Denotational Semantics of PCF in Agda

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Abstract

In synthetic domain theory, all sets are predomains, domains are pointed sets, and functions are implicitly continuous. The denotational semantics of PCF (Plotkin's version) presented here illustrates how it might look if synthetic domain theory can be implemented in Agda. As a work-around, the code presented here uses unsatisfiable postulates to allow Agda to type-check the definitions.

The (currently illiterate) Agda source code used to generate this document can be downloaded from https://github.com/pdmosses/xds-agda, and browsed with hyperlinks and highlighting at https://pdmosses.github.io/xds-agda/.

{-# OPTIONS --rewriting --confluence-check #-}

module PCF.All where

import PCF.Domain-Notation

import PCF. Types

import PCF.Constants

import PCF. Variables

import PCF.Environments

import PCF.Terms

import PCF.Checks

```
module PCF.Domain-Notation where
open import Relation. Binary. Propositional Equality. Core
  using (_≡_; refl) public
Domain = Set
variable D E : Domain
-- Domains are pointed
postulate
               : \{D : Domain\} \rightarrow D
-- Fixed points of endofunctions on function domains
postulate
                : \{D : Domain\} \rightarrow (D \rightarrow D) \rightarrow D
  fix
   -- Properties
  fix-fix : \forall \{D\} (f : D \rightarrow D) \rightarrow
                       fix f \equiv f (fix f)
  fix-app: \forall \{P D\} (f : (P \rightarrow D) \rightarrow (P \rightarrow D)) (p : P) \rightarrow (P \rightarrow D)
                       fix f p \equiv f (fix f) p
-- Lifted domains
postulate
  \mathbb{L}
                : Set \rightarrow Domain
                \colon \{\mathsf{P} : \mathsf{Set}\} \to \mathsf{P} \to \mathbb{L} \; \mathsf{P}
                : \{P : Set\} \{D : Domain\} \rightarrow (P \rightarrow D) \rightarrow (\mathbb{L} P \rightarrow D)
   -- Properties
  elim^{-\sharp}-\eta: \forall \{PD\} (f: P \rightarrow D) (p: P) \rightarrow
                    (f^{\sharp})(\eta p) \equiv f p
  elim-^{\sharp}-\bot: \forall {PD} (f: P \rightarrow D) \rightarrow
                     (f <sup>♯</sup>) ⊥ ≡ ⊥
-- Flat domains
\_+\bot : \mathsf{Set} \to \mathsf{Domain}
\mathsf{S} + \bot = \mathbb{L} \mathsf{S}
-- McCarthy conditional
-- t \longrightarrow d_1 , d_2 : D (t : Bool +\perp ; d_1, d_2 : D)
open import Data.Bool.Base
  using (Bool; true; false; if then else ) public
   \_ \_ , \_ : {D : Domain} \rightarrow Bool +\bot \rightarrow D \rightarrow D
   -- Properties
                     : \forall \{D\} \{d_1 d_2 : D\} \rightarrow (\eta \text{ true} \longrightarrow d_1, d_2) \equiv d_1
  true-cond
   false-cond : \forall \{D\} \{d_1 d_2 : D\} \rightarrow (\eta \text{ false} \longrightarrow d_1, d_2) \equiv d_2
  bottom-cond : \forall \{D\} \{d_1 d_2 : D\} \rightarrow (\bot \longrightarrow d_1, d_2)
```

```
module PCF. Types where
open import Data.Bool.Base
