

# The University of Melbourne

## Practice Exam Paper

School of Computing and Information Systems

COMP90038 Algorithms and Complexity

Leo's solution

### Question 1

A. Give the names of two *stable* sorting algorithms, together with their worst-case time complexities. Write the names and complexities in the box:

1. Insertion sort,  $\Theta(n^2)$
2. Merge sort,  $\Theta(n \log n)$

B. Give the names of two *unstable* sorting algorithms, together with their worst-case time complexities. Write the names and complexities in the box:

1. Selection sort,  $\Theta(n^2)$
2. Quick sort,  $\Theta(n^2)$

### Question 2

(4 marks)

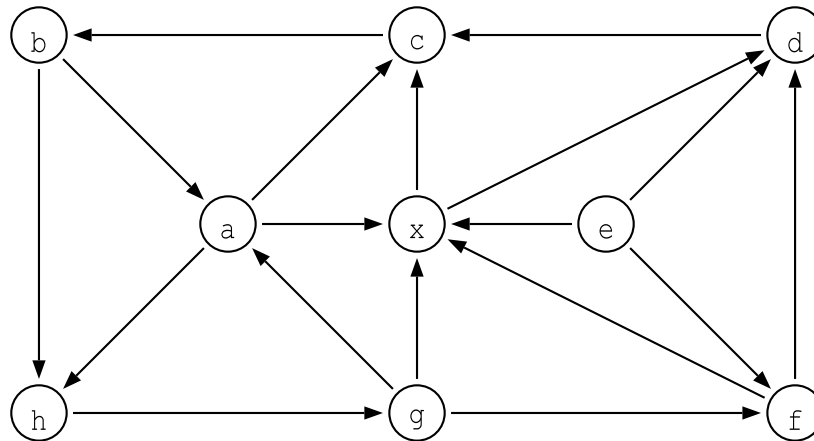
We are given an array  $A$  holding  $n$  integers, for some large  $n$ . **The array is sorted**, and the values in  $A$  range from -2147483648 to 2147483647, evenly distributed. Give  $\Theta$  expressions for the following tasks:

- A. Running the insertion sort algorithm on the array  $A$ :  $\Theta(n)$
- B. Running the selection sort algorithm on the array  $A$ :  $\Theta(n^2)$
- C. Performing binary search for integer  $k$  which is not in  $A$ :  $\Theta(\log n)$
- D. Performing interpolation search for integer  $k$  not in  $A$ :  $\Theta(1)$

### Question 3

For the directed graph below, list the order in which the nine nodes are visited during a depth-first (DFS) traversal, as well as the order in which they are visited during a breadthfirst (BFS) traversal. As always, assume that any **ties are resolved** by taking nodes in alphabetical order. Write the answers in the boxes given.

(4 marks)



DFS sequence:

a c b h g f d x e

BFS sequence:

a c h x b g d f e

## Question 4

(4 marks)

Given the pattern A T G A and the text

T C A T C A T C C A T G C A C A A T G A C T T T

how many character comparisons will Horspool's algorithm make before locating the pattern in the text? Write the number in the box:

$$1+2+1+1+2+4(\text{successful searching}) = 11$$

## Question 5

Assume the array  $A$  holds the keys 77, 64, 15, 43, 28, 91, 80, 32, 56 in index positions 1 to 9. Show the heap that results after application of the linear-time bottom-up heap construction algorithm. You may show the heap as a tree or as an array.

[ 91 64 80 56 28 15 77 32 43 ]

(4 marks)

## Question 6

The functions  $A$ – $D$  are defined recursively as follows (all divisions round down, to the closest integer):

$$A(n) = 2 A(n/3) + 2, \quad \text{with } A(1) = 1$$

$$B(n) = B(n/2) + n/2, \quad \text{with } B(1) = 1$$

$$C(n) = 512 C(n/8) + 4n^2, \quad \text{with } C(1) = 4$$

$$D(n) = 4 D(n/2) + n^2, \quad \text{with } D(1) = 2$$

In the following table, for each of the four functions, tick the most precise correct statement about the function's rate of growth:

	$O(n)$	$\Theta(n)$	$O(n \log n)$	$\Theta(n^2)$	$O(n^2 \log n)$	$\Theta(n^2 \sqrt{n})$	$O(n^3)$
$A$	✓						
$B$		✓					
$C$							✓
$D$					✓		

## Question 7

For each of **A–D** below, answer yes or no, and, in each case, briefly explain your reasoning (just a justification of your answer, rather than detailed calculations). A yes/no answer that is not justified will not attract marks, even if correct.

