

# Access Control Models

SIO

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# Access types

- Physical access *Equipmentos expostos*
  - Physical contact between a subject and the object of interest
    - Facility, room, network, computer, storage device, authentication token, etc.
  - Out of scope of this course ...
- Informatic or electronic access
  - Information-oriented contact between a subject and the object of interest
    - Contact through request-response dialogs
  - Contact is mediated by
    - Computers and networks
    - Operating systems, applications, middleware, devices, etc.

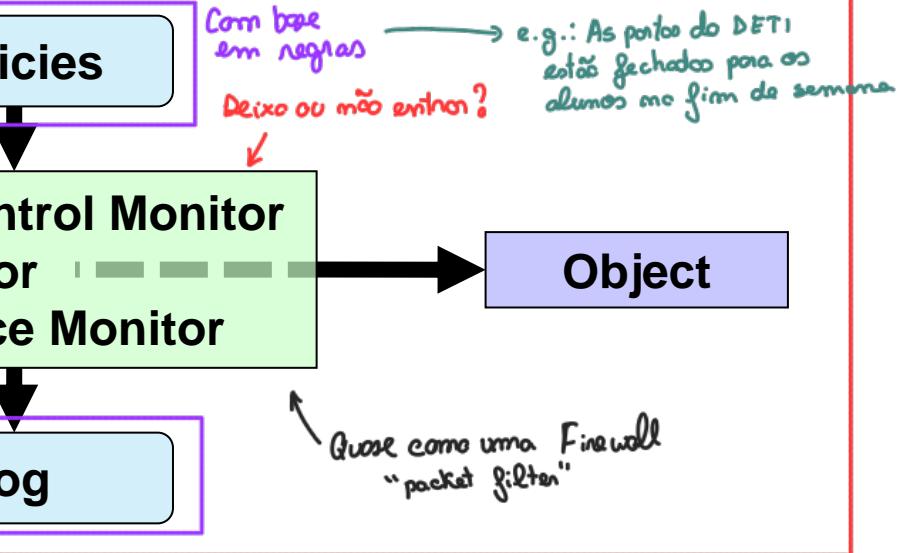
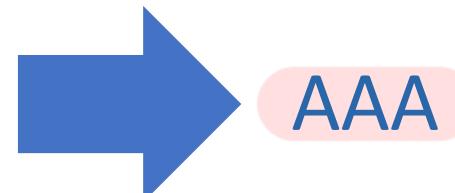
# Access control

- Definition

- Policies and mechanisms that mediate the access of a subject to an object

- Normal requirements

- Authentication ↗ Quem está a fazer o pedido!
  - With some Level of Assurance (LoA)
  - Authorization policies
  - Accountability → logging



Regras: São a passagem para algo escrito (concreto) de uma política!  
Políticas: decisões de alto nível, não define como fazer!

# Access control: Subjects and Objects

- Both are digital entities
- Subjects are something exhibiting activity:
  - Processes
  - Computers
  - Network elements
- Objects are targets of actions (resources):
  - Stored data
  - CPU time
  - Memory
  - Processes
  - Computers
  - Networks
- An entity can be both subject & object

Fork bomb já não manda o PC a baixar ...

# Least privilege principle

↳ Princípio do privilégio mínimo

"Every program and every user of the system should operate using the least set of privileges necessary to complete the job" → Nunca devemos trabalhar com privilégios superiores aos que precisamos!

J. H. Saltzer, M. D. Schroeder, Proc. of The protection of information in computer systems, IEEE, 63(9) 1975

- Privilege:
  - Authorization to perform a given task
  - Similar to access control clearance
- Subjects should have, at any time, the exact privileges required to their assigned tasks
  - Less privileges than the required create unsurpassable barriers
  - More privileges than the required create vulnerabilities
    - Damage resulting from accidents or errors
    - Potential interactions among privileged programs
    - Misuse of a privileges
    - Unwanted information flows
    - "need-to-know" military restrictions

↳ e.g.: Trabalhando sempre com o **root** é perigoso...  
é possível acabar com a máquina...

↳ Não deve saber coisas que não precisa...

# Access control models

## • Access control matrix

- Matrix with all access rights for subjects relatively to objects
- Conceptual organization

Apenas o conceito! // (matrix conceptual)

"por objetos": ACLs  
"por sujeitos": Capacidades

	O1	O2	...	Om-1	Om
S1		Access rights			
S2					
...					
Sn-1					
Sn					

OBJETOS

Normalmente os direitos associados a um objeto são uma ACL Access Control List

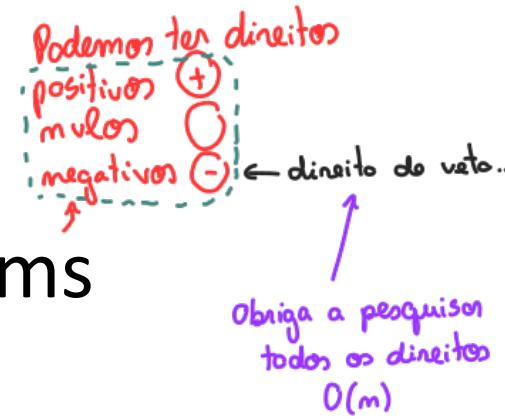
DIREITOS positivos / nulos / negativos

Capacidade! (guardada no token)  
no Auth 2.0 vai no token ... SIO

# Access control models

- **ACL-based mechanisms**

- ACL: Access Control List
- Matrix column



- List of access rights for specific subjects
  - Access rights can be positive or negative
  - Default subjects may often be used
- Usually, ACLs are stored along with objects
  - e.g., for file system objects

	O1	O2	...	Om-1	Om
S1		Access rights			
S2					
...					
Sn-1					
Sn					

↑  
Para 1 objeto as  
permissões para os sujeitos

# Access control models

Este sujeito, tem vários permissões...

