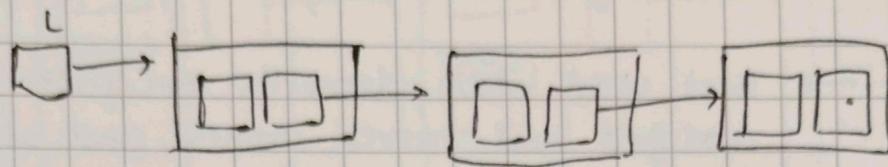
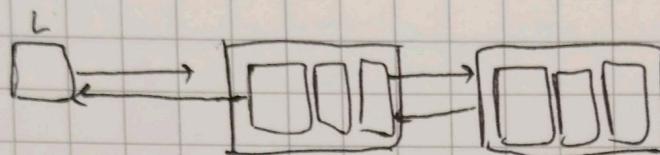


* LINKED LIST IMPLEMENTATION

SINGLY - LINKED



DOUBLY - LINKED



& CIRCULAR (Which I wont draw because screw dat.)

Characteristics :-

- ① Elements are not stored physically the same as logical representation
- ② There is overhead storage (extra storage needed to contain the link)

Eg.
one char elem = 1 byte

one char elem = 9 bytes

→ TRAVERSAL

Has to begin from 1st node to the last node
(In the case of singly-linked list)

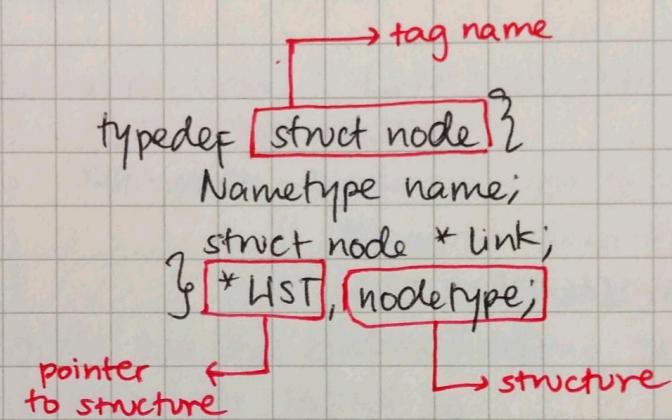
WHEN TO USE ARRAY VS. LINKED LIST?

Array → Best used when number of elements are known and changes are rarely anticipated.

Linked list → Best used when elements are often modified and moved around, and when the maximum number of elements is not known or defined.

PRACTICE

```
typedef struct {  
    char LName [16];  
    char FName [24];  
    char MI;  
} Nametype;
```



3 data types generated :-

`struct node`, `LST`, `nodetype`

* Note:

`struct node` ≈ `nodetype`
`nodetype` + ≈ `LST`
`struct node*` ≈ `LST`

(Sideline, for reference)
Internet version
`struct node {
 struct Name name;
 struct node * link;
};`

PASS BY COPY :