Question	Answer/explanation
<b>A.</b>  Is $\sqrt{n} \in \Omega(n)$ ?	No, because $n/(\text{root } n) = \text{root } n$

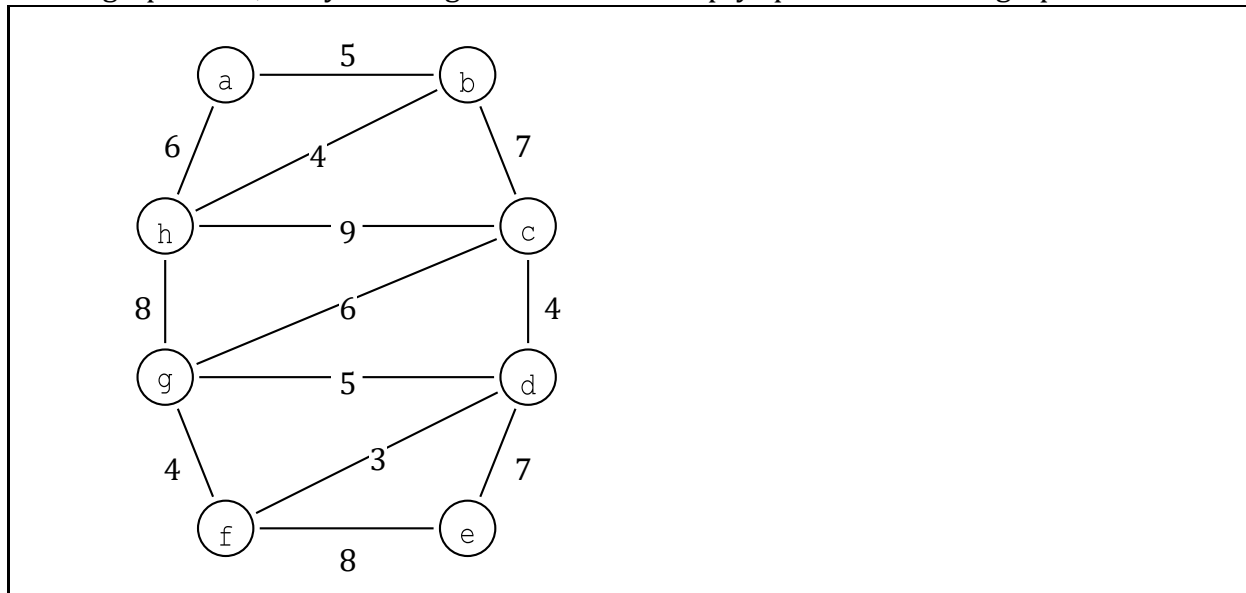
(4 marks)

<p><b>B.</b></p> <p>Is <math>n^{(2\log n)} \in O(n^{(\log n)})</math>?</p>	<p>No, because</p> $\frac{n^{(2\log n)}}{n^{(\log n)}} = n^{(\log n)}$
<p><b>C.</b></p> <p>Is <math>\Theta(\log(n^{2\log n})) = \Theta(\log(n^{\log n}))</math>?</p>	<p>Yes, because</p> $\frac{\log(n^{2\log n})}{\log(n^{\log n})} = \frac{2\log n}{\log n} = 2$
<p><b>D.</b></p> <p>Is <math>\Theta(\log(2^n)) = \Theta(\log(3^n))</math>?</p>	<p>Yes, because</p> $\frac{\log(2^n)}{\log(3^n)} = \frac{(\log(3^n) * \log(2)(3))}{\log(3^n)} = \log(2)(3)$

## Question 8

(6 marks)

A. The box below contains a weighted undirected graph with eight nodes. Give a minimum spanning tree for the graph. You may do that either by outlining a minimum spanning tree on the graph itself, or by drawing the tree in the empty space next to the graph.



These edges were chosen: a-b, b-h, b-c, c-d, d-f, f-g, d-e

B. Given a weighted graph  $G = \langle V, E \rangle$ , a subgraph  $\langle V, E' \rangle$  (that is,  $E' \subseteq E$ ) which is a tree with *minimum* weight is a *maximum spanning tree* for  $G$ .

We want a transformation of the graph  $G$  so that we can run Prim's algorithm on the transformed graph  $G'$ , and the algorithm will find a maximum spanning tree for  $G$ .

Describe such a (systematic) transformation from  $G$  to  $G'$ .

(I think the red word should be "**maximum**", otherwise this question is unreasonable.)

