Updating Views

How can I insert a tuple into a table that doesn’t exist?

```
CREATE VIEW  JoeBarSells AS
  SELECT beer, price
  FROM  Sells
  WHERE bar = ‘joe bar’;
```

If we make the following insertion:

```
INSERT INTO  JoeBarSells
VALUES(“bud special”, 3.5)
```

It becomes:

```
INSERT INTO  Sells
VALUES(NULL, ‘bud special’, 3.5)
```

Q: Is the new tuple in table Sells? In the view JoeBarSells?

Non-Updatable Views

```
CREATE VIEW  Champaign-view  AS

        SELECT  name, product, store
        FROM     Person, Purchase
        WHERE   Person.city = “Champaign”    AND
                    Person.name = Purchase.buyer
```

How can we add the following tuple to the view?

(“Joe”, “Shoe Model 12345”, “Nine West”)

We need to add “Joe” to Person first. One copy ? More copies ?

CS411
Database Systems

view

06d: SQL-4
Constraints and Triggers

Why Do We Learn This?

Expected behavior of program

"Assertion"

↓

Expected behavior of data

"health"

part of <u>schema</u>

① ②

Constraints and Triggers

View updates

rel.

tables"

attr.

values

①

- A *constraint* is a relationship among data elements that the DBMS is required to enforce.
  - Example: key constraints.

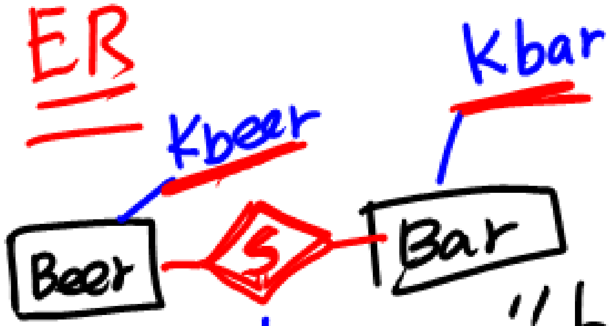
②

- *Triggers* are only executed when a specified condition occurs, e.g., insertion of a tuple.
  - Easier to implement than many constraints.

Kinds of Constraints

- Keys. = unique null
- Foreign-key, or referential-integrity.
- a) Value-based constraints.
  - Constrain values of a particular attribute.
- b) Tuple-based constraints.
  - Relationship among components.
- c) Assertions: any SQL boolean expression.

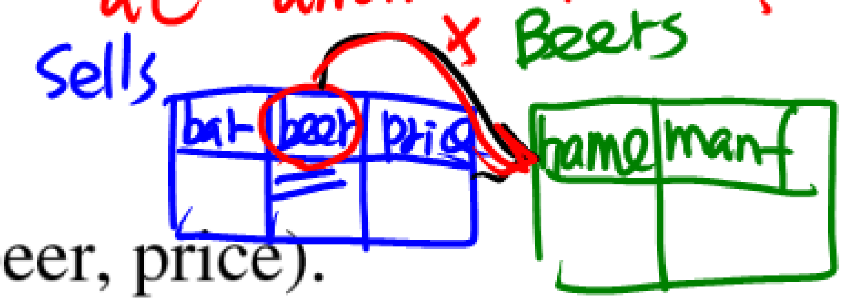
  table-based,

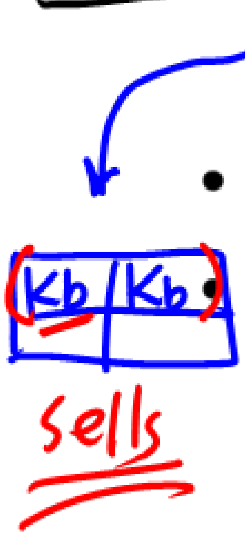


Foreign Keys

a key that is at another table.

"beer" = key at
at Sells Beers.





- Consider Relation Sells(bar, beer, price).
- We might expect that a beer value is a real beer --- something appearing in Beers.name .
- A constraint that requires a beer in Sells to be some key in Beers is called a *foreign -key* constraint.

constraint ?

- must be in beers table.
- has to be a key.

Expressing Foreign Keys

- Use the keyword REFERENCES, either:
  1. Within the declaration of an attribute, when only one attribute is involved.
  2. As an element of the schema, as:

  FOREIGN KEY (<list of attributes>)
    REFERENCES <relation> (<attributes>)

- Referenced attributes must be declared PRIMARY KEY or UNIQUE.

