Irrationality of $\sqrt{2}$

Is the length of the diagonal whose side is 1 expressible as a ratio of two integers? [1], p. 6

This side's length is $\sqrt{2}$ because of the Pythagorean theorem:

$$a^2 = 1^2 + 2^2$$

In other words, do there exist two integers a,b with no divisor in common such that $\left(a/b\right)^2=2$

"No divisor in common": an integer would defeat the thesis of irrationality as well, but we're looking specifically for a rational.

Let's try to give geometric interpretations of the square of $\frac{a}{b}$ equaling a number.

$$\left(rac{a}{b}
ight)^2 = 2 \Rightarrow \ rac{a^2}{b^2} = 2 \Rightarrow \ a^2 - 2b^2$$

 $2b^2$ is even, because the definition of evenness is divisibility by 2, or equivalently, being a multiple of 2, which $2b^2$ clearly is.

The square of an odd integer being odd necessitates defining "squaring" geometrically.

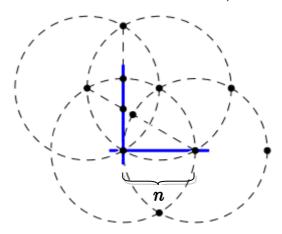
Squaring is multiplying a number by itself. Geometrically, translating a number into a new dimension creates a new number which is one higher power of the number.

- Opening a height of 3 from a width of 3 produces an **area** of 3×3 , from power 1 to 2.
- Opening a depth of 3 from a height and width both of 3 produces a **volume** of $3\times 3\times 3$, or power from 2 to 3.

We are concerned with the power 2.

But how do we open such new dimension for a number n geometrically?

- 1. Given a segment of length n;
- 2. Draw a perpendicular line, that is, forming a right angle, from one of its end points;
- 3. Intersect the perpendicular line with the arc from the other end point.



Construction of a square of a number n

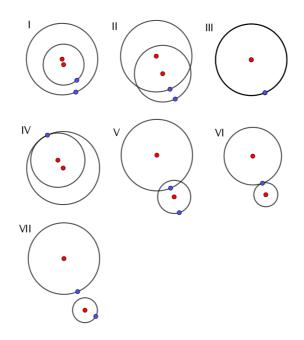
 $Procedural\ construction\ of\ this\ scene\ is\ done\ in\ {\tt Irrationality_of_Square_Root_of_Two.nb}\ .$

Object oriented construction is done in the package pslab.geom in PSLab.

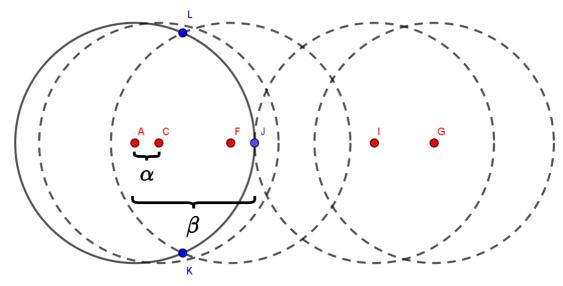
Intersection of circles A and B?

Given two circles with specified center and radius, what is their intersection?

Cases for intersection



Same-radius circles



Where:

- $oldsymbol{lpha} \in \mathbb{R}$ is the distance between the centers
- $\,eta>0\in\mathbb{R}$ is the radius
- $\alpha, \beta \in \mathbb{R}$,

Then:

- If $\alpha=0$ (concentric), the circles' intersection is the circle (they're the same) (case III)
- If $\alpha=2\,\beta$ (for example |AI|), the circles' intersection is one point (in this case, J) (case VI)
- If $0<\alpha<2\,\beta$ (for example |AF|), the circles' intersection is two points (in this case, $\{L,K\}$) (case II)
- If $\alpha>2$ β (for example |AG|), the circles' intersection is null (case VII)
- · In this scenario, once circle can't contain another

Distinct radius circles

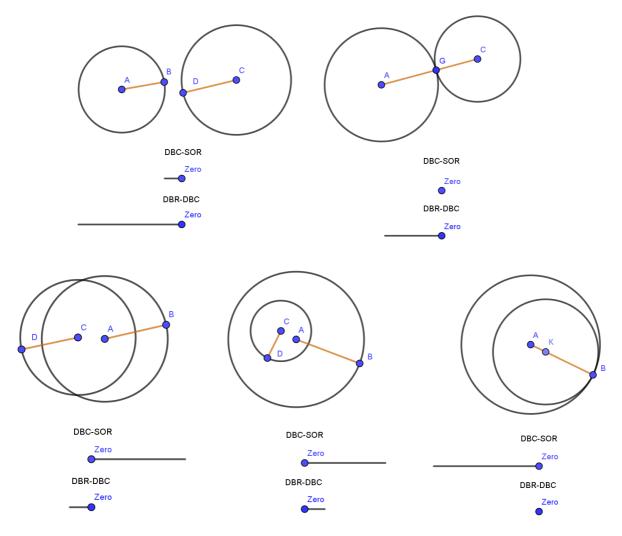
Let's define a circle from two points, for center and radius point, plus a measure α of revolution of r around c, measured from 0 (where the circle reduces to a segment) to 1 (one full revolution). So, a record $\{c, r, \alpha\}$, where:

- c the center point
- ullet radius point
- lpha the revolution, from 0 to 1

How to decide intersection of two circles in this form?

