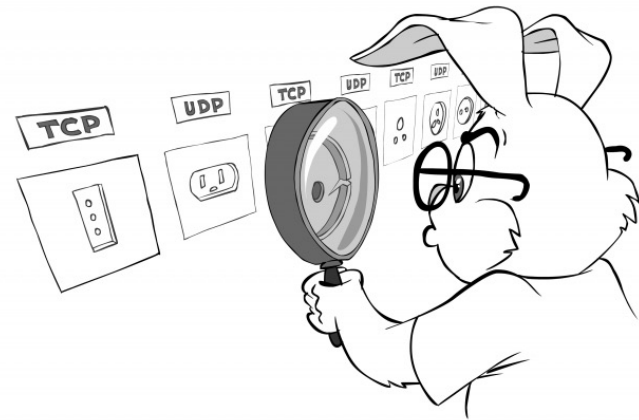


# Outline

- Network Security Threads
- Securing end-to-end connections
  - Transport Layer Security (TLS)
- Security at Transport Layer
  - TCP, UDP, QUIC
- Security Network Layer
  - IP/ICMP, IPSec, VPNs, IPv6 Security



# Transport Layer Attacks and Mitigation

## 1. TCP SYN Flood (DoS/DDoS Attack)

- attackers exploit the TCP three-way handshake by sending a **large number of SYN** requests to a target server **without completing the handshake**, consuming server resources and causing DoS
- Mitigation:
  - **SYN cookies**: the **server encodes connection information\*** in the **sequence # field of SYN-ACK TCP header**; if the client completes the handshaking (w/ ACK) and the cookie received is verified, resources are allocated for the connection
    - currently implemented in most recent OSs
    - e.g., setting the feature in Linux: `echo 1 > /proc/sys/net/ipv4/tcp_syncookies`
  - **Backlog, Rate limiting**: limits the backlog queue size (pending conn.) or the # of SYN requests a server accepts from a single source

\* a cryptographic hash of the connection information, such as IP addresses, port numbers, and a secret value

# Transport Layer Attacks and Mitigation

## 2. TCP Reset Attack

- attackers **inject forged TCP RST** flag (reset) packets to forcefully terminate a TCP session between two parties, if they can guess or sniff the current sequence #s
- Mitigation:
  - use **sequence number randomization** to make sequence numbers more unpredictable
  - implement **TCP/IP stack hardening** on the server side, which ignores RST packets that do not match the sequence numbers of the active session
  - **TCP RST limiting**
  - firewalls and intrusion prevention systems
  - note: TLS protection does not impair a malicious TCP connection reset

# Transport Layer Attacks and Mitigation

## 3. TCP Session Hijacking

- an attacker intercepts and takes over an active TCP session by guessing **sequence numbers** and injecting malicious **data** into the session
  - altering TCP's **sequence numbers** may confuse the TCP state machine, causing desynchronization between the client and server, which may result in connection stalls, unnecessary retransmissions, or dropped segments
  - injecting malicious **data** will be detected and never reach the app, if TLS is in use
- Mitigation:
  - use **TLS to encrypt and authenticate session traffic**, which prevents an attacker from modifying or injecting data into an active session
  - Initial Sequence Number (ISN) randomization
  - ensure enhanced session management, namely, through:

# Transport Layer Attacks and Mitigation

## 3. TCP Session Hijacking (cont.)

- **TCP-AO** – cryptographically authenticates TCP segments using MAC over TCP header for every packet; protects the TCP state machine
  - TCP Authentication Option, not widely deployed (RFC5925)
- **Connection Token Validation** – adds an extra layer of session-specific validation at the transport or application layer
  - e.g., token negotiated during handshaking, not part of standard TCP
- **IPSec (Transport Mode)** – provides encryption and integrity protection at the network layer, securing TCP sessions
- **Middlebox Support** – firewalls, IDS/IPS can monitor and protect the TCP session state