Example: With Attribute

```
CREATE TABLE Beers (
  name     CHAR(20) PRIMARY KEY,  / Unique
  manf     CHAR(20) );

CREATE TABLE Sells (
  bar      CHAR(20),
  beer     CHAR(20) REFERENCES Beers(name),
  price    REAL );
```

Example: As Element

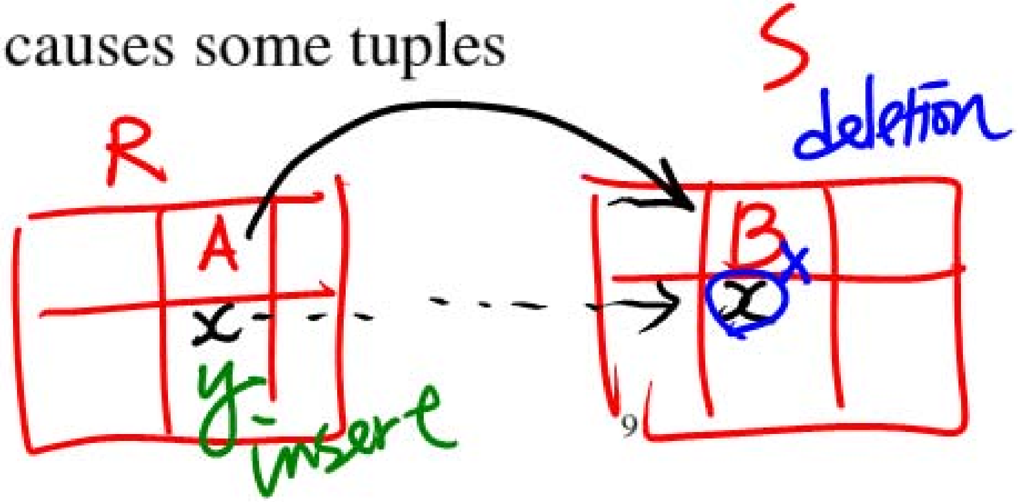
```
CREATE TABLE Beers (
  name    CHAR(20) PRIMARY KEY,
  manf    CHAR(20) );

CREATE TABLE Sells (
  bar     CHAR(20),
  beer    CHAR(20),
  price   REAL,
  FOREIGN KEY(beer) REFERENCES
     Beers(name));
```

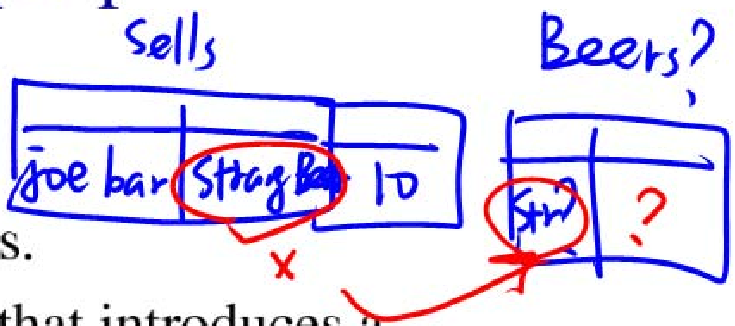
Enforcing Foreign-Key Constraints

- If there is a foreign-key constraint from attributes of relation *R* to the primary key of relation *S*, two violations are possible:

1. An insert or update to *R* introduces values not found in *S*.
2. A deletion or update to S causes some tuples of *R* to "dangle."



Actions Taken -- 1



- Suppose *R* = Sells, *S* = Beers.
- An insert or update to Sells that introduces a nonexistent beer must be rejected. X

at R

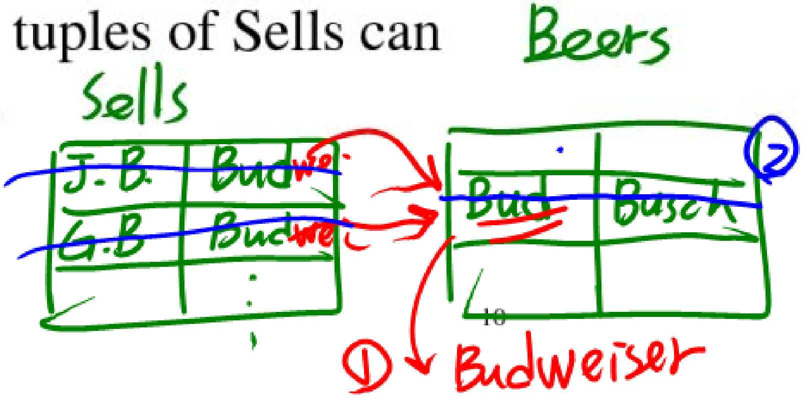
- A deletion or update to Beers that removes a beer value found in some tuples of Sells can be handled in three ways.

