Mark: 4/20 (total score: 5.8/29)



+100/1/48+

Computer Security Foundations 11/11/2022	Duration: 1H LEIC Midterm
28 227 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	be corrected automatically. Please follow these rules:
<ul> <li>Mark your answer using only blue a:</li> </ul>	and black pen. No pencil or light-coloured pens.
<ul> <li>Check only inside each box and be gen</li> </ul>	perous on ink. Erased boxes will not be detected automatically
<ul> <li>Only the boxes matter for the automa</li> </ul>	tic correction. You may underline text or take notes on the sid
	here are 30 questions in total, each with 4 options. Each can check one or two choices per question. The scoring for
$\bullet$ One checked correct answer (100%).	<ul> <li>Two checked answers, one correct (50%).</li> </ul>
• One checked incorrect answer (-20%)	<ul> <li>Two checked answers, none correct (-20%).</li> </ul>
<ul> <li>No checked answers (0%).</li> </ul>	$\bullet$ More than two checked answers (-20%).
_0	Code your 9-digit upYYYYXXXXX student number horizontally on the left, and replicate it below. Write also your first and last name below.  Student Number:  up 2 0 2 0 0 7 5 3 4  First and Last Name:  Pedro Macedo
Group 1 Introduction (4 Question 1.1 • Which of the following	questions) ng concepts is not associated with the risk management

0.5/1

Question 1.2 \ Which of the following is not a common reason for an attacker to compromise a server that answers to requests from a large number of clients/users?

0.5/1

To make the machine part of a botnet. Geo-political and strategic motivation. 🗶

Security proof.

Cost/benefit analysis. ✗



As part of a supply-chain attack.

To cause a data breach.

Threat analysis.

Mitigation. X

	Question 1.3 . Pick the incorrect statement or indicate that all are correct.
	A security goal is stated in terms of avoiding loss of value to an asset.
V/4	A security policy describes how security mechanisms are used to realize a security model.
)/1	X All are correct.
	A trust model describes which assumptions we can rely on to build security.
	Question 1.4 \$\&\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \
	Message that triggers an array access out of bounds.
)/1	Text filled into edit box that contains JavaScript.
)/ I	JavaScript performing heap spraying.
	➤ Overflow in an integer computation. ✓
	Group 2 Software Security (10 questions)
	Question 2.1 - One can bypass Address Space Layout Randomization (ASLR) protections by:
	The oddresses are always Error
	Jumping to system code using Return Oriented Programming.  Extracting addresses using other vulnerabilities.  Trial and error. X  When it is used ASLR. The ROP
0.2/1	Extracting addresses using other vulnerabilities.
	Trial and error. X
	NORG C + DE EXPEL - COL
	Question 2.2 & Certain format strings can make a program crash (with very high probability).  Which of the following printf commands is very likely to crash the program?
(2)	#DO ARTHUR FOR COLOR AND STUDIES AND A STUD
0.2/1	int n=0; printf("%f%d",n);
	Question 2.3 Recall the buffer overflow lab. You explored a buffer overflow vulnerability by
	creating a malicious input that overwrites the return address of the vulnerable function with an address, within your buffer where you stored shellcode to be executed. Which security technique had to be
14	disabled during compilation to allow you to overwrite the return address of the vulnerable function?  Data Execution Prevention  Stack Canary  Always verifies it  Value of the core  Address Space Layout Randomisation was everwrite  Ouestion 2.4. A capary built from string terminating bytes offers better security than a random
/1	None of the other choices Address Space Layout Randomisation X was overwriting
	Question 2.4 h Canaly built from String terminating Dyage office Security built a failten
	In my opinion, it a carry,
2	It is harder to guess than a random canary. has a sti terminating bytes,
/1 ( -)	They can be generated much more efficiently.
	They can be generated much more efficiently.  It makes it more difficult to overwrite the stack.  Random bytes can be all 0.  Random bytes can be all 0.  They can be generated much more efficiently.  Figure out prints ("/.s"), se read will stop in the canary, and overwriting the return address.
	Random bytes can be all 0.
	not overwriting

