

# BhTSL, Behavior Trees Specification and Processing

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## Abstract

In the context of game development, there is always the need for describing behaviors for various entities, whether NPCs or even the world itself. That need requires a formalism to describe properly such behaviors. As the gaming industry has been growing, many approaches were proposed. First, finite state machines were used and evolved to hierarchical state machines. As that formalism was not enough, a more powerful concept appeared. Instead of using states for describing behaviors, people started to use tasks. This concept was incorporated in behavior trees. This paper focuses in the specification and processing of Behavior Trees. A DSL designed for that purpose will be introduced. It will also be discussed a generator that produces  $\text{\LaTeX}$  diagrams to document the trees, and a Python module to implement the behavior described. Additionally, a simulator will be presented. These achievements will be illustrated using a concrete game as a case study.

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## 1 Introduction

At some point in the video-game history, NPCs (Non-Playable Characters) were introduced. With them came the need to describe behaviors. And with these behaviors came the need of the existence of a formalism so that they can be properly specified.

As time passed by, various approaches were proposed and used, like finite and hierarchical state machines. These are state-based behaviors, that is, the behaviors are described through states. Although this is a clear and simplistic way to represent and visualize small behaviors, it becomes unsustainable when dealing with bigger and more complex behaviors. Some time later, a new and more powerful concept was introduced: using tasks instead of states to describe behaviors. This concept is incorporated in what we call behavior trees.



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Behavior Trees (BT for short) were first used in the videogame industry in the development of the game *Halo 2*, released in 2004 [5]. The idea is that people create a complex behavior by only programming actions (or tasks) and then design a tree structure whose leaf nodes are actions and the inner nodes determine the NPC's decision making. Not only these provide an easy and intuitive way of visualizing and designing behaviors, they also provide a good way to work with scalability through modularity, solving the biggest issue from state-based design. Since then, multiple gaming companies adopted this concept and, in recent years, behavior trees are also being used in different areas like Artificial Intelligence and Robotics.

In this context, we felt that it could be useful to have a DSL to specify BTs independently of application area and the programming language chosen for the implementation. The language must be compact and easy to use but it should be expressive enough to be applied to real situations. In that sense a new kind of node was included, as will be described.

This paper will introduce the DSL designed and the compiler implemented to translate it to a programming language, in this case Python. Additionally, the compiler also generates L<sup>A</sup>T<sub>E</sub>X diagrams to produce graphical documentation for each BT specified.

A small example will be described in our language as a case study to illustrate all the achievements attained.

The paper is organized as follow: Concepts and State of the Art frameworks are presented in Section 2. Architecture and language specification are proposed in Section 3. Compiler development is discussed in Section 4. An illustrative example is presented in Section 5, before concluding the paper in Section 6. The paper also includes two appendices, one that contains the tokens table, and another to show the complete specification of TGame used as an example.

## 2 Concepts

This section will be built based on references [1, 4, 3].

Formally, a BT is a tree whose internal nodes are called control flow nodes and leafs are called execution nodes.

A behavior tree executes by periodically sending ticks to its children, in order to traverse the entire tree. Each node, upon a tick call, returns one of the following three states to its parent: **SUCCESS** if the node was executed with success; **FAILURE** if the execution failed; or **RUNNING** if it could not finish the execution by the end of the tick. In the last case, the next tick will traverse the tree until it reaches the running execution node, and will try again to run it.

### 2.1 Control Flow Nodes

Control flow nodes are structural nodes, that is, they don't have any impact in the state of the system. They only control the way the subsequent tree is traversed. In the classical formulation, there are 4 types of control flow nodes: **Sequence**, **Selector**, **Parallel** and **Decorator**.

A sequence node (figure 1a) visits its children in order, starting with the first, and advancing for the next one if the previous succeeded. Returns:

- **SUCCESS** - if all children succeed;
- **FAILURE** - if a child fails;
- **RUNNING** - if a child returns **RUNNING**.

Like the sequence, the selector node (figure 1c) also visits its children in order, but it only advances if the child that is being executed returns **FAILURE**. Returns:

- **SUCCESS** - if a child succeeds;
- **FAILURE** - if all children fails;
- **RUNNING** - if a child returns **RUNNING**.

