Data Structures

Introduction – 1/4

- A language processor makes frequent use of the search operation over its data structures
- Data structures used in language processing can be classified on basis of the following criteria
 - Nature of a data structure :- whether linear or non linear
 - Purpose of a data structure :- whether search or allocation
 - Lifetime of a data structure :- whether used during language processing or during target program execution

Introduction – 2/4

Linear data structure

- Consists of linear arrangement of elements in memory
- Requires contiguous area of memory
- Designer is forced to overestimate the memory requirements
- Leads to wastage of memory

Non linear data structure

- Accessed using pointers
- Elements do not occupy contiguous area of memory
- Leads to lower search efficiency

Introduction – 3/4

Search data structure

- Used during language processing to maintain attribute information concerning different entities in SP
- Entry for entity is created only once
- Searched for large no of times
- Search efficiency is very important
- Constitute various tables of information

Introduction – 4/4

Allocation data structure

- Address of memory area allocated to entity is known to user of that entity
- No search operations are conducted
- Speed of allocation or deallocation and efficiency of memory utilization are important
- Used to handle programs with nested structures

TP rarely uses search data structures

Search data structures

Set of entries, each entry accommodating the information concerning one entity

Each entry has a key field which forms basis of search

Key field is symbol field containing name of entity

Entry formats – 1/2

- Each entry in search structure is a set of fields
- Entry consists of two parts
 - Fixed part
 - Variant part
- Each part consists of set of fields
- Fields of fixed part exist in each entry of search structure

Entry formats – 2/2

- Value in tag field of fixed part determines the information to be stored in variant part of entry
- Fixed and variant part may be nested (i.e. variant part may itself consist of fixed and variant parts)

Fixed and Variable length entries – 1/3

- An entry may be declared as a record or a structure of language in which the language processor is being implemented
- In **fixed length entry format**, each record is defined to consist of the following fields
 - Fields in the fixed part of entry
 - U_{vi} SF_{vi}, i.e. set of fields in all variant parts of entry
- Records in search structure have identical format which enables use of homogeneous linear data structures like arrays

Fixed and Variable length entries – 2/3

 Use of linear search organizations enables use of efficient search procedures

- In variable length entry format, record consists of following fields
 - Fields in fixed part of entry, including tag field
 - $\{ f_j \mid f_j \in SF_{vj} \text{ if tag} = v_j \}$
- No memory wastage occurs

Fixed and Variable length entries – 3/3

- When variable length entry format is used, search method may require knowledge of length of entry
- Hence, record might consists of fields
 - A length field
 - Fields in fixed part of entry, including tag field

$$- \{ f_j \mid f_j \in SF_{vj} \text{ if tag} = v_j \}$$

Depict as follows

length

Hybrid entry formats – 1/2

- Compromise between fixed and variable entry formats to combine access efficiency of fixed entry format with memory efficiency of variable entry format
- Split into two halves
 - Fixed part
 - Variant part
- Pointer is added to fixed part
- It points to variable part of entry

Hybrid entry formats – 2/2

- Fixed part of all entries are organized into an efficient search structure eg. Linear structure
- Variable part does not need to be located through a search
- Hence it is put into an allocation data structure
- Hybrid entry format is depicted below

Fixed part	pointer		length	variable part
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Operations on search structures

- Add :- Add entry of a symbol
- Search :- Search and locate entry of a symbol
- Delete :- Delete entry of a symbol

- Entry for a symbol is created only once, but can be searched many no of times
- Deletion operation is not common

Generic search procedure – 1/2

- 1. Make a prediction concerning entry of search data structure which symbol s may be occupying. Let this entry be e
- Let s_e be symbol occupying eth entry. Compare s with s_e. Exit with success if two match

3. Repeat steps 1 and 2 till it can be concluded that the symbol does not exist in search data structure

Generic search procedure – 2/2

- Nature of prediction varies with organization of search data structure
- Each comparison is called a probe
- Efficiency is determined by no of probes performed in search procedure
- Following notation to represent no of probes in a search
 - $-p_s$:- no of probes in successful search
 - $-p_{ij}$:- no of probes in unsuccessful search

Other Topics

- Table organization
- Sequential search organization
- Binary search organization
- Hash table organization
 - Hashing functions
 - Collision handling methods
- Linked list and tree structured organization
- Allocation data structures
 - Stacks
 - Heaps