# Transport Layer Attacks and Mitigation

## 1. UDP Flood (DoS/DDoS Attack)

- attackers flood a target with large volumes of UDP packets, overwhelming the network or server, leading to resource exhaustion
- Mitigation:
  - **Rate limiting** – implement rate limits on incoming UDP traffic
  - **DDoS protection services** – filter out malicious traffic using
    - cloud-based (e.g. Cloudflare, Akamai Prolexic, AWS Shield)
    - network-based DDoS mitigation solutions (e.g., Verizon, AT&T,...)
    - open source (sw, tools): iptables/NetFilter; DDoS Deflate; Snort; Suricata

# Transport Layer Attacks and Mitigation

## 2. UDP Amplification

- attackers exploit UDP-based services (such as DNS or NTP) that respond to small queries with large responses
- by sending forged requests with the target's IP address, the attacker **amplifies** the attack traffic, overwhelming the target with answers to queries it didn't ask ("reflected attack")
- Mitigation:
  - configure UDP services (e.g., DNS, NTP) to **prevent amplification**, using DNS Response Rate Limiting (DNS RRL) and ensuring that NTP servers are patched
    - RRL detects patterns in arriving queries, and when it finds a pattern that suggests abuse, it can reduce the rate at which the replies are sent (from BIND 9.10 on)
  - complement with **egress filtering** to prevent traffic from leaving your network with a spoofed source address (attacker inside the network domain)

# Transport Layer Attacks and Mitigation

## 3. UDP and TCP Port Scanning

- Mitigation:
  - use a firewall to block unnecessary or unused ports
  - enable IDS/IPS to detect and block suspicious port scanning activity
- **Slow Port Scanning** – spread over hours and days; hard to detect
  - time-distributed probing (eludes rate-limiting security approaches)
  - low-traffic volume (looks like legitimate traffic)

difficulties:

- detection is resource intensive (long-term traffic analysis and continuous mon.)
- even harder if slow scanning rate is adjusted dynamically
- risk of false positives



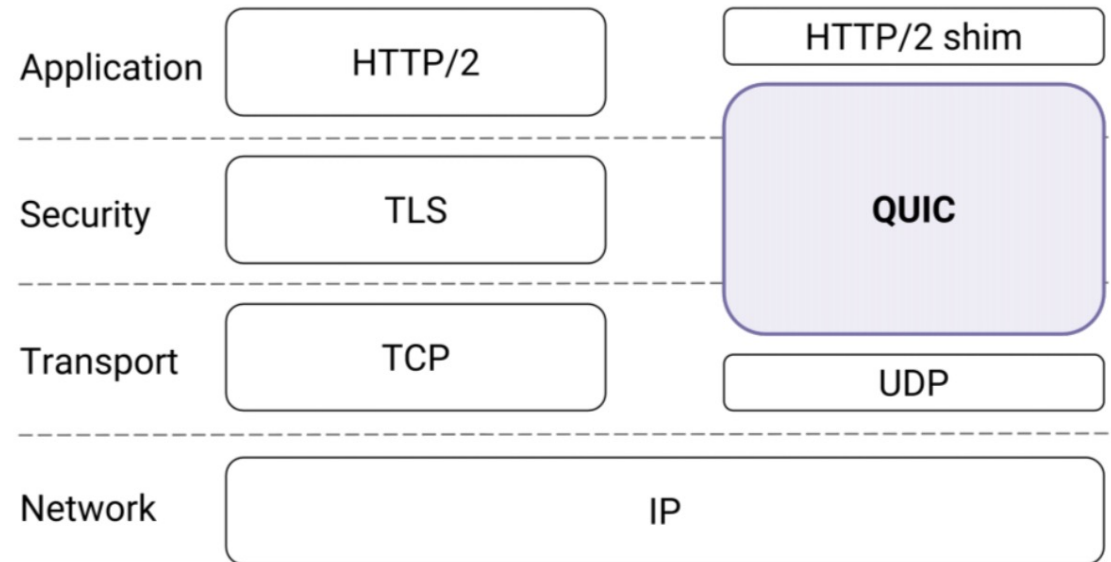
# Transport Layer - QUIC

## QUIC - UDP-Based Multiplexed and Secure Transport

- secure general-purpose transport protocol, no native
- introduced by Google (2012), now a IETF standard RFC 9000 (2021)
- **motivation**: overcome TCP limitations (single stream, HOL blocking, unsecure and slow handshaking)
- runs over UDP for flexibility and user space control (instead of kernel)
- **core features**: secure handshaking, flow control, congestion control, error control, reliable delivery, 0-RTT, stream-based (ordered sequence of bytes), multiple streams per connection (multiplexing)
- integrated as a Chromium component, supported by major browsers

