Example

```
int A[16][128]={all values};
int B[16][128]={all values};
int C[16];
int main() {
       int tmp,i,j;
       for (i=0;i<16;i++) {
              tmp=0;
              for (j=0; j<128; j++)
                     tmp+=A[i][j]*B[i][j];
              C[i]=tmp;
return 0;
```



Example: 1) Memory accesses

```
int main() {
  int tmp,i,j;
  for (i=0;i<16;i++) {
       tmp=0;
       for (j=0; j<128; j++)
       tmp+=A[i][j] * B[i][j];
       C[i]=tmp;
  return 0:
               Memory reads:
               128+128
    Memory write
                    Memory accesses
```

```
ldr
         r9, [r8, r1, asl 2]
        r4, [r7, r1, asl 2]
ldr
         r0, r4, r9, r0
mla
add
         r1, r1, #1
         r1, #128
cmp
bne
         Ini-j
         r5, r5, #1
add
         r5, #16
cmp
str
         r0, [r6, #4]!
         Ini-i
bne
```

Accesses for fetching instructions are not considered in this example.

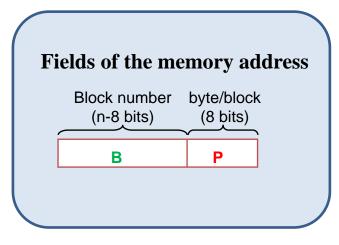
Example: 2) Memory layout

```
A LANGE OF THE PARTY OF THE PAR
```

```
int A[16][128]={values};
int B[16][128]={values};
                          Address
                                        Memory
int C[16];
                         0x0C000000
                                      A[0][0]
int main() {
                                      A[0][1]
      return 0; }
                           + 0x2000
                                              16x128x4 B
                                      A[1][0]
Linker script
SECTIONS {
                         0x0C002000
                                      B[0][0]
.data : {
                                      C[0]
                         0x0C004000
*(.data)
                                                  16x4 B
                                      C[1]
.bss : {
*(.bss)
.text : {
                                      main()
  *(.text) }
```

Example: 3) Memory blocks

Assume a cache memory with block of 256B



Address	Memory	В	lock number
0x0C0000 <mark>00</mark>	A[0][0]		C0000
	A[0][1]		C0000
	•••	256 B	
0x0C0000 F0	A[0][62]		
	A[0][63]		
0x0C0001 <mark>00</mark>	A[0][64]		00004
	•••		C0001
0x0C000200	A[1][0]		C0002
its/block			_
ocks			

 $\frac{256 \ bytes/block}{4 \ bytes/element} = 64 \ elements/block$ Matriz A needs 32 memory blocks

Example: 4) Cache structure

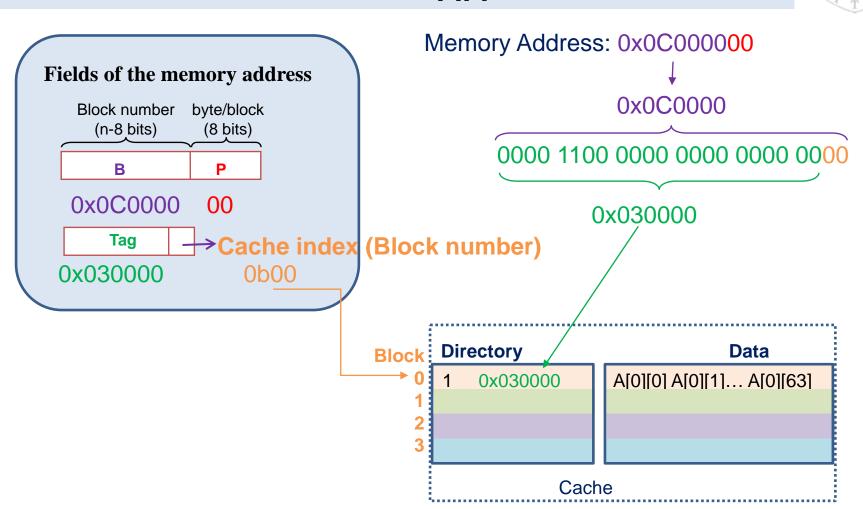
Assume a cache memory of 1KB with block of 256B

$$\frac{1024 \ bytes}{256 \ bytes/block} = 4 \ blocks$$

Block	Directory	Data	
0	V Tag	Byte-0 byte-1	byte-255
1			
2			
3			
	Cache		

Example: 5) Direct mapped cache

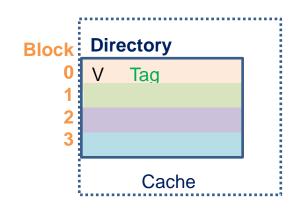
Assume a cache memory of 1KB with block of 256B Where in the cache is A[0][0] stored?