	f() 分子	100/3/46+
Question 2.5 $\clubsuit$ Recall what you have studied managed by a function $f$ that is called by a function correct for the positioning of the following data p of $g$ , (3) parameters passed by $g$ , (4) return a	about the classical oction $g$ . Indicate wieces: (1) local vari	
	High address  Low address	A B Power ters C D.local vor- of  for program analysis techniques
beyond safety. The C-like program on the right ill  Dynamic taint analysis.  Control-flow integrity.	printf untaintented char *fgets	char *fmt,);
Question 2.7  Consider the code below, when An attacker that controls the value of nresp can control of the value of nresp can control of the can lead to allocating too much memory.  None of the other choices.	Xnresp = packet_g if (nresp > 0) { response = xma for (i = 0; i < response [i] =	et_int(); ulloc(nresp*sizeof(char*));
Question 2.8 $\clubsuit$ Suppose you intend to us instructions of functions $f_1$ , $f_2$ , $f_3$ , $f_4$ , $f_5$ (in so the stack as illustrated below. Which of the followstrategy would still work without any changes to	ome order) and you wing functions could	take parameters and the ROP
<ul> <li>☐ All functions could take parameters.</li> <li>☐ Function f<sub>2</sub>.</li> <li>☐ It is not possible to pass any parameters.</li> <li>✗ Function f<sub>3</sub>.</li> </ul>	High ad	$egin{array}{c} \&f_1 \ \&f_5 \ \&f_4 \ \end{array}$ ,
Question 2.9 A Recall the buffer overflow lab variable receives an untested input. For the explois placed at the end of the input bytes, you should	itation strategy whe	ormed, where a buffer in a local ere the injected code (shellcode)
In the place of the vulnerable function's return place in the stack.  In the place of the vulnerable function's rehigher place in the stack.	eturn address, write	another address pointing to a
<ul><li>In the place of the vulnerable function's return to the other options.</li></ul>	irn address, write th	e address of the buffer. X

0.5/1

0.5/1

0/1

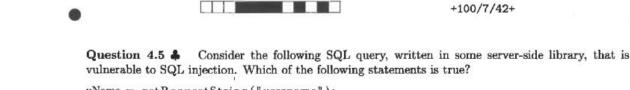
0.5/1



	Question 3.1  Virtual machines are commonly used as a security mechanism. Which of the following sentences is not true?
	An advantage of virtual machines is that software may be transparently executed as in a non-virtualised operating system.
0/1	An advantage of virtual machines is that acquiring root in the virtual machine does not grant root in a host operating system.
	A disadvantage of virtual machines is that malware may exploit the virtualisation layer to remain undetected.
	A disadvantage of virtual machines is that malware may detect that is is being virtualised.
	Question $3.2 \clubsuit$ Which of the above is <b>not</b> a systems security principle that we have studied in the classes?
0.54	
0.5/1	☐ Separation of privilege ★ Compromise recording
	Question 3.3  The Android operating system is a particular distribution built on top of Linux. Which is not an additional domain-specific restriction that Android implements?
	There is no root user, to prevent applications from escalating privileges.
0.5/1	It isolates different applications by registering each application with a different UNIX user.
0.5/1	It implements a form of Mandatory Access Control, so that no application can change its permissions.       ✓
	It enforces Manifest Permissions per application, to restrict its system capabilities.
	Question 3.4  Many modern operating systems require system support for an hardware component called a Trusted Platform Module (TPM). Which is the main rationale for that requirement?
	☐ To store user's cookies when browsing the web.≺
0/1	To guarantee that the user only installs official software.
	To protect the user login process.
	To prevent malicious bootloaders from compromising the operating system's boot process.
	Question 3.5 . Which fundamental mechanism does a UNIX operating system have in place to enforce isolation?
	System call interposition, by virtualising the address space of different user processes, monitoring all virtual address translations.
0.5/1	Software fault isolation, by virtualising the address space of different user processes, monitoring all virtual address translations.
	Software fault isolation, by ensuring that the control-flow of user processes never accesses invalid virtual memory regions.
	System call interposition, by preventing different user processes to communicate via system calls.

