-slice, the process goes to down (except for bottom queue)

• Example New process (A, B, C, D)

Top Queue (1 tick)

A
B
C
D

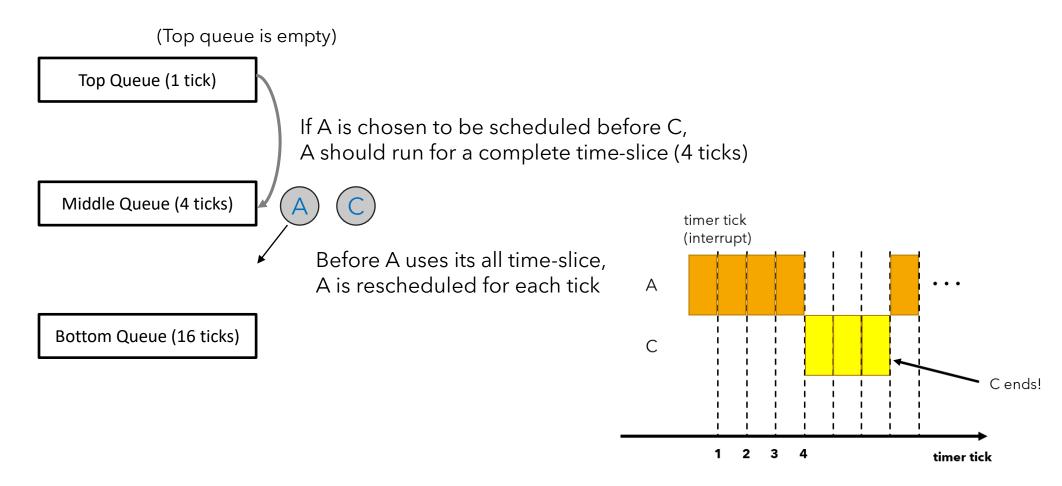
In the top queue, all processes should be scheduled at each timer tick

Middle Queue (4 ticks)

If a top process does not end in 1 tick, the process goes to the middle queue

Bottom Queue (16 ticks)

Example



• Example

(Top queue is empty)

Top Queue (1 tick)

Because A doesn't end after using its time-slice,
A goes to the bottom queue

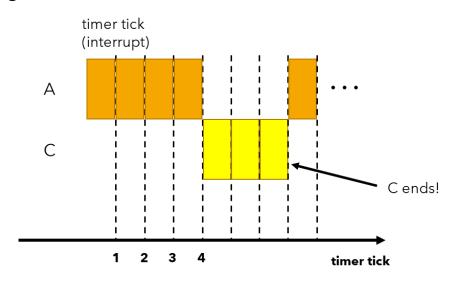
timer tick
(interrupt)

A

Bottom Queue (16 ticks)

A

C



- Small tips
 - Don't need too many changes in original code understanding the scheduler flow is important
 - Don't need linked-list queue recommend to use fixed-sized arrays to represent each priority level
 - To understand more detail about xv6 scheduler, study Chapter 5 in the xv6 book (https://pdos.csail.mit.edu/6.828/2018/xv6/book-rev11.pdf)
 - To study MLFQ scheduler, see OSTEP Chapter 8 (https://pages.cs.wisc.edu/~remzi/OSTEP/cpu-sched.pdf)

• Build xv6 with 'CPUS = 1' flag to test easier (In Makefile)

```
ifndef CPUS
CPUS := 1
endif
```

- Hints
 - per-process structure: struct proc in proc.h

```
37 // Per-process state
38 struct proc {
39    uint sz;
40    pde_t* pgdir;
41    char *kstack;
    rocess
42    enum procstate state;
43    int pid;
44    struct proc *parent;
45    struct trapframe *tf;
46    struct context *context;
47    void *chan;
48    int killed;
49    struct file *ofile[NOFILE];
50    struct inode *cwd;
51    char name[16];
52 };
```

You can add member variables for per-process variable
 e.g.,) allowed ticks, level, ...

• Hints

New processes are initialized in allocproc()

```
static struct proc*
allocproc(void)
  struct proc *p;
  char *sp;
  acquire(&ptable.lock);
  for(p = ptable.proc; p < &ptable.proc[NPROC]; p++)</pre>
    if(p->state == UNUSED)
      goto found;
  release(&ptable.lock);
  return 0;
 ound:
  p->state = EMBRYO;
  p->pid = nextpid++;
  release(&ptable.lock);
```

• Scan ptable to find an empty space

(Don't need to be modified)

- Good place to add codes for initialization
 - e.g.,) find empty space in top queue and add p

• Hints

```
void scheduler(void) {
 struct proc *p;
 struct cpu *c = mycpu();
 c \rightarrow proc = 0;
 for(;;){
   sti();
   acquire(&ptable.lock);
    for(p = ptable.proc; p < &ptable.proc[NPROC]; p++){</pre>
     if(p->state != RUNNABLE)
      c->proc = p;
      switchuvm(p);
      p->state = RUNNING;
      swtch(&(c->scheduler), p->context);
      switchkvm();
      c->proc = 0;
   release(&ptable.lock);
```

- In scheduler(), we need to find which process to schedule
 - Vanilla xv6 scans ptable, but we need to scan queues
 - When scanning queues, we should look at the top level to the bottom level in order
- After swtch(), we can distinguish if the process used just before is over or not
 - We need to modify # of ticks remaining in the process and manage queues according to the case

Testing

- Run simple bench programs to test whether implemented scheduler works well
 - bench 1~3: fork 10 child processes with a fixed execution time (1/4/16 ticks)
 - bench 4: fork 10 child processes with different execution times ([0, 5, 10, 15, 20, 0, ...] ticks)
 - bench 5: fork 50 child processes with I fixed + 49 different times (200 + [0,3,6,9,12,0, ...] ticks)
- To compile bench programs, you need to modify Makefile
 - e.g.,) add _bench I \, ... under UPROGS=\
 - You can make & test any programs if you want
- You need to modify param.h to test bench programs
 - #define NPROC 16 → 64

```
1 #define NPROC 64 // maximum number of processes
2 #define KSTACKSIZE 4096 // size of per-process kernel stack
3 #define NCPU 8 // maximum number of CPUs
```

```
170
        cat
        echo\
        forktest
        _grep
        init\
        kill'
        _ln\
177
        _ls\
        _mkdir\
        _rm\
        sh∖
        _stressfs\
        _usertests
        _wc\
        zombie'
        mytest\
        swtchtest
        realtest'
         bench1
        bench2\
```

- Deadline
 - ~ 2022.11.02 (Wed) 23:59
- Hand-in procedure
 - projl_201812345.patch
 - Run the following command and upload proj I_201812345.patch
 - git diff > proj1_201812345.patch
 - Check the patch file with Notepad and confirm your modifications are in the patch file
 - Report
 - Submit an I-2 page report
 - Free format (Korean/English)
 - Description of your implementation
 - Analysis of benchmark programs including comparison with xv6 default scheduler

Finally...

Do NOT hesitate to ask questions!

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Thank you!