Seminar Report: Namy

Pedro Garcia Mota, Pau Espin Pedrol May 15, 2013

1 Introduction

Introduce in a couple of sentences the seminar and the main topic related to distributed systems it covers.

The aim of this seminar is to implement a distributed name server in erlang similar to DNS. Instead of addresses we will store process identifiers of hosts.

2 Work done

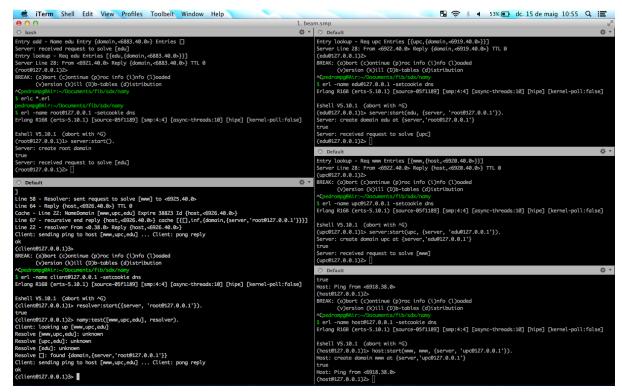
Present the new code that you have added or how you have implemented a required functionality by using small Erlang code snippets (you do not need to copy&paste all the code).

2.1 entry.erl

```
-module(entry).
-\mathbf{export}(\lceil \mathbf{lookup}/2, \ \mathrm{add}/3, \ \mathrm{remove}/2 \rceil).
lookup (Req, Entries) ->
      case lists: keyfind (Req, 1, Entries) of
      \{ -, \text{Entry} \} \rightarrow
         Entry;
      false ->
         unknown
   end.
add(Name, Entry, Entries) ->
   lists: keystore (Name, 1, Entries, {Name, Entry}).
remove (Name, Entries) ->
   lists: keydelete (Name, 1, Entries).
2.2
       cache.erl
-module(cache).
-\mathbf{export}(\lceil \operatorname{lookup}/2, \operatorname{new}/0, \operatorname{add}/4, \operatorname{remove}/2 \rceil).
```

```
lookup (Name, Cache) ->
    case lists: keyfind (Name, 1, Cache) of
     \{-, \text{ Expire }, \text{ Id}\} \rightarrow
         case time: valid (Expire, time:now()) of
         true ->
            {ok, Id};
          false ->
            invalid
       end;
     false ->
       unknown
  end.
\text{new}() \rightarrow
  [].
add(NameDomain, Expire, Reply, Cache) ->
  T = \{NameDomain, Expire, Reply\},\
  lists: keystore (NameDomain, 1, Cache, T).
remove (Name, Cache) ->
  lists: keydelete (Name, 1, Cache).
```

3 Experiments



Working version of the vanilla set up with TTL of 0 seconds.

4 Open questions

• i) In the vanilla set-up the time-to-live (TTL) is zero seconds. What happens if we extend this to two or four seconds? Try to quantify the amount of message traffic reduced.

If we enable the cache, changing the TTL value for a number greater than 0, all name resolutions are cached, so the amount of message traffic is reduced.

Let's do, for all tests, Np pings from host A to host B ('www').

Let's assume we have 1 Root nameserver node, an 'edu' nameserver node, a 'upc' nameserver node, which sums a total of up to *NameServers*=3 nameservers involved in the name lookup.

Name reslutions have to be redone every TTL seconds, thus the amount of name resolutions to be done for Np=1000 pings at PingPeriod=1 seconds per ping is:

$$\label{eq:totalTime} \begin{split} & \text{TotalTime} = \frac{Np}{PingPeriod} \\ & \text{NsResolutions} = \frac{TotalTime}{TTL} = \frac{Np}{PingPeriod} \end{split}$$

It can be easily probed that as NsResolutions increases, the amount of total messages increases by the amount of messages needed to resolve the name. In this case we have 3 nameservers which interacts with the host doing the ping, so each time we need to do the name resolution we have to request and get the response at each level, at a total cost of Nameservers*2 messages.

That makes Nmessages = (Np ·2) + (NsResolutions · NsServers · 2) = (Np · 2) +
$$(\frac{Np}{PingPeriod} \cdot NsServers \cdot 2)$$

Note: Np is multiplied by 2 because we count request and reply.

So, in our case, depending only on TTL: Nmessages = $(2000) + (\frac{6000}{TTL})$

Which for TTL=2 makes a total of 5000 messages, and for TTL=4 the amount is of 3500.

- ii) Extend TTL to a minute and analyze the impact (e.g. how many nodes need to know about the change? what happens with cached information?) of the following changes:
 - a) shutdown a host and start it up registered under the same name;

After shutting down the 'www' host and starting it up again with the same name the client does not reply any petitions. This happen because the cache entry for the 'www' host is not invalidated and the resolver is trying to contact a process that does not exist anymore.

```
(client@127.0.0.1)4> namy:test([www,upc,edu], resolver).
Client: looking up [www,upc,edu]
Resolve [www,upc,edu]: found {host,<6926.40.0>}
Client: sending ping to host [www,upc,edu] ... Client: no reply ok
```

− b) shutdown a host and start it up registered under a new name.

After shutting down the 'www' host and starting it up again with the name 'ftp' the resolver uses the cache to ask the [upc,edu] server for the new host, locates it and replies. After the first petition the complete the [ftp,upc,edu] entry is already cached.

```
(client@127.0.0.1)19> namy:test([ftp,upc,edu], resolver).
Client: looking up [ftp,upc,edu]
Resolve [ftp,upc,edu]: unknown
Resolve [upc,edu]: found {domain,<6925.40.0>}
Client: sending ping to host [ftp,upc,edu] ... Client: pong reply ok
(client@127.0.0.1)20> namy:test([ftp,upc,edu], resolver).
Client: looking up [ftp,upc,edu]
Resolve [ftp,upc,edu]: found {host,<6926.67.0>}
Client: sending ping to host [ftp,upc,edu] ... Client: pong reply ok
```

• iii) Our cache also suffers from old entries that are never removed. Invalid entries are removed and updated but if we never search for the entry we will not remove it. What can we do to solve this issue?

We could set a timer or a separate process on the resolver to check for invalid entries in the cache and remove them.

5 Personal Opinion

Your personal opinion of the seminar, including whether it should be included in next year's course or no.

The seminar should be included in the next course as it expands the concepts of name servers and gives the opportunity to implement a naming distributed system similar to DNS, having a better idea of how these systems are implemented.