

CPSC 335- Algorithm Engineering

Fall 2020

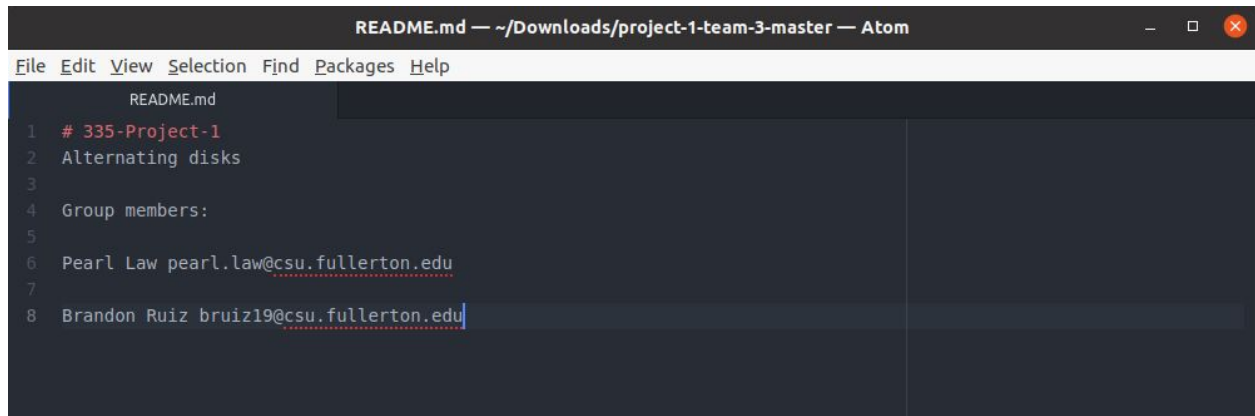
Instructor: Mike Peralta

Project 1: Team 3

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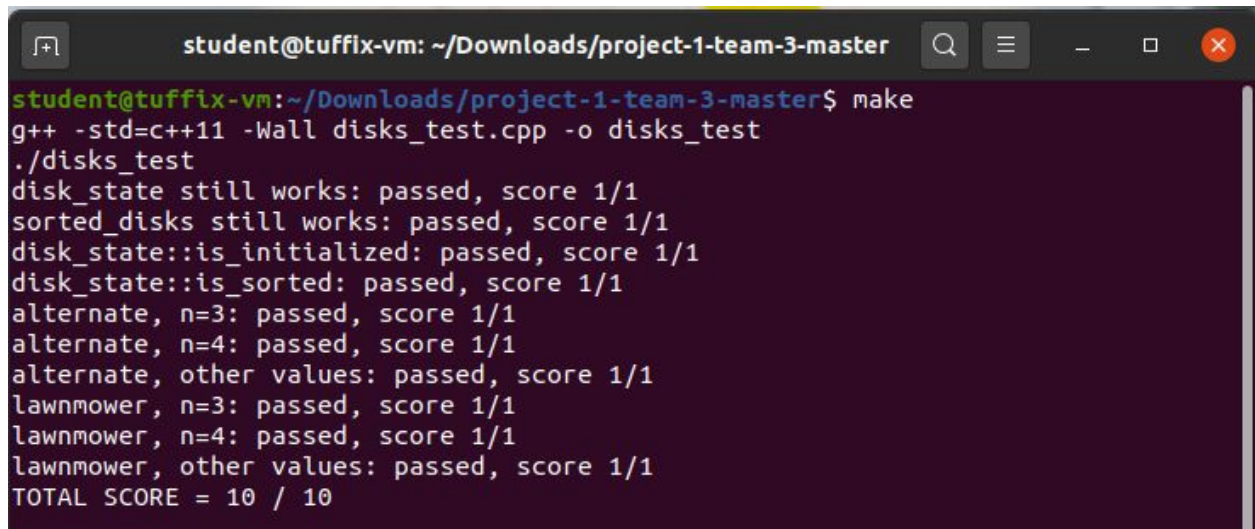
Brandon Ruiz (bruiz19@csu.fullerton.edu)

Tuffix Screenshot of README.md inside Atom editor:



```
README.md — ~/Downloads/project-1-team-3-master — Atom
File Edit View Selection Find Packages Help
README.md
1 # 335-Project-1
2 Alternating disks
3
4 Group members:
5
6 Pearl Law pearl.law@csu.fullerton.edu
7
8 Brandon Ruiz bruiz19@csu.fullerton.edu
```

Screenshot of code executing make command:



```
student@tuffix-vm: ~/Downloads/project-1-team-3-master
student@tuffix-vm:~/Downloads/project-1-team-3-master$ make
g++ -std=c++11 -Wall disks_test.cpp -o disks_test
./disks_test
disk_state still works: passed, score 1/1
sorted_disks still works: passed, score 1/1
disk_state::is_initialized: passed, score 1/1
disk_state::is_sorted: passed, score 1/1
alternate, n=3: passed, score 1/1
alternate, n=4: passed, score 1/1
alternate, other values: passed, score 1/1
lawnmower, n=3: passed, score 1/1
lawnmower, n=4: passed, score 1/1
lawnmower, other values: passed, score 1/1
TOTAL SCORE = 10 / 10
```

Problem name: Disk sorter

Input: a positive integer n and a list of $2n$ disks of alternating colors light-dark, starting with light

Output: a list of $2n$ disks, the first n disks are dark, the next n disks are light, and an integer m representing the number of swaps to move the light ones after the dark ones