# Transport Layer - QUIC

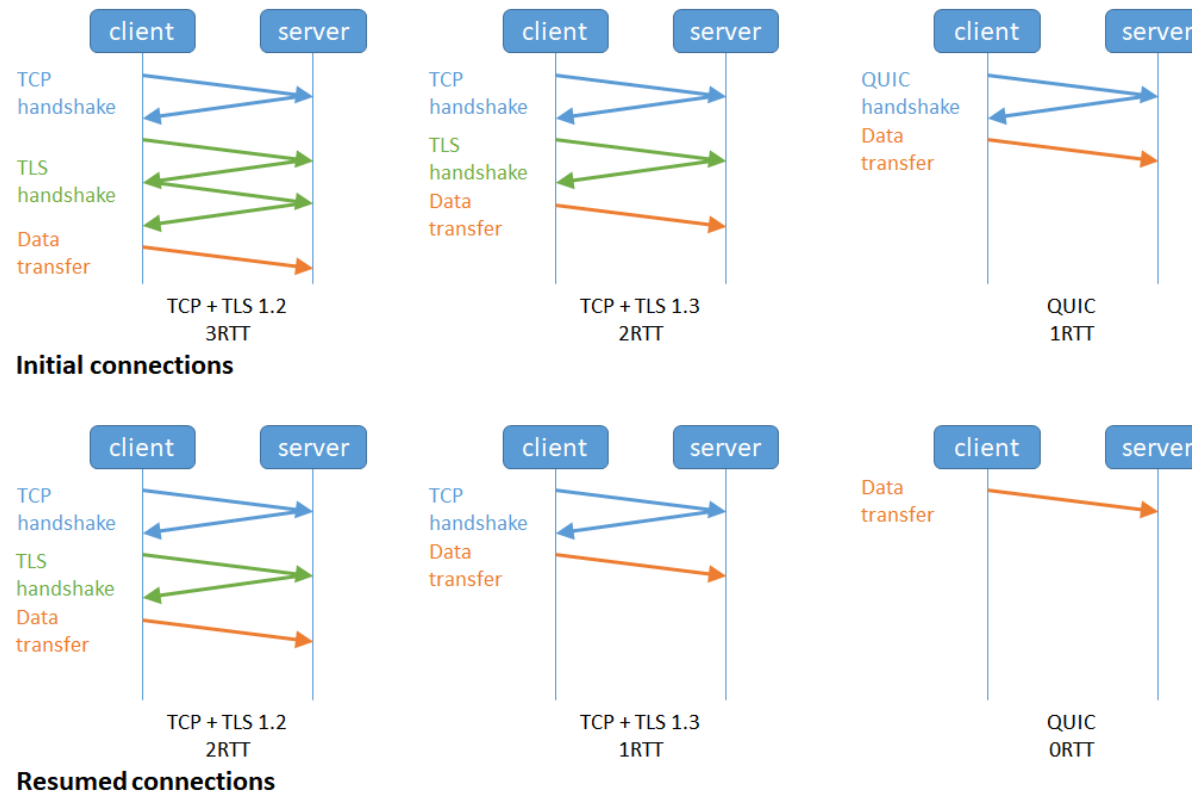
- HTTP/3 built on top of IETF QUIC



**Figure 1: QUIC in the traditional HTTPS stack.**

# Transport Layer - QUIC

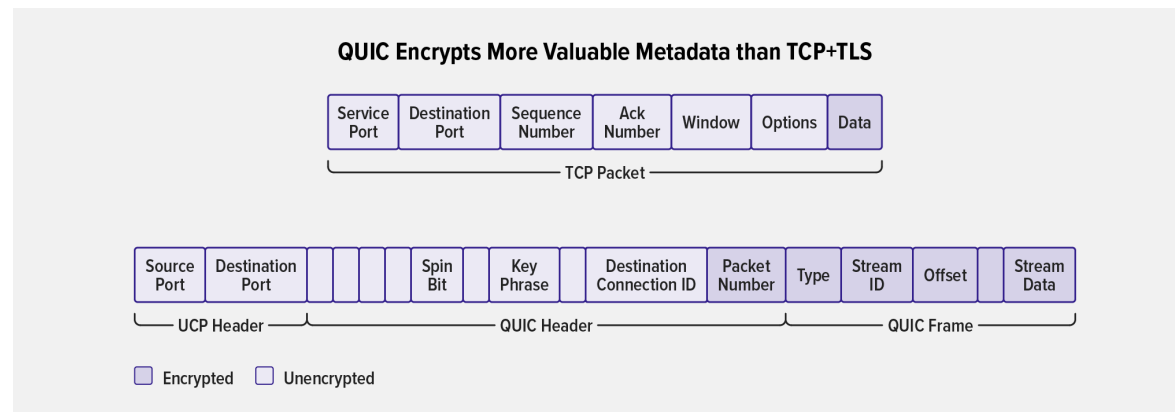
- handshaking – initial latency further reduced



# Transport Layer - QUIC

## QUIC - UDP-Based Multiplexed and Secure Transport

- each endpoint generates connection IDs (decoupling semantics from 5-tuple)
- stateful, connection-oriented protocol
- end-to-end transport protocol info is hidden from the network
- QUIC is not a replacement for TCP, app may use TCP if QUIC finds a fatal error
- better response time, security and privacy in detriment of interoperability



# Transport Layer - QUIC

## Security Features

- **Built-in encryption**
  - integrates TLS 1.3 natively, making encryption mandatory for all connections
  - **forward secrecy** and **0-RTT encryption**: on session/connections resumption
- **Reduced attack surface**
  - (almost) no cleartext handshaking
  - **encryption of metadata**: unlike TCP, where metadata (seq #s, flags, etc.) is exposed, QUIC encrypts its headers, offering greater privacy
- **Protecting against connection migration**
  - **connection migration**: ability to maintain a session in presence of network changes (e.g., switching from Wi-Fi to mobile data), using Connection IDs
  - stateless reset: “silent” mechanism to detect and handle connection failures, e.g. if server state is lost

# Transport Layer - QUIC

## Security Features (more...)

- **Resistance to replay attacks**
  - although QUIC supports 0-RTT (more vulnerable to replay attacks than 1-RTT) for faster session resumption, it restricts the type of data (e.g., only idempotent requests) that can be sent during the 0-RTT phase
