

Subtrees with small branching number

Pedro Marun

Carnegie Mellon University

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- ▶ A tree T is *finitely branching* if it is $<\aleph_0$ -branching. It is *infinitely branching* if $|I_T(x)| \geq \aleph_0$ for every $x \in T$.
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If b is a cofinal branch, then b is an uncountable 1-branching subtree.



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Definition (Nyikos)

Let T be a tree. The *fine wedge topology* on T is generated by all sets of the form $\uparrow x$ and their complements, where

$\uparrow x = \{y \in T : x \leq y\}$ and $x \in T$.

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If $X \subseteq T$, write $\uparrow X = \{y \in T : \exists x \in X (x \leq y)\}$.

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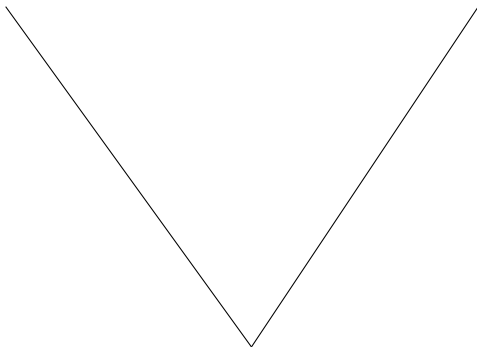
Let $x \in W = (\uparrow x_0) \cap \cdots \cap (\uparrow x_{n-1}) \cap (\uparrow y_0)^c \cap \cdots \cap (\uparrow y_{m-1})^c$.

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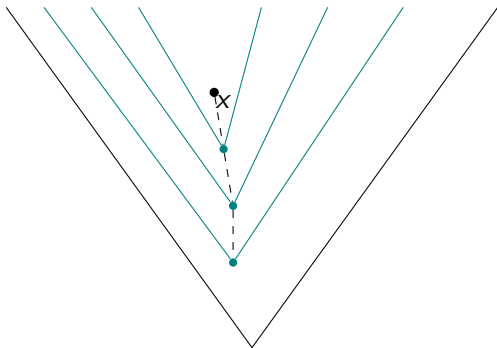
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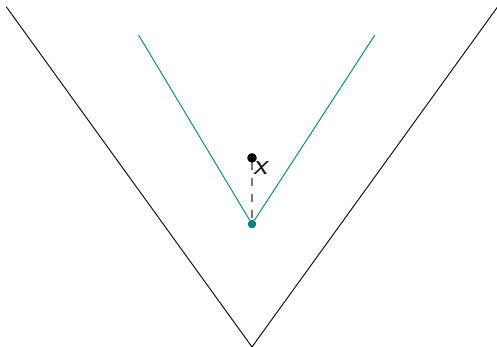


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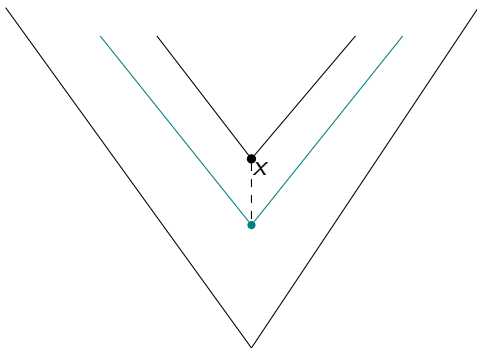


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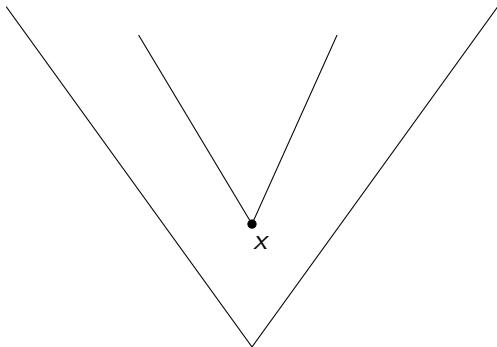


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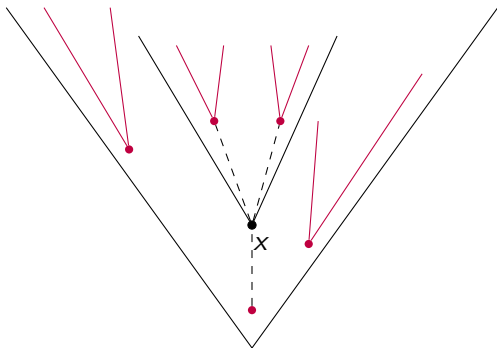


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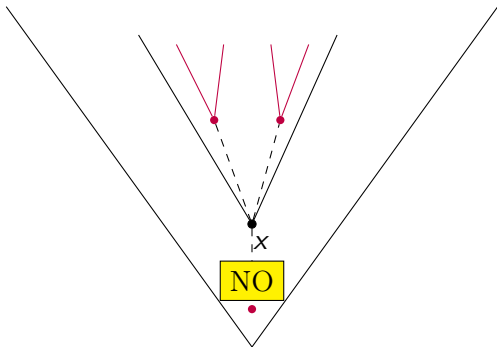


