

1 (text) Dijkstra [15 points] Find the shortest path from E to all other nodes using the Dijkstra Algorithm (Show your work and intermediate steps)

From to Cost

Visited: $[E]$ $E \rightarrow E : 0 \rightarrow$ Set all to ∞ ,
 $E \rightarrow C : 4$ go to the cheapest cost (D)
 $E \rightarrow K : 5$
 $E \rightarrow D : 1$

then $[E, D]$

$ED \rightarrow H : 3$
 $ED \rightarrow C : 3$ choose C

$\Rightarrow [E, D, C]$

$EDC \rightarrow B : 4 \rightarrow [E, D, C, B]$
 $EDCB \rightarrow A : 1 \Rightarrow [E, D, C, B, A]$

\Rightarrow Choose $H \rightarrow [E, P, H]$

$EDH \rightarrow G : 8$ choose G
 $EDH \rightarrow M : 8$

$\Rightarrow EDCBA \rightarrow P : 15$

$EDCBA \rightarrow G : 13$

$EDHG \rightarrow P : 11$
 $EDHG \rightarrow O : 9$

From all this, the shortest path is:

$E \rightarrow D : 1$
 $E \rightarrow C : 3$ (E, D, C)
 $E \rightarrow K : 5$ (E, D, C, K)

$EDHG \rightarrow O : 9$ Done

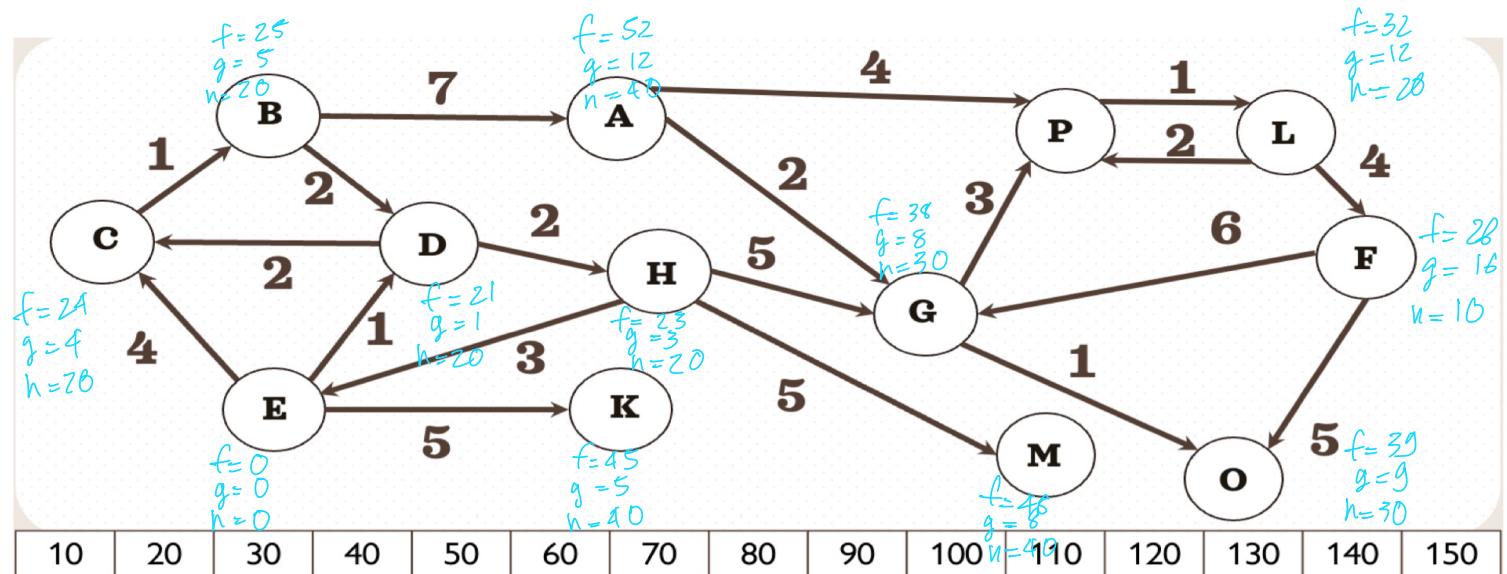
$E \rightarrow H : 3$ EDH
 $E \rightarrow B : 4$ $EDCB$
 $E \rightarrow G : 8$ $EDHG$
 $E \rightarrow M : 8$ $EDHGM$
 $E \rightarrow P : 11$ $EDHGP$
 $E \rightarrow O : 9$ $EDHGO$
 $E \rightarrow L : 12$ $EDHEPL$
 $E \rightarrow F : 16$ $EDHEPLF$
 $E \rightarrow A : 11$ $EDCBA$

$[EDHGP]$
 $EDHEP \rightarrow L : 12 \Rightarrow [EDHEPL]$

$EDHEPL \rightarrow F : 16$

$\Rightarrow [EDHGPLF]$

2 (text) A* [20 points] Find the shortest path from E to all other nodes using the A* Algorithm (Show your work and intermediate steps) Hint: Use the geographical location values to derive the heuristic to use.



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3(text) Comparison

We can see that Dijkstra and A* both locate node "H" from E. They both take the same # of iterations, even though A* is known to be faster. This isn't the case since they run the same amount of time.

7. Text

Time Complexity for Graph Direct check is $O(V^3)$ because it takes $n \times n$ matrix, which we can look at how it looks in picture

$$n \begin{bmatrix} \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot \end{bmatrix} \Rightarrow n \times n \Rightarrow O(V^2)$$

Space is also $O(V^2)$ because it needs to save it down as a matrix $n \times n$

For Path Finder:

Time Complex: $O(V+E)$ where n is # of vertices. Because that we use DFS to explore and worst case of DFS is $O(V+E)$

Space also $O(V+E)$ since path list take up V space and visit set would store all the V and they all visited E in total.