Na linha

- Capability-based mechanisms
  - Capability: unforgeable authorization token
  - Matrix row
  - Contains object references and access rights
- Access granting
  - Transmission of capabilities between subjects
  - Mediated / non-mediated
- Usually, capabilities are kept by subjects
  - e.g., OAuth 2.0 access tokens

	O1	O2	...	Om-1	Om
S1		Access rights			
S2					
...					
Sn-1					
Sn					

Pode ir no token :

# Access control models: MAC and DAC

- **Mandatory access control (MAC)**
  - Fixed access control policy implemented by the access control monitor
  - Access control rights cannot be tailored by subjects or object owners
    - ↳ Nunca mudam,
- **Discretionary access control (DAC)**
  - Some subjects can update rights granted or denied to other subjects for a given object
  - Usually this is granted to object owners and system administrators
    - ↳ Podem mudar consecutivamente { estao fora do Monitor,

# Access control models: Role-Based Access Control (RBAC)

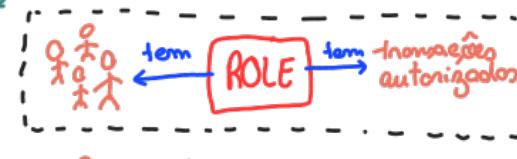
Role-Based Access Control

D.F. Ferraiolo and D.R. Kuhn, "Role Based Access Control",  
15th National Computer Security Conference, Baltimore, Oct. 1992

- Not DAC or MAC

- Roles are dynamically assigned to subjects
- For access control what matters is the role played by the subject
  - And not the subject's identity
  - Identity is mostly relevant for role access and logging

- Groups: (many subjects) união de pessoas <sup>FIXO</sup> pertence a ponto.
- Roles: união de direitos
  - ↳ Podemos assumir um role qualquer, escolher quais quero POR SESSÃO



- Os sujeitos assumem roles por sessão!



Quero distribuir os direitos por sujeitos diferentes  
Não por tabela!  
Associo direitos a transações e  
não direitos a objetos

- Access control binds roles to (meaningful) operations

- Operations are complex, meaningful system transactions !
  - Not the ordinary, low-level read/write/execute actions on individual objects
  - Operations can involve many individual, lower-level objects

Não queremos redundância!

# Access control models: RBAC role assignment

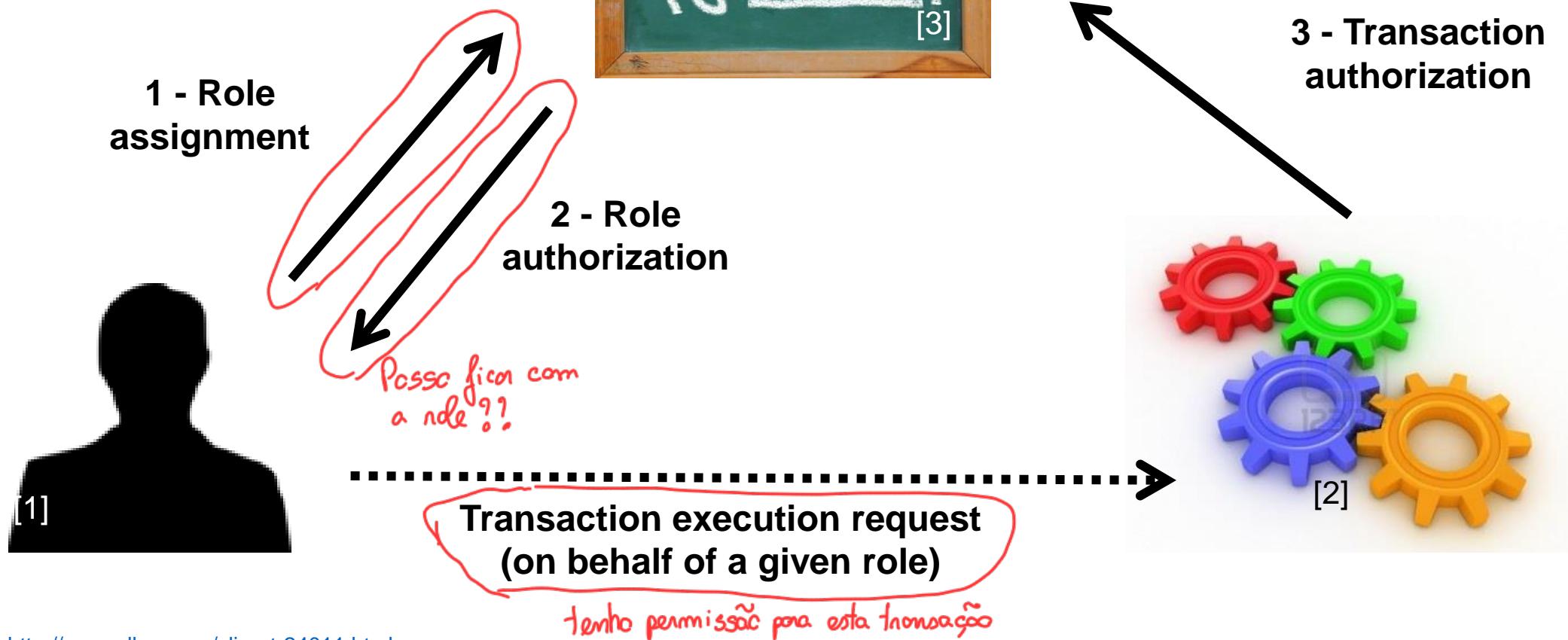
- All subject activity on the system is conducted through transactions
  - And transactions are allowed to specific roles
  - Thus, all active subjects are required to have some active role
- A subject can execute a transaction iff it has selected or been assigned a role which can use the transaction

# Access control models: RBAC role authorization

Se ele pertencer à role ...  
já

- A subject's active role must be authorized for the subject
  - In that case, the subject may assume the role
- Transaction authorization:
  - A subject can execute a transaction iff
    - the transaction is authorized through the subject's role memberships
  - and
    - there are no other constraints that may be applied across subjects, roles, and permissions