  using (Bool)
open import Ágda.Builtin.Nat
  using (Nat)
open import PCF.Domain-Notation
  using (Domain; \_+\bot)
-- Syntax
data Types : Set where
 variable \sigma \tau : Types
\mathsf{infixr}\ 1 \ \_{\Rightarrow}\_
-- Semantics \mathcal D
\mathcal{D}: \mathsf{Types} \to \mathsf{Domain}
\mathcal{D} \iota
            = Nat + \bot
          = Bool + \perp
\mathcal{D} o
\mathcal{D}\left(\sigma \Rightarrow \tau\right) = \mathcal{D} \ \sigma \to \mathcal{D} \ \tau
variable x y z : \mathcal{D} \sigma
```

```
module PCF. Constants where
open import Data.Bool.Base
    using (Bool; true; false; if then else )
open import Agda.Builtin.Nat
    using (Nat; _+_; _-_; _==_)
open import PCF.Domain-Notation
    using (\eta; _{\sharp}; fix; \bot; _{\longrightarrow}, _{]})
open import PCF. Types
    using (Types; o; \iota; \_\Rightarrow_; \sigma; \mathcal{D})
-- Syntax
data \mathcal{L}: Types \rightarrow Set where
    tt : \mathcal{L} o
    ff: \mathcal{L} o
    \supset_i : \mathcal{L} (o \Rightarrow \iota \Rightarrow \iota \Rightarrow \iota)
    \supset_0: \mathcal{L} (0 \Rightarrow 0 \Rightarrow 0 \Rightarrow 0)
   Y: \{\sigma : \mathsf{Types}\} \to \mathcal{L} ((\sigma \Rightarrow \sigma) \Rightarrow \sigma)
k: (\mathsf{n} : \mathsf{Nat}) \to \mathcal{L} \iota
    +1': \mathcal{L}(\iota \Rightarrow \iota)
   -1': \mathcal{L}(\iota \Rightarrow \iota)
    Z : \mathcal{L} (\iota \Rightarrow 0)
variable c : \mathcal{L} \sigma
-- Semantics
\mathcal{A}[\![\![\ \_\!]\!]: \mathcal{L} \ \sigma \to \mathcal{D} \ \sigma
\mathcal{A}[\![ \ \mathsf{tt} \qquad ]\!] = \eta \ \mathsf{true}
\mathcal{A}[\![ff]] = \eta \text{ false}
\mathcal{A}[[\mathbf{k} \, \mathbf{n} \, ]] = \eta \, \mathbf{n}
\mathcal{A}[\![+1']\!] = (\lambda \mathsf{n} \to \eta (\mathsf{n} + 1))^{\sharp}
\mathcal{A}[\![-1']\!] = (\lambda n \rightarrow \text{if } n == 0 \text{ then } \perp \text{else } \eta (n-1))^{\sharp}
\mathcal{A}[\![\![ \ \mathsf{Z} \ ]\!]\!] = (\lambda \ \mathsf{n} \to \eta \ (\mathsf{n} == 0))^{\sharp}
```

```
module PCF. Environments where
open import Data.Bool.Base
   using (Bool; if then else )
open import Data.Maybe.Base
   using (Maybe; just; nothing)
open import Agda.Builtin.Nat
   using (Nat; \_==\_)
open import Relation.Binary.PropositionalEquality.Core
   using (\equiv ; refl; trans; cong)
open import PCF.Domain-Notation
   using (\bot)
open import PCF. Types
   using (Types; \iota; o; \_\Rightarrow\_; \mathcal{D})
open import PCF. Variables
   using (\mathbf{V}; var; Env)
-- \rho \perp is the initial environment
\rho \bot: Env
\rho \perp \alpha = \perp
-- (\rho [ x / \alpha ]) \alpha' = x when \alpha and \alpha' are identical, otherwise \rho \alpha'
[ / ]: \{\sigma : \mathsf{Types}\} \to \mathsf{Env} \to \mathcal{D} \ \sigma \to \mathcal{V} \ \sigma \to \mathsf{Env}
\rho \left[ \times / \alpha \right] = \lambda \alpha' \rightarrow h \rho \times \alpha \alpha' (\alpha == V \alpha') where
