



Propositional logic 1 of 2

Definition:

Let PV be a set of propositional variables. The set of propositional formulas, denoted F (or F_{PV}), is the smallest set s.t.

Basis: if $x \in PV$, then $x \in F$

I.S.: If $P_1, P_2 \in F$, then

$$\neg P_1, (P_1 \wedge P_2), (P_1 \vee P_2), (P_1 \rightarrow P_2), (P_1 \leftrightarrow P_2) \in F$$

Ex 1:

$$((x \leftrightarrow y) \vee \neg(x \rightarrow z))$$

Parse Tree:



Definition:

A truth assignment (t.a.) is a function $\tau: PV \rightarrow \{0, 1\}$

Annotations: τ points to the function symbol, False points to 0, True points to 1.

Definition: Extended Truth Assignment

$$\tau^*: F \rightarrow \{0, 1\}$$

$$\tau^*(P) = \begin{cases} \tau(x) & \text{if } P = x \in PV \\ \tau^*(P_1) \cdot \tau^*(P_2) & \text{if } P = (P_1 \wedge P_2) \\ \max(\tau^*(P_1), \tau^*(P_2)) & \text{if } P = (P_1 \vee P_2) \end{cases}$$

Annotation: $\min(\tau^(P_1), \tau^*(P_2))$ also works*

$$\tau^*((P_1 \rightarrow P_2)) = \begin{cases} 1 & , \text{ if } \tau^*(P_1) = 0 \text{ or } \tau^*(P_2) = 1 \\ 0 & , \text{ o/w} \end{cases}$$

$$\tau^*((P_1 \leftrightarrow P_2)) = \begin{cases} 1 & , \text{ if } \tau^*(P_1) = \tau^*(P_2) \\ 0 & , \text{ o/w} \end{cases}$$

Conventions:

Precedence (high to low)

(i) \neg

(ii) \wedge

(iii) \vee

(iv) $\rightarrow, \leftrightarrow$ (right associate)

Omit outermost parentheses

Definition:

2 formulas P_1, P_2 are **logically equivalent** iff for any t.a. τ , $\tau^*(P_1) = \tau^*(P_2)$.

Notation: $P_1 \text{ LEQV } P_2$, or $(P_1 \leftrightarrow P_2)$

Definition:

A formula P_1 **logically implies** a formula P_2 iff for any t.a. τ , if $\tau^*(P_1) = 1$, then $\tau^*(P_2) = 1$.

Definition:

τ satisfies P means $\tau^*(P) = 1$

τ falsifies P means $\tau^*(P) = 0$

Definition:

P is a **tautology** (valid formula) means every t.a. satisfies P .

P is **satisfiable** means some t.a. satisfies P .

P is **unsatisfiable** (contradiction) means every t.a. falsifies P .

P is **falsifiable** means some t.a. falsifies P .

Theorem: (5.7, 5.9)

For any formula P_1, P_2 ,

(i) $P_1 \text{ LEQV } P_2$ iff $P_1 \leftrightarrow P_2$ is a tautology

(ii) P_1 logically implies P_2 iff $P_1 \rightarrow P_2$ is a tautology

Truth Tables

e.g.

| x | y | z | $(x \leftrightarrow y) \vee \neg(x \rightarrow z)$ | | |
|---|---|---|--|---|---|
| 0 | 0 | 0 | 1 | 1 | 0 |
| 0 | 0 | 1 | 1 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 1 | 1 |
| 1 | 0 | 1 | 0 | 0 | 0 |
| 1 | 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 1 | 1 | 1 | 0 |

Definition: Normal Forms (5.9)

- (i) A **literal** is a variable or the negation of a variable.
- (ii) A **term** is a literal or the conjunction (AND) of 2 or more literals.
- (iii) A **clause** is a literal or the disjunction (OR) of 2 or more literals.
- (iv) A **disjunctive normal form** (DNF) formula is a **term** or the **disjunction** of 2 or more **terms**.
- (v) A **conjunctive normal form** (CNF) formula is a **clause** or the **conjunction** of 2 or more **clauses**.