Algorithm HW2 PS1

Chien-I, Tseng

B11901072

1 Problem Statement

This program solves the Maximum Planar Subset problem by finding the maximum number of non-intersecting chords in a circle and listing the chords in the solution.

2 Algorithm Design

2.1 Dynamic Programming Approach

2.1.1 States Definition

- 1. dp[i][j] represents the maximum number of the non-intersecting chords in the region between i and j.
- 2. path[i][j] represents the endpoints of the chords chosen for the solution.

2.1.2 Base Cases

- 1. when $i \ge j$, dp[i][j] = 0
- 2. if dp[i][j] has calculated, return the value

2.1.3 Recursion

case1: if there is a chord starting at i and ending at "end", dp[i][j] = 1 + dp[i+1][end-1] + dp[end+1][j] case2: skip i, dp[i][j] = dp[i+1][j]

chose the maximum value above these two cases.

2.2 Space Complexity and Time Complexity

We have $O(n^2)$ dp table and $O(n^2)$ path table to fill, so the space complexity is $O(n^2)$. For each space, there are O(n) possible chords to consider

3 Data Structure

3.1 Chord

Written in the file ./src/maxPlanarSubset.h

```
7 struct Chord {
8    int start, end;
9    Chord(int s, int e) : start(s), end(e) {}
10 };
```

3.2 maxPlanarSubset

4 Function

4.1 Solve

Written in the file ./src/maxPlanarSubset.cpp

This function contains the main logic of dynamic programming and store the endpoints in path[i][j] for the solution reconstruction.

4.2 Solution Reconstruction

According to path[i][j], we can reconstruct the solution chord set.

```
20 void MaxPlanarSubset::reconstructSolution(int i, int j, std::vector<Chord>& result) {
21    if (i >= j) return;
22
23    if (path[i][j] == -1) {
24        reconstructSolution(i + 1, j, result);
25    } else {
26        result.push_back(Chord(i, path[i][j]));
27        reconstructSolution(i + 1, path[i][j] - 1, result);
28        reconstructSolution(path[i][j] + 1, j, result);
29    }
30 }
```