

Dicionários / Tabelas de Dispersão II

22/11/2023

Ficheiro ZIP

- Está disponível no Moodle um **ficheiro ZIP** de suporte aos tópicos de hoje
- O tipo abstrato **Hash Table** usando **Separate Chaining**
- **Versão “simples”**, que permite trabalho autónomo de desenvolvimento e teste


Sumário

- Recap
- Hash Tables – Representação usando **Separate Chaining**
- Análise detalhada do TAD Hash Table – Separate Chaining
- Desempenho computacional

Let's
RECAP

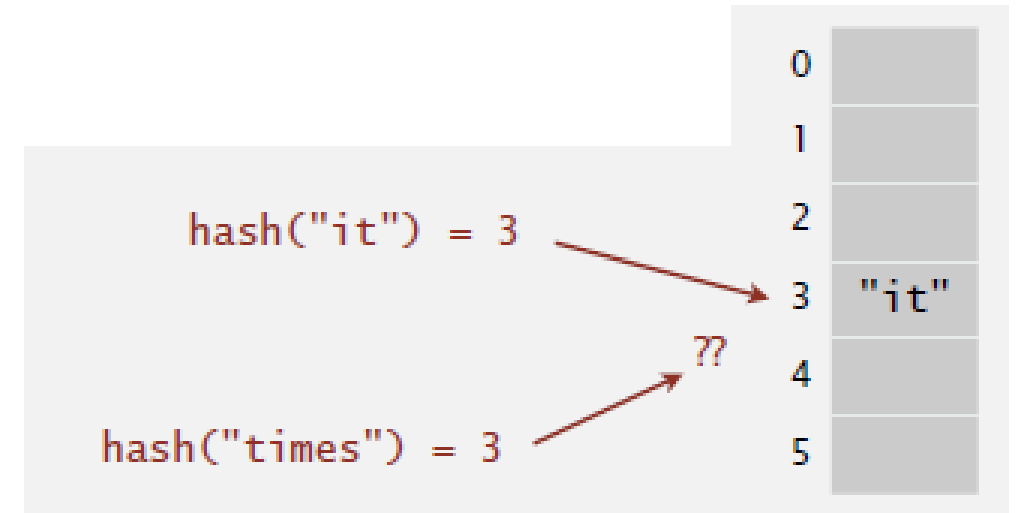
Recapitulação

Tabelas de Dispersão – Hash Tables

- Estrutura de dados para armazenar pares (**chave**, **valor**)
- Sem **chaves duplicadas**
- **Sem limite** de tamanho
- **Sem** uma **ordem** implícita
- MAS, com **acesso rápido !!** 

Hash Tables – Open Addressing

- Armazenar itens numa tabela/array indexada pela chave
 - Índice é função da chave
- Função de **Hashing**: para calcular o **índice** a partir da **chave**
 - Rapidez !!
- **Colisão**: 2 **chaves diferentes** originam o **mesmo índice** da tabela




[Sedgewick & Wayne]

Inserir na tabela – Linear Probing

- Guardar na **posição i**, se estiver disponível
- Caso contrário, tentar $(i + 1) \% M$, $(i + 2) \% M$, etc.
- Inserir **L** -> índice = 6
- **Colisão !!**
- ...

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
st[]	P	M			A	C	S	H			E				R	X
M = 16	L															

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
st[]	P	M			A	C	S	H	L		E				R	X
M = 16																



[Sedgewick & Wayne]

Resolução de colisões – Linear Probing

- Aceder à **posição i**
- Se necessário, tentar em $(i + 1) \% M$, $(i + 2) \% M$, etc.

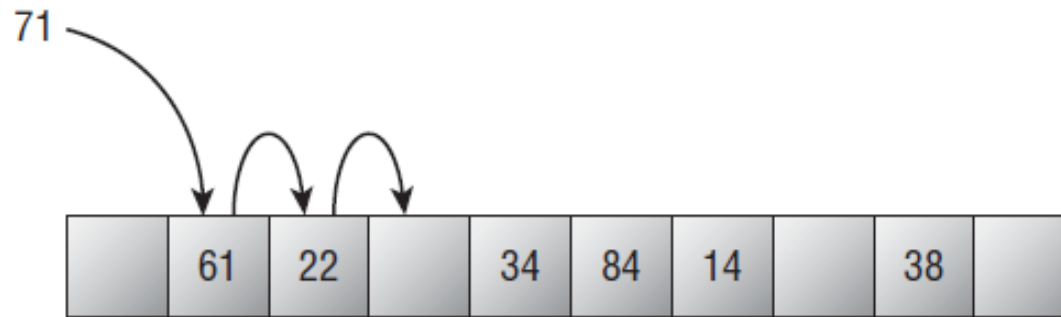


Figure 8-2: In linear probing, the algorithm adds a constant amount to locations to produce a probe sequence.