Let's define the following variables and map out the same cases.

- DBC: the distance between centers
- DBR: the difference between radiuses (subtraction in absolute)
- SOR: the sum of radiuses



Each case for the values DBC-SOR, DBR-DBC above is exclusive. This defines the following table.

	$DBC - SOR < 0 \Leftrightarrow DBC < SOR$	$DBC - SOR = 0 \Leftrightarrow DBC = SOR$	$DBC - SOR > 0 \Leftrightarrow DBC > SOR$
$\mathrm{DBR} - \mathrm{DBC} < 0 \Leftrightarrow \mathrm{DBR} < \mathrm{DBC}$	$\begin{array}{l} DBR < DBC < SOR \\ \text{Circle in proper complement; null} \\ \text{intersection} \end{array}$	$\begin{array}{l} DBR < DBC \land DBR < SOR \\ \text{Circle in improper complement; point} \\ \text{intersection} \end{array}$	$\begin{array}{l} DBR < DBC \wedge SOR < DBC \\ \text{Not contained or in complement; pair of} \\ \text{points intersection} \end{array}$
$DBR - DBC = 0 \Leftrightarrow DBR = DBC$	$\begin{array}{l} DBR < SOR \land DBC < SOR \\ \text{Circle improperly contained; point} \\ \text{intersection} \end{array}$	DBR = DBC = SOR	
$DBR - DBC > 0 \Leftrightarrow DBR > DBC$			$\begin{aligned} SOR &< DBC < DBR \\ \text{Circle properly contained; null intersection} \end{aligned}$

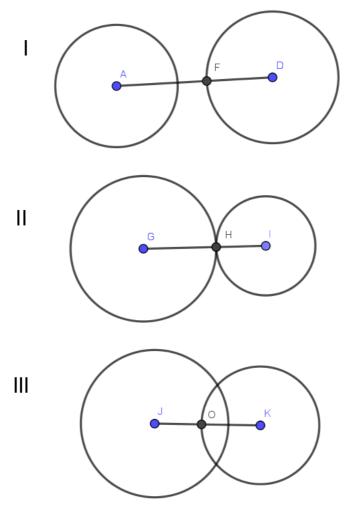
The missing cases didn't appear in our visual analysis. Let's analyze in more detail each of our variables and their relations.

- DBC
 - If DBC is zero, the circles are concentric
 - o Otherwise, DBC must be positive
- SOR
 - If SOR is zero, both circles are points; this is not permitted
 - SOR is positive
- DBR
 - DBR is taken to be in absolute; so it either is zero or positive
 - If DBR is zero, the circles' radiuses are equal
- DBC and SOR
 - If the **DBC** equals **SOR**, we know the circles intersect in one point
 - If DBC is larger than SOR, there is a gap between the circles, or the circles are in each other's complement
 - If **DBC** is smaller than **SOR**, there is a negative gap between the circles, meaning the circles are **not** in each other's complements.
- DBR and DBC
 - \circ The relationship between **DBR** and **DBC** seems (if any) unclear. We can't be positive DBR DBC is relevant to the result, or that it is not a coincidence.

Geometrically

The above inferences use numbers and the order relation. What can be observed geometrically?

- If two circles have one point of intersection, then their centers must be collinear with the point of intersection.
 - Collinear centers and intersection algorithm
 - 1. Draw a segment between the two circles' centers
 - 2. Intersect one of the circles with this segment
 - 3. If the point of intersection is in the complement of the other circle, then the intersection is null (case I below)
 - 4. If the point of intersection is in the other circle itself, then the intersection is that point (case II below)
 - 5. Else (the point is in the interior of the circle), the intersection is a pair of points (case III below)

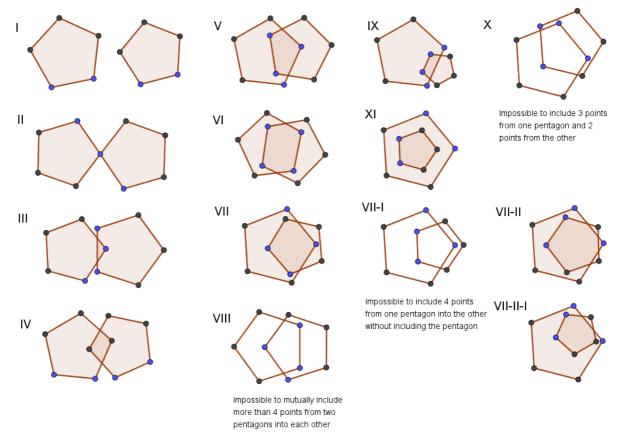


This algorithm necessitates identifying a point being on a circle vs. it being on the interior of it.