① *cascade

update

② Delete,

③ set Null



Actions Taken -- 2

- The three possible ways to handle beers that suddenly cease to exist are:

1. *Default* : Reject the modification.
2. *Cascade* : Make the same changes in Sells.
   - Deleted beer: delete Sells tuple.
   - Updated beer: change value in Sells.

X 3. *Set NULL* : Change the beer to NULL.

beer@Sells if beer is key

Example: Cascade

- Suppose we delete the Bud tuple from Beers.
  - Then delete all tuples from Sells that have beer = ’Bud’.
- Suppose we update the Bud tuple by changing ’Bud’ to ’Budweiser’.
  - Then change all Sells tuples with beer = ’Bud’ so that beer = ’Budweiser’.

Example: Set NULL

- Suppose we delete the Bud tuple from Beers.
  - Change all tuples of Sells that have beer = ’Bud’ to have beer = NULL.
- Suppose we update the Bud tuple by changing ’Bud’ to ’Budweiser’.
  - Same change.

Choosing a Policy

- When we declare a foreign key, we may choose policies SET NULL or CASCADE independently for deletions and updates.
- Follow the foreign-key declaration by:

ON [UPDATE, DELETE][SET NULL CASCADE]

- Two such clauses may be used.
- Otherwise, the default (reject) is used.

Example

```
CREATE TABLE Sells (
  bar    CHAR(20),
  beer   CHAR(20),
  price REAL,
  FOREIGN KEY(beer)
     REFERENCES Beers(name)
     ON DELETE SET NULL
     ON UPDATE CASCADE );
```

Check (not in current MySQL)

- Attribute-based

- Tuple-based

Attribute-Based Checks

- Put a constraint on the value of a particular attribute.

a conditon that must hold.

- CHECK(<condition>) must be added to the declaration for the attribute.
- The condition may use the name of the attribute, but any other relation or attribute name must be in a subquery.

Example

① ~ attr check

② ~ any condition

#

Foreign key?

scope

name

subquery

```
CREATE TABLE Sells (
  bar    CHAR(20),
  beer   CHAR(20)   CHECK ( beer IN
        (SELECT name FROM Beers)),
  price REAL CHECK ( price <= 5.00 )
);
```

③ when to "check"

⇒ [insert / update this attr.

i) name in Beers ≠ key

ii) "one-side" check.

Q: How is Check different from Foreign Key?

```
… beer CHAR(20) CHECK ( beer IN
        (SELECT name FROM Beers))
… price REAL CHECK (price <= 5.00)
```

1. The kind of conditions to enforce?

2. The timing/actions of enforcing?

Timing of Checks

- An attribute-based check is checked only when a value for that attribute is inserted or updated.
  - Example: CHECK (price <= 5.00) checks every new price and rejects it if it is more than $5.
  - Example: CHECK (beer IN (SELECT name FROM Beers)) not checked if a beer is deleted from Beers (unlike foreign-keys).

Tuple-Based Checks

- CHECK (<condition>) may be added as another element of a schema definition.
- The condition may refer to any attribute of the relation, but any other attributes or relations require a subquery.
- Checked on insert or update only.

Example: Tuple-Based Check

- Only Joe's Bar can sell beer for more than $5:

```
CREATE TABLE Sells (
    bar ①      CHAR(20),
    beer       CHAR(20),
    price ②    REAL,
    CHECK (bar = 'Joe''s Bar' OR
                price <= 5.00)
);
```

Element

select avg(price)
from . . .

subquery to ref any attr/table.

only Joe Bar can be exp

Q: Why do we need tuple-level check?

- We can do attribute-based check, why tuple level?

- Give examples that you need to use tuple check.

attr-level
tuple-level
table

Assertions (not in current MySQL)

- These are database-schema elements, like relations or views.