# RBAC rules



[1] From <http://www.clker.com/clipart-24011.html>

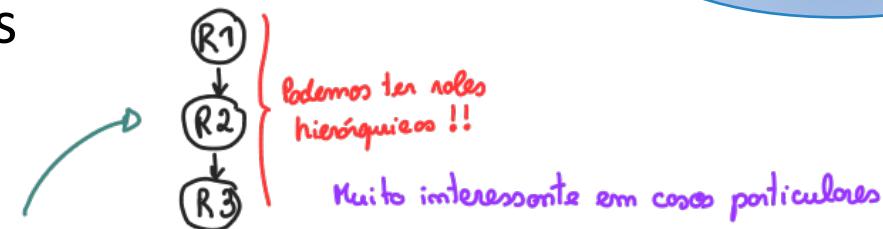
[2] From [http://www.123rf.com/photo\\_12115593\\_three-dimensional-colored-toothed-wheels.html](http://www.123rf.com/photo_12115593_three-dimensional-colored-toothed-wheels.html)

[3] From <http://www1.yorksolutions.net/Portals/115255/images/MyRoleIs.jpg>

# RBAC variants

Role-Based Access Control

- RBAC 0
  - No role hierarchies
  - No role constraints



- RBAC 1
  - RBAC 0 w/ role hierarchies (privilege inheritance)

- RBAC 2
  - RBAC 0 w/ role constraints (separation of duties)

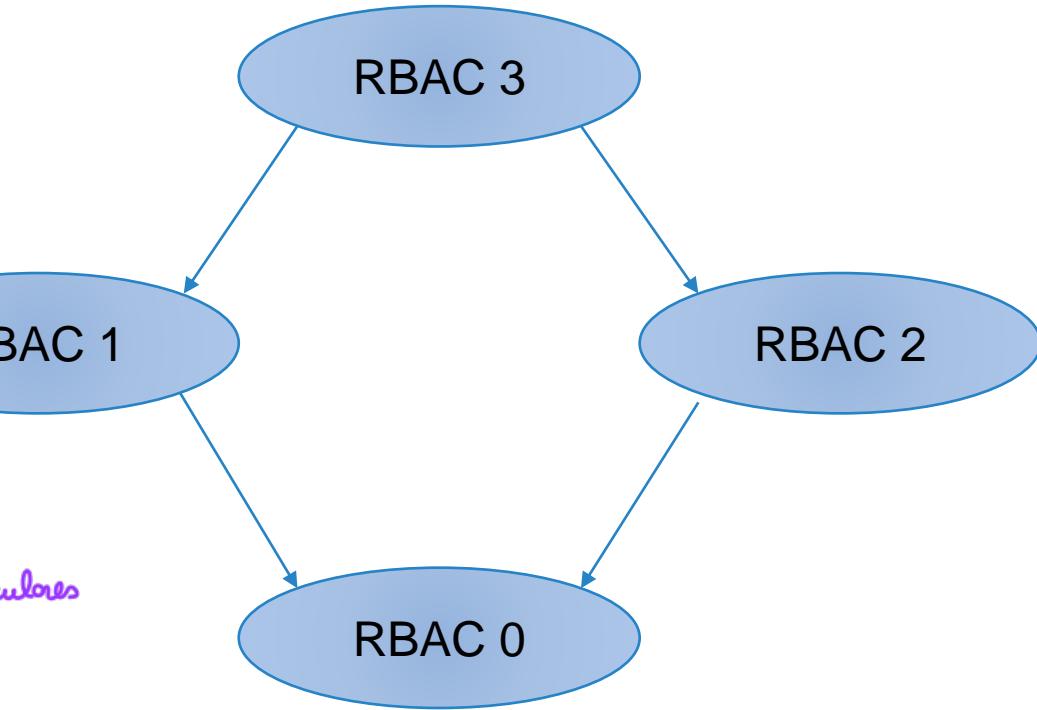
*separação de deveres*

Isolam os papéis entre todos os roles!  
Evitar uma entidade a manipular a outra

Quem tem uma role  
não tem outra!

- RBAC 3
  - RBAC 1 + RBAC 2

Hierarquia + Separação de deveres



"quase como criar uma classe  
para cada conjunto de permissões  
e garantir condições"

Cashier  
Auditor

→ A user could have both the "Cashier" and "Auditor" roles.

# NIST RBAC model

TESTE  
?

- **Flat RBAC**

- Simple RBAC model w/ user-role review

Poder de verificar qual é role que um utilizador tem, e quais pode assumir! //

Revisão de Roles

- **Hierarchical RBAC**

- Flat RBAC w/ role hierarchies (DAG or tree)
- General and restricted hierarchies

- **Constraint RBAC**

- RBAC w/ role constraints for separation of duty

separação de deveres

- **Symmetric RBAC**

- RBAC w/ permission-role review

Dada uma permissão quais roles têm, e quais permisões tem uma role

Revisão de Permissão

"utilizador - papel"

## User-role review

Which users can have a role?

Role → users

Which roles can a user have?

User → roles

"direito - papel"

## Permission-role review

Which permissions has a role?

Role → permissions

Which roles have a permission?

Permission → roles

# Access control models: Context-Based Access Control (CBAC)

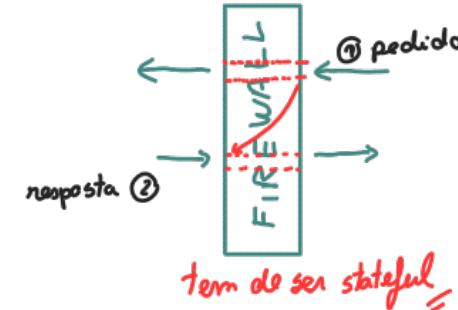
- Access rights have an historical context

Baseado em histórico

- The access rights cannot be determined without reasoning about past access operations

Com base no que fiz no passado!,,

- Example: Stateful packet filter firewall



- Example: Chines Wall policy

ou está na chino ou  
ma monoglia

- Conflict groups

Monogis não podem  
passar a moralha...  
(conflictos passados)

lo o que ferg mo  
porada afecta o  
futuro...