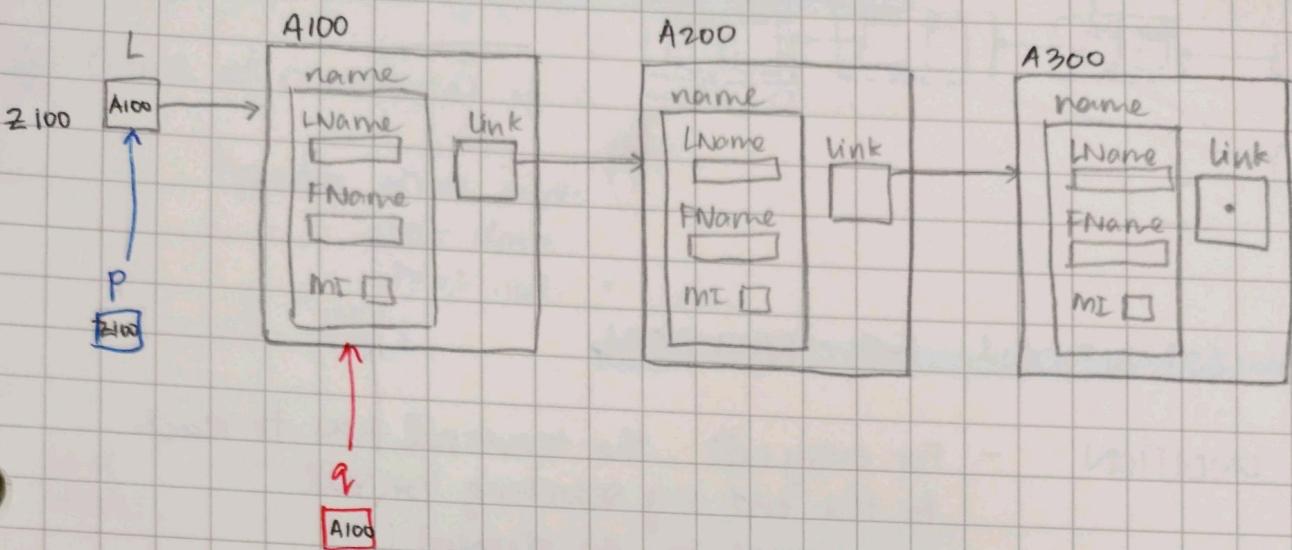
`void displayList (LST A)`

→ `void displayList (struct node * A)`

PASS BY ADDRESS :

`void insertFirst (LST *A, Nametype N)`

→ `void insertFirst (struct node **A,
 struct Name N)`



`LIST *P;` // declaring `*P.` `*P` exists but not pointing
`P = &L;` // making `*P` point to `L`

`LIST q;` // `q` not pointing to anything yet
`q = L;` // `q` pointing to first node

Q. ACCESS THE VARIABLE "MI" VIA `L`, `q` & `p`.

- ANS.
- a) via `L` $L \rightarrow \text{name.MI}$
 - b) via `q` $q \rightarrow \text{name.MI}$
 - c) via `p` $(*p) \rightarrow \text{name.MI}$

Q. Statements to move `q` and `p` point to the next node.

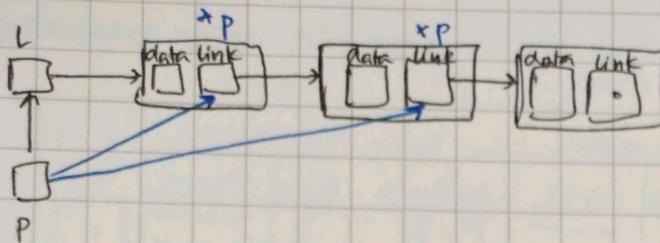
`q = q → link`

`p = &(*p) → link`

lets break this down

Pointer of a pointer i.e the pointer holds the address of another pointer. The "link" portion of the node is a pointer. `*p` is an alias for the pointer pointing to the node with link. `p` holds $(*p) \rightarrow \text{link}$'s address. i.e `p` is pointing to $(*p) \rightarrow \text{link}$.

Illustrated next page.



LOOP CONDITIONS FOR TRAVERSAL

1 CONDITION

- For every call, the traversal has to reach to the end e.g searching for all occurrences of an element.