Example: 5) Direct mapped cache (2)

Assume a cache memory of 1KB with block of 256B

Cache block for each data

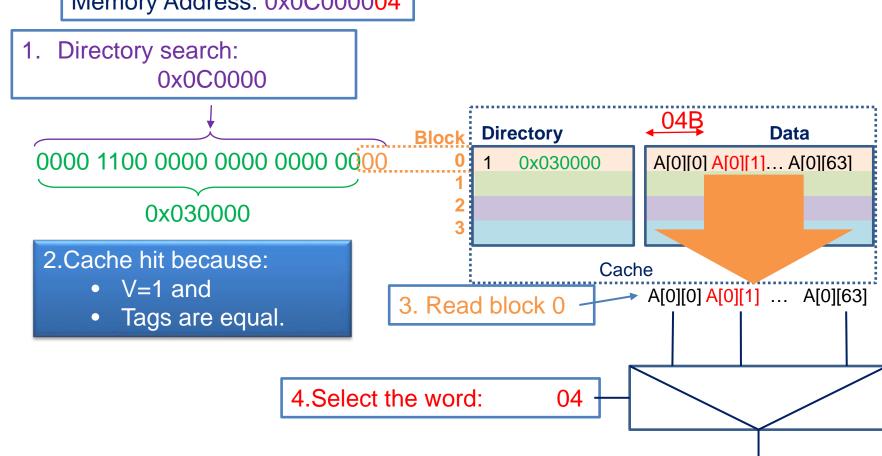


Address	Memory
0x0C000000	A[0][0] A[0][1] A[0][63]
0x0C000400	A[2][0] A[2][1] A[2][63]
0x0C002000	B[0][0] B[0][1] B[0][63]
0x0C0040 <mark>00</mark>	C[0] C[1] C[15]

Example: 6) Cache hit

Assume we want to read A[0][1], and the cache contents are:

Memory Address: 0x0C000004



Example: 7) Cache miss

Assume we want to read B[0][0], and the cache contents are:

Memory Address: 0x0C002000

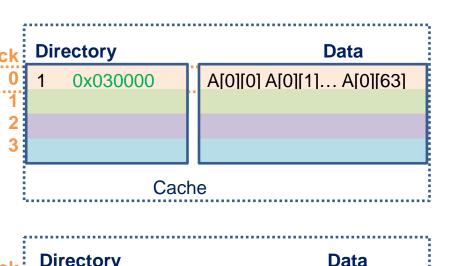
1. Directory search: 0x0C0020

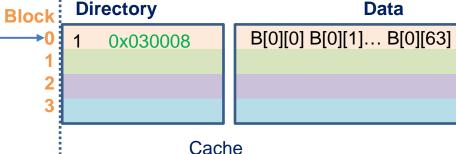
0000 1100 0000 0000 0010 00<mark>00.</mark>

0x030008

- 2. Cache miss
- 3. The block has to be fetched from memory
- 4. Read block 0

5.Select the word: 00





Example: accesses in a direct mapped cache (1)



i=0, j=0

1. Read A[0][0]

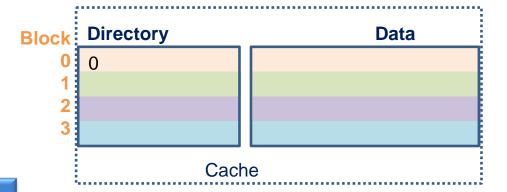
Memory Address: 0x0C000000

Block: 0

Tag: 0x030000

Cache miss (v=0)

```
for (i=0;i<16;i++) {
  tmp=0;
  for (j=0;j<128;j++)
      tmp+=A[i][j] * B[i][j];
  C[i]=tmp; }</pre>
```



Cold miss or compulsory miss: First Access to the block

Example: accesses in a direct mapped cache (2)



i=0, j=0

- 1. Read A[0][0]
- 2. Read B[0][0]

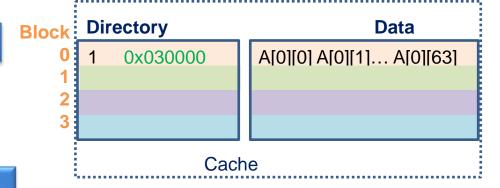
Memory Address: 0x0C000000

Block: 0

Tag: 0x030008

Cache miss (different tags)

```
for (i=0;i<16;i++) {
   tmp=0;
   for (j=0;j<128;j++)
       tmp+=A[i][j] * B[i][j];
   C[i]=tmp; }</pre>
```