- **Obfuscation of traffic characteristics**
  - encrypted packet numbers
  - minimal exposure of connection metadata impairs traffic analysis and MitM
- **Improved privacy**
  - reduced latency and fingerprinting: 0-RTT and reduced amount of unencrypted handshake information (harder for external entities to fingerprint users)
  - connection ID randomization and renegotiation (NEW\_CONNECTION\_ID frames)

# Transport Layer - QUIC

## Security Features (more...)

- DoS protection
  - **stateless server resumption**: support for stateless server resumption via token sent during the handshake, reducing server-side memory for connection state
  - **client-validated handshake**: prevention of amplification attacks during the handshake by requiring clients to validate their IP addresses before the server commits significant resources to the connection
- **Resistance to network surveillance and censorship**
  - encrypts headers and control information, making it harder to block or censor
  - obfuscation of protocol versions, impairing protocol fingerprinting

# Transport Layer Attacks and Mitigation

## QUIC

### 1. Amplification attacks

- an attacker may send small initial packets with spoofed source addresses (claiming to be the victim) to a QUIC server, which then responds with larger handshake packets
  - although QUIC includes anti-spoofing mechanisms like stateless resets and address validation, improperly configured servers or excessive retries can make this attack feasible
- Mitigation:
  - implementing robust address validation mechanisms (e.g., QUIC retry packets), stateless resets, rate limiting