D.F.C. Brewer and M.J. Nash,  
**"The Chinese Wall Security Policy"**,  
IEEE Symposium on Security and Privacy, 1989

- Access control policies need to address past accesses to objects from members of conflict groups

# Access control models: Attribute-Based Access Control (ABAC)

- Access control decisions are made based on attributes associated with relevant entities

*escreve em yaml*

- OASIS XACML architecture

- Policy Administration Point (PAP)

- Where policies are managed

- Policy Decision Point (PDP)

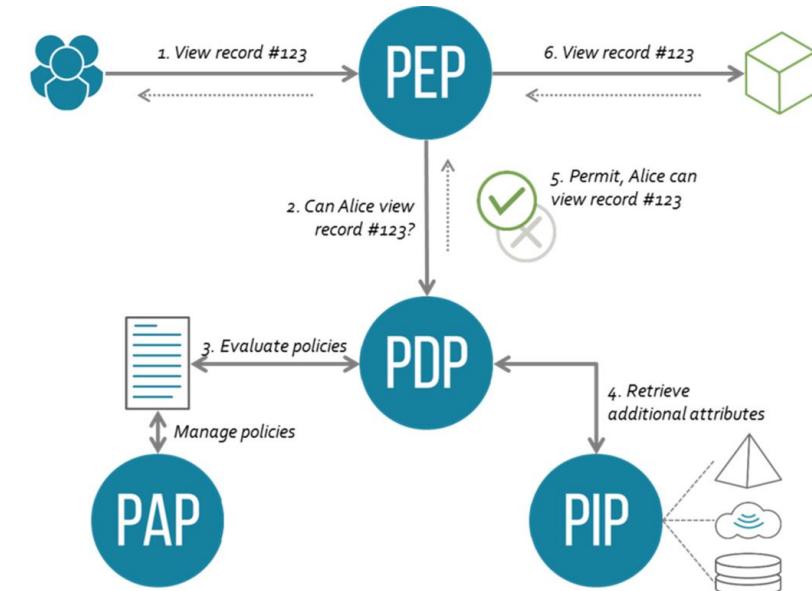
- Where authorization decisions are evaluated and issued

- Policy Enforcement Point (PEP)

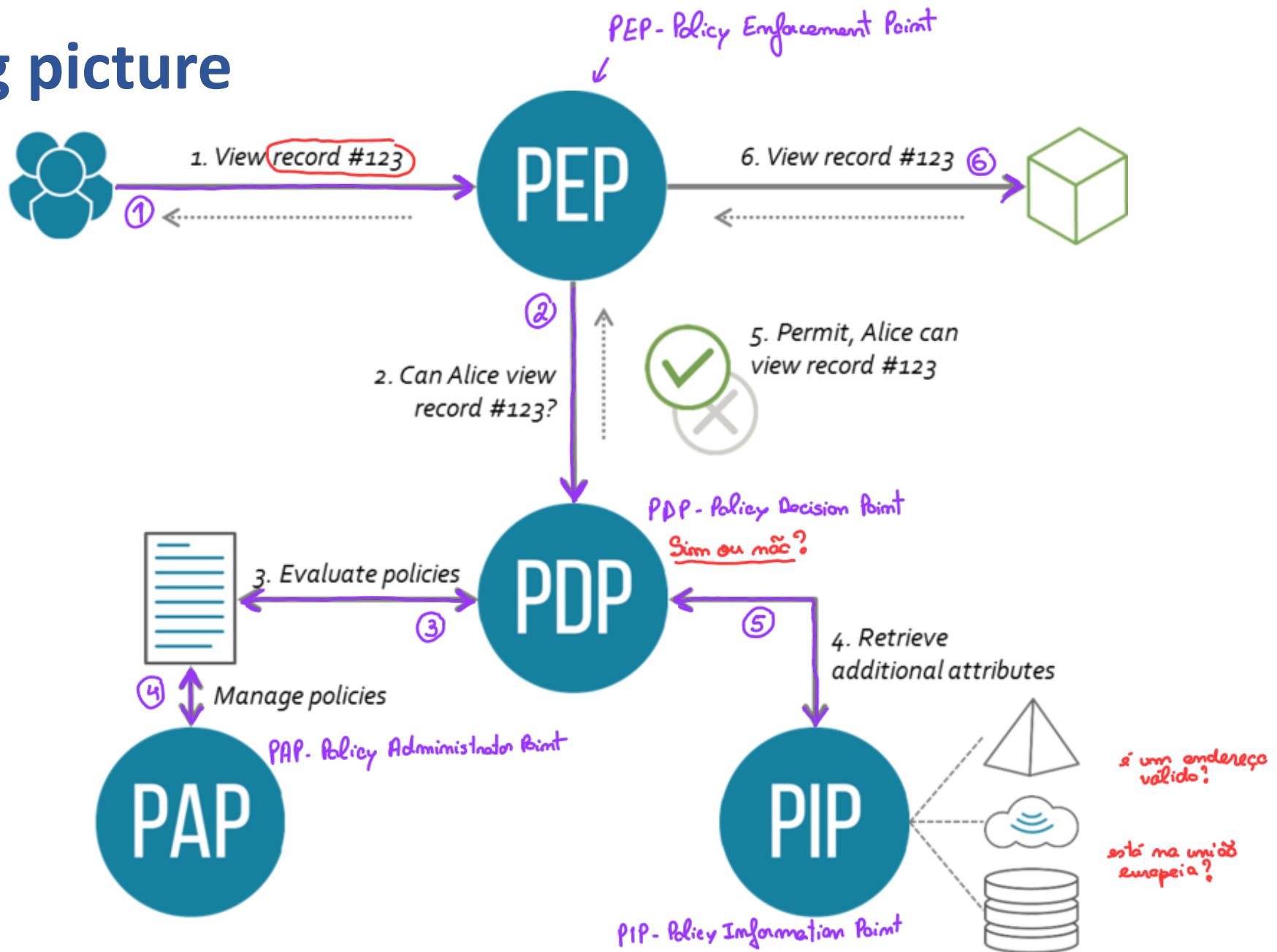
- Where resource access requests are intercepted and confronted with PDP's decisions

- Policy Information Point (PIP)

- Where the PDP gets external information



# XACML big picture



# XACML: Access control with PEP and PDP

- A subject sends a request
  - Which is intercepted by the Policy Enforcement Point (PEP)
- The PEP sends an authorization request to the Policy Decision Point (PDP)
  - With some subject's attributes
- The PDP evaluates the authorization request against its policies and reaches a decision
  - Which is returned to the PEP
  - Policies are retrieved from a Policy Retrieval Point (PRP)
  - Useful attributes are fetched from Policy Information Points (PIP)
  - Policies are managed by the Policy Administration Point (PAP)

# Access control models: Break-the-glass

Emergência!