[Stephens]

Resolução de colisões – Quadratic Probing

- Aceder à **posição i**
- Se necessário, tentar em $(i + 1) \% M$, $(i + 4) \% M$, $(i + 9) \% M$, etc.

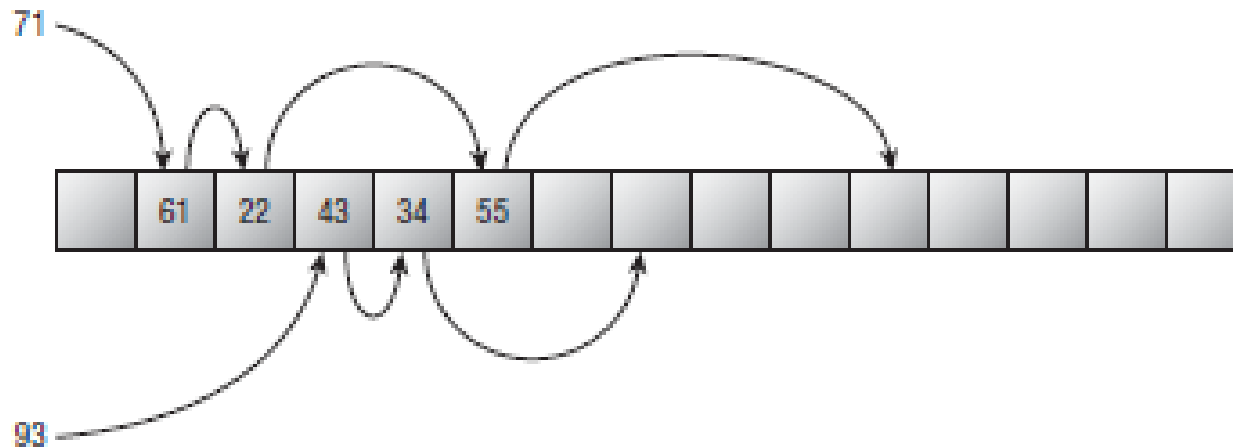


Figure 8-4: Quadratic probing reduces primary clustering.

[Stephens]

Análise – Linear Probing – Knuth, 1963

- Fator de carga – Load Factor

$$\lambda = N / M$$



- Nº médio de tentativas para encontrar um item

$$1/2 \times (1 + 1/(1 - \lambda))$$

-> 1.5, se $\lambda = 50\%$

-> 3, se $\lambda = 80\%$

- Nº médio de tentativas para inserir um item ou concluir que não existe

$$1/2 \times (1 + 1/(1 - \lambda)^2)$$

-> 2.5, se $\lambda = 50\%$

-> 13, se $\lambda = 80\%$



Resizing + Rehashing

- **Objetivo** : fator de carga $< 1/2$
- **Duplicar o tamanho** do array quando fator de carga $\geq 1/2$
- **Reduzir para metade** o tamanho do array quando fator de carga $\leq 1/8$
- Criar a nova tabela e **adicionar**, um a um, todos os **itens**

before resizing

	0	1	2	3	4	5	6	7
keys[]		E	S			R	A	
vals[]		1	0			3	2	

after resizing

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
keys[]					A		S				E				R	
vals[]					2		0				1				3	

[Sedgewick & Wayne]

Exemplo

- Hash Table (String, String)

Exemplo – $M = 17$ – $N = 12$

```
size = 17 | Used = 12 | Active = 12
0 - Free = 0 - Deleted = 0 - Hash = 68, 1st index = 0, (December, The last month of the year)
1 - Free = 1 - Deleted = 0 -
2 - Free = 0 - Deleted = 0 - Hash = 70, 1st index = 2, (February, The second month of the year)
3 - Free = 1 - Deleted = 0 -
4 - Free = 1 - Deleted = 0 -
5 - Free = 1 - Deleted = 0 -
6 - Free = 0 - Deleted = 0 - Hash = 74, 1st index = 6, (January, 1st month of the year)
7 - Free = 0 - Deleted = 0 - Hash = 74, 1st index = 6, (June, 6th month)
8 - Free = 0 - Deleted = 0 - Hash = 74, 1st index = 6, (July, 7th month)
9 - Free = 0 - Deleted = 0 - Hash = 77, 1st index = 9, (March, 3rd month)
10 - Free = 0 - Deleted = 0 - Hash = 77, 1st index = 9, (May, 5th month)
11 - Free = 0 - Deleted = 0 - Hash = 79, 1st index = 11, (October, 10th month)
12 - Free = 0 - Deleted = 0 - Hash = 78, 1st index = 10, (November, Almost at the end of the year)
13 - Free = 1 - Deleted = 0 -
14 - Free = 0 - Deleted = 0 - Hash = 65, 1st index = 14, (April, 4th month)
15 - Free = 0 - Deleted = 0 - Hash = 65, 1st index = 14, (August, 8th month)
16 - Free = 0 - Deleted = 0 - Hash = 83, 1st index = 15, (September, 9th month)
```