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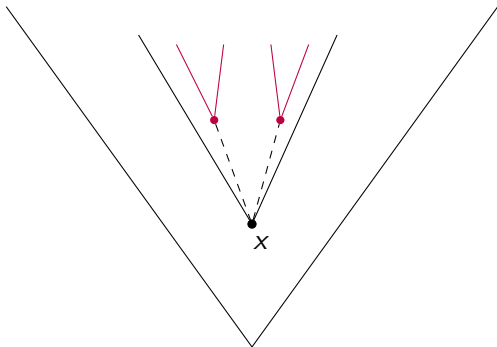


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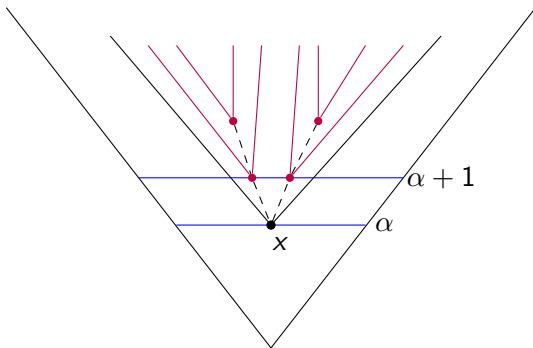


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3. The set of $\{ht(x) : x \text{ is safe}\}$ is uncountable.

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Let T be an infinitely branching \aleph_1 -tree. Then T is a Lindelöf tree if and only if it has the Lindelöf property with respect to the fine wedge topology.

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If $x \in S$ and $y < x$, then $x \in \uparrow f(y)$, so x is safe for S . So

$S \subseteq \{x \in T : x \text{ is safe}\}$, hence the latter is uncountable. □

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Lemma (folklore [2])

Let T be a Suslin tree. The poset \mathbb{P}_T has the ccc and is countably distributive. Moreover, if $D \subseteq \mathbb{P}_T$ is dense and open, then there exists $\alpha < \omega_1$ such that $T \restriction [\alpha, \omega_1) \subseteq D$.

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Let T be a Suslin tree in the universe V . If W is an outer model and $b \in W$ is a cofinal branch through T , then b is \mathbb{P}_T -generic over V .

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$$\{\text{Suslin}\} \subseteq \{\text{Lindelöf}\} \subseteq \{\text{Aronszajn}\}$$

A non-Lindelöf Aronszajn tree

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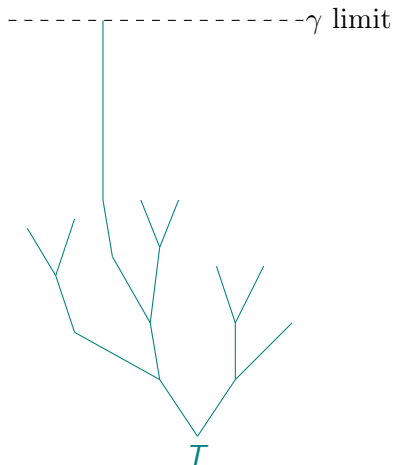
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Now recursively build, level by level, an infinitely branching tree U which has T as a subtree.

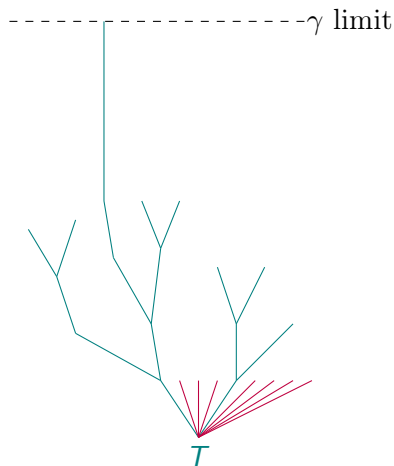


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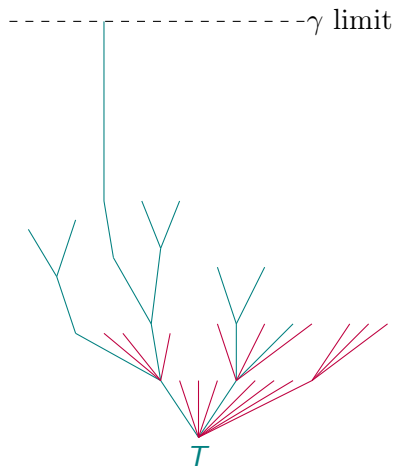


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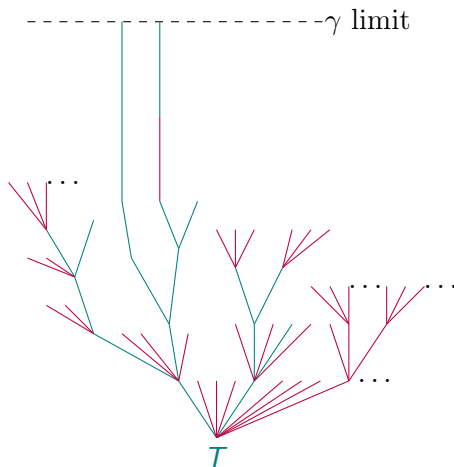