CREATE Table Sells (
)

- Defined by:

```
        CREATE ASSERTION <name>
                CHECK ( <condition> );
```

- Condition may refer to any relation or attribute in the database schema.

schema = rel. + assertion

Example: Assertion

? why use assert can I use ?

- attr check
- tuple check

- In Sells(bar, beer, price), no bar may charge an average of more than $5.

```
CREATE ASSERTION NoRipoffBars CHECK
   (
   NOT EXISTS (
        SELECT bar FROM Sells
        GROUP BY bar
        HAVING 5.00 < AVG(price)
   ));
```

Bars with an average price above $5

hold. if all bars cheap

avg of price by bar,

Example: Assertion

- In Drinkers(name, addr, phone) and Bars(name, addr, license), there cannot be more bars than drinkers.

```
CREATE ASSERTION FewBar CHECK (
  (SELECT COUNT(*) FROM Bars) <=
  (SELECT COUNT(*) FROM Drinkers)
);
```

Timing of Assertion Checks

- In principle, we must check every assertion after every modification to any relation of the database.
- A clever system can observe that only certain changes could cause a given assertion to be violated.
  - Example: No change to Beers can affect FewBar. Neither can an insertion to Drinkers.

Triggers: Motivation

- Attribute- and tuple-based checks have limited capabilities.
- Assertions are sufficiently general for most constraint applications, but they are hard to implement efficiently.
  - The DBMS must have real intelligence to avoid checking assertions that couldn't possibly have been violated.

Triggers: Solution

programmer.

- A trigger allows the user to specify when the check occurs.
- Like an assertion, a trigger has a general-purpose condition and also can perform any sequence of SQL database modifications.

Event-Condition-Action Rules

Why more eff. than assertion?

- Another name for "trigger" is *ECA rule*, or event-condition-action rule.
- *Event* : typically a type of database → insert on sells modification, e.g., "insert on Sells."
- *Condition* : Any SQL boolean-valued expression. newValue & Beers.name
- *Action* : Any SQL statements.

→ insert new val ~~to~~ Beers.

Example: A Trigger

- There are many details to learn about triggers.
- Here is an example to set the stage.
- Instead of using a foreign-key constraint and rejecting insertions into Sells(bar, beer, price) with unknown beers, a trigger can add that beer to Beers, with a NULL manufacturer.

schema element => create table R (...)

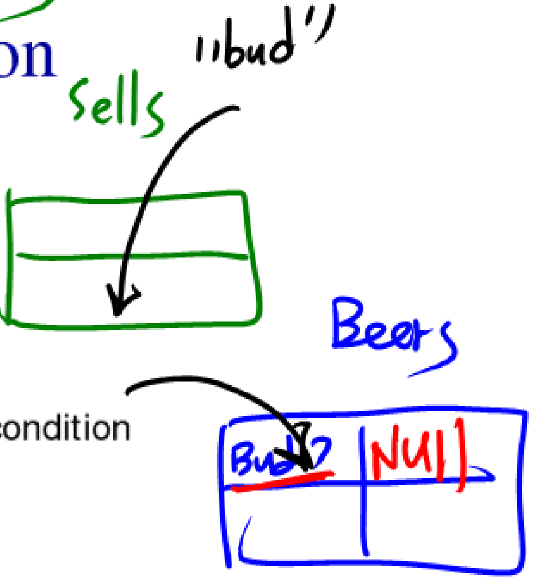
Example: Trigger Definition

```
CREATE TRIGGER BeerTrig
    AFTER INSERT ON Sells
    REFERENCING NEW ROW AS NewTuple
    FOR EACH ROW
    WHEN (NewTuple.beer NOT IN
        (SELECT name FROM Beers))
    INSERT INTO Beers(name)
        VALUES(NewTuple.beer);
```

The event

The condition

The action



What happen?

if no, => reject.

Options: The Event

- AFTER can be BEFORE.
  - Also, INSTEAD OF, if the relation is a view.
    - A great way to execute view modifications: have triggers translate them to appropriate modifications on the base tables.
- INSERT can be DELETE or UPDATE.
  - And UPDATE can be UPDATE . . . ON a particular attribute.

Options: FOR EACH ROW

- Triggers are either *row-level* or *statement-level*.
- FOR EACH ROW indicates row-level; its absence indicates statement-level.
- Row level triggers are executed once for each modified tuple.
- Statement-level triggers execute once for an SQL statement, regardless of how many tuples are modified.