TAD Hash Table

```
HashTable* HashTableCreate(unsigned int capacity, hashFunction hashF,  
| | | | | | | probeFunction probeF, unsigned int resizeIsEnabled);
```

```
void HashTableDestroy(HashTable** p);
```

```
int HashTableContains(const HashTable* hashT, const char* key);
```

```
char* HashTableGet(HashTable* hashT, const char* key);
```

```
int HashTablePut(HashTable* hashT, const char* key, const char* value);
```


```
int HashTableReplace(const HashTable* hashT, const char* key,  
| | | | | | | const char* value);
```

```
int HashTableRemove(HashTable* hashT, const char* key);
```




Hash Table Header + Hash Table Bin

```
struct _HashTableHeader {  
    unsigned int size;  
    unsigned int numActive;  
    unsigned int numUsed;  
    hashFunction hashF;  
    probeFunction probeF;  
    unsigned int resizeIsEnabled;  
    struct _HashTableBin* table;  
};
```

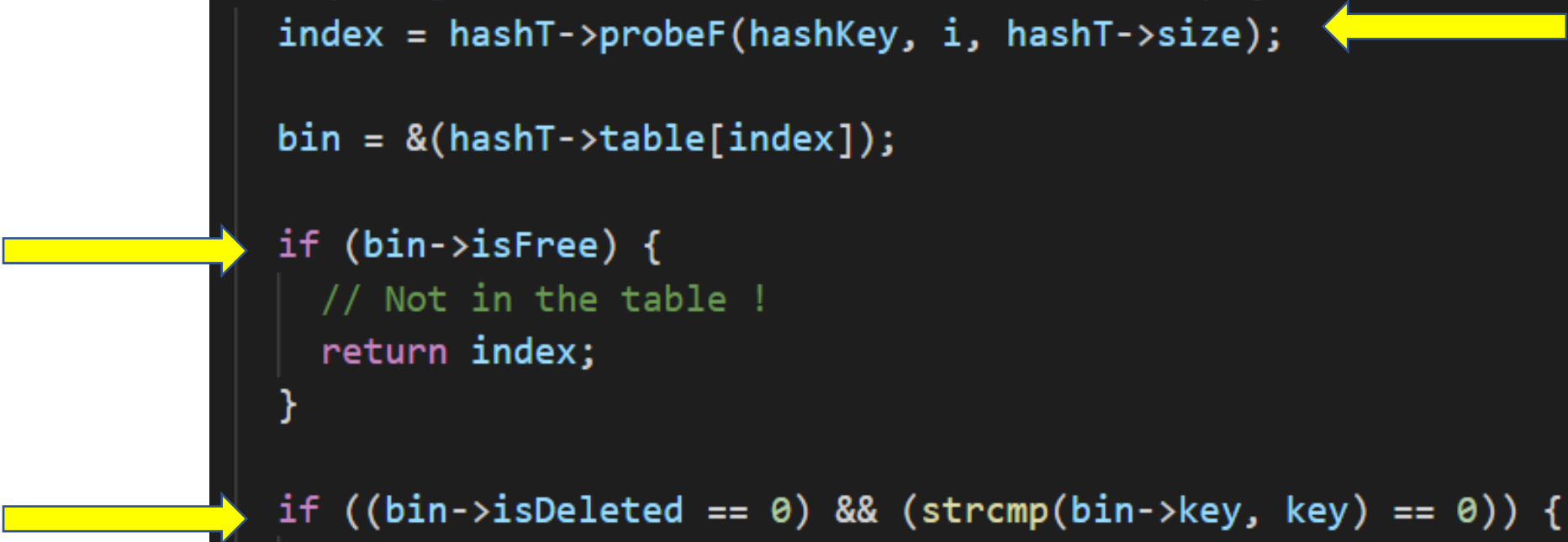


```
struct _HashTableBin {  
    char* key;  
    char* value;  
    unsigned int isDeleted;  
    unsigned int isFree;  
};
```



Procura de uma chave

```
for (unsigned int i = 0; i < hashT->size; i++) {  
    index = hashT->probeF(hashKey, i, hashT->size);  
  
    bin = &(hashT->table[index]);  
  
    if (bin->isFree) {  
        // Not in the table !  
        return index;  
    }  
  
    if ((bin->isDeleted == 0) && (strcmp(bin->key, key) == 0)) {  
        // Found it !  
        return index;  
    }  
}
```