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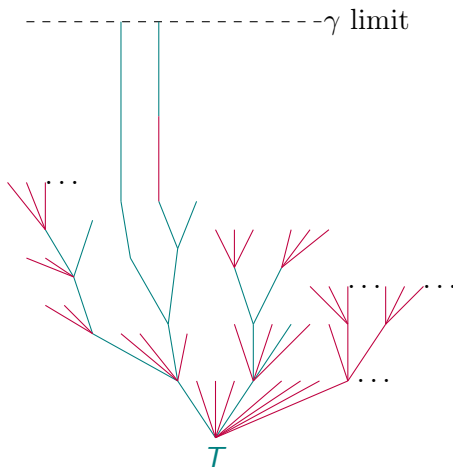


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$$\{\text{Suslin}\} \subseteq \{\text{Lindelöf}\} \subsetneq \{\text{Aronszajn}\}$$

(\diamond) A special Lindelöf tree

Let $\vec{f} = \langle f_\alpha : \alpha \in \text{lim}(\omega_1) \rangle$ be such $f_\alpha : \alpha \rightarrow [\alpha]^{<\omega}$ and \vec{f} guesses any $f : \omega_1 \rightarrow [\omega_1]^{<\omega}$ stationarily often.

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At stage $\alpha = \omega \cdot \alpha$, for each pair $(x, q) \in (T \restriction \alpha) \times \mathbb{Q}$ with $\varphi(x) < q$, choose a cofinal branch b through $T \restriction \alpha$ such that $x \in b$, $\sup(\varphi \restriction b) = q$ and **the unique point of b immediately above x is not in $f_\alpha(x)$** . Then put a new node $y \in T_\alpha$ above b and let $\varphi(y) = q$.

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Let $\vec{f} = \langle f_\alpha : \alpha \in \text{lim}(\omega_1) \rangle$ be such $f_\alpha : \alpha \rightarrow [\alpha]^{<\omega}$ and \vec{f} guesses any $f : \omega_1 \rightarrow [\omega_1]^{<\omega}$ stationarily often.

Build an infinitely branching tree T with underlying set ω_1 together with a specializing function $\varphi : T \rightarrow \mathbb{Q}$ by recursion on levels, maintaining that if $\varphi(x) < q$ then there is some $y > x$ with $\varphi(y) = q$.

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Therefore, T is Lindelöf and so, under \diamond ,

$$\{\text{Suslin}\} \subsetneq \{\text{Lindelöf}\} \subsetneq \{\text{Aronszajn}\}$$

Definition

Let T be a normal infinitely branching \aleph_1 -tree. Consider the following poset \mathbb{P} : conditions are functions $p \in \prod_{x \in F} [I(x)]^{<\omega}$, where $F \in [T]^{<\omega}$, such that:

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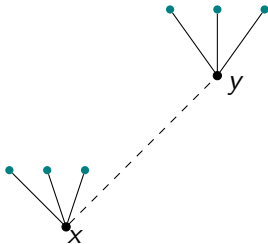
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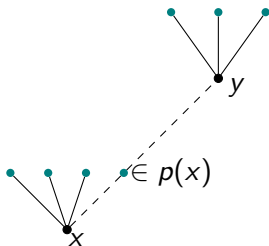


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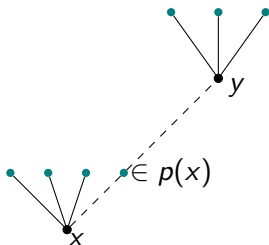


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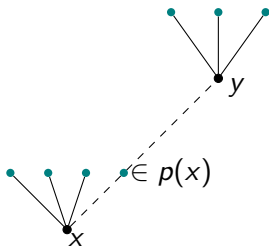
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If $p \in \mathbb{P}$ and $x \in \text{dom}(p)$, then p is a promise that $p(x) = I_{\dot{S}}(x)$.

Theorem (M.)

If T is a normal, infinitely branching Aronszajn tree, then \mathbb{P} has the ccc and $\Vdash_{\mathbb{P}} \dot{S}$ is a finitely branching normal subtree of T .

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Corollary

If MA_{\aleph_1} holds, then there are no Lindelöf trees.

Other ways of adding subtrees?

Theorem (M.)

Let T be an infinitely branching \aleph_1 tree. Suppose \mathbb{P} is a poset, G is \mathbb{P} -generic over V and $S \in V[G]$ is a finitely branching subtree of T .

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Let T be an infinitely branching \aleph_1 tree. Suppose \mathbb{P} is a poset, G is \mathbb{P} -generic over V and $S \in V[G]$ is a finitely branching subtree of T . Suppose that \mathbb{P} is either

- ▶ countably closed*
- ▶ strongly proper for a stationary set of countable elementary substructures of some (large) H_λ .*

Then $S \in V$.

Thank you :)

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