"eu quero ultrapassar a limitação"

"uma razão de força maior"

Políticos break - the - glass

- In some scenarios it may be required to overcome the established access limitations
  - e.g., in a life-threatening situation
    - ← um médico encontra uma pessoa desmaiada, precisa de ver o seu relatório clínico,
- In those cases, the subject may be presented with a break-the-glass decision upon a deny
  - Can overcome the deny at their own responsibility
  - Logging is fundamental to prevent abuses

Ignorar a não permissão  
em casos excepcionais...

# Separation of Duties

R.A. Botha, J.H.P. Eloff, “Separation of duties for access control enforcement in workflow environments”, IBM Systems Journal, 2001

- Fundamental security requirement for fraud and error prevention
  - Dissemination of tasks and associated privileges for a specific business process among multiple subjects
  - Often implemented with RBAC RBAC 2
- Damage control
  - Segregation of duties helps reducing the potential damage from the actions of one person
  - Some duties should not be combined into one position

# Segregation of duties

## ISACA (Inf. Systems Audit and Control Ass.) matrix guideline

↑  
Matriz de segregação  
de deveres

Exhibit 2.9—Segregation of Duties Control Matrix													
	Control Group	Systems Analyst	Application Programmer	Help Desk and Support Manager	End User	Data Entry	Computer Operator	Database Administrator	Network Administrator	Systems Administrator	Security Administrator	Systems Programmer	Quality Assurance
Control Group		X	X	X		X	X	X	X	X		X	
Systems Analyst	X			X	X		X				X		X
Application Programmer	X			X	X	X	X	X	X	X	X	X	X
Help Desk and Support Manager	X	X	X		X	X		X	X	X		X	
End User		X	X	X			X	X	X			X	X
Data Entry	X		X	X			X	X	X	X	X	X	
Computer Operator	X	X	X		X	X		X	X	X	X	X	
Database Administrator	X			X	X	X	X		X	X			X
Network Administrator	X			X	X	X	X	X					
System Administrator	X			X	X		X	X				X	
Security Administrator			X				X	X				X	
Systems Programmer	X		X	X	X	X	X	X		X	X		X
Quality Assurance		X	X		X						X		

X—Combination of these functions may create a potential control weakness.

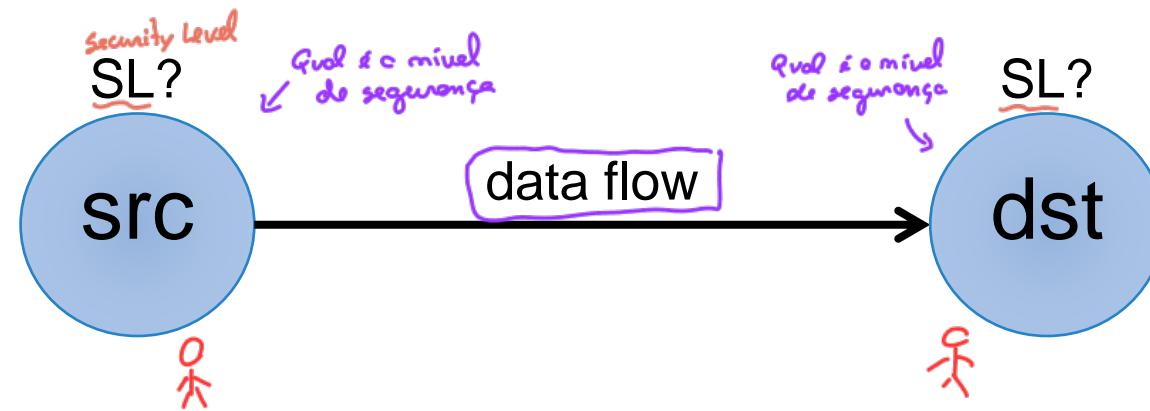
Pode

Não pode  
ser! //

# Information flow models

Quando falamos de **MAC - Mandatory Access control** pensamos misto:

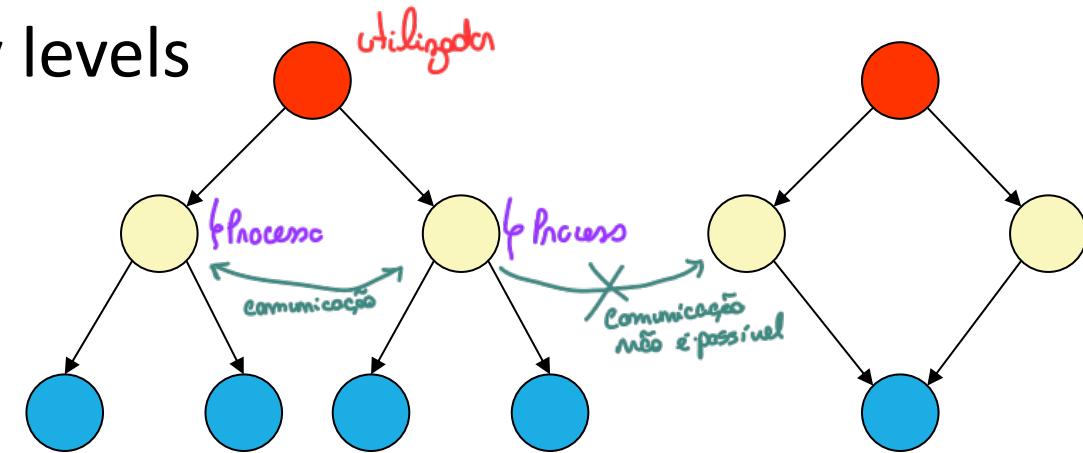
- Authorization is applied to data flows { → e.g.: Materia classificada em Portugal  
→ Só podem ser acedidos por pessoas autorizadas  
• Só pessoas com um determinado nível podem aceder!,,}
- Considering the data flow source and destination
- Goal: avoid unwanted/dangerous information flows