A parallel node (figure 1e), as the name implies, visits its children in parallel. Additionally, it has a parameter  $M$  that acts as a success rate. For  $N$  children and  $M \leq N$ , it returns:

- **SUCCESS** - if  $M$  children succeed;
- **FAILURE** - if  $N - M + 1$  children fail;
- **RUNNING** - otherwise.

A decorator (figure 1b) is a special node that has an only one child, and uses a policy (set of rules) to manipulate the return status of its child, or the way it ticks it. Some examples of decorator nodes are:

1. **Inverter** - inverts the **SUCCESS/FAILURE** return status of the child;
2. **Max- $N$ -Times** - the child can only fail  $N$  times. After that it only returns **FAILURE** without ticking the child.

## 2.2 Execution Nodes

Execution nodes are the simplest, yet the more powerful. They are the ones that have access to the state of the system, and can update it. There are two types of execution nodes: **Action** and **Condition**.

Upon the execution of a tick, an action node (figure 1d) runs a chunk of code that can return either **SUCCESS**, **FAILURE** or **RUNNING**

The condition node (figure 1f) verifies a proposition, returning **SUCCESS/FAILURE** if the proposition is/is not valid. This node never returns **RUNNING**.

## 2.3 Control Flow Nodes with memory

Sometimes, when a node returns **RUNNING**, we want it to remember which nodes he already executed, so that the next tick doesn't execute them again. We call this nodes with memory. And they are represented by adding a  $\_*$  to the symbols mentioned previously. This is only sintatic sugar because we can also represent these nodes with a non-memory BT, but that will not be discussed here.

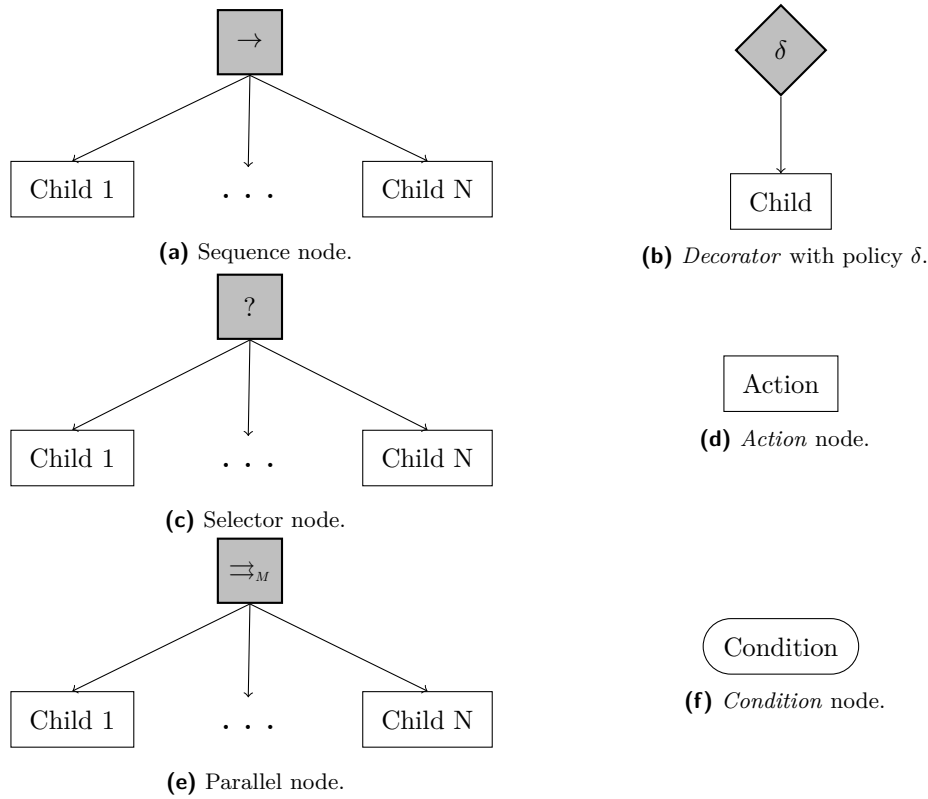
Please note that, while we avoid the re-execution of nodes with this type of node, we also lose the reactivity that this re-execution provides.