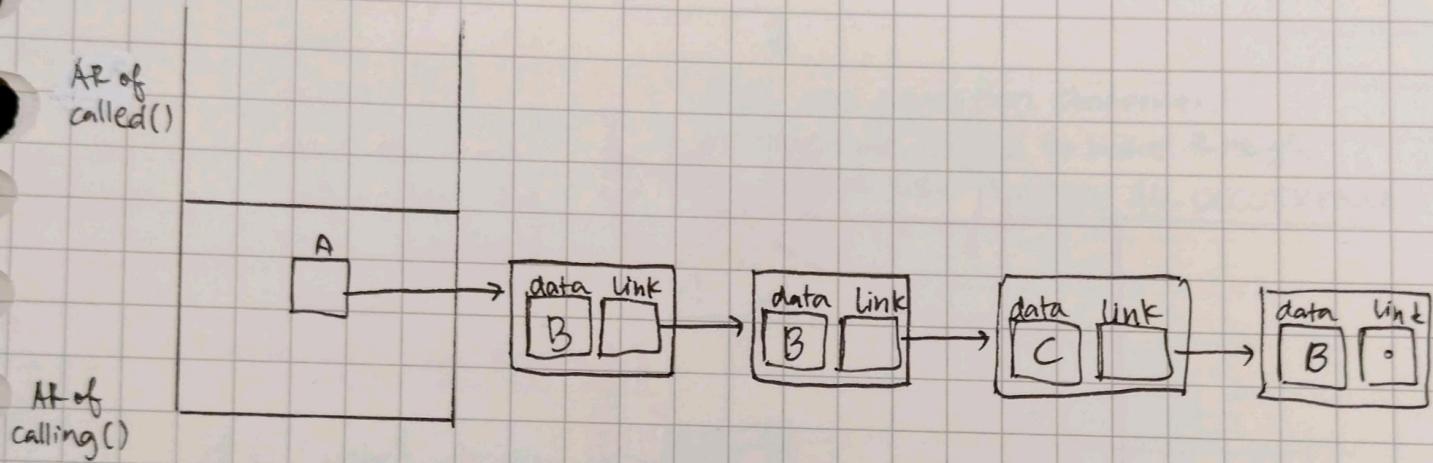
2 CONDITIONS

- When the traversal does not require to traverse the whole list to reach the end.
Eg. delete the first instance of an element.

PRACTICE EXERCISES & SIMULATIONS

Q. Delete all occurrences of a given element.

```
typedef struct node {  
    char data;  
    struct node * link;  
}* LIST;
```



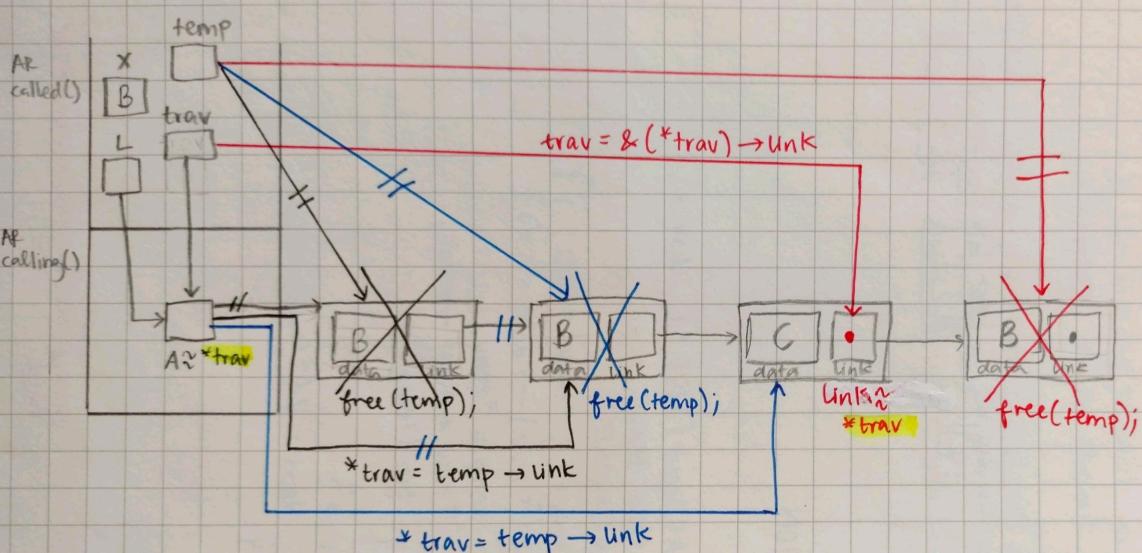
ANS. & SIMULATION

```
void deleteAllOccur (LIST * L, char x) {  
    LIST * trav, temp;  
    if (*trav == NULL) {  
        for (trav = L; *trav != NULL; ) {  
            if ((*trav) → data == x) {  
                temp = *trav;  
                *trav = temp → link;  
                free (temp);  
            }  
            else {  
                trav = &(*trav) → link;  
            }  
        }  
    }  
}
```

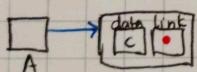
SIMULATION & DRAWING ON NEXT PAGE

deleting all occurrences

SIMULATION



Resulting LIST



CODE

```
void deleteAllOccur(LIST * L, char x) {
```

 deletion involved, pass by address

LIST * trav, temp;
pointer
Pointer
node

```
for (trav = L; *trav != NULL; ) {
```

Note:

In the execution
Stack, trav & L
hold the same
value; they are
both pointers to
pointer.

trav = L, not trav = &L.