- Src and Dst security-level attributes
  - Information flows should occur only between entities with given security level (SL) attributes
  - Authorization is given based on the SL attributes

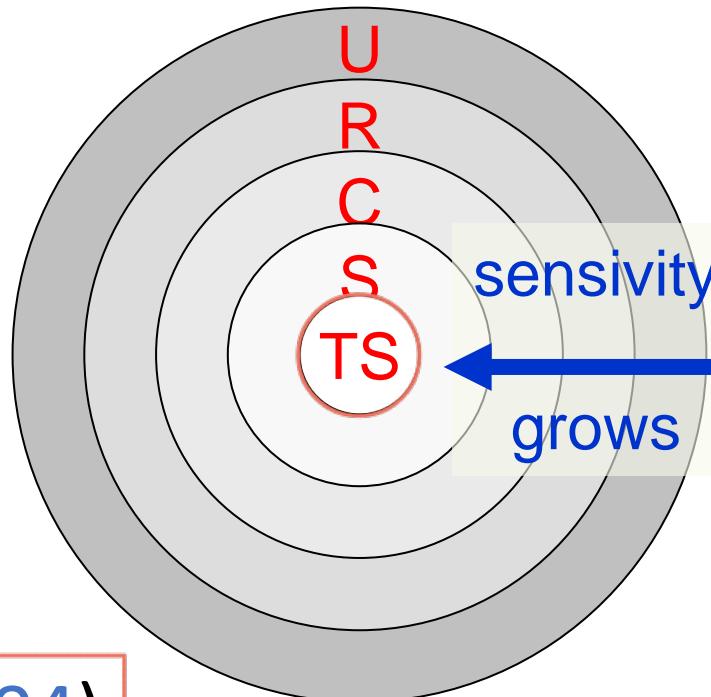
# Multilevel security

- Subjects (or roles) act on different security levels
  - Levels do not intersect themselves
  - Levels have some partial order
    - Hierarchy
    - Lattice
- Levels are used as attributes of subjects and objects
  - Subjects: security level clearance → credenciação
  - Objects: security classification  
Documentos físicos
- Information flows & security levels
  - Same security level → authorized
    - Still, restricted to a “need to know”
  - Different security levels → controlled



# ■ MS levels: Military / Intelligence organizations

- Typical levels
  - Top secret
  - Secret
  - Confidential
  - Restricted
  - Unclassified



- **Portugal** (NTE01, NTE04)
    - Muito Secreto
    - Secreto
    - Confidencial
    - Reservado
- Quanto mais o dono do país.
- Marca

- **EU** example
  - EU TOP SECRET
  - EU SECRET
  - EU CONFIDENTIAL
  - EU RESTRICTED
  - EU COUNCIL / COMMISSION
- **NATO** example
  - COSMIC TOP SECRET (CTS)
  - NATO SECRET (NS)
  - NATO CONFIDENTIAL (NC)
  - NATO RESTRICTED (NR)

# Security categories (or compartments)

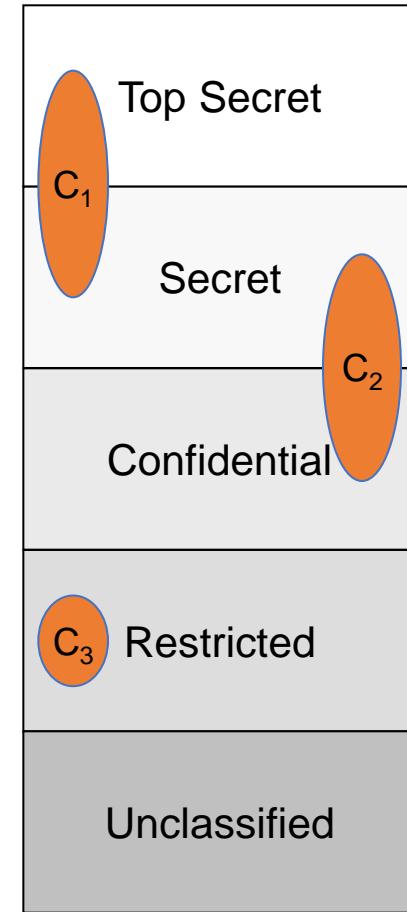
*sub-divisão dos níveis de segurança*

- Self-contained information environments
  - May span several security levels
- Military environments
  - Military branches, military units
- Civil environments
  - Departments, organizational units
- An object can belong to different compartments and have a different security classification in each of them
  - (top-secret, crypto), (secret, weapon)

