

# Scm.index

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# 1 Scm.Abstract-Syntax

```
module Scm.Abstract-Syntax where

open import Data.Integer.Base renaming (Z to Int) public
open import Data.String.Base using (String) public

data Con : Set -- constants, *excluding* quotations
variable K : Con
Ide = String -- identifiers (variables)
variable I : Ide
data Exp : Set -- expressions
variable E : Exp
data Exp : Set -- expression sequences
variable E : Exp

data Body : Set -- body expression or definition
variable B : Body
data Body+ : Set -- body sequences
variable B+ : Body+
data Prog : Set -- programs
variable P : Prog

-----
-- Literal constants

data Con where -- basic constants
  int : Int → Con -- integer numerals
  #t : Con -- true
  #f : Con -- false

-----
-- Expressions

data Exp where
  con : Con → Exp -- expressions
  ide : Ide → Exp -- K
  ( _ ∘ _ ) : Exp → Exp → Exp -- I
  (λ λ _ ∘ _ ) : Ide → Exp → Exp -- (lambda I E)
  (if _ ∘ _ ∘ _ ) : Exp → Exp → Exp → Exp -- (if E E1 E2)
  (set! _ ∘ _ ) : Ide → Exp → Exp -- (set! I E)

data Exp where
  _ ∘ _ : Exp -- empty sequence
  _ ∘ _ ∘ _ : Exp → Exp → Exp -- prefix sequence E E
```

```

-----  

-- Definitions and Programs  

  

data Body where
  uu_ : Exp → Body          -- side-effect expression E
  (define _ uu_) : Ide → Exp → Body  -- definition (define I E)
  (begin _) : Body+ → Body        -- block (begin B+)  

  

data Body+ where
  uu_ : Body → Body+         -- body sequence
  _ uu_ : Body → Body+ → Body+ -- single body sequence B
  _ _ uu_ : Body → Body+ → Body+ → Body+ -- prefix body sequence B B+  

  

data Prog where
  uuu : Prog                -- programs
  uu_ : Body+ → Prog         -- empty program
  uu_ : Body+ → Prog         -- non-empty program B+  

  

infix 30 uu_
infixr 20 _ uu_

```

## 2 Scm.Auxiliary-Functions

```

module Scm.Auxiliary-Functions where

open import Scm.Notation
open import Scm.Abstract-Syntax
open import Scm.Domain-Equations

-- Environments  $\rho : \mathbf{U} = \mathbf{Ide} \rightarrow \mathbf{L}$ 

postulate _ ==_ : Ide → Ide → Bool
 $\underline{\_}[\underline{\_}/\underline{\_}] : \mathbf{U} \rightarrow \mathbf{L} \rightarrow \mathbf{Ide} \rightarrow \mathbf{U}$ 
 $\rho[\alpha/\mathbb{I}] = \lambda \mathbb{I}' \rightarrow \eta(\mathbb{I} ==^L \mathbb{I}') \longrightarrow \alpha, \rho \mathbb{I}'$ 

postulate unknown : L
--  $\rho \mathbb{I} = \text{unknown}$  represents the lack of a binding for  $\mathbb{I}$  in  $\rho$ 

postulate initial-env : U
-- initial-env shoud include various procedures and values

-- Stores  $\sigma : \mathbf{S} = \mathbf{L} \rightarrow \mathbf{E}$ 

 $\underline{\_}[\underline{\_}/\underline{\_}]' : \mathbf{S} \rightarrow \mathbf{E} \rightarrow \mathbf{L} \rightarrow \mathbf{S}$ 
 $\sigma[\epsilon/\alpha]' = \lambda \alpha' \rightarrow (\alpha ==^L \alpha') \longrightarrow \epsilon, \sigma \alpha'$ 

assign : L → E → C → C
assign =  $\lambda \alpha \epsilon \theta \sigma \rightarrow \theta(\sigma[\epsilon/\alpha]')$ 

hold : L → (E → C) → C
hold =  $\lambda \alpha \kappa \sigma \rightarrow \kappa(\sigma \alpha) \sigma$ 

postulate new : (L → C) → C
-- new  $\kappa \sigma = \kappa \alpha \sigma'$  where  $\sigma \alpha = \text{unallocated}$ ,  $\sigma' \alpha \neq \text{unallocated}$ 

alloc : E → (L → C) → C
alloc =  $\lambda \epsilon \kappa \rightarrow \text{new}(\lambda \alpha \rightarrow \text{assign} \alpha \epsilon (\kappa \alpha))$ 
-- should be  $\perp$  when  $\epsilon \mid \text{-M} == \text{unallocated}$ 

initial-store : S
initial-store =  $\lambda \alpha \rightarrow \eta \text{ unallocated M-in-E}$ 

postulate finished : C
-- normal termination with answer depending on final store

truish : E → T
truish =
 $\lambda \epsilon \rightarrow (\epsilon \in \neg T) \longrightarrow$ 
 $((\epsilon \mid \neg T) ==^T \eta \text{ false}) \longrightarrow \eta \text{ false}, \eta \text{ true} ,$ 
 $\eta \text{ true}$ 