HashTableGet

```
char* HashTableGet(HashTable* hashT, const char* key) {  
    int index = _searchHashTable(hashT, key);  
    if (index == -1 || hashT->table[index].isFree == 1) {  
        // NOT FOUND  
        return NULL;  
    }  
  
    struct _HashTableBin* bin = &(hashT->table[index]);  
    char* result = (char*)malloc(sizeof(char) * (1 + strlen(bin->value)));  
    strcpy(result, bin->value);  
  
    return result;  
}
```



Tarefas

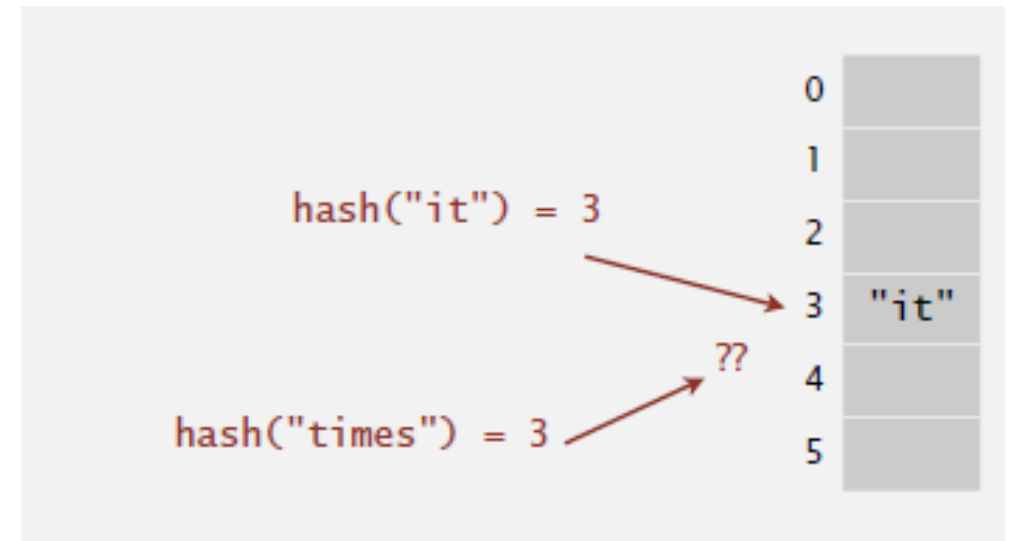
- Analisar as funções desenvolvidas
- E o simples programa de teste

Hash Tables

– Separate Chaining

Colisões – Como proceder ?

- Duas **chaves distintas** são mapeadas no **mesmo índice** da tabela !!
- Como gerir de modo eficiente ?
- Sem usar “demasiada” memória !!
- **Alternativa** ao Open Addressing ?