*↳ Pode ser as 2 categorias*

*eu tenho um nível por categoria:*

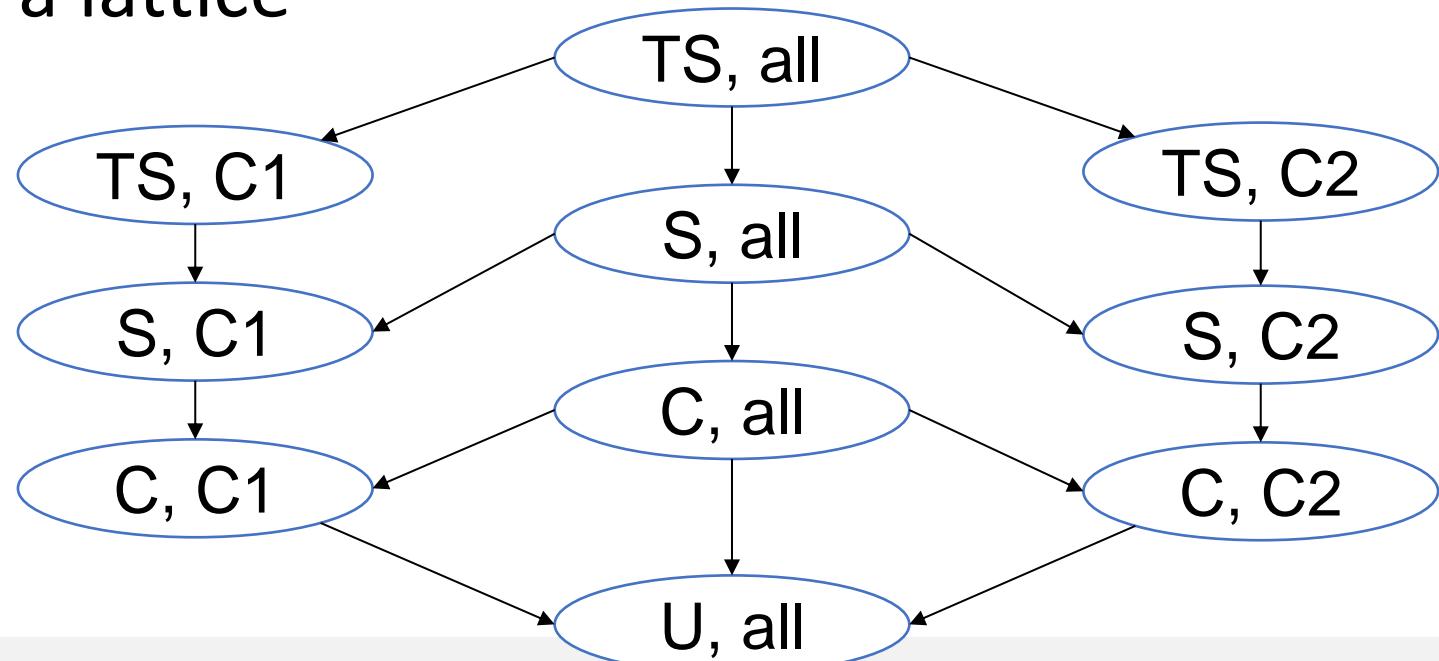
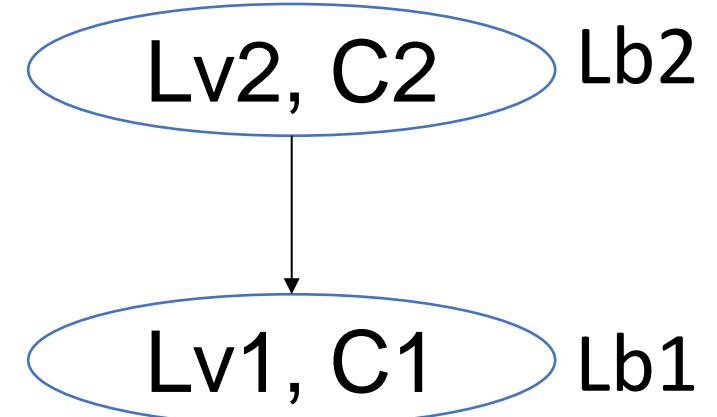
*TOP SECRET MARINHA*  
*SECRET CYBER ↑*  
*(...)* *t security label!*



# Security labels

$$\textcircled{S} \xrightarrow[c_s \subseteq c_d]{SL_s \leq SL_d} \textcircled{D}$$

- Label = Category + Level
- Relative order between labels
  - $Lb1 \leq Lb2 \Rightarrow C1 \subseteq C2 \wedge Lv1 \leq Lv2$   
contida ou igual       nível de sequência ou menor ou igual
- Labels form a lattice



# Bell-La Padula MLS Model

*Confidencialidade*

D. Elliott Bell, Leonard J. La Padula, "Secure Computer Systems: Mathematical Foundations", MITRE Tech. Report 2547, Vol. I, 1973

- Access control policy for controlling information flows
  - Addresses data confidentiality and access to classified information
  - Addresses disclosure of classified information
    - Object access control is not enough
    - One needs to restrict the flow of information from a source to authorized destinations
- Combines access control matrixes with MLS

IDEA

Consideremos o  
SECRETISMO

- Os níveis menos seguros podem produzir informação para os níveis mais seguros
- Os níveis mais seguros NÃO podem produzir informação para os níveis mais seguros

# Bell-La Padula MLS Model: secure state transitions

- Simple security condition (no read up)

— S can read O iff  $L(S) \geq L(O)$

- \*-property (no write down)