   h: \{\sigma \tau : \mathsf{Types}\} \to \mathsf{Env} \to \mathcal{D} \ \sigma \to \mathcal{V} \ \sigma \to \mathcal{V} \ \tau \to \mathsf{Maybe} \ (\sigma \equiv \tau) \to \mathcal{D} \ \tau
   h \rho \times \alpha \alpha' (just refl) = \times
   h \rho \times \alpha \alpha' nothing = \rho \alpha'
      \underline{\phantom{a}} = T_{\underline{\phantom{a}}} : (\sigma \ \tau : \mathsf{Types}) \to \mathsf{Maybe} \ (\sigma \equiv \tau)
   (\sigma \Rightarrow \tau) = T (\sigma' \Rightarrow \tau') = f (\sigma = T \sigma') (\tau = T \tau') where
                f: Maybe (\sigma \equiv \sigma') \rightarrow \text{Maybe} (\tau \equiv \tau') \rightarrow \text{Maybe} ((\sigma \Rightarrow \tau) \equiv (\sigma' \Rightarrow \tau'))
                f = \lambda \{ (just p) (just q) \rightarrow just (trans (cong (<math>\_\Rightarrow \tau) p) (cong (\sigma' \Rightarrow \_) q) \}
                          ; \_ \_ \rightarrow nothing \}
   \iota == T \iota = \text{just refl}
   o == T o = just refl
   \_ == T \_ = nothing
    ==\mathsf{V}_{-}:\{\sigma\ \tau:\,\mathsf{Types}\}\to\mathcal{V}\ \sigma\to\mathcal{V}\ \tau\to\mathsf{Maybe}\ (\sigma\equiv\tau)
   var i \sigma == V var i' \tau =
      if i == i' then \sigma == T \tau else nothing
```

```
module PCF. Terms where
open import PCF. Types
  using (Types; \_\Rightarrow\_; \sigma; \mathcal{D})
open import PCF.Constants
using (L; \mathcal{A}[\![\_]\!]; c) open import PCF.Variables
   using (V; Env; [[]])
open import PCF.Environments
  using (_[_/_])
-- Syntax
data Terms : Types \rightarrow Set where
   V : {\sigma : \mathsf{Types}} \to \mathcal{V} \ \sigma \to \mathsf{Terms} \ \sigma
                                                                                         -- variables
   L \qquad : \{\sigma : \mathsf{Types}\} \to \mathcal{L} \ \sigma \to \mathsf{Terms} \ \sigma
                                                                                       -- constants
  variable M N : Terms \sigma
infixl 20 _⊔_
-- Semantics
\mathcal{A}' \llbracket \quad \rrbracket : \mathsf{Terms} \ \sigma \to \mathsf{Env} \to \mathcal{D} \ \sigma
\mathcal{A}' \llbracket V \alpha
                  ]\!] \ \rho = \rho \ [\![ \ \alpha \ ]\!]
\mathcal{A}' \llbracket L \mathsf{c} \qquad \rrbracket \rho = \mathcal{A} \llbracket \mathsf{c} \rrbracket
```

```
{-# OPTIONS --rewriting --confluence-check #-}
open import Agda.Builtin.Equality
open import Agda. Builtin. Equality. Rewrite
module PCF.Checks where
open import Data.Bool.Base
open import Agda.Builtin.Nat
open import Relation.Binary.PropositionalEquality.Core
  using (\equiv ; refl)
open import PCF.Domain-Notation
open import PCF. Types
open import PCF.Constants
open import PCF. Variables
open import PCF. Environments
open import PCF. Terms
  {-# REWRITE fix-app elim-^{\sharp}-\eta elim-^{\sharp}-\bot true-cond false-cond #-}
-- Constants
pattern N n = L (k n)
pattern succ = L + 1'
pattern pred\perp = L - 1'
                = L \supset_i
