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# C pointers

Pointers

Dynamic memory

Geoff's self-checklist:

☒ ~~Start iClicker Cloud~~

☐ Record lecture

# Announcements

- Lab quizzes this week! Please attend your registered section
  - arrays
  - data types
  - loops
  - file input using `fgets` and `scanf`
- See Piazza post @124 for procedure
  - Send a private message to your Zoom TA in case of technical issues

# Modifying parameters with call-by-reference

- Consider these two “swap” functions and their calls
  - Trace the call stack to see what is printed after each call

```
void swap_v(int a, int b) {  
    int temp = a;  
    a = b;  
    b = temp;  
}
```

```
...  
int x = 3;  
int y = 9;  
swap_v(x, y);  
printf("x: %d, y: %d", x, y);
```

```
void swap_r(int* a, int* b) {  
    int temp = *a;  
    *a = *b;  
    b = temp;  
}
```

```
...  
int p = 3;  
int q = 9;  
swap_r(&p, &q);  
printf("p: %d, q: %d", p, q);
```

# Pointer to a pointer?

```
int main() {  
    int x = 5;  
    int* p = &x;  
    *p = 6;  
    int** q = &p;  
    int*** r = &q;  
  
    printf("%d\n", *p);  
    printf("%d\n", *q);  
    printf("%d\n", *(*q));  
}
```

- "You can keep adding levels of pointers until your brain explodes or the compiler melts – whichever happens soonest"
  - stackoverflow user JeremyP

# Back to call-by-reference

- Consider the following function that adds two parameters supplied by reference

```
int add(int* num1, int* num2) {  
    int sum = *num1 + *num2;  
    return sum;  
}  
  
int main() {  
    int a = 2;  
    int b = 4;  
    int c = add(&a, &b);  
    printf("sum = %d", c);  
}
```

- Can we modify the add function so that it uses a pointer to return the answer?

# Returning pointers

- Will it work if we just change the return type to pointer, and return the sum variable's address?

```
int* add(int* num1, int* num2) {  
    int sum = *num1 + *num2;  
    return &sum;  
}
```

```
int main() {  
    int a = 2;  
    int b = 4;  
    int* c = add(&a, &b);  
    printf("sum = %d", *c);  
}
```

This will have problems!  
Think about what is (or was)  
on the call stack

# Passing array elements as parameters

- Arrays are passed by reference by default

```
double getMaximum(double data[], int size); // prototype
double getMaximum(double* data, int size); // equivalent prototype

double answer = getMaximum(myarr, length); // function call
```

- Note that we do not need to provide "&" when specifying the address of the entire array (i.e. the address of the first element)
- If we want to specify the address of an individual element of the array, we would need the address operator
  - e.g. &data[4]

# Pointer arithmetic

- If we know the address of the first element of an array, we can compute the addresses of the other array elements
  - (or whatever comes after, if it is meaningful)

```
int A[5];
int* q = &A[0];
printf("q address: %p\n", q);
A[0] = 2;
A[1] = 4;

printf("value of q: %d\n", *q);
printf("value of q+1: %d\n", *(q + 1));
```

```
int x = 5;
int* p = &x;
printf("p address: %p\n", p);
printf("value of p: %d\n", *p);
printf("value of p+1: %d\n", *(p + 1));
```





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# Dynamic memory

# Dynamic memory

- Arrays declared as local variables must have a known size at compile time
  - but sometimes we don't know how much space we need until runtime
- Suppose we expect users to only need to store up to 1000 values, so we hard-code "1000" as an array size
  - What if user needs more? Change code and recompile
  - What if user only needs 5? Wastes memory
- If the value 1000 is hard-coded, this is hard to find and change
  - especially if used in multiple locations
- If hardcoded as a symbolic constant, still cannot change without recompiling

# Memory management in C

- We have already seen how locally-declared variables are placed on the function call stack
  - allocation and release are managed automatically
- The available stack space is extremely limited
  - placing many large variables or data structures on the stack can lead to stack overflow
- Stack variables only exist as long as the function that declared them is running

# Dynamic memory allocation

- At run-time, we can request extra space on-the-fly, from the *memory heap*
- Request memory from the heap – "allocation"
- Return allocated memory to the heap (when we no longer need it) – "deallocation"
- Unlike stack memory, items allocated on the heap must be explicitly freed by the programmer

# Dynamic memory allocation

- Function `malloc` returns a pointer to a memory block of at least `size` bytes:

```
ptr = (cast-type*) malloc(byte-size);
```