```

```

-- Lists

cons : F
cons =

$$\lambda \epsilon \kappa \rightarrow (\# \epsilon == \perp 2) \longrightarrow \text{alloc}(\epsilon \downarrow 1) (\lambda \alpha_1 \rightarrow \text{alloc}(\epsilon \downarrow 2) (\lambda \alpha_2 \rightarrow \kappa ((\alpha_1, \alpha_2) \text{-in-}\mathbf{E})) , \perp$$


list : F
list = fix  $\lambda$  list'  $\rightarrow$ 

$$\lambda \epsilon \kappa \rightarrow (\# \epsilon == \perp 0) \longrightarrow \kappa (\eta \text{ null M-in-}\mathbf{E}) ,$$


$$\text{list}' (\epsilon \dagger 1) (\lambda \epsilon \rightarrow \text{cons} \langle (\epsilon \downarrow 1) , \epsilon \rangle \kappa)$$


car : F
car =

$$\lambda \epsilon \kappa \rightarrow (\# \epsilon == \perp 1) \longrightarrow \text{hold}((\epsilon \downarrow 1) \dashv \downarrow^2 1) \kappa , \perp$$


cdr : F
cdr =

$$\lambda \epsilon \kappa \rightarrow (\# \epsilon == \perp 1) \longrightarrow \text{hold}((\epsilon \downarrow 1) \dashv \downarrow^2 2) \kappa , \perp$$


setcar : F
setcar =

$$\lambda \epsilon \kappa \rightarrow (\# \epsilon == \perp 2) \longrightarrow \text{assign}((\epsilon \downarrow 1) \dashv \downarrow^2 1)$$


$$(\epsilon \downarrow 2)$$


$$(\kappa (\eta \text{ unspecified M-in-}\mathbf{E})) , \perp$$


setcdr : F
setcdr =

$$\lambda \epsilon \kappa \rightarrow (\# \epsilon == \perp 2) \longrightarrow \text{assign}((\epsilon \downarrow 1) \dashv \downarrow^2 2)$$


$$(\epsilon \downarrow 2)$$


$$(\kappa (\eta \text{ unspecified M-in-}\mathbf{E})) , \perp$$


```

### 3 Scm.Domain-Equations

```
module Scm.Domain-Equations where

open import Scm.Notation
open import Scm.Abstract-Syntax using (Ide; Int)

-- Domain declarations

postulate L : Domain -- locations
variable α : L
N : Domain -- natural numbers
T : Domain -- booleans
R : Domain -- numbers
P : Domain -- pairs
M : Domain -- miscellaneous
variable μ : M
F : Domain -- procedure values
variable φ : F
postulate E : Domain -- expressed values
variable ε : E
S : Domain -- stores
variable σ : S
U : Domain -- environments
variable ρ : U
C : Domain -- command continuations
variable θ : C
postulate A : Domain -- answers

E = E
variable ε : E

-- Domain equations

data Misc : Set where null unallocated undefined unspecified : Misc

N = Nat ⊥
T = Bool ⊥
R = Int + ⊥
= L × L
M = Misc + ⊥
F = E → (E → C) → C
-- E = T + R + P + M + F
S = L → E
U = Ide → L
C = S → A
```

```
-- Injections, tests, and projections
```

```
postulate
```