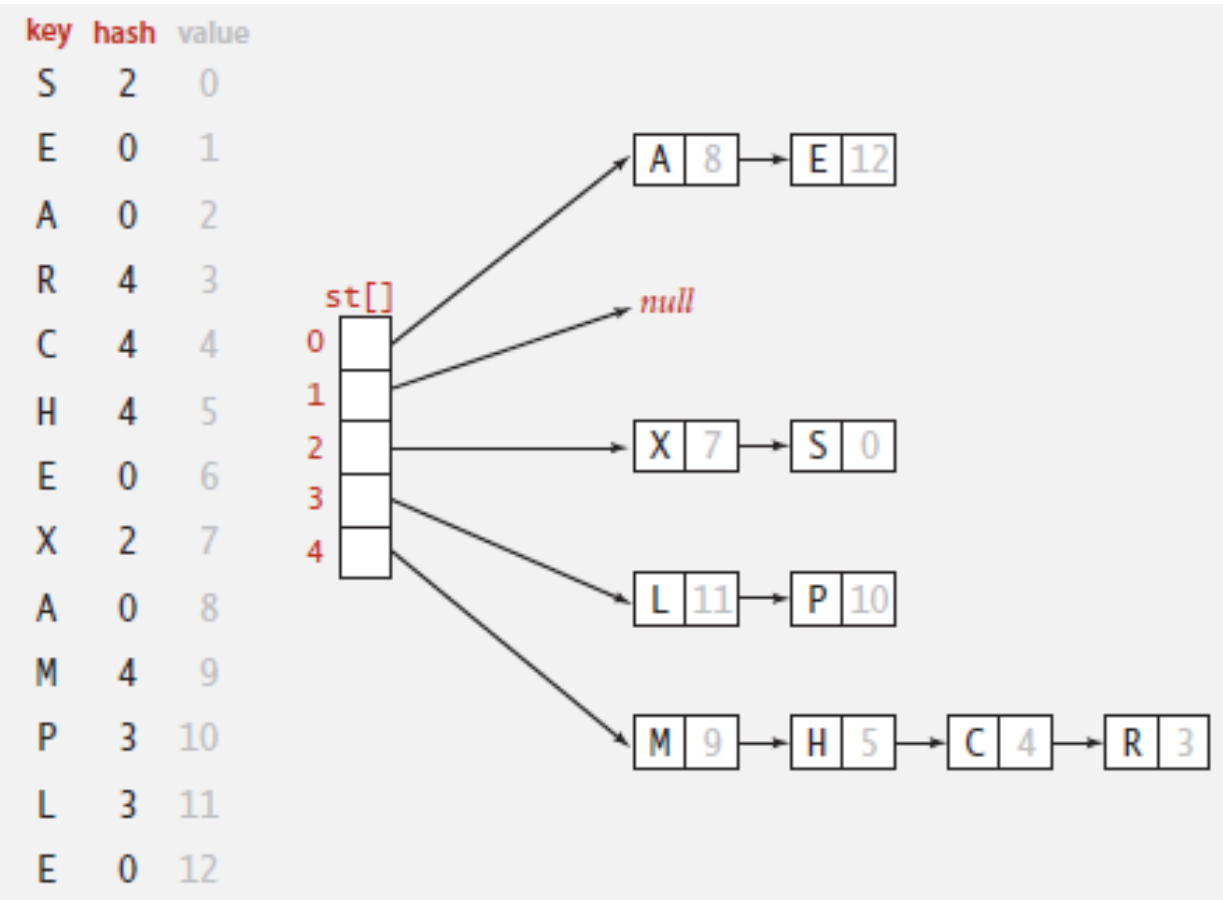


[Sedgewick & Wayne]

Separate Chaining (IBM, 1953)


- Array de **M ponteiros** ($M < N$)
- Mapear a chave em $[0..(M - 1)]$
- Inserir no **início** de uma lista, se não existir
- Procurar **apenas numa lista**

[Sedgewick & Wayne]




Hash Table Header + Hash Table Bin

```
struct _HashTableHeader {  
    unsigned int size;  
    unsigned int numBins;  
    hashFunction hashF;  
    List** table;  
};
```



```
struct _HashTableBin {  
    char* key;  
    char* value;  
};
```

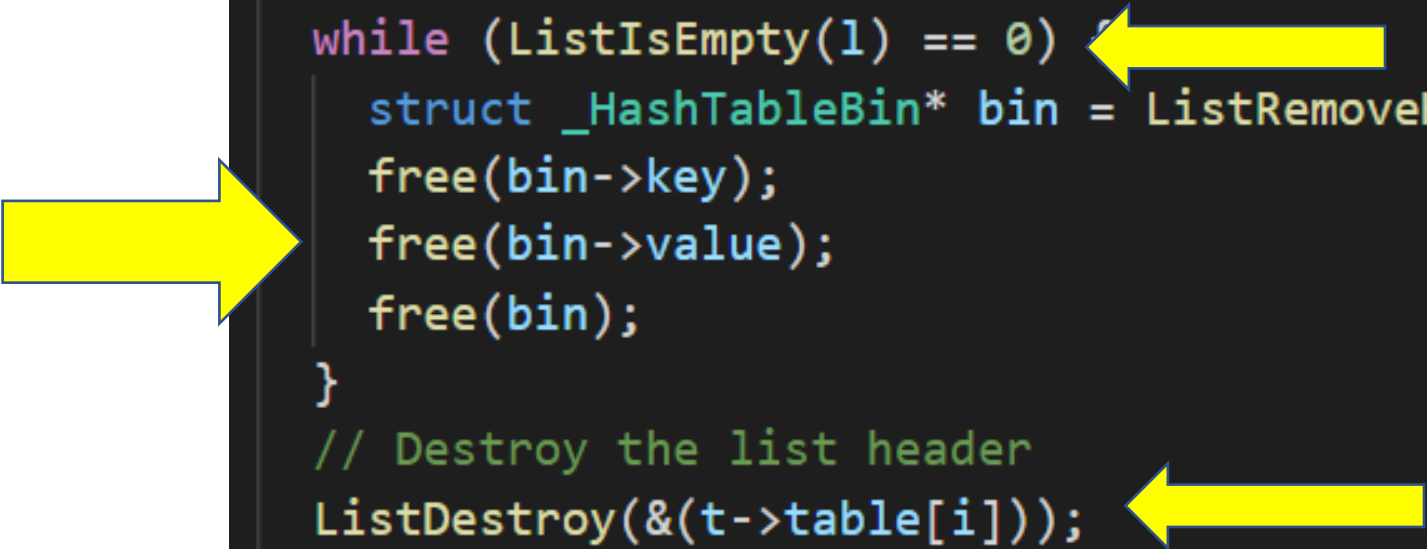


HashTableCreate

```
HashTable* hTable = (HashTable*)malloc(sizeof(struct _HashTableHeader));  
assert(hTable != NULL);  
hTable->table = (List**)malloc(size * sizeof(List*));  
assert(hTable->table != NULL);  
  
hTable->size = size;  
hTable->numBins = 0;  
hTable->hashF = hashF;  
  
for (int i = 0; i < size; i++) {  
    hTable->table[i] = ListCreate(comparator);  
}
```