		Question 3.6 ♣ The UNIX filesystem can be seen as an instance of:	
	1	Capability lists, since each file enumerates the permissions for all the users who may read, write or execute it.	
0.2/1	/3	Role-based access control, since each file has an access control list for three fixed roles (owner,group,others) and the association between users and groups may be modified independently.	
		Attribute-based access control, since users may change their group membership without the need to change file permissions. ⋈	
		$\square$ Access control lists, since only the owner of each file (or root) may read, write or execute it. $\times$	
		Question 3.7 . For efficiency reasons, the address space of the Linux kernel is not completely independent from that of user processes due to kernel mapping. Which of the following is correct?	
		Part of each user process's memory is mapped into the kernel memory space, to reduce the size of the address translation tables.	
)/1		Part of the kernel memory space is mapped into the memory space of each user process, to reduce the number of cache misses.	
		Since the discovery of speculative execution attacks such as Meltdown and Spectre, Linux no longer uses kernel mapping, as it violates the principle of separation of privilege.	
		Part of the kernel memory space is mapped into the memory space of each user process, to speed up system calls. X	
		Question 3.8 • In the environment variable and setuid program lab we have experimented with environment variables and setuid programs. Which of the following sentences is not true?	
		☐ If the shell detects that it is being run under a setuid process, it may drop its privilege.	
)/1		When a process forks a child process, it passes on its environment, excluding some critical variables if the fork is a system call.	
		When the shell forks a child process, it passes on its environment, excluding some critical variables if it has a different effective user id.	
	The system function allows calling a shell function within a program with the program's environment.		
		Question 3.9 & Remember the environment variable and setuid program lab. The Linux file system associates a owner, a group, and 12 permission bits with each entry. For a file with the 12th bit (the Set-UID bit) active, it can be executed:	
(1 (7)	Only by a user belonging to the group.		
	Only by users whose set (owner, group, others) in the permission bits, has the executable bit active.		
	0	Only by the owner.	××
		Only by the owner or by a user belonging to the group.	lher

Question 3.10 . Under Linux, Docker crucially relies on seccomp-bpf to confine containers. Select the incorrect choice or mark that all choices are correct. If an attacker has access to certain system calls inside a container, it can exploit Docker to acquire root in the host OS. 🕝 Users can manually configure seccomp-bpf filters to block various system calls inside the con--0.2/1tainer. All choices are correct. Using the default Docker seccomp-bpf filters, an attacker that acquires root in the container does not directly acquire root in the host OS. Group 4 Web Security (6 questions) Question 4.1 . Cookies may be used for various purposes. Name the incorrect one below. Encrypting user communication. Managing user sessions. -0.2/1Personalising user experience. Tracking user activity. Question 4.2 & Which sentence best characterises a Cross-Site Request Forgery (CSRF) attack? "When a malicious origin can send requests to a server in another origin ...": ... where the user is already logged in. ... and trigger a side-effect on the server. 0/1... and measure side-channels such as re-... and read its response. sponse time. X Question 4.3 A classical protection against Cross-Site Scripting (XSS) attacks is for the site to adopt a Content Security Policy (CSP). Which of the following statements is true? Stored XSS attacks cannot be prevented with a CSP because the malicious payloads are already stored in the server. Preventing reflected XSS attacks requires both CSP and SRI (Subresource Integrity). 0/1XSS attacks using inline scripts are blocked by a default CSP. DOM-based XSS attacks cannot be prevented with a CSP because they occur on the server-side when processing a user request. Question 4.4 & Which assignment can be seen as a valid analogy in the table below? Systems Security Web Security (1)=Pages,(2)=HTML,(3)=JavaScript,(4)=Popups Processes (1)=DOM,(2)=Cookies,(3)=iFrames,(4)=Fetch ② ③ Files -0.2/1X ①=Pages,②=Cookies,③=HTTP,④=Frames Sockets Sub-processes (1)=Frames,(2)=DOM,(3)=Images,(4)=Sub-frames



vulnerable to SQL injection. Which of the following statements is true? uName = getRequestString("username"); uPass = getRequestString("userpassword"); sql = 'SELECT\_\*\_FROM\_Users\_WHERE\_Name\_="' + uName + '"\_AND\_Pass\_="' + uPass + '"' A SQL injection attack would not be possible if the clauses for the Name and Pass fields did not enclose strings with double quotes ("). A SQL injection attack may be able to delete the USERS table. A SQL injection attack will only be able to read existing data from the USERS table. A SQL injection attack that bypasses authentication is only possible because the clause for the Name field appears before the clause for the Pass field. X Question 4.6 According the Same-Origin Policy (SOP), which of the following is allowed? HTML code in a page from origin A can send a POST request to a page from origin B. HTML code in a page from origin A can send and read the response of a GET request to a page from origin B.

JavaScript code in a frame from origin A can exchange data with a frame from origin B. JavaScript code in a page from origin A can inspect the HTML code of a frame from origin B.

0/1

0/1