$$\begin{array}{l} \underline{\text{T-in-E}} : \mathbf{T} \rightarrow \mathbf{E} \\ \underline{\in_{\mathbf{T}}} : \mathbf{E} \rightarrow \text{Bool} + \perp \\ \underline{|_{\mathbf{T}}} : \mathbf{E} \rightarrow \mathbf{T} \\ \\ \underline{\mathbf{R}\text{-in-}\mathbf{E}} : \mathbf{R} \rightarrow \mathbf{E} \\ \underline{\in_{\mathbf{R}}} : \mathbf{E} \rightarrow \text{Bool} + \perp \\ \underline{|_{\mathbf{R}}} : \mathbf{E} \rightarrow \mathbf{R} \\ \\ \underline{\text{-in-}\mathbf{E}} : \mathbf{E} \rightarrow \mathbf{E} \\ \underline{\in_{-}} : \mathbf{E} \rightarrow \text{Bool} + \perp \\ \underline{|_{-}} : \mathbf{E} \rightarrow \\ \\ \underline{\mathbf{M}\text{-in-}\mathbf{E}} : \mathbf{M} \rightarrow \mathbf{E} \\ \underline{\in_{\mathbf{M}}} : \mathbf{E} \rightarrow \text{Bool} + \perp \\ \underline{|_{\mathbf{M}}} : \mathbf{E} \rightarrow \mathbf{M} \\ \\ \underline{\mathbf{F}\text{-in-}\mathbf{E}} : \mathbf{F} \rightarrow \mathbf{E} \\ \underline{\in_{\mathbf{F}}} : \mathbf{E} \rightarrow \text{Bool} + \perp \\ \underline{|_{\mathbf{F}}} : \mathbf{E} \rightarrow \mathbf{F} \end{array}$$

```
-- Operations on flat domains
```

```
postulate
```

$$\begin{array}{l} \underline{==^L} : \mathbf{L} \rightarrow \mathbf{L} \rightarrow \mathbf{T} \\ \underline{==^M} : \mathbf{M} \rightarrow \mathbf{M} \rightarrow \mathbf{T} \\ \underline{==^R} : \mathbf{R} \rightarrow \mathbf{R} \rightarrow \mathbf{T} \\ \underline{==^T} : \mathbf{T} \rightarrow \mathbf{T} \rightarrow \mathbf{T} \\ \underline{<^R} : \mathbf{R} \rightarrow \mathbf{R} \rightarrow \mathbf{T} \\ \underline{+^R} : \mathbf{R} \rightarrow \mathbf{R} \rightarrow \mathbf{R} \\ \underline{\wedge^T} : \mathbf{T} \rightarrow \mathbf{T} \rightarrow \mathbf{T} \end{array}$$

## 4 Scm.Notation

```

module Scm.Notation where

open import Data.Bool.Base  using (Bool; false; true) public
open import Data.Nat.Base   renaming (N to Nat) using (suc) public
open import Data.String.Base using (String) public
open import Data.Unit.Base  using (T)
open import Function         using (id; _○_ ) public

Domain = Set -- unsound!

variable
A B C      : Set
D E F      : Domain
n          : Nat

-----  

-- Domains

postulate
⊥ : D           -- bottom element
fix : (D → D) → D -- fixed point of endofunction

-----  

-- Flat domains

postulate
_+_     : Set → Domain      -- lifted set
η       : A → A +_
_ SHARP : (A → D) → (A +_ → D) -- Kleisli extension

Bool_    = Bool +_
Nat_    = Nat +_
String_ = String +_

postulate
_==_    : Nat_ → Nat → Bool_ -- strict numerical equality
_>=_    : Nat_ → Nat → Bool_ -- strict greater or equal
_→_,_   : Bool_ → D → D → D -- McCarthy conditional

-----  

-- Sum domains

postulate
_+_     : Domain → Domain → Domain      -- separated sum
inj₁   : D → D + E                      -- injection
inj₂   : E → D + E                      -- injection
[_,_]  : (D → F) → (E → F) → (D + E → F) -- case analysis