HashTableDestroy

```
for (int i = 0; i < t->size; i++) {  
    List* l = t->table[i];  
    // Free the HT bins of each list  
    while (ListIsEmpty(l) == 0) {  
        struct _HashTableBin* bin = ListRemoveHead(l);  
        free(bin->key);  
        free(bin->value);  
        free(bin);  
    }  
    // Destroy the list header  
    ListDestroy(&(t->table[i]));  
}  
free(t->table);  
free(t);
```





Procurar

```
// Search for the key
// If found, the list current node is updated
//
static int _searchKeyInList(List* l, char* key) {
    if (ListIsEmpty(l)) {
        return 0;
    }

    // Needed for the comparator
    // Shallow copy of the key: just the pointer
    struct _HashTableBin searched;
    searched.key = key;

    ListMoveToHead(l);
    return ListSearch(l, &searched) != -1;
}
```



Inserir

```
int HashTablePut(HashTable* hashT, char* key, char* value) {
    unsigned int index = hashT->hashF(key) % hashT->size;
    List* l = hashT->table[index];

    if (_searchKeyInList(l, key) == 1) {
        // FOUND, cannot be added to the table
        return 0;
    }

    // Does NOT BELONG to the table
    // Insert a new bin in the list

    struct _HashTableBin* bin = (struct _HashTableBin*)malloc(sizeof(*bin));
    bin->key = (char*)malloc(sizeof(char) * (1 + strlen(key)));
    strcpy(bin->key, key);
    bin->value = (char*)malloc(sizeof(char) * (1 + strlen(value)));
    strcpy(bin->value, value);

    ListInsert(l, bin);
    hashT->numBins++;

    return 1;
}
```

Substituir



```
int HashTableReplace(const HashTable* hashT, char* key, char* value) {
    unsigned int index = hashT->hashF(key) % hashT->size;
    List* l = hashT->table[index];

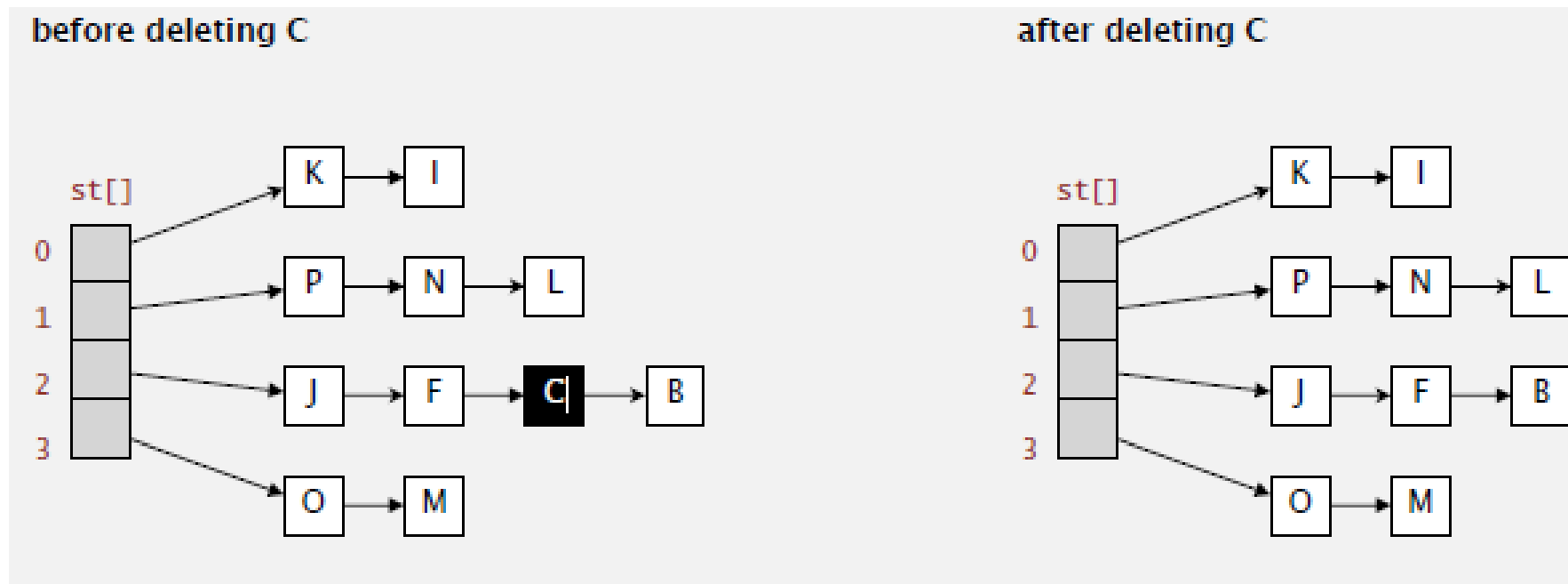
    // Search and update current, if found
    if (_searchKeyInList(l, key) == 0) {
        return 0;
    }

    struct _HashTableBin* bin = ListGetCurrentItem(l);

    free(bin->value);
    bin->value = (char*)malloc(sizeof(char) * (1 + strlen(value)));
    strcpy(bin->value, value);

    return 1;
}
```

