



# B24 July 16 Lec 2 Notes

**Proof:** (of Gram-Schmidt orthogonalization algorithm)

Induct on  $n$  (length of  $x_1, \dots, x_n$ )

Base case: ( $n=1$ ) This is trivial

I.H: Assume algorithm works for lists up to length  $n$ .

I.S: Consider  $x_1, \dots, x_n$  L.I.

$$\begin{aligned}\langle v_{n+1}, v_j \rangle &= \left\langle x_{n+1} - \sum_{k=1}^n \frac{\langle x_{n+1}, v_k \rangle}{\|v_k\|^2} v_k, v_j \right\rangle \\ &= \langle x_{n+1}, v_j \rangle - \sum_{k=1}^n \frac{\langle x_{n+1}, v_k \rangle}{\|v_k\|^2} \underbrace{\langle v_k, v_j \rangle}_{\substack{\text{By I.H.,} \\ \text{this is } 0 \\ \text{unless } k=j}}\end{aligned}$$

$$= \langle x_{n+1}, v_j \rangle - \frac{\langle x_{n+1}, v_j \rangle}{\|v_j\|^2} \langle v_j, v_j \rangle$$

$$= 0$$

Remains to show that:

$$\text{Span} \{x_1, \dots, x_{n+1}\} = \text{Span} \{v_1, \dots, v_{n+1}\}$$

We know from I.H.:

$$\text{Span} \{x_1, \dots, x_n\} = \text{Span} \{v_1, \dots, v_n\}$$

**Proof:**  $\text{Span} \{x_1, \dots, x_{n+1}\} \subseteq \text{Span} \{v_1, \dots, v_{n+1}\}$

$$x_{n+1} = \underbrace{x_{n+1} - \sum_{k=1}^n \frac{\langle x_{n+1}, v_k \rangle}{\|v_k\|^2} v_k}_{= v_{n+1}} + \underbrace{\sum_{k=1}^n \frac{\langle x_{n+1}, v_k \rangle}{\|v_k\|^2} v_k}_{\in \text{Span}(v_1, \dots, v_n)}$$

$$\in \text{Span}(v_1, \dots, v_{n+1}) \quad \square$$

**Proof:**  $\text{Span} \{x_1, \dots, x_{n+1}\} \supseteq \text{Span} \{v_1, \dots, v_{n+1}\}$

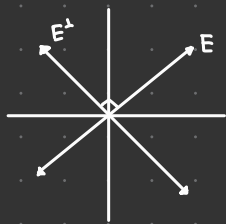
$$v_{n+1} = \underbrace{x_{n+1} - \sum_{k=1}^n \frac{\langle x_{n+1}, v_k \rangle}{\|v_k\|^2} v_k}_{\in \text{Span}(v_1, \dots, v_n) = \text{Span}(x_1, \dots, x_n)}$$

$\square$

### Definition:

If  $E \subset V$  is a subspace ( $V$  is IPS), we define the **orthogonal complement** of  $E$  by

$$E^\perp := \{x \in V : x \perp E\}$$



### Theorem:

If  $E \subset V$  is a subspace ( $V$  is IPS),  $v \in V$ , there exists unique  $v_1 \in E$ ,  $v_2 \in E^\perp$  s.t.

$$v = v_1 + v_2$$

### Proof:

Existence: Let  $v_1 = P_E v$ , and  
 $v_2 = v - P_E v$

$v_2 \perp E$  and  $v_1 \in E$  so that existence is proven, since  $v = v_1 + v_2$

Uniqueness:

If  $v_1 \in E$ ,  $v_2 \in E^\perp$  are s.t.  $v = v_1 + v_2$ , then  $v_1 = v - v_2 \in E$  and  $v - v_1 = v_2 \perp E$

$$\Rightarrow v_1 = P_E v$$

$$\Rightarrow v_2 = v - P_E v$$

□

### Question:

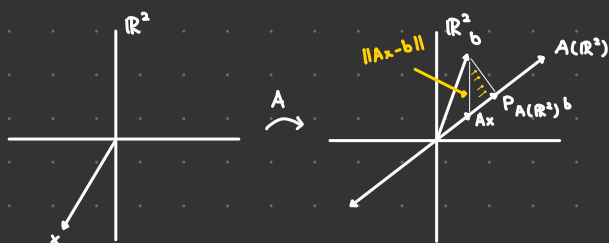
Given  $A: V \rightarrow W$  and  $b \in W$ , how do we minimize  $\|Ax - b\|$  ( $x \in V$ )?

Distance between  
 $Ax$  and  $b$

(for instance if  $A$  is a "cost" function and  $b=0$ , this is minimizing cost.)

If  $Ax=b$  has a solution  $x_0$ , then  $\|Ax-b\|$  is minimized at  $\|Ax_0-b\| = 0$ .

Otherwise ...