```

```

-----  

-- Product domains  

postulate
   $\_ \times \_ : \text{Domain} \rightarrow \text{Domain} \rightarrow \text{Domain}$  -- cartesian product
   $\_, \bar{\_} : D \rightarrow E \rightarrow D \times E$  -- pairing
   $\downarrow^2_1 : D \times E \rightarrow D$  -- 1st projection
   $\downarrow^2_2 : D \times E \rightarrow E$  -- 2nd projection
   $\downarrow^3_1 : D \times E \times F \rightarrow D$  -- 1st projection
   $\downarrow^3_2 : D \times E \times F \rightarrow E$  -- 2nd projection
   $\downarrow^3_3 : D \times E \times F \rightarrow F$  -- 3rd projection
-----  

-- Tuple domains  

 $\_ \wedge \_ : \text{Domain} \rightarrow \text{Nat} \rightarrow \text{Domain}$  --  $D \wedge n$  n-tuples
 $D \wedge 0 = T$ 
 $D \wedge 1 = D$ 
 $D \wedge \text{suc } (\text{suc } n) = D \times (D \wedge \text{suc } n)$ 
-----  

-- Finite sequence domains  

postulate
   $\langle \rangle : \text{Domain} \rightarrow \text{Domain}$  --  $D$  domain of finite sequences
   $\langle \_ \rangle : D \rightarrow \text{Domain}$  -- empty sequence
   $\langle \_ \rangle : (D \wedge \text{suc } n) \rightarrow D$  --  $\langle d_1, \dots, d_{n+1} \rangle$  non-empty sequence
   $\# : D \rightarrow \text{Nat} \perp$  --  $\# d$  sequence length
   $\_ \S \_ : D \rightarrow D \rightarrow D$  --  $d \S d$  concatenation
   $\_ \downarrow \_ : D \rightarrow \text{Nat} \rightarrow D$  --  $d \downarrow n$  nth component
   $\_ \dagger \_ : D \rightarrow \text{Nat} \rightarrow D$  --  $d \dagger n$  nth tail
-----  

-- Grouping precedence  

infixr 1       $\_ + \_$ 
infixr 2       $\_ \times \_$ 
infixr 4       $\_, \_$ 
infixr 8       $\_ \wedge \_$ 
infixr 20      $\_ \longrightarrow \_ , \_$   

 $\llbracket \_ \rrbracket = \text{id}$ 

```

## 5 Scm.Semantic-Functions

```

module Scm.Semantic-Functions where

open import Scm.Notation
open import Scm.Abstract-Syntax
open import Scm.Domain-Equations
open import Scm.Auxiliary-Functions

 $\mathcal{K}[\_]$  : Con  $\rightarrow \mathbf{E}$ 
 $\mathcal{E}[\_]$  : Exp  $\rightarrow \mathbf{U} \rightarrow (\mathbf{E} \rightarrow \mathbf{C}) \rightarrow \mathbf{C}$ 
 $\mathcal{E}[\_]$  : Exp  $\rightarrow \mathbf{U} \rightarrow (\mathbf{E} \rightarrow \mathbf{C}) \rightarrow \mathbf{C}$ 

 $\mathcal{B}[\_]$  : Body  $\rightarrow \mathbf{U} \rightarrow (\mathbf{U} \rightarrow \mathbf{C}) \rightarrow \mathbf{C}$ 
 $\mathcal{B}^+[\_]$  : Body+  $\rightarrow \mathbf{U} \rightarrow (\mathbf{U} \rightarrow \mathbf{C}) \rightarrow \mathbf{C}$ 
 $\mathcal{P}[\_]$  : Prog  $\rightarrow \mathbf{A}$ 

-- Constant denotations  $\mathcal{K}[\ K \ ] : \mathbf{E}$ 
 $\mathcal{K}[\text{int } Z \ ] = \eta Z \text{ R-in-}\mathbf{E}$ 
 $\mathcal{K}[\#t \ ] = \eta \text{ true T-in-}\mathbf{E}$ 
 $\mathcal{K}[\#f \ ] = \eta \text{ false T-in-}\mathbf{E}$ 

-- Expression denotations
 $\mathcal{E}[\text{con } K \ ] \rho \kappa = \kappa (\mathcal{K}[ K \ ])$ 
 $\mathcal{E}[\text{ide } I \ ] \rho \kappa = \text{hold } (\rho I) \kappa$ 

$$\mathcal{E}[(\lambda E \sqcup E)] \rho \kappa = \\ \mathcal{E}[E] \rho (\lambda \epsilon \rightarrow \\ \mathcal{E}[E] \rho (\lambda \epsilon \rightarrow \\ (\epsilon \mid\!\!-\mathbf{F}) \epsilon \kappa))$$


$$\mathcal{E}[(\lambda l \sqcup E)] \rho \kappa = \\ \kappa (\\ (\lambda \epsilon \kappa' \rightarrow \\ \text{list } \epsilon (\lambda \epsilon \rightarrow \\ \text{alloc } \epsilon (\lambda \alpha \rightarrow \\ \mathcal{E}[E] (\rho [\alpha / l]) \kappa')) \\ ) \mathbf{F}\text{-in-}\mathbf{E})$$


$$\mathcal{E}[(\text{if } E \sqcup E_1 \sqcup E_2)] \rho \kappa = \\ \mathcal{E}[E] \rho (\lambda \epsilon \rightarrow \\ \text{truish } \epsilon \longrightarrow \mathcal{E}[E_1] \rho \kappa , \mathcal{E}[E_2] \rho \kappa)$$


$$\mathcal{E}[(\text{set! } l \sqcup E)] \rho \kappa = \\ \mathcal{E}[E] \rho (\lambda \epsilon \rightarrow \\ \text{assign } (\rho l) \epsilon ( \\ \kappa (\eta \text{ unspecified M-in-}\mathbf{E})))$$

--  $\mathcal{E}[\_]$  : Exp  $\rightarrow \mathbf{U} \rightarrow (\mathbf{E} \rightarrow \mathbf{C}) \rightarrow \mathbf{C}$ 

$$\mathcal{E}[\text{uuu}] \rho \kappa = \kappa \langle \rangle$$


$$\mathcal{E}[E \text{ uu } E] \rho \kappa = \\ \mathcal{E}[E] \rho (\lambda \epsilon \rightarrow \\ \mathcal{E}[E] \rho (\lambda \epsilon \rightarrow \\ \kappa (\langle \epsilon \rangle \S \epsilon)))$$