= L Y
pattern if
pattern Y
pattern Z = L Z
-- Variables
f = var 0 \iota
g = var 1 (\iota \Rightarrow \iota)
h = var 2 (\iota \Rightarrow \iota \Rightarrow \iota)
a = var 3 \iota
b = var 4 \iota
-- Arithmetic
check-41+1 : \mathcal{A}' \llbracket \text{ succ } \sqcup N \text{ 41 } \rrbracket \rho \bot \equiv \eta \text{ 42}
check-41+1 = refl
check-43-1 : \mathcal{A}' \llbracket \text{ pred} \bot \bigsqcup N \text{ 43} \rrbracket \rho \bot \equiv \eta \text{ 42}
check-43-1 = refl
-- Binding
check-id : \mathcal{A}' \| (\mathbf{1} \times \mathbf{1} \times \mathbf{1} \times \mathbf{1} \times \mathbf{1}) \| \mathbf{1} \times \mathbf{1} \times \mathbf{1} \| \rho \| = \eta 42
check-id = refl
check-k : \mathcal{A}' [ (\hbar a \perp \hbar b \perp V a) \perp N 42 \perp N 41 ]] \rho \perp \equiv \eta 42
check-k = refl
check-ki : \mathcal{A}' [ (^{\uparrow} a _{\square} ^{\uparrow} b _{\square} ^{\lor} b) _{\square} ^{\lor} 41 _{\square} ^{\lor} 42 ]] \rho \bot \equiv \eta 42
check-ki = refl
```

```
check-suc-41 : \mathcal{A}' \parallel (\hat{x} \text{ a} \sqcup (\text{succ} \sqcup V \text{ a})) \sqcup N \text{ 41} \parallel \rho \perp \equiv \eta \text{ 42}
check-suc-41 = refl
check-pred-42 : \mathcal{A}' \llbracket ( \mathring{\Lambda} \text{ a} \sqcup (\text{pred} \bot \sqcup V \text{ a})) \sqcup N \text{ 43 } \rrbracket \rho \bot \equiv \eta \text{ 42}
check-pred-42 = refl
check-if-zero : \mathcal{A}' [ if \sqcup (Z \sqcup N 0) \sqcup N 42 \sqcup N 0 ] \rho \bot \equiv \eta 42
check-if-zero = refl
check-if-nonzero : \mathcal{A}' [ if \sqcup (Z \sqcup N 42) \sqcup N 0 \sqcup N 42 ] \rho \bot \equiv \eta 42
check-if-nonzero = refl
-- fix (\lambda f. 42) \equiv 42
check-fix-const:
   \mathcal{A}' \llbracket Y \sqcup (\lambda f \sqcup N 42) \rrbracket \rho \bot
   \equiv \eta 42
check-fix-const = fix-fix (\lambda \times \rightarrow \eta 42)
-- fix (\lambdag. \lambdaa. 42) 2 \equiv 42
check-fix-lambda:
   \mathcal{A}' \llbracket Y \sqcup (\lambda g \sqcup \lambda a \sqcup N 42) \sqcup N 2 \rrbracket \rho \bot
   \equiv n 42
check-fix-lambda = refl
-- fix (\lambdag. \lambdaa. ifz a then 42 else g (pred a)) 101 \equiv 42
check-countdown:
   \mathcal{A}' \llbracket Y \sqcup (\lambda g \sqcup \lambda a \sqcup
                             (\mathsf{if} \mathrel{\sqcup} (Z \mathrel{\sqcup} V \mathsf{a}) \mathrel{\sqcup} N \mathsf{42} \mathrel{\sqcup} (V \mathsf{g} \mathrel{\sqcup} (\mathsf{pred} \bot \mathrel{\sqcup} V \mathsf{a}))))
       \rho \perp
   \equiv \eta 42
check-countdown = refl
-- fix (\lambdah. \lambdaa. \lambdab. ifz a then b else h (pred a) (succ b)) 4 38 \equiv 42
check-sum-42:
   \mathcal{A}' \llbracket (Y \sqcup (\lambda h \sqcup \lambda a \sqcup \lambda b \sqcup
                             (if \sqcup (Z \sqcup V a) \sqcup V b \sqcup (V h \sqcup (pred \perp \sqcup V a) \sqcup (succ \sqcup V b)))))
       \equiv \eta 42
check-sum-42 = refl
-- Exponential in first arg?
```