Lawnmower Algorithm Pseudocode:

```
def lawnmower(integer n, list of disks L):
    swaps = 0

    if n == 0:
        return Nothing

    else:
        for i = 0 to 2n do
            for j = 0 to 2n - 1 do          // left to right movement
                if L[j] > L[j + 1]:        // swap if light disk (L) > dark disk (D)
                    then
                        temp = L[j]
                        L[j] = L[j + 1]
                        L[j + 1] = temp
                        swaps += 1

            for k = 2n - 1 to 1 do          // right to left movement
                if L[k] < L[k - 1]:        // swap if dark disk (D) < light disk (L)
                    then
                        temp = L[k]
                        L[k] = L[k - 1]
                        L[k - 1] = temp
                        swaps += 1

        return sorted disk list and number of swaps needed
```

Alternate Algorithm Pseudocode:

```
define alternate (integer n, list of disks L):
    swaps = 0

    if n = 0;
        return no value (0)

    else
        for i = 0 to n - 1 do
            for j = 0 to n - 1 do
                if L[j + 1] < L[j]
                    then
                        temp = L[j]
                        L[j] = L[j + 1]
                        L[j + 1] = tmp
                        swap += L[j + 1]

|
return disk list after sorted with executed swaps
```

Mathematical Analysis:

*Abbreviations: s.c = step count

1. Lawnmower algorithm:

a. **Evaluate if statement** (if $n == 0$) = **1 step**

b. **Evaluate else statement** = **1 step**

1. s.c_{1st inner for loop} = # loop iterations * s.c loop block
loop iterations = $((2n - 1) - 0 / 1) + 1 = 2n$

s.c loop block => evaluate if statement:

if statement = 1 step

then block = 4 steps

else = 0 steps (no else statement to execute)

s.c_{if statement} = $1 + \max(4, 0) = 1 + 4 = 5$ steps

s.c_{1st inner for loop} = $2n * 5 = \underline{10n}$

2. s.c_{2nd inner for loop} = # loop iterations * s.c loop block
loop iterations = $((1 - (2n - 1)) / 1) + 1 = -2n + 2 + 1 = |-2n + 3| = 2n + 3$
(take absolute value because we cannot have negative iterations/steps)

s.c loop block => evaluate if statement:

if statement = 1 step

then block = 4 steps

else = 0 steps (no else statement to execute)

s.c_{if statement} = $1 + \max(4, 0) = 1 + 4 = 5$ steps

s.c_{2nd inner for loop} = $(2n+3) * 5 = \underline{10n + 15}$

3. s.c_{outer for loop} = else overhead + $\sum_{i=0}^{2n} \text{outer for loop} * (\text{s.c}_{1st \text{ inner for loop}} + \text{s.c}_{2nd \text{ inner for loop}})$
 $= 1 + \sum_{i=0}^{2n} (10n + 10n + 15)$
 $= 1 + \sum_{i=0}^{2n} (20n + 15)$
 $= 1 + \sum_{i=0}^{2n} (20n) + \sum_{i=0}^{2n} (15)$
 $= 1 + 0 + \sum_{i=1}^{2n} (20n) + (15 * (2n + 1))$
 $= 1 + 20 * \sum_{i=1}^{2n} (n) + 30n + 15$
 $= 1 + 20 * (n * 2n) + 30n + 15$

$$= 1 + 20 * (2n^2) + 30n + 15$$

$$\text{s.c}_{\text{outer for loop}} = 40n^2 + 30n + 16$$

c. Evaluate function statements-

initializing variable swaps = 1 step

return statement = 1 step

$$\text{Entire function s.c} = 4 + 40n^2 + 30n + 16 = 40n^2 + 30n + 20$$

The step count for the entire algorithm, after dropping the constant, dominated term, and multiplicative constant, is n^2 . This corresponds to the time complexity efficiency class $O(n^2)$.

2. Alternate algorithm:

d. **Evaluate if statement** (if $n == 0$) = 1 step

e. **Evaluate else statement** = 1 step

01. Step count (1st for loop): number of loop executed * step count block
number of loop executed: $(((n-1)-0) / 1) + 1 = n$

step count block: evaluating the if statement

if statement = 1 step

step count block = 0 steps (does not execute anything)

else = 0 steps (does not follow up on execution)

step count (1st loop): $1 + \max(0, 0) = 1 + 0 = 1$ step

Step count (1st for loop): $n * 1 = n$

02. Step count (2nd for loop): number of loop executed * step count block
number of loop executed: $(((n-1)-0) / 1) + 1 = n$

step count block: evaluating the if statement

if statement = 1 step

step count block = 3 steps

else = 0 steps (does not follow up on execution)

step count (2nd loop): $1 + \max(3, 0) = 1 + 3 = 4$ steps

Step count (2nd for loop): $n * 4 = 4n$

03. Step count (outside for loop): outside else + Σ (outside for loop) * (1st for loop + 2nd for loop)

$$= 1 + \sum_{i=0}^n (n + 4n)$$

$$\begin{aligned}
&= 1 + \sum_{i=0}^n (5n) \\
&= 1 + 0 + 5 * \sum_{i=1}^n (n) \\
&= 1 + 0 + 5 * (n * n)
\end{aligned}$$

Step count (outside for loop) = $5n^2 + 1$

f. **Evaluate whole algorithm:** $5n^2 + 1 + 3 = 5n^2 + 4$

The step count for the whole algorithm after removing the values unaffected by n is n^2 . This falls under the category of time complexity efficiency of the class: $O(n^2)$.