## 2.4 State of The Art

In the gaming industry there is some interesting projects that use tools based on Behavior trees as the main focus to describe NPCs behaviors. Unreal Engine [2] and Unity<sup>1</sup> are two examples of major game engines that use them. In their case, instead of a language, they offer a graphical user interface (GUI) to specify the BTs, through a drag and drop tactic. Upon the creation of an execution node, the programmer needs to specify the action or condition that will be executed. The nodes mentioned before are all implemented in these engines,

<sup>1</sup> <https://unity.com>

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■ **Figure 1** BT Nodes structure.

126 along with some extensions. All the nodes that were mentioned before are implemented in  
 127 both of these engines, along with some extensions.

128 In addition to game engines, there are also frameworks like Java Behavior Trees<sup>2</sup> for Java  
 129 and Owl<sup>3</sup> for Python that implement BTs. In this case, they work as a normal library.

### 130 3 Architecture and Specification

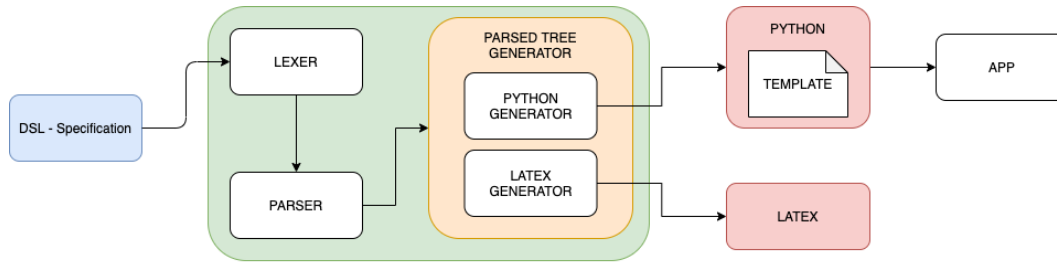
131 In this section, it will be explained the general architecture of our system to process BTs,  
 132 that is depicted in Figure 2. After introducing its modules, one subsection is devoted to the  
 133 BhTSL domain specific language design.

134 The input for our system, **DSL - Specification**, is a text file describing the behavior  
 135 which should follow the language syntax. The compiler, which is represented as the green  
 136 rounded rectangle in the diagram of Figure 2, is composed of the following modules: a **Lexer**,  
 137 a **Parser** and a **Code Generator**.

138 This generator has two sub-generators. The **Latex Generator**, that is responsible for the  
 139 generation of the  $\text{\LaTeX}$  code to draw the tree diagram representing the behavior specified.  
 140 And the **Python Generator**, that produces the fragment of Python code that implements  
 141 the desired behavior according to a template predefined by us in the context of this project;  
 142 that code fragment can be later imported by any Python application that aims to.

<sup>2</sup> <https://github.com/gaia-ucm/jbt>

<sup>3</sup> <https://github.com/eykd/owl1>

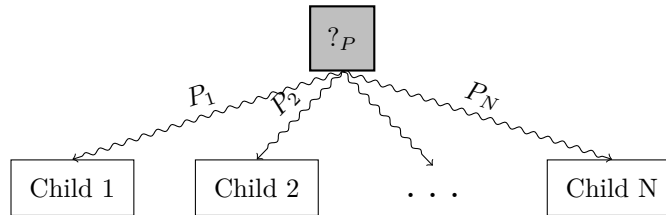


■ **Figure 2** System Architecture.

### 3.1 BhTSL

Before we start describing the DSL, we will introduce a new node, called **Probability Selector** (Figure 3 depicts that concept), that provides us with a relevant extension to the standard formalism for a more powerful behavior specification. This extension improves the expressiveness of BhTSL language.

A probability selector node is like a normal selector node, but instead of visiting its children from left to right, it visits them randomly, taking into account that each child has a probability, defined by the user, of being chosen first.



■ **Figure 3** Probability Selector node.

#### 3.1.1 Syntax

In our language, each specification represents one and only one behavior. An input file, containing the behavior specification text, is divided into 3 components:

- *Behavior* - main behavior tree;
- *Definitions* (optional) - node definitions that can be referenced in other nodes or in the main BT;
- *Code* - Python block that contains the code fragments described the execution nodes, and other code that the programmer wishes to add.