Apagar é fácil !



[Sedgewick & Wayne]

Apagar

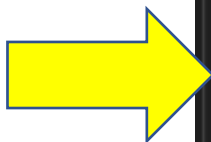
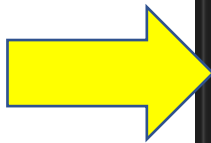
```
int HashTableRemove(HashTable* hashT, char* key) {
    unsigned int index = hashT->hashF(key) % hashT->size;
    List* l = hashT->table[index];

    // Search and update current, if found
    if (_searchKeyInList(l, key) == 0) {
        return 0;
    }

    // Get rid of the bin
    struct _HashTableBin* bin = ListGetCurrentItem(l);
    free(bin->key);
    free(bin->value);
    free(bin);

    // Get rid of the list node
    ListRemoveCurrent(l);
    hashT->numBins--;

    return 1;
}
```

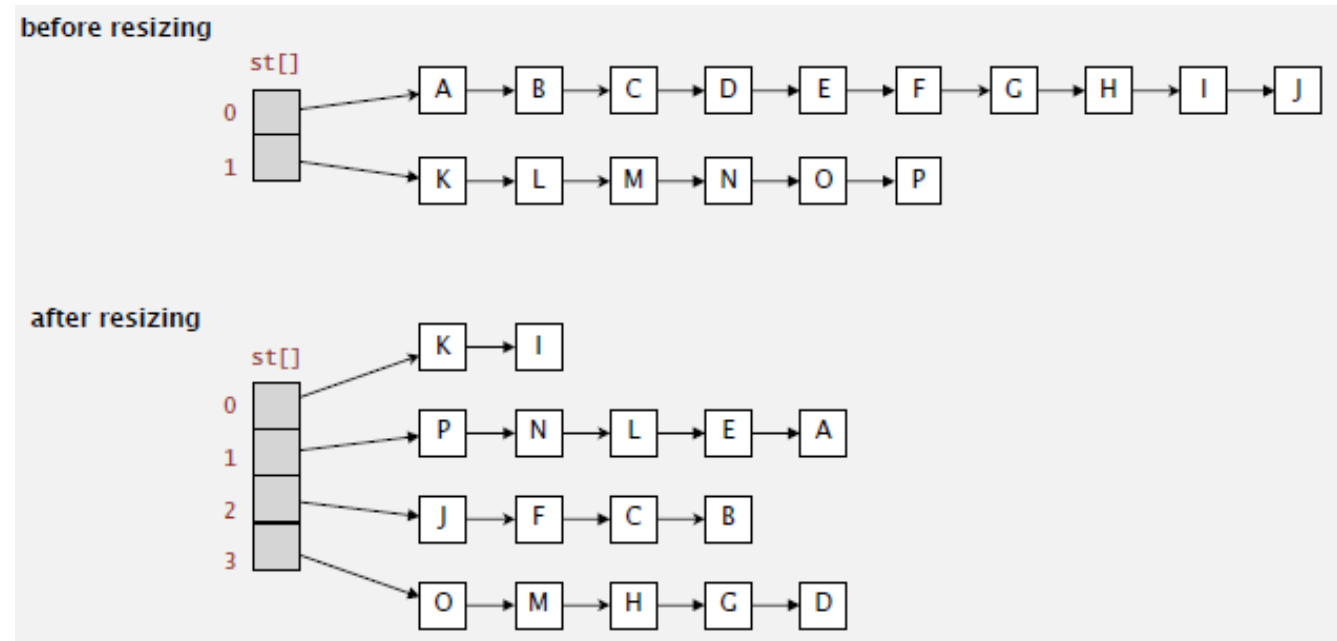


Eficiência Computacional

- Em média, N/M elementos em cada lista – Load Factor
- Procurar / inserir \rightarrow nº de comparações é proporcional a N/M
 - M vezes mais rápido que na procura sequencial
- M muito grande \rightarrow demasiadas listas vazias
- M demasiado pequeno \rightarrow listas muito longas
- Escolha habitual : $M \approx N/4 \rightarrow O(1)$