Only one condition statement
because we expect to move through
the entire list. "delete All OCCURRENCES"

```
if ((*trav) -> data == x) { // checking if data == x, value to be  
    deleted.
```

```
    temp = *trav; // let temp hold the address of the node to be  
    deleted;
```

*trav = temp -> link;

the link
pointing to the
current node

 let it point to the node
 after the current node (the current node is
 the one that needs to be deleted)

```
    free(temp); // deallocate the node that needs to be deleted.
```

} else {

```
    trav = &(*trav) -> link;
```

 moving to the next node if current node
 doesn't need to be deleted.

}

Q. Same structure, write the code of the function that will delete the first two occurrences

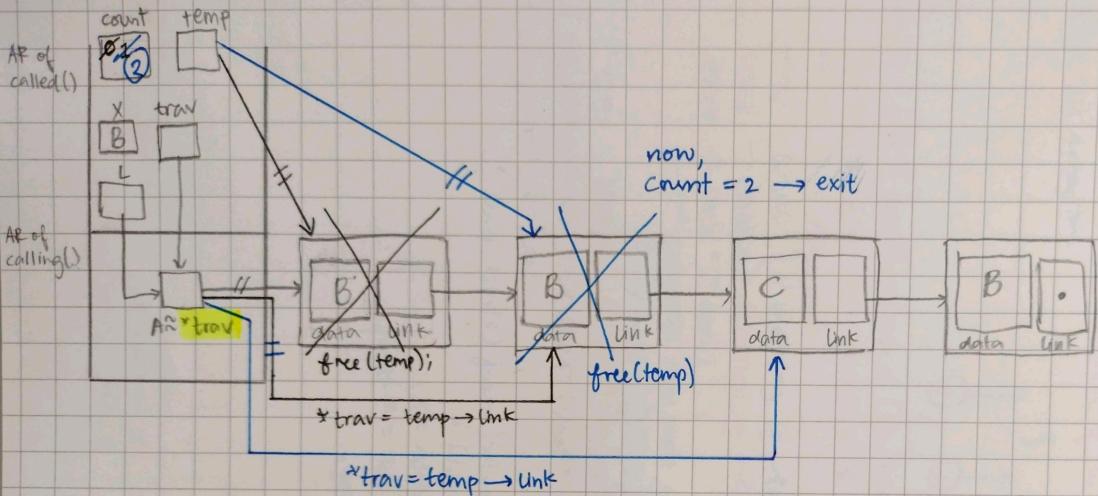
```
void deleteFirstTwo (LIST * L, char x) {  
    LIST * trav, temp;  
    int count = 0;  
    for (trav = L; *trav != NULL && count != 2;) {  
        if ((*trav) → data == x) {  
            temp = *trav;  
            *trav = temp → link;  
            free (temp);  
            count++;  
        } else {  
            trav = &(*trav) → link;  
        }  
    }  
}
```

CODE

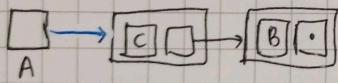
```
void deleteFirstTwo (LIST * L, char x) {  
    LIST * trav, temp;  
    int count = 0; // counter to count no. of occurrences  
    for (trav = L; *trav != NULL && count != 2;) {  
        // two conditions because we don't expect to  
        // traverse the whole list  
        if ((*trav) → data == x) {  
            temp = *trav;  
            *trav = temp → link;  
            free (temp);  
            count++; // increment count if occurrence is spotted  
        } else {  
            trav = &(*trav) → link; // traverse if data ≠ x  
        }  
    }  
}
```

deleting 2 occurrences

simulation



RESULTING LIST



★ PRACTICE

WITH TEST CASES, delete B

A B A C

A C B B

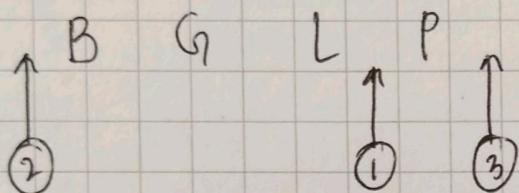
Q. WITH THE SAME DATASTRUCTURE, INSERT SORTED FUNCTION

LIST CONTAINS

B G L P

TEST CASES

- ① INSERT N
- ② INSERT A
- ③ INSERT S



```
typedef struct node {  
    char data;  
    struct node * link;  
} LIST;
```