We want to avoid using numbers, so we must avoid using arbitrary Euclidean distances. We could try defining a point as being in the **interior** of a circle as that which is in a circle the intersection of which with the first circle results in a segment, instead of a point.

This way, we define interactions between objects as generating new objects, and evaluate the generated objects.

Let's evaluate the same kind of interaction definition for another case of regular polygon, the pentagon.



There is an algorithm being followed in this construction. The algorithm is that with each step, we include one more point from one polygon into the other, symmetrically, that is, alternating from which polygon is the point included (ignoring step II for now).

Some observations.

- The polygon formed by the inclusion has n+2 sides, where n is how many points are included in the intersection (from whichever pentagon).
 - VII-II is an exception: 3 points in the intersection, 7 sides
 - VII-II is different from V in that there are 3 points from one pentagon and 0 from the other, while V has 2 from one and 1 from the other; that is, the difference between number of points included from each pentagon, or inclusion asymmetry (what's both subtracted from one and added to the other), is 1 greater
 - ullet This ((2,1) o (3,0)) is the first increase in inclusion asymmetry that's meaningful, where meaningful means the resulting intersection polygon has a different number of sides

Let's list the inclusion asymmetry chains for pentagons.

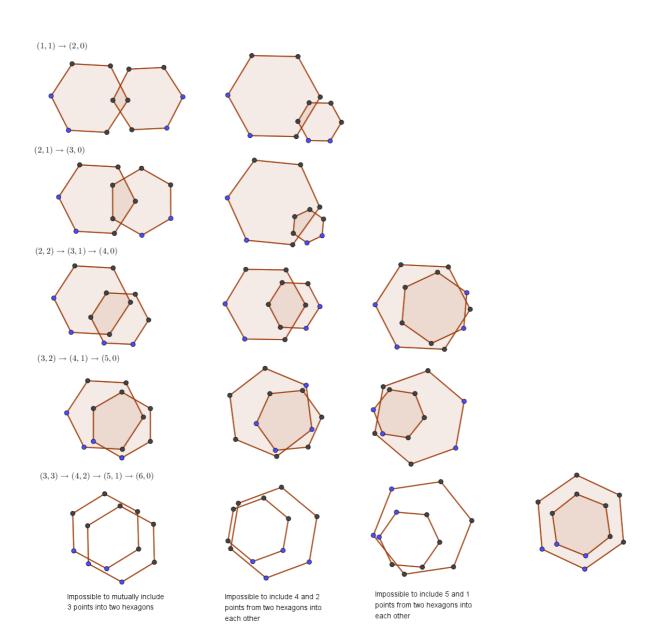
- $(1,1) \to (2,0)$: $4 \to 4$ sides
- ullet $(2,1)
 ightarrow (3,0) \colon 5
 ightarrow 7$ sides
- (2,2)
 ightarrow (3,1)
 ightarrow (4,0)
 vert 6
 ightarrow 6 sides
- (5,0): 5 sides (proper inclusion)

The max inclusion asymmetry chain for a polygon of 5 sides is the chain with sum of included sides 4.

The (1,0) inclusion, obviously symmetric, has always 3 sides.

Let's look at polygons with more sides.

Hexagon



- · Chains from pentagon
 - \circ $(1,1) \rightarrow (2,0)$: $4 \rightarrow 4$ sides
 - ullet $(2,1)
 ightarrow (3,0) \colon 5
 ightarrow 5$ sides
 - ullet $(2,2)
 ightarrow (3,1)
 ightarrow (4,0) \colon 6
 ightarrow 6
 ightarrow 8$ sides
- (3,2)
 ightarrow (4,1)
 ightarrow (5,0): 7
 ightarrow 7
 ightarrow 7 sides
- (6,0): 6 sides (proper inclusion)

Once again, the max (improper) chain is of size sides - 1.

Polygons with more sides

- (4,3) o (5,2) o (6,1) o (7,0): ?
- (4,4) o (5,3) o (6,2) o (7,1) o (8,0):
 (5,4) o (6,3) o (7,2) o (8,1) o (9,0):

References

1. Mathematics and Logic, Mark Cac and Stanislaw Ulam, Dover Publications, 1968