Data-Driven Night

Thessaloniki Java Meetup Group

GENERAL

What is C* (I)

- Fast Distributed DB
- High Availability
- Near-Linear Horizontal Scalability

C*

What is C* (II)

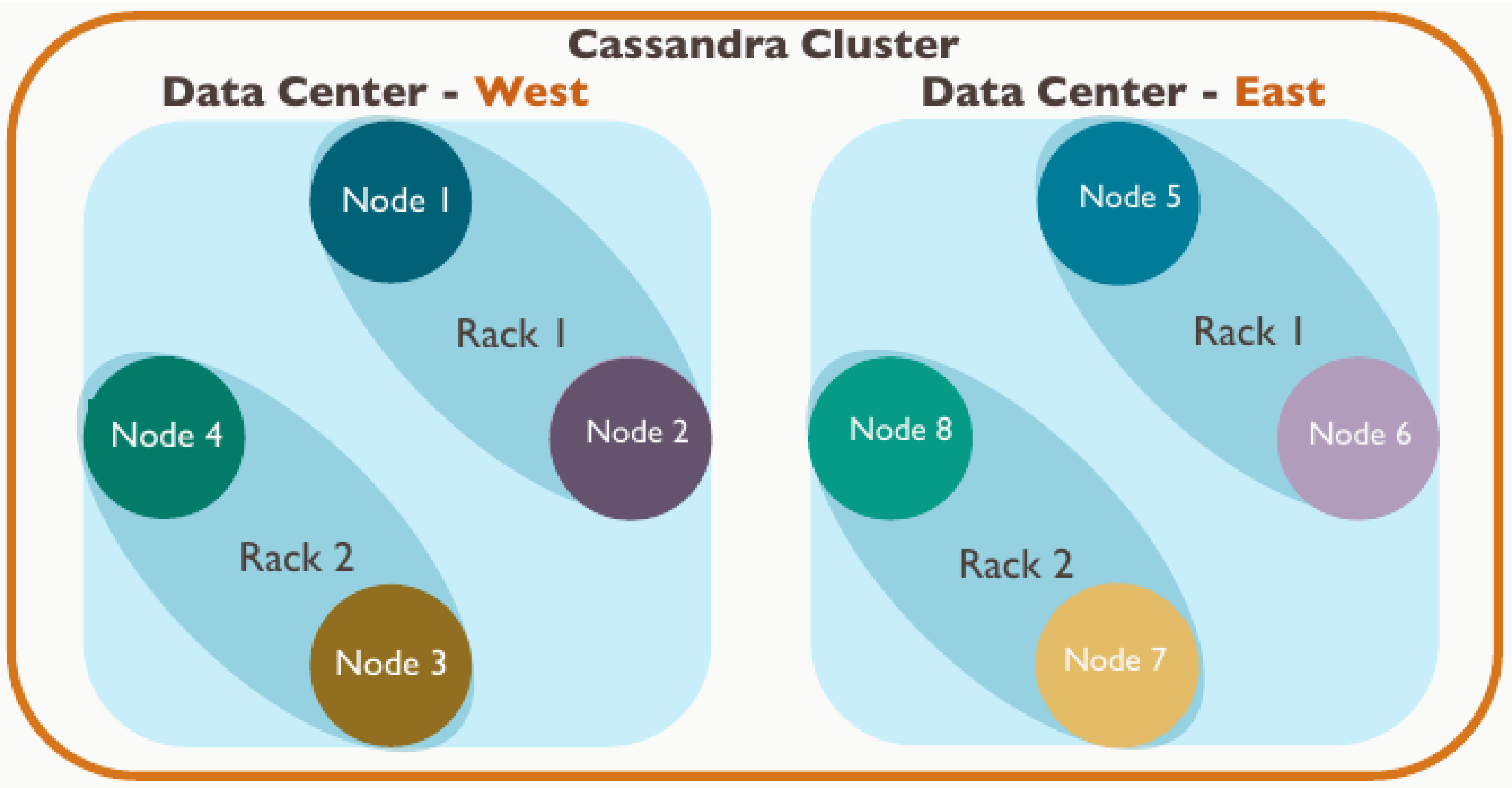
- Predictable Performance
- Fault Tolerance (P2P)
- Cannot replace RDBMS ad hoc
  → Data Model is different!

  → Transactions
    - Kind of native in RDBMS
    - C* lightweight transaction mechanism exist but should be generally avoided
    - Transactions can be implemented in the Application

Origins

- Google Big Table
  - Storage Model
- Amazon Dynamo
  - Distribution backbone
- Facebook integrated these two (2008)
  - Later released as Cassandra
  - Nowadays an Apache project