Options: REFERENCING

- INSERT statements imply a new tuple (for row-level) or new set of tuples (for statement-level).
- DELETE implies an old tuple or table.
- UPDATE implies both.
- Refer to these by

[NEW OLD][TUPLE TABLE] AS <name>

Options: The Condition

- Any boolean-valued condition is appropriate.
- It is evaluated before or after the triggering event, depending on whether BEFORE or AFTER is used in the event.
- Access the new/old tuple or set of tuples through the names declared in the REFERENCING clause. (or fixed by "OLD", "NEW" in MySQL.)

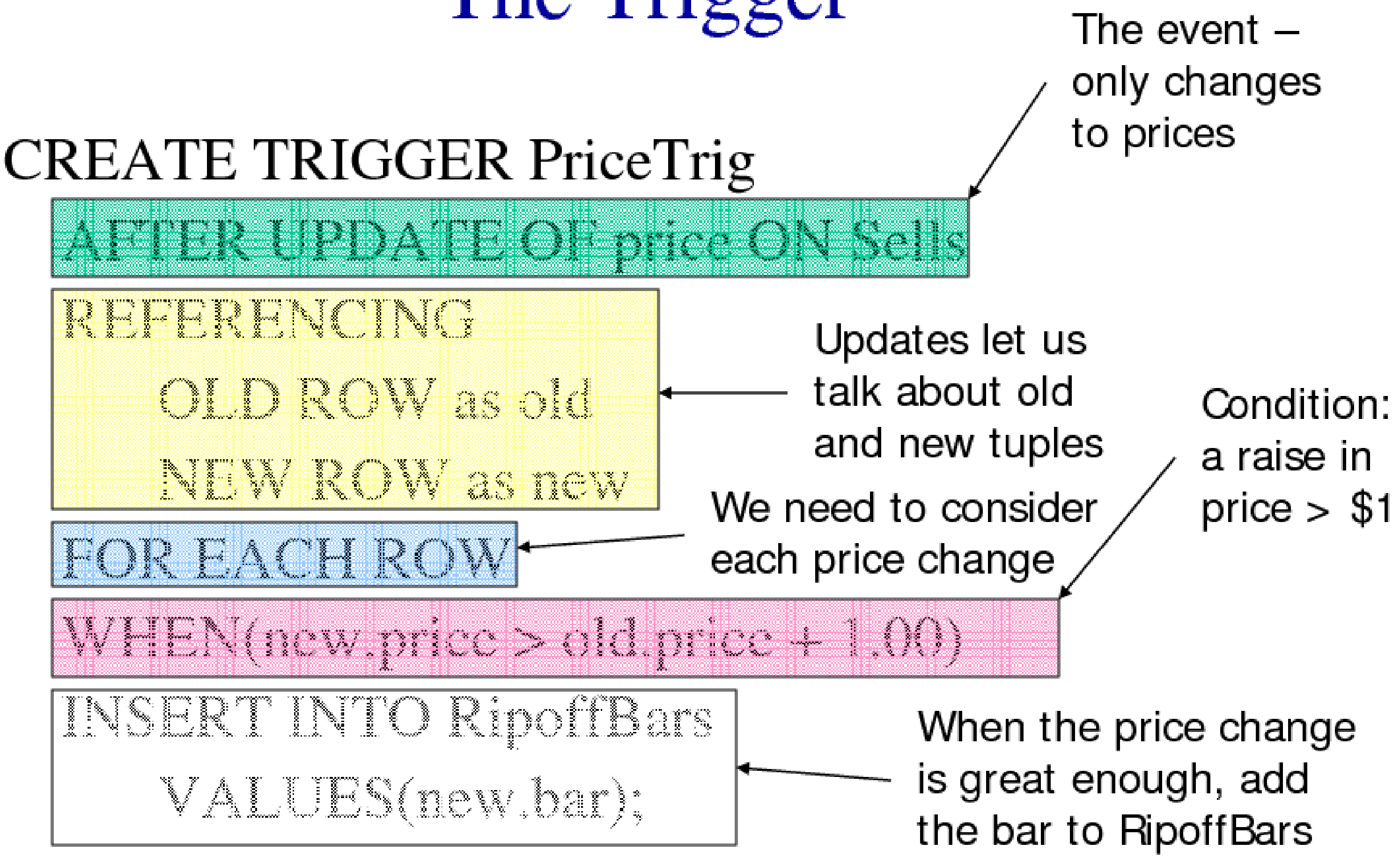
Options: The Action

- There can be more than one SQL statement in the action.
  - Surround by BEGIN . . . END if there is more than one.
- But queries make no sense in an action, so we are really limited to modifications.