We suspect  $Ax = P_{A(R^2)} b$  will minimize  $\|Ax-b\|$ , and indeed:

### Proposition:

Let  $V, W$  be IPS,  $b \in W$  and  $A: V \rightarrow W$  a L.T. Then

$$\inf_{x \in V} \|Ax - b\| = \|P_{\text{range}(A)} b - b\|$$

### Proof:

Minimizing  $\|Ax - b\|$  is equivalent to minimizing  $\|Ax - b\|^2$ , and

$$\begin{aligned} \|Ax - b\|^2 &= \left\| \underbrace{Ax - P_{\text{range}(A)} b}_{\substack{Ax \in \text{range } A \\ P_{\text{range}(A)} b \in \text{range } A}} + \underbrace{P_{\text{range}(A)} b - b}_{\in (\text{range } A)^\perp} \right\|^2 \\ &= \|Ax - P_{\text{range}(A)} b\|^2 + \|P_{\text{range}(A)} b - b\|^2 \quad \text{Pythagorean thm.} \\ &\geq \|P_{\text{range}(A)} b - b\|^2 \end{aligned}$$

### Remark:

Because minimizing  $\|Ax - b\|$  is equivalent to minimizing  $\|Ax - b\|^2$ , and

$Ax = P_{\text{range}(A)} b$  is called the "least squares solution".

### Question:

How do we find  $P_{\text{range}(A)} b$ ?

### Method 1:

Find an orthogonal basis for  $\text{range}(A)$  (by GS-alg.) and use a formula for  $P_{\text{range}(A)} b$ .

### Definition:

Let  $A$  be an  $m \times n$  matrix:

$$A = \begin{bmatrix} a_{11} & \cdots & a_{1n} \\ \vdots & \ddots & \vdots \\ a_{m1} & \cdots & a_{mn} \end{bmatrix}$$

Then the Hermitian adjoint or adjoint of  $A$  is defined by the  $n \times m$  matrix:

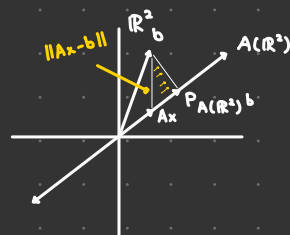
$$A^* := \begin{bmatrix} \overline{a_{11}} & \cdots & \overline{a_{m1}} \\ \vdots & \ddots & \vdots \\ \overline{a_{1n}} & \cdots & \overline{a_{mn}} \end{bmatrix}$$

**Method 2:** ( $A$  is an  $m \times n$  matrix)

$$A = \begin{bmatrix} a_{11} & \dots & a_{1n} \\ \vdots & \ddots & \vdots \\ a_{m1} & \dots & a_{mn} \end{bmatrix}$$

Since  $Ax \in \text{range } A \quad \forall x \in \mathbb{F}^n$ , we have:

$$Ax = P_{\text{range } A} b \quad \text{iff} \quad b - Ax \perp \text{range } A$$



Recall the columns of  $A$  span  $\text{range } A$ , so  $b - Ax \perp \text{range } A$  iff

$$\langle b - Ax, \begin{bmatrix} a_{11} \\ \vdots \\ a_{m1} \end{bmatrix} \rangle = \dots = \langle b - Ax, \begin{bmatrix} a_{1n} \\ \vdots \\ a_{mn} \end{bmatrix} \rangle = 0$$

and

$$\langle b - Ax, \begin{bmatrix} a_{1k} \\ \vdots \\ a_{mk} \end{bmatrix} \rangle = \underbrace{\begin{bmatrix} \overline{a_{1k}} & \dots & \overline{a_{mk}} \end{bmatrix}}_{1 \times m} \cdot \underbrace{(b - Ax)}_{\in \mathbb{F}^m}$$

So

$$\langle b - Ax, \begin{bmatrix} a_{11} \\ \vdots \\ a_{m1} \end{bmatrix} \rangle = \dots = \langle b - Ax, \begin{bmatrix} a_{1n} \\ \vdots \\ a_{mn} \end{bmatrix} \rangle = 0$$

iff

$$\begin{bmatrix} \overline{a_{1k}} & \dots & \overline{a_{mk}} \end{bmatrix} \cdot (b - Ax) = 0 \quad \text{for } 1 \leq k \leq n$$

iff

$$A^* (b - Ax) = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}$$

iff

$$\underline{A^* A x = A^* b}$$

This is called the normal equation,  
and its solution  $x$  satisfies  $Ax = P_{\text{range } A} b$   
i.e. gives the least squares solution.