Conflict:

2 memory blocks are mapped to the same cache block

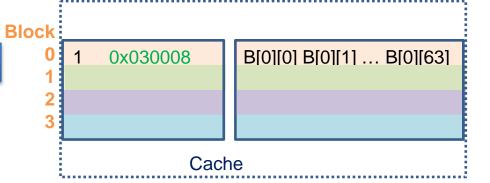
Example: accesses in a direct mapped cache (3)



```
i=0,j=1
1. Read A[0][1]
Memory Address: 0x0C000004
Block: 0
Tag: 0x030000
```

```
for (i=0;i<16;i++) {
  tmp=0;
  for (j=0;j<128;j++)
     tmp+=A[i][j] * B[i][j];
  C[i]=tmp; }</pre>
```

Cache miss (different tags)



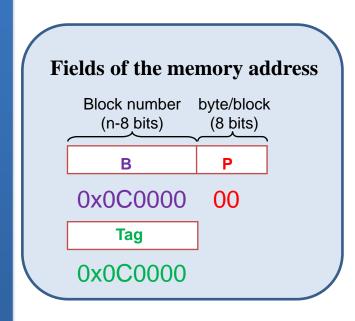
Internal loop

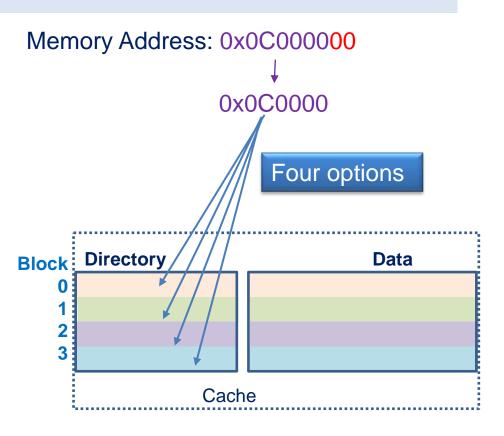
Due to conflicts, all reads are misses: 128x2 read misses

Example: 8) Fully associative cache

Assume a cache memory of 1KB with block of 256B and fully associative Where in the cache is A[0][0] stored?

ANYWHERE





Example: accesses in a fully associative cache (1)



i=0, j=0

1. Read A[0][0]

Memory Address: 0x0C000000

Tag: 0x0C0000

Cache miss (V=0)

2. Read B[0][0]

Memory Address: 0x0C002000

Tag: 0x0C0020

Cache miss (V=0)

i=0, j=1

1. Read A[0][1]

Memory Address: 0x0C000004

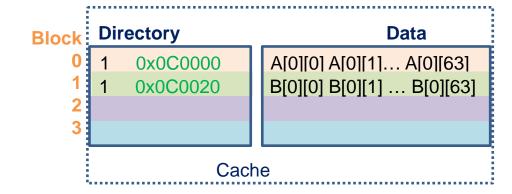
2. Read B[0][1]

Cache hit

Memory Address: 0x0C002004

Cache hit

```
for (i=0;i<16;i++) {
  tmp=0;
  for (j=0;j<128;j++)
     tmp+=A[i][j] * B[i][j];
  C[i]=tmp; }</pre>
```



Example: accesses in a fully associative cache



```
i=0, j=64
```

1. Read A[0][64]

Memory Address: 0x0C000100

Tag: 0x0C0001

Cache miss (V=0)

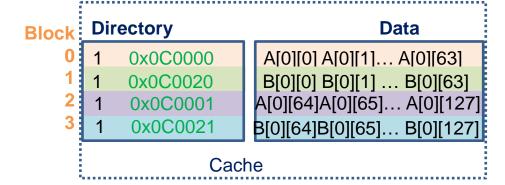
2. Read B[0][64]

Memory Address: 0x0C002100

Tag: 0x0C0021

Cache miss (V=0)

```
for (i=0;i<16;i++) {
   tmp=0;
   for (j=0;j<128;j++)
       tmp+=A[i][j] * B[i][j];
   C[i]=tmp; }</pre>
```



Internal loop

Only compulsory misses: 4 read misses, one for each memory block

Example: 9) Set associative cache

Assume a cache memory of 1KB with block of 256B and 2 ways
Where in the cache is A[0][0] stored?

Number of sets =
$$\frac{Number \ of \ blocks}{Blocks/Set} = \frac{2}{2}$$