For each edge in  $G$ , we change the weight from  $(x)$  to  $(-x)$ . For example, the weight of "a to b" is 5 originally, now we let it become -5.

Then we can run Prim's algorithm to find the maximum tree.

## Question 9

(6 marks)

Consider the function  $F$  below. The function takes as input an integer array  $A$ , and the size  $n$  of  $A$ . The array indices run from 1 to  $n$ . The division used is integer division, that is, it rounds down to the closest smaller (or equal) integer value.

In the box, give a  $\Theta$  expression for the function's time complexity.

**function**  $F(A[\cdot], n)$

$s \leftarrow 0$

$m \leftarrow n$

**while**  $m > 0$  **do**

**for**  $i \leftarrow 1$  **to**  $m$  **do**

$s \leftarrow s + A[i]$

$m \leftarrow m/2$

$$n + (n/2) + (n/4) + \dots + 1 = 2n - 1 = \Theta(n)$$

## Question 10

Using pseudo-code, give an algorithm for deleting the smallest element of a binary search tree (a BST). Assume a non-empty binary tree  $T$  has attributes  $T.left$ ,  $T.right$ , and  $T.root$  which denote  $T$ 's left sub-tree, right sub-tree, and the key of  $T$ 's root node, respectively. You can use these tests if they seem useful:  $\text{isLeaf}(T)$  tests whether the binary tree  $T$  is a leaf, and  $\text{isEmpty}(T)$  tests whether it is empty.

**Function**  $\text{deleteSmallest}(T)$

**if**  $\text{isEmpty}(T.left)$

$T = T.right$       *//Delete T*

**else**

$\text{deleteSmallest}(T.left)$

(10 marks)

## Question 11

Consider an array  $A$  of  $n$  distinct integers (that is, all elements are different). It is known that  $A$  was originally sorted in ascending order, but  $A$  was then right-rotated  $r$  places, where  $0 < r < n$ . In other words, the last  $r$  elements were moved from the end of the array to the beginning, with all other elements being pushed  $r$  positions to the right. For example, for  $n = 7$  and  $r = 3$ , the result may look like this:

[43,46,58,12,20,29,34]

For  $r = 5$ , the result, based on the same original array, would be

[29,34,43,46,58,12,20]

You know that the given  $A[0..n-1]$  has this particular form, that is, for some  $r$ , the sequence  $A[r], \dots, A[n-1], A[0], \dots, A[r-1]$  is in ascending order, but you do not know what  $r$  is. Design an algorithm to find the largest integer in  $A$ . Full marks are given for an algorithm that works in time  $O(\log n)$ ; half marks are given for a solution that is correct, but less efficient.

```
Function findMax(A[])
    return find(A, 0, A.length-1)
```

```
Function find(A[], start, end)
    mid = (start+end)/2
    if A[mid]>A[mid+1]
        return A[mid]
    if A[mid]>A[0]
        return find(A, mid+1, A.length-1)
    return find(A, 0, mid-1)
```

(10 marks)

## Question 12

Two programmers face the following problem. Given an array containing  $n$  random integers in random order, find the largest integer. The integers are placed in cells  $A[1] \dots A[n]$ .

Programmer X has come up with the code shown below, on the left. (In the programming language used, arrays are indexed from 0, but X's method does not use  $A[0]$ .)

```
function X( $A[\cdot], n$ )
   $max \leftarrow A[1]$ 
   $i \leftarrow 2$ 
  while  $i \leq n$  do
    if  $A[i] > max$  then
       $max \leftarrow A[i]$ 
     $i \leftarrow i + 1$ 
  return  $max$ 
```

```
function Y( $A[\cdot], n$ )
   $i \leftarrow n$ 
  while  $i > 0$  do
     $A[0] \leftarrow A[i]$ 
     $i \leftarrow i - 1$ 
    while  $A[0] > A[i]$  do
       $i \leftarrow i - 1$ 
  return  $A[0]$ 
```

Programmer Y has solved the same problem differently, as shown above on the right.

Compare the two solutions using three criteria: Correctness, time complexity class, and the number of comparisons performed. Write your analysis in the box:

	Left one	Right one
Correctness	Correct	Correct
Time complexity	$O(n)$	$O(n)$
Number of comparisons	$(n-1) + n$	$n + x$ (where $x$ may be 2 to $n+1$ )

\*The green part is for “if  $A[i] > max$ ” (left) or “while  $A[0] > A[i]$ ” (right).

\*The red part is the number of comparisons for “while  $i \leq n$ ” (left) or “while  $i > 0$ ” (right).