To illustrate our idea about the DSL we plan to design (formally defined by a grammar in section 4), we present below an example of a specification written in the intended language.

```

behavior : [
  sequence : [
    condition : $cond1,
    condition : $cond2
    memory selector : [
      parallel : $par1,
      prob_selector : $prob1
    ]
  ]
]
  
```

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```
170     ]
171 ]
172
173 parallel par1 : 10 [
174     action : $action1,
175     action : $action2
176 ]
177
178 prob_selector prob1 : [
179     $e1 -> decoraror : INVERTER [
180         action : $action1
181     ],
182     $e2 -> action : $action2
183 ]
184
185 %%
186
187 def action1(entity):
188     pass
189
190 def action2(entity):
191     pass
192
193 def cond1(entity):
194     pass
195
196 def cond2(entity):
197     pass
198
199 def e1(entity):
200     pass
201
202 def e2(entity):
203     pass
204
```

## 4 Tool development

In the next subsection the implementation of the BhTSL processor will be detailed, as well as the language specification will be presented.

### 4.1 Lexical analysis

The first step in the development of a compiler is the lexical analysis, that converts a char sequence into a token sequence. The tokens table can be seen in Appendix.

### 4.2 Syntatic analysis

Syntatic analysis, or parsing, it the process of analyzing a string of symbols conforming the rules of a grammar.

Below we list the context free grammar that formally specifies BhTSL syntax:

```
root : behavior CODE
      | behavior definitions CODE
```

```

218         | definition behavior CODE
219
220     behavior : BEHAVIOR ':' '[' node ']'
221
222     node : SEQUENCE ':' '[' nodes ']'
223           | SEQUENCE ':' VAR
224           | MEMORY SEQUENCE ':' '[' nodes ']'
225           | MEMORY SEQUENCE ':' VAR
226           | SELECTOR ':' '[' nodes ']'
227           | SELECTOR ':' VAR
228           | MEMORY SELECTOR ':' '[' nodes ']'
229           | MEMORY SELECTOR ':' VAR
230           | PROBSELECTOR ':' '[' prob_nodes ']'
231           | PROBSELECTOR ':' VAR
232           | MEMORY PROBSELECTOR ':' '[' prob_nodes ']'
233           | MEMORY PROBSELECTOR ':' VAR
234           | PARALLEL ':' INT '[' nodes ']'
235           | PARALLEL ':' VAR
236           | DECORATOR ':' INVERTER '[' node ']'
237           | DECORATOR ':' VAR
238           | CONDITION ':' VAR
239           | ACTION ':' VAR
240
241     nodes : nodes ',' node
242            | node
243
244     prob_nodes : prob_nodes ',' prob_node
245                | prob_node
246
247     prob_node : VAR RIGHTARROW node
248
249     definitions : definitions definition
250                 | definition
251
252     definition : SEQUENCE NODENAME ':' '[' nodes ']'
253                | SELECTOR NODENAME ':' '[' nodes ']'
254                | PROBSELECTOR NODENAME ':' '[' prob_nodes ']'
255                | PARALLEL NODENAME ':' INT '[' nodes ']'
256                | DECORATOR NODENAME ':' INVERTER '[' node ']'
257

```

### 4.3 Semantic analysis

As usual, from a static semantics perspective, the compiler will check the source text for non-declared variables and variable redeclaration.

A variable can only be accessed if it is declared, either in the *definitions* section (if it represents a control flow node), or in the *code* section (execution node). Additionally, a variable can only be declared one time, to avoid ambiguity in the memory access by the processor.

### 4.4 Code generator

The compiler can generate two different outputs: a  $\text{\LaTeX}$  file, that contains the  $\text{\LaTeX}$  commands to draw a diagram for the BT specified; and a Python file, that contains the functions that implement the specified behavior.

## 268 4.5 BT Simulator

269 In order to test the Python code generated and to help the BT specifier to debug the behavior  
 270 he is willing to describe, we also developed an interpreter that imports the Python behavior  
 271 file and simulates its execution.

272 This additional tool proved to be useful.

## 273 4.6 Implementation

274 All the project was develop using Python.

275 To generate automatically the compiler from the tokens and grammar specifications, we used  
 276 the PLY (Python Lex-Yacc)<sup>4</sup> library, which is an implementation of the `lex` and `yacc` lexer  
 277 and parser generator tools for Python.

278 To implement the Code Generator module, we resorted to the well-known tree-traversal  
 279 approach to produce the output visiting in the appropriate sequence the parse tree nodes;  
 280 some standard libraries were used for that purpose.

281 Additionally, we created a L<sup>A</sup>T<sub>E</sub>X library, `behaviortrees.sty`, to draw the trees specified.

282 This library is used in the L<sup>A</sup>T<sub>E</sub>X generator.

## 283 5 Example

284 In this section it will be presented an toy-example to illustrate how to specify a behavior in  
 285 BhTSL language.

286 **Game description:** Suppose that in some game, called `TGame`, it is intended to have a  
 287 *guard that patrols a house*. The guard has the following behavior: while he is patrolling, if  
 288 he sees the player, activates an alarm and then, depending on the level of courage he has,  
 289 decides (based on probabilities) whether he runs away or fights the player. In case of running  
 290 away, he constantly checks if he still sees the player, returning to patrolling in case he doesn't.  
 291 If he still sees it, he keeps running. The same thing happens when he chooses to fight the  
 292 player, only this time he checks if the player is already dead or not.

293 **Game specification:** The complete specification in BhTSL for the behavior described  
 294 above can be seen in Appendix B. Figure 4 shows the tree diagram of this specification after  
 295 compiling it to L<sup>A</sup>T<sub>E</sub>X.

## 296 6 Conclusion

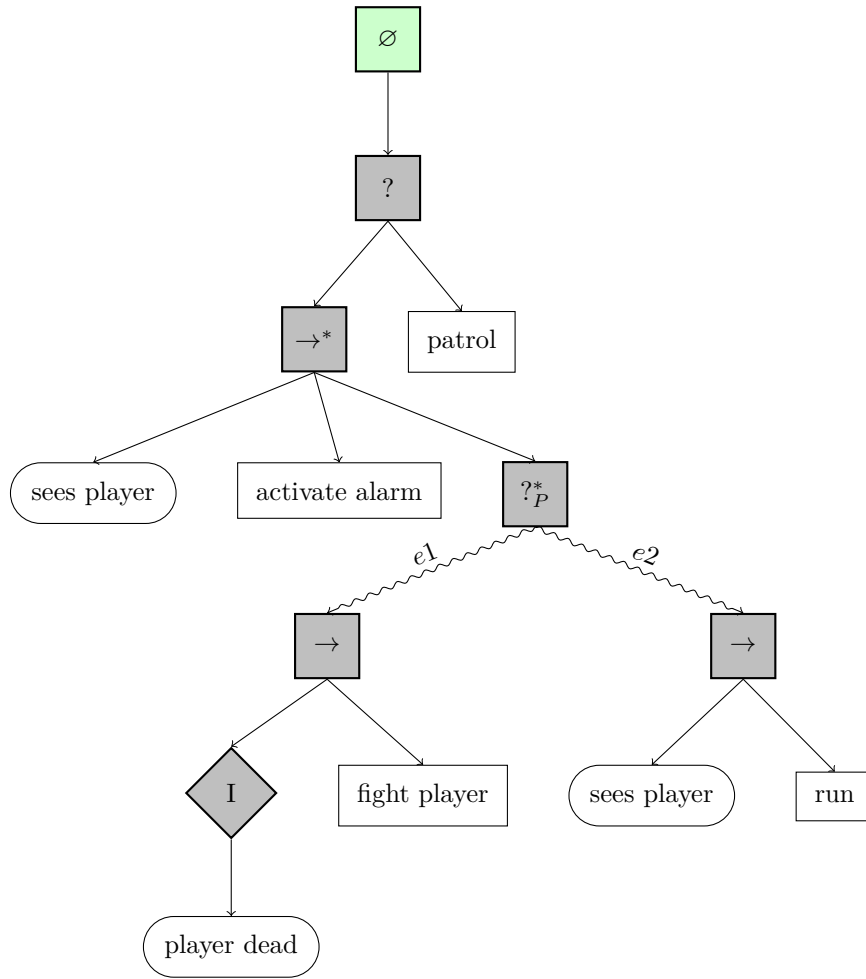
297 As games industry is growing significantly every day, the need for a formal way to describe  
 298 behaviors is also increasing requiring more and more expressiveness keeping it easy to learn,  
 299 to use and to understand. After some initial attempts not powerful enough, a new approach  
 300 called Behavior Trees (BT) appeared. This paper describes a project in which we are working  
 301 on, aimed at designing a DSL to write BT and developing the respective compiler to generator  
 302 Python functions to be incorporated in final Python programs created to implement games  
 303 or other kind of applications.

304 Along the paper the DSL designed, called BhTSL, was introduced by example and  
 305 specified by a context free grammar. The architecture of the BhTSL processor was depicted  
 306 and discussed, and the development of the compiler that produces the Python code library

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<sup>4</sup> <https://www.dabeaz.com/ply/>





■ **Figure 4** Example: TGame behavior tree diagram automatically generated by our tool

was described. In that context an example of a game specification was presented and the  $\text{\LaTeX}$  fragment that is generated to draw the BT was shown.

Although not detailed or exemplified, the simulator developed to help on debugging the BT specified in BhTSL language was mentioned along the paper.

## References

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### A Tokens Table

The following table displays the full set of tokens of BhTSL language defined in terms of regular expressions (REs) as utilized in our compiler.

<i>Tokens</i>	
Name	Value
literals	([]),:%
RIGHTARROW	->
BEHAVIOR	\bbehavior\b
SEQUENCE	\bsequence\b
SELECTOR	\bselector\b
PROBSELECTOR	\bprobselector\b
PARALLEL	\bparallel\b
DECORATOR	\bdecorator\b
CONDITION	\bcondition\b
ACTION	\baction\b
INVERTER	\bINVERTER\b
MEMORY	\bmemory\b
INT	\d+
VAR	\$_w+
NODENAME	\b\$_w+\b
CODE	%%(. \n)+

**Table 1** BhTSL Tokens Table for Lexical Analysis.

### B Example Specification

This appendix lists the TGame full specification that is an example of a game formal description in BhTSL.