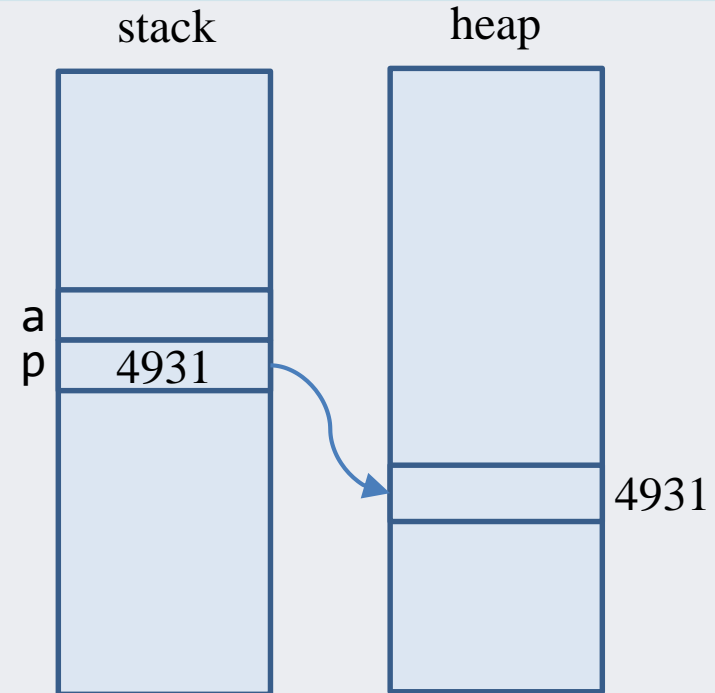
- Function `free` returns the memory block (previously allocated with `malloc`) and pointed to by `ptr` to the memory heap:

```
free(ptr);
```

- The system knows how many bytes need to be freed, provided we supply the correct address `ptr`

# Heap example

```
int main() {  
    int a;  
    int *p = (int*) malloc(sizeof(int));  
    *p = 10;  
}
```



- If there is no free memory left on the heap, **malloc** will return a null pointer
- Note: **malloc** only allocates the space but does not initialize the contents.
  - Use **calloc** to allocate and clear the space to binary zeros

# Allocating dynamic arrays

- Suppose we want to allocate space for exactly 10 integers in an array

```
#include <stdio.h>
#include <stdlib.h>

int main() {
    int* i;
    i = (int*) malloc(10*sizeof(int));
    if (i == NULL) {
        printf("Error: can't get memory...\n");
        exit(1); // terminate processing
    }

    i[0] = 3; // equivalent: *(i+0) = 3;
    i[1] = 16; // *(i+1) = 16;
    printf("%d", *i);
    ...
}
```

# Allocating dynamic arrays

From user input, variable array size

```
#include <stdio.h>
#include <stdlib.h>

int main() {
    int employees, index;
    double* wages;
    printf("Number of employees? ");
    scanf("%d", &employees);

    wages = (double*) malloc(employees * sizeof(double))
    if (!wages) { // equivalent: if (wages == NULL)
        printf("Error: can't get memory...\n");
    }

    printf("Everything is OK\n");
    ...
}
```

See `dma_examples.c`



# Dangling pointers

- When we are done with an allocated object, we free it so that the system can reclaim (and later reuse) the memory

```
int main() {  
    int* i = (int*) malloc(sizeof(int));  
    *i = 5;  
    free(i);  
  
    printf("%d", *i);  
}
```

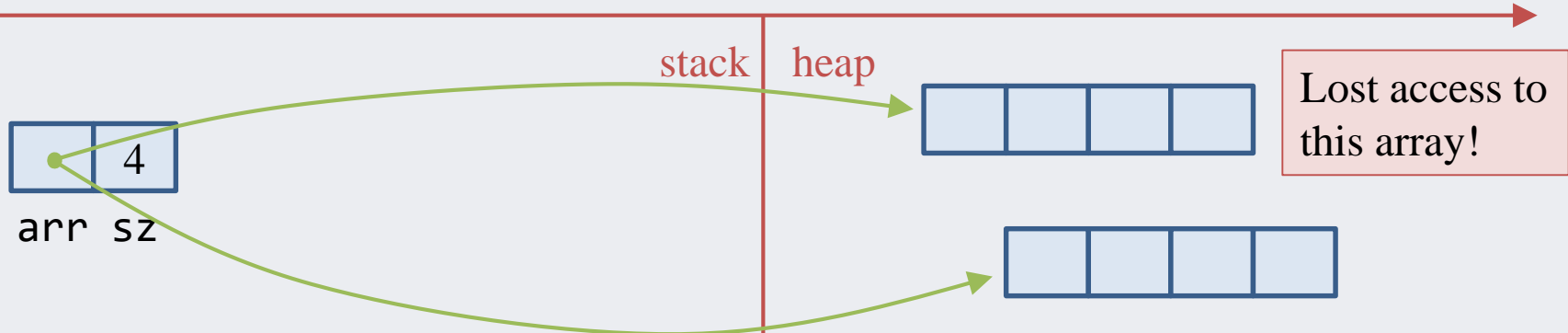
The space is marked as free, but the value remains until it is overwritten