Another Example

- Using Sells(bar, beer, price) and a unary relation RipoffBars(bar) created for the purpose, maintain a list of bars that raise the price of any beer by more than $1.

The Trigger



```
CREATE TRIGGER PriceTrig
  AFTER UPDATE OF price ON Sells
  REFERENCING
       OLD ROW as old
       NEW ROW as new
  FOR EACH ROW
  WHEN(new.price > old.price + 1.00)
  INSERT INTO RipoffBars
       VALUES(new.bar);
```

The event – only changes to prices

Updates let us talk about old and new tuples

We need to consider each price change

Condition: a raise in price > $1

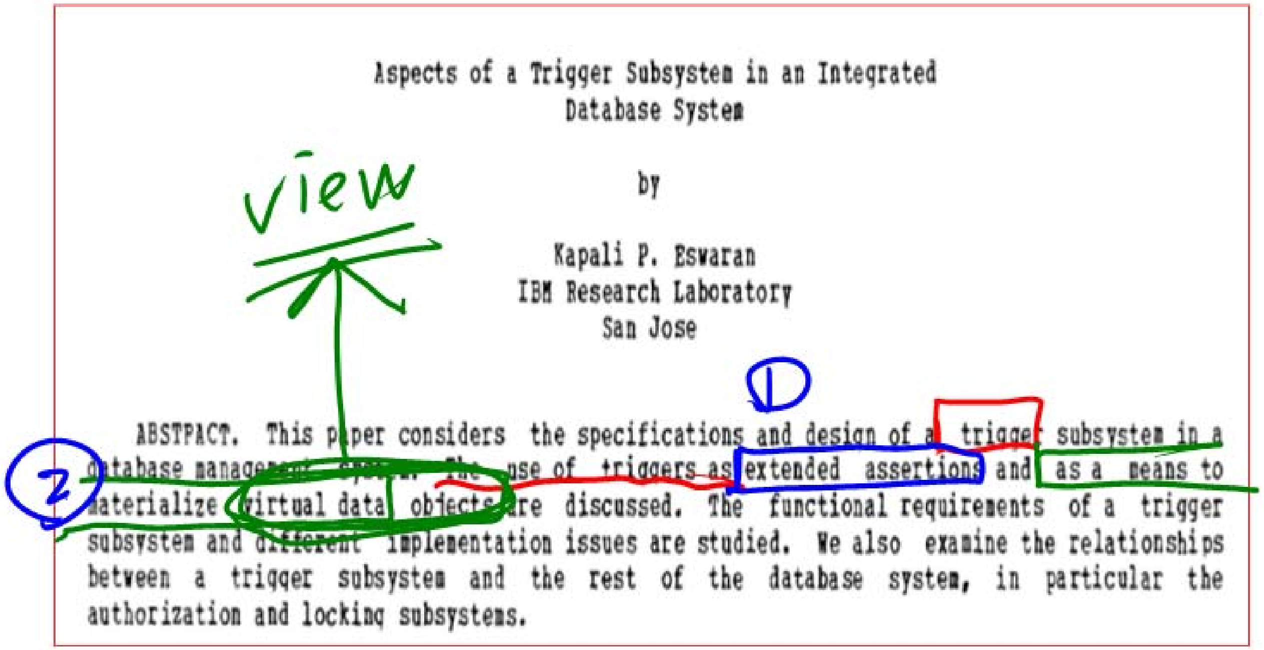
When the price change is great enough, add the bar to RipoffBars

Behind the Scene: Why Trigger was invented?

Aspects of a trigger subsystem in an integrated database system. Proceedings of the 2nd international conference on Software engineering. 1976.

1. Extended assertions. (why?)
2. ??

Behind the Scene: This is why…



Aspects of a Trigger Subsystem in an Integrated
Database System

by

Kapali P. Eswaran
IBM Research Laboratory
San Jose

ABSTPACT. This paper considers the specifications and design of a trigger subsystem in a database management system. The use of triggers as extended assertions and as a means to materialize virtual data objects are discussed. The functional requirements of a trigger subsystem and different implementation issues are studied. We also examine the relationships between a trigger subsystem and the rest of the database system, in particular the authorization and locking subsystems.