```
behavior : [  
  selector : [  
    memory sequence : [  
      condition : $sees_player,  
      action : $activate_alarm,  
      memory prob_selector : [  
        $e1 -> sequence : [  
          decorator : INVERTER [  
            condition : $player_dead  
          ],  
          action : $fight_player  
        ],  
        $e2 -> sequence : [  
          condition : $sees_player,  
          action : $run  
        ]  
      ]  
    ],  
    action : $patrol  
  ]  
]
```

```
346 ]
347
348 %%
349 def sees_player(patroller):
350     from math import sqrt
351     player_x = patroller['player']['x']
352     player_y = patroller['player']['y']
353
354     patroller_x = patroller['x']
355     patroller_y = patroller['y']
356
357     if sqrt(patroller_x ** 2 - player_x** 2 +
358         patroller_y**2 - player_y**2) <= patroller['vision_radius']:
359         return SUCCESS
360
361     return FAILURE
362
363
364 def activate_alarm(patroller):
365     print("ALARM ACTIVATED!!")
366     patroller['alarm_activated'] = True
367     return SUCCESS
368
369
370 def player_dead(patroller):
371     if patroller['player']['hp'] == 0:
372         return SUCCESS
373     return FAILURE
374
375
376 def fight_player(patroller):
377     patroller['player']['hp'] -= 10
378     return RUNNING
379
380
381 def run(patroller):
382     patroller['x'] = patroller['x'] - 10
383     patroller['y'] = patroller['y'] - 10
384
385     print("Running away from player!")
386     return RUNNING
387
388
389 def patrol(patroller):
390     patroller['x'] = patroller['x'] + 10
391     patroller['y'] = patroller['y'] + 10
392     return RUNNING
393
394
395 def e1(patroller):
396     return patroller['guts'] / 10
397
398 def e2(patroller):
399     return 1 - patroller['guts'] / 10
400
```