Resizing + Rehashing

- **Objetivo** : fator de carga aprox. constante
- **Duplicar o tamanho** do array quando $N/M \geq 8$
- **Reduzir para metade** o tamanho do array quando $N/M \leq 2$
- Criar a nova tabela e **adicionar**, um a um, todos os **itens**



[Sedgewick & Wayne]

Tarefa

- Implementar uma função para fazer **Resizing + Rehashing**
- Adaptar o tamanho da tabela à evolução do **fator de carga**

Separate Chaining

```
size = 17 | Active = 12
0 -
  | Hash = 68, (December, 12th month)
1 -
2 -
  | Hash = 70, (February, 2nd month of the year)
3 -
4 -
5 -
6 -
  | Hash = 74, (January, 1st month of the year)
  | Hash = 74, (July, 7th month)
  | Hash = 74, (June, 6th month)
7 -
8 -
9 -
  | Hash = 77, (March, 3rd month)
  | Hash = 77, (May, 5th month)
10 -
  | Hash = 78, (November, 11th month)
11 -
  | Hash = 79, (October, 10th month)
12 -
13 -
14 -
  | Hash = 65, (April, 4th month)
  | Hash = 65, (August, 8th month)
15 -
  | Hash = 83, (September, 9th month)
16 -
```

Open Addressing + Linear Probing

```
size = 17 | Used = 12 | Active = 12
0 - Free = 0 - Deleted = 0 - Hash = 68, 1st index = 0, (December, The last month of the year)
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4 - Free = 1 - Deleted = 0 -
5 - Free = 1 - Deleted = 0 -
6 - Free = 0 - Deleted = 0 - Hash = 74, 1st index = 6, (January, 1st month of the year)
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```

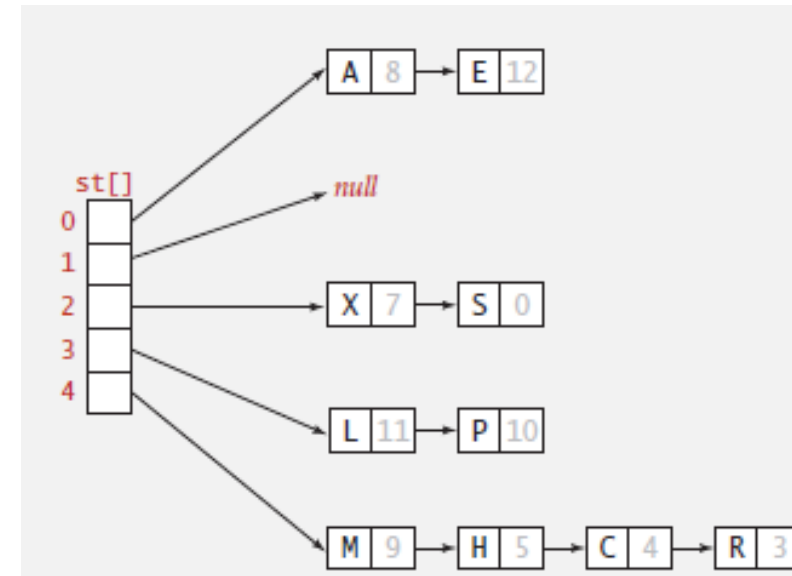
Eficiência Computacional

implementation	guarantee			average case			ordered ops?	key interface
	search	insert	delete	search hit	insert	delete		
separate chaining	N	N	N	$3-5^*$	$3-5^*$	$3-5^*$		<code>equals()</code> <code>hashCode()</code>
linear probing	N	N	N	$3-5^*$	$3-5^*$	$3-5^*$		<code>equals()</code> <code>hashCode()</code>
* under uniform hashing assumption								

[Sedgewick & Wayne]

Separate Chaining **vs** Linear Probing

- **Separate Chaining**
 - Desempenho não se degrada abruptamente
 - Pouco sensível a funções de hashing menos boas
-
- **Linear Probing**
 - Menos espaço de memória desperdiçado



	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
keys[]	P	M			A	C	S	H	L		E				R	X
vals[]	10	9			8	4	0	5	11		12				3	7

[Sedgewick & Wayne]

Hash Tables **vs** Balanced Search Trees

- Tabelas de Dispersão
- Código mais simples
- Melhor alternativa se **não** pretendermos **ordem**
- Mais **rápidas**, para chaves simples
- Árvores Binárias Equilibradas
- Pior caso : **$O(\log N)$** vs **$O(N)$**
- Suportam **ordem**
- **compareTo()** **vs** **equals()** + **hashCode()**