```

```

-- Body denotations  $\mathcal{B}[\![\ B\ ]\!]$  :  $\mathbf{U} \rightarrow (\mathbf{U} \rightarrow \mathbf{C}) \rightarrow \mathbf{C}$ 

$$\mathcal{B}[\![\ \sqcup\ E\ ]\!] \rho \kappa = \mathcal{E}[\![\ E\ ]\!] \rho (\lambda \epsilon \rightarrow \kappa \rho)$$


$$\begin{aligned} \mathcal{B}[\![\ (\text{define } l \sqcup E)\ ]\!] \rho \kappa = \\ \mathcal{E}[\![\ E\ ]\!] \rho (\lambda \epsilon \rightarrow (\rho l ==^L \text{unknown}) \longrightarrow \\ \text{alloc } \epsilon (\lambda \alpha \rightarrow \kappa (\rho [\alpha / l])), \\ \text{assign } (\rho l) \epsilon (\kappa \rho)) \end{aligned}$$


$$\mathcal{B}[\![\ (\text{begin } B^+)\ ]\!] \rho \kappa = \mathcal{B}^+[\![\ B^+\ ]\!] \rho \kappa$$

-- Body sequence denotations  $\mathcal{B}^+[\![\ B^+\ ]\!]$  :  $\mathbf{U} \rightarrow (\mathbf{U} \rightarrow \mathbf{C}) \rightarrow \mathbf{C}$ 

$$\mathcal{B}^+[\![\ \sqcup\ B\ ]\!] \rho \kappa = \mathcal{B}[\![\ B\ ]\!] \rho \kappa$$


$$\mathcal{B}^+[\![\ B \sqcup B^+\ ]\!] \rho \kappa = \mathcal{B}[\![\ B\ ]\!] \rho (\lambda \rho' \rightarrow \mathcal{B}^+[\![\ B^+\ ]\!] \rho' \kappa)$$

-- Program denotations  $\mathcal{P}[\![\ \Pi\ ]\!]$  :  $\mathbf{A}$ 

$$\mathcal{P}[\![\ \sqcup\sqcup\ ]\!] = \text{finished initial-store}$$


$$\mathcal{P}[\![\ \sqcup\ B^+\ ]\!] = \mathcal{B}^+[\![\ B^+\ ]\!] \text{initial-env } (\lambda \rho \rightarrow \text{finished}) \text{ initial-store}$$


```

## 6 Scm.index

```
module Scm.index where
  import Scm.Notation
  import Scm.Abstract-Syntax
  import Scm.Domain-Equations
  import Scm.Semantic-Functions
  import Scm.Auxiliary-Functions
```