— S can write O iff  $L(S) \leq L(O)$

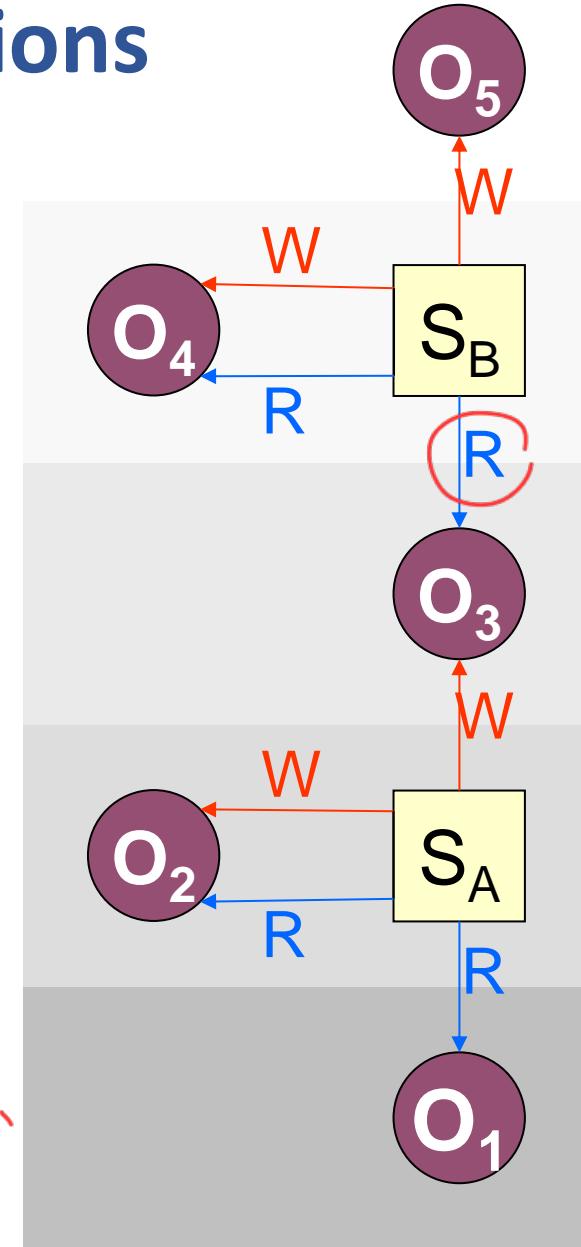
— aka confinement property

- Discretionary Security Property

— DAC-based access control within the same security level

Top Secret  $\rightarrow$  Secret  
Top Secret  $\leftarrow$  Secret

Só podemos fugir a informação para cima  
Mas posso ler de baixo...



# Biba Integrity Model

↳ Consideremos a INTEGRIDADE

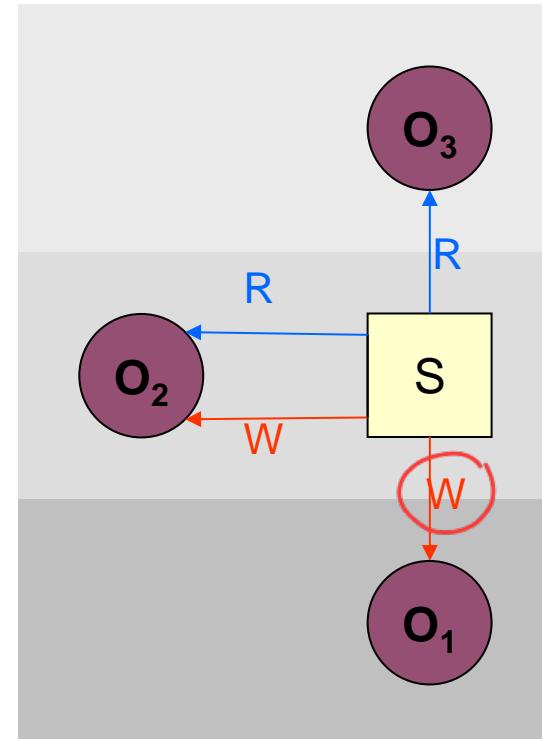
oposto da anterior ...

Se confiamos nos níveis  
de integridade superiores //

K. J. Biba, "Integrity Considerations for Secure Computer Systems", MITRE Technical Report 3153, The Mitre Corporation, April 1977

- Access control policy for enforcing integrity control over data flows
  - Uses integrity levels, not security levels
  - Similar to Bell-La Padula, with inverse rules
- Simple Integrity Property (no read down)
  - S can read O iff  $I(S) \leq I(O)$
- Integrity \*-Property (no write up)
  - S can write O iff  $I(S) \geq I(O)$

Não considera  
o SECRETISMO



# Windows mandatory integrity control

MAC

- Allows mandatory (priority and critical) access control enforcement prior to evaluate DACLs → os negros MACs são superiores à negros DACs
  - If access is denied, DACLs are not evaluated
  - If access is allowed, DACLs are evaluated

- Em windows podemos ter MAC e DAC

→ MAC tem direito de voto

- Windows tem integrity level!

Só confiamos em integrity level superior

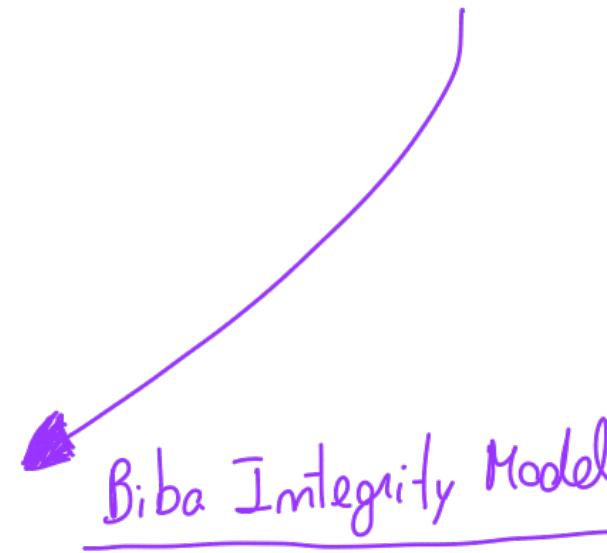
- untrusted
- low
- ...
- system
- ...

→ um processo de baixa  
Confiança não consegue virar dono!

# Windows mandatory integrity control: integrity labels

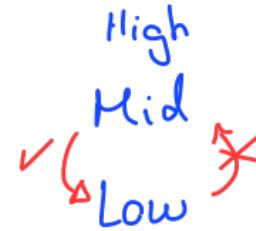
- Untrusted
- Low (or AppContainer)
- Medium
- Medium Plus
- High
- System
- Protected Process

↳ No gestor de tarefas do windows conseguimos ver isto ...



# Windows mandatory integrity control: users and processes

MAC



- **User integrity level**

- Medium: standard users *(mínimo)*
- High: elevated users *(Alto)*

- **Process integrity level**

- The minimum associated to the owner and the executable file
- User processes usually are Medium or High
  - Except if executing Low-labeled executables
- Service processes: High

# Windows mandatory integrity control

- Securable objects mandatory label
  - NO\_WRITE\_UP (default) ↗ Biba Integrity Model //
  - NO\_READ\_UP
  - NO\_EXECUTE\_UP