ANS.

```
void insertSorted (LIST * L, char x) {  
    LIST temp, * trav;  
    for (trav = L; * trav != NULL && (*trav) -> data < x;  
         trav = &(*trav) -> link) {}  
    temp = (LIST) malloc(sizeof(struct node));  
    if (temp != NULL) {  
        temp -> data = x;  
        temp -> link = * trav;  
        * trav = temp;  
    }  
}
```

CODE

```
void insertSorted ( LIST * L, char x ) {  
    LIST temp, * trav;  
    for( trav = L; *trav != NULL && (*trav) -> data < x;  
        trav = &(*trav) -> link ) { }
```

two conditions,
we're not expected to
traverse through the
whole list everytime

(*trav) -> data < x

trav = &(*trav) -> link

compares ASCII
value of the characters

Butterfly notation
loop, once the
conditions are no
longer met, loop exists.

```
temp = (LIST) malloc (size of (struct node));
```

// dynamically allocate temp so that it
points to a new node

```
if (temp != NULL) { // checking if allocation is successful
```

temp -> data = x; // assigning values to the new node

temp -> link = *trav;

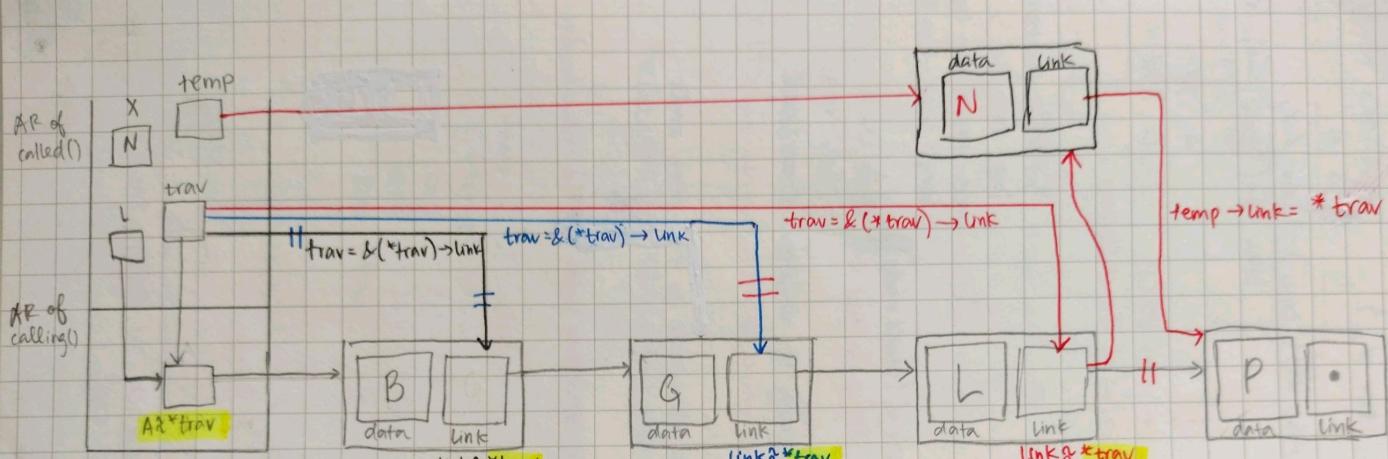
↳ connects new node to the next node
in the list (OR NULL if last node)

*trav = temp;

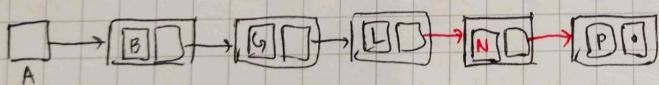
↳ connects previous node to the new node

}

SIMULATION ON NEXT PAGE
TEST CASE 1



RESULTING LIST



* PRACTICE
WITH OTHER
REMAINING TEST CASES