- If the pointer continues to refer to the deallocated memory, it will behave unpredictably when dereferenced (and the memory is reallocated) – a **dangling pointer**
  - Leads to bugs that can be subtle and brutally difficult to find
  - So, set the pointer to **NULL** after freeing `i = NULL;`

# Memory leaks

- If you lose access to allocated space (e.g. by reassigning a pointer), that space can no longer be referenced, or freed
  - And remains marked as allocated for the lifetime of the program

```
int* arr;  
int sz = 4;  
arr = (int*) malloc(sz*sizeof(int));  
arr[2] = 5;  
arr = (int*) malloc(sz*sizeof(int));  
arr[2] = 7;
```



See `memory_leak.c`

# Exercise

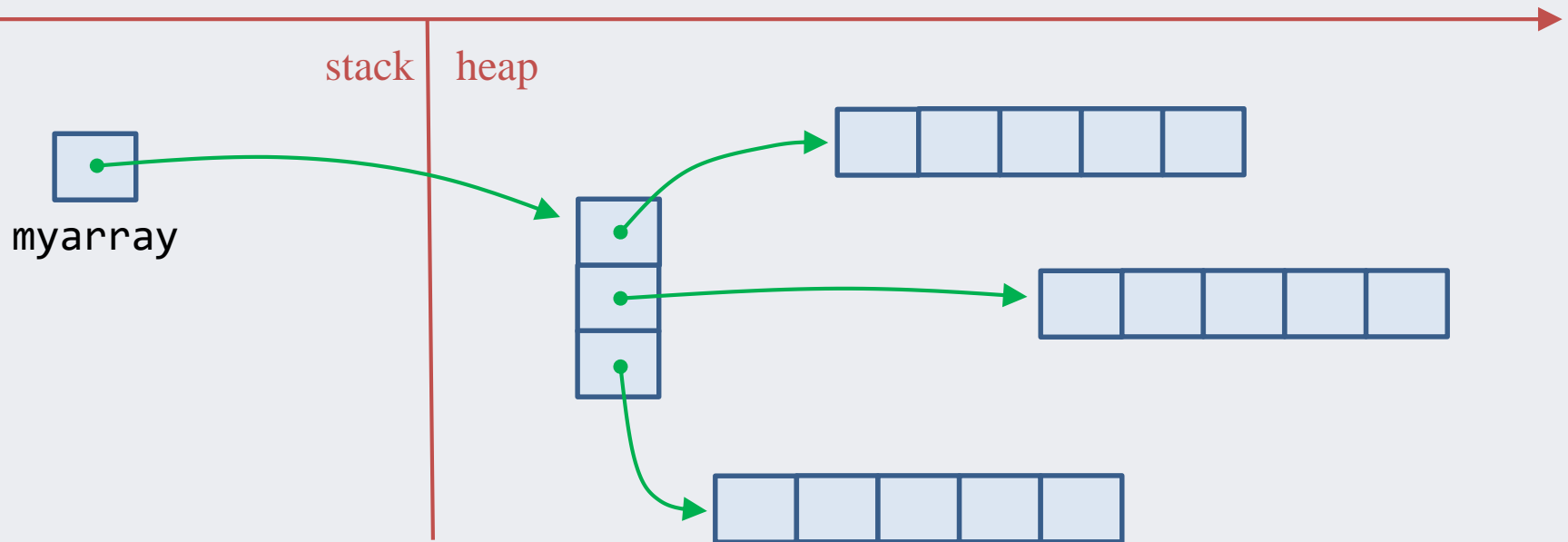
- What is printed to the screen?
  - Also clearly identify memory leaks and dangling pointers

```
int w;  
int z;  
int* t = (int*) malloc(sizeof(int));  
int* y = (int*) malloc(sizeof(int));  
int* x = (int*) malloc(sizeof(int));  
  
*x = 3;  
*y = 5;  
z = *x + *y;  
w = *y;  
*x = z;  
free(x);  
*t = 2;  
y = &z;  
x = y;  
free(t);  
printf("*x=%d, *y=%d, z=%d, w=%d\n", *x, *y, z, w);
```

# Dynamic allocation of a 2D array

See dma\_2d.c for details

```
int dim_row = 3;  
int dim_col = 5;  
int** myarray; // pointer to a pointer
```



# Stack memory vs heap memory

- Stack
  - fast access
  - allocation/deallocation and space automatically managed
  - memory will not become fragmented
  - local variables only
  - limit on stack size (OS-dependent)
  - variables cannot be resized
- Heap
  - variables accessible outside declaration scope
  - no (practical) limit on memory size
  - variables can be resized
  - (relatively) slower access
  - no guaranteed efficient use of space
  - memory management is programmer's responsibility

# Readings for this lesson

- Thareja
  - Appendices A, B, E
- Next class:
  - Thareja, Chapters 4 – 5