If  $A^* A$  is invertible, then

$$A^* A x = A^* b \Rightarrow x = (A^* A)^{-1} A^* b$$

So  $Ax = A(A^* A)^{-1} A^* b = P_{\text{range } A} b$

Note:  $A$  is not square

Since this holds for arbitrary  $b \in \mathbb{F}^m$ ,  $P_{\text{range}(A)} = A(A^* A)^{-1} A^*$  (if  $A^* A$  is invertible)

### Theorem:

For an  $m \times n$  matrix  $A$ ,

$$\text{Ker}(A) = \text{Ker}(A^*A)$$

Proof:

If  $x \in \text{Ker}(A)$ , then  $A^*Ax = A^*0 = 0$

i.e.  $x \in \text{Ker}(A^*A)$ .

If  $x \in \text{Ker}(A^*A)$ , consider the identity

$$\begin{aligned}\|Ax\|^2 &= \langle Ax, Ax \rangle \\ &= \langle A^*Ax, x \rangle, \forall x \in \mathbb{F}^n\end{aligned}$$

$$\text{So } A^*Ax = 0 \Rightarrow \langle A^*Ax, x \rangle = 0$$

$$\Rightarrow Ax = 0$$

$$\Rightarrow x \in \text{Ker}(A)$$

□

### Corollary:

For an  $m \times n$  matrix  $A$ ,  $A^*A$  is invertible iff  $\text{rank}(A) = n$ .

Proof:

$A^*A$  is invertible iff  $\text{Ker}(A^*A) = \{0\}$

iff  $\text{Ker}(A) = 0$  By prev. thm.

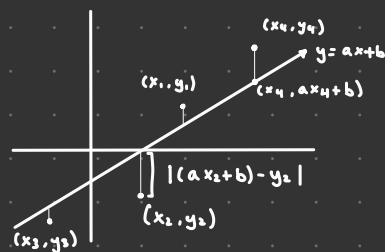
iff  $\text{rank}(A) = n$

□

Suppose we conjecture that two quantities  $y, x$  are related linearly i.e.

$y = ax + b$  for some  $a, b \in \mathbb{R}$ , and we want to find  $a, b$  from some data points

$(x_1, y_1), \dots, (x_n, y_n)$



In general, there won't be any  $a, b \in \mathbb{R}$  s.t.  $y_i = ax_i + b$  for  $1 \leq i \leq n$  (for instance due to error in our experiment which produces the data points  $(x_i, y_i)$ ), but to find the "best" line approximation, we can aim to minimize:

$$\begin{aligned} \sum_{k=1}^n |(ax_k + b) - y_k|^2 &= \left\| \begin{bmatrix} ax_1 + b - y_1 \\ \vdots \\ ax_n + b - y_n \end{bmatrix} \right\|^2 \\ &= \left\| \underbrace{\begin{bmatrix} x_1 & 1 \\ \vdots & \vdots \\ x_n & 1 \end{bmatrix}}_A \underbrace{\begin{bmatrix} a \\ b \end{bmatrix}}_x - \underbrace{\begin{bmatrix} y_1 \\ \vdots \\ y_n \end{bmatrix}}_b \right\|^2 \end{aligned}$$

Ex 1:

If our data points are  $(1, 1), (2, 3), (3, 2)$

We want to minimize

$$\begin{aligned} &= \left\| \begin{bmatrix} x_1 & 1 \\ \vdots & \vdots \\ x_n & 1 \end{bmatrix} \begin{bmatrix} a \\ b \end{bmatrix} - \begin{bmatrix} y_1 \\ \vdots \\ y_n \end{bmatrix} \right\|^2 \\ &= \left\| \begin{bmatrix} 1 & 1 \\ 2 & 1 \\ 3 & 1 \end{bmatrix} \begin{bmatrix} a \\ b \end{bmatrix} - \begin{bmatrix} 1 \\ 3 \\ 2 \end{bmatrix} \right\|^2 \end{aligned}$$

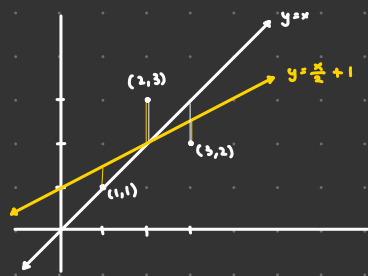
Method 2 tells us to solve  $A^*Ax = A^*b$  or

$$\begin{bmatrix} 1 & 2 & 3 \\ 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 2 & 1 \\ 3 & 1 \end{bmatrix} \begin{bmatrix} a \\ b \end{bmatrix} = \begin{bmatrix} 1 & 2 & 3 \\ 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 3 \\ 2 \end{bmatrix}$$

$$\Rightarrow \begin{bmatrix} 14a + 6b \\ 6a + 3b \end{bmatrix} = \begin{bmatrix} 13 \\ 6 \end{bmatrix}$$

$$\Rightarrow a = \frac{1}{2}, b = 1$$

Ex 1 continued...



$y = \frac{x}{2} + 1$  has error of  $(\frac{1}{2})^2 + 1^2 + (\frac{1}{2})^2 = \frac{3}{2}$

$y = x$  has error of  $0^2 + 1^2 + 1^2 = 2$