# Transport Layer Attacks and Mitigation

## QUIC

### 2. Reflection attacks

- like amplification attacks, reflection attacks involve sending requests to a server with a spoofed source IP address (that of a victim). The server then reflects its response back to the victim, overwhelming them with traffic
  - although QUIC includes client address validation before sending larger responses, servers that improperly validate source addresses or allow excessive retries could be vulnerable
- Mitigation:
  - QUIC servers should ensure proper **client address validation**, **rate limit** large responses, and implement **packetretry tokens** during the handshake to prevent reflection

# Transport Layer Attacks and Mitigation

## QUIC

### 3. Connection hijacking

- the use of encryption and Connection IDs make it harder to hijack sessions compared to TCP, however, an attacker could still attempt to intercept Connection IDs and guess seq# (very unlikely) and inject malicious packets into an active session
- Mitigation:
  - very unlikely to occur; impossible to spoof traffic without encryption keys

# Transport Layer Attacks and Mitigation

## QUIC

### 4. Packet reordering and delay attacks

- attackers may reorder packets or introduce artificial delays to **degrade the performance** of a QUIC connection (timeout and retries)
- Mitigation:
  - QUIC is resilient to packet reordering, but detecting abnormal levels of reordering or delays, on stable networks, could help identify potential attacks

### 5. Side-channel attacks

- attackers could observe patterns in encrypted traffic, such as packet sizes, timing, or metadata, to infer information about the connection
- Mitigation:
  - padding QUIC packets for obfuscating traffic patterns, randomizing Connection IDs, use time-constant processing on specific QUIC packet fields

# Transport Layer Attacks and Mitigation

## QUIC

### 6. Downgrade attacks

- QUIC is designed to enforce TLS 1.3; there's a possibility of attempting a downgrade to an insecure version of TLS or a less secure cipher
- Mitigation:
  - QUIC is resistant to downgrade attacks; proper configuration of cipher suites and protocols on both client and server sides can further prevent downgrades

### 7. Resource exhaustion attacks

- an attacker could initiate many uncomplete handshake requests, causing the server to allocate resources without establishing actual connections (resource exhaustion)
- Mitigation:
  - stateless retries and client address validation prevent servers from being overwhelmed
  - servers limit resource allocation for incomplete handshakes

# Transport Layer Attacks and Mitigation

## QUIC

### 8. Replay attacks (0-RTT)

- an attacker could capture and replay a client's 0-RTT data to the server, potentially causing the server to process the same request multiple times
- Mitigation:
  - sensitive operations (e.g., transactions) **should not** be processed in 0-RTT mode. QUIC servers store session tickets to detect and prevent replayed data

### 9. Traffic correlation attacks

- an attacker monitoring network traffic might correlate connections based on Connection ID patterns, packet sizes, and timing, even if the content is encrypted, which may lead to deanonymization
- Mitigation:
  - Connection ID rotation and traffic obfuscation techniques reduces risks

# Transport Layer Attacks and Mitigation

## QUIC

### Real concern?

- Amplification and reflection attacks: **yes**, if misconfigured or improperly deployed
- Resource exhaustion (DoS) attacks: **yes**, especially for high-traffic services
- Replay attacks (in 0-RTT): **low-moderate**, especially for critical services or financial transactions
- Side-channel and traffic correlation attacks: **low**, especially for privacy-sensitive applications.
- Implementation bugs and misconfigurations: **yes**
- In summary:
  - if QUIC is **properly configured** and **regularly maintained**, the risk of attacks is very low compared to older protocols like TCP with TLS 1.2

# Transport Layer Attacks and Mitigation

## QUIC

Additional info:

- QUIC IETF WG
  - <https://datatracker.ietf.org/wg/quic/about/>
- Current implementations (2025):
  - <https://github.com/quicwg/quicwg.github.io/blob/main/implementations.md>
- Further reading:
  - Adam Langley, et al. “*The QUIC Transport Protocol: Design and Internet-Scale Deployment*”. In ACM SIGCOMM '17. ACM, New York, NY, USA, 183–196. DOI: <https://doi.org/10.1145/3098822.3098842>
  - IETF RFC 9000, May 2021, <https://datatracker.ietf.org/doc/html/rfc9000>