



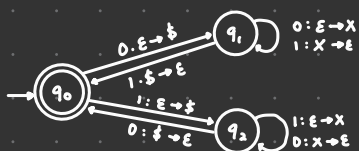
Context Free Languages 3 of 3

Definition: Deterministic PDA (DPDA)

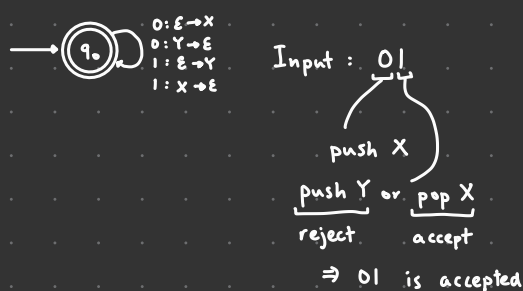
A PDA for which on every input, there's exactly one computation

Ex 1: (8.5 on pg. 253) (of a DPDA)

$$L_1 = \{x \in \{0,1\}^* : \#_0(x) = \#_1(x)\}$$

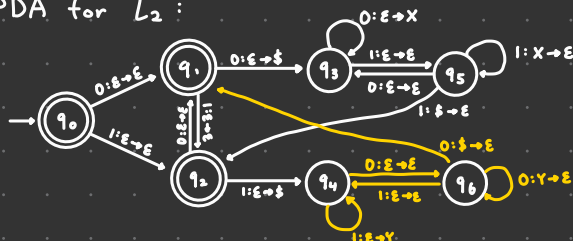


Ex 2: Of a non-deterministic / standard PDA for L_1

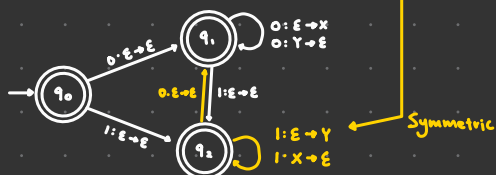


Ex 3: $L_2 = \{x \in (0,1)^* : \#_{00}(x) = \#_{11}(x)\}$

A DPDA for L_2 :



A Standard PDA for L_2 :



Theorem: The BIG Result Part 2

Let L be a language. Then:

$$L = \mathcal{L}(G) \text{ for some CFG } G \text{ iff } L = \mathcal{L}(M) \text{ for some PDA } M.$$

There are two variations of a PDA: Accept-by-state / Accept-by-stack / Two stack PDA.