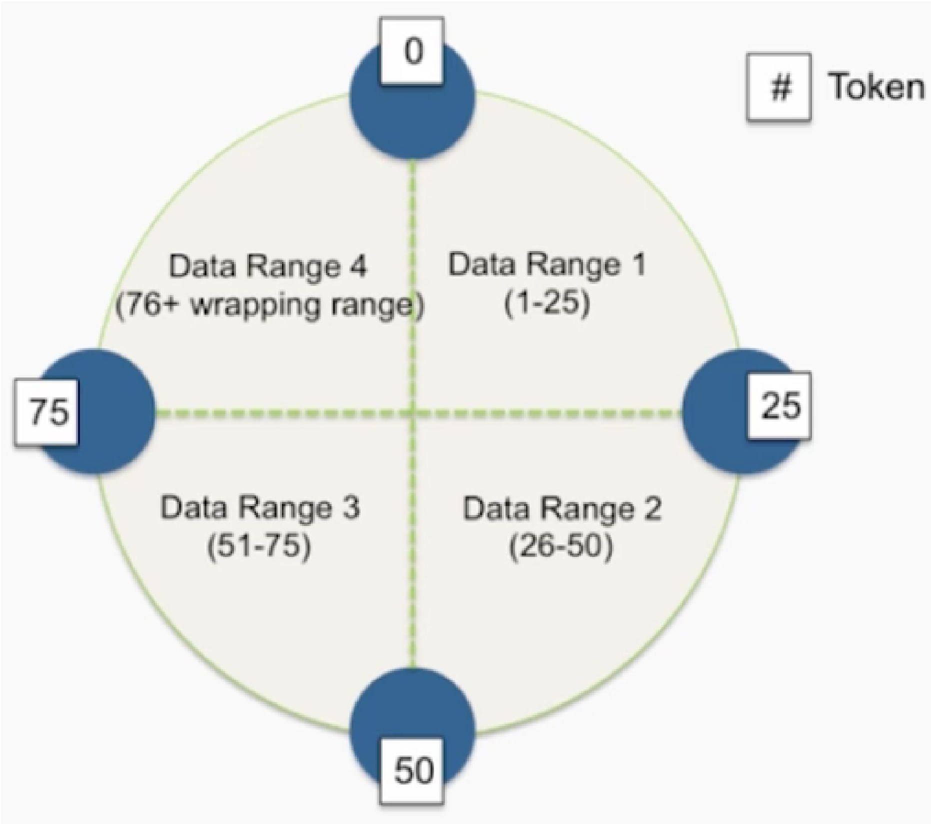
Structure (I)



Cassandra Cluster
Data Center - West
Data Center - East
Node 1
Rack 1
Node 2
Node 4
Rack 2
Node 3
Node 5
Rack 1
Node 6
Node 8
Rack 2
Node 7

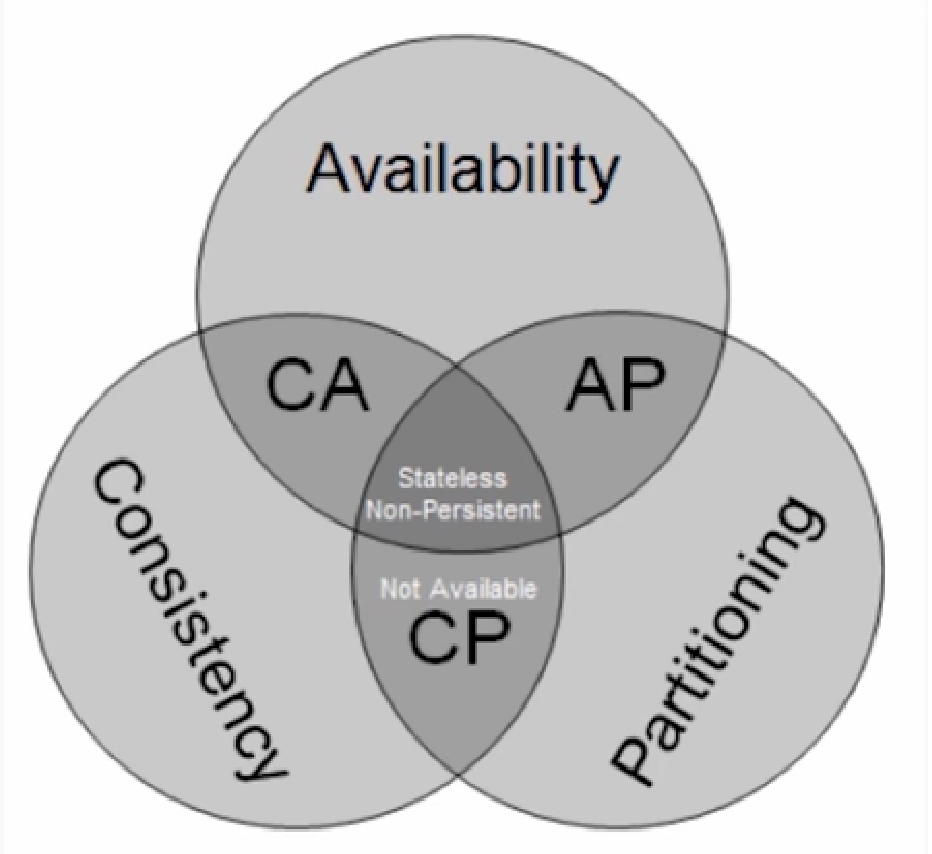
Structure (II)

- Hash Ring
- P2P
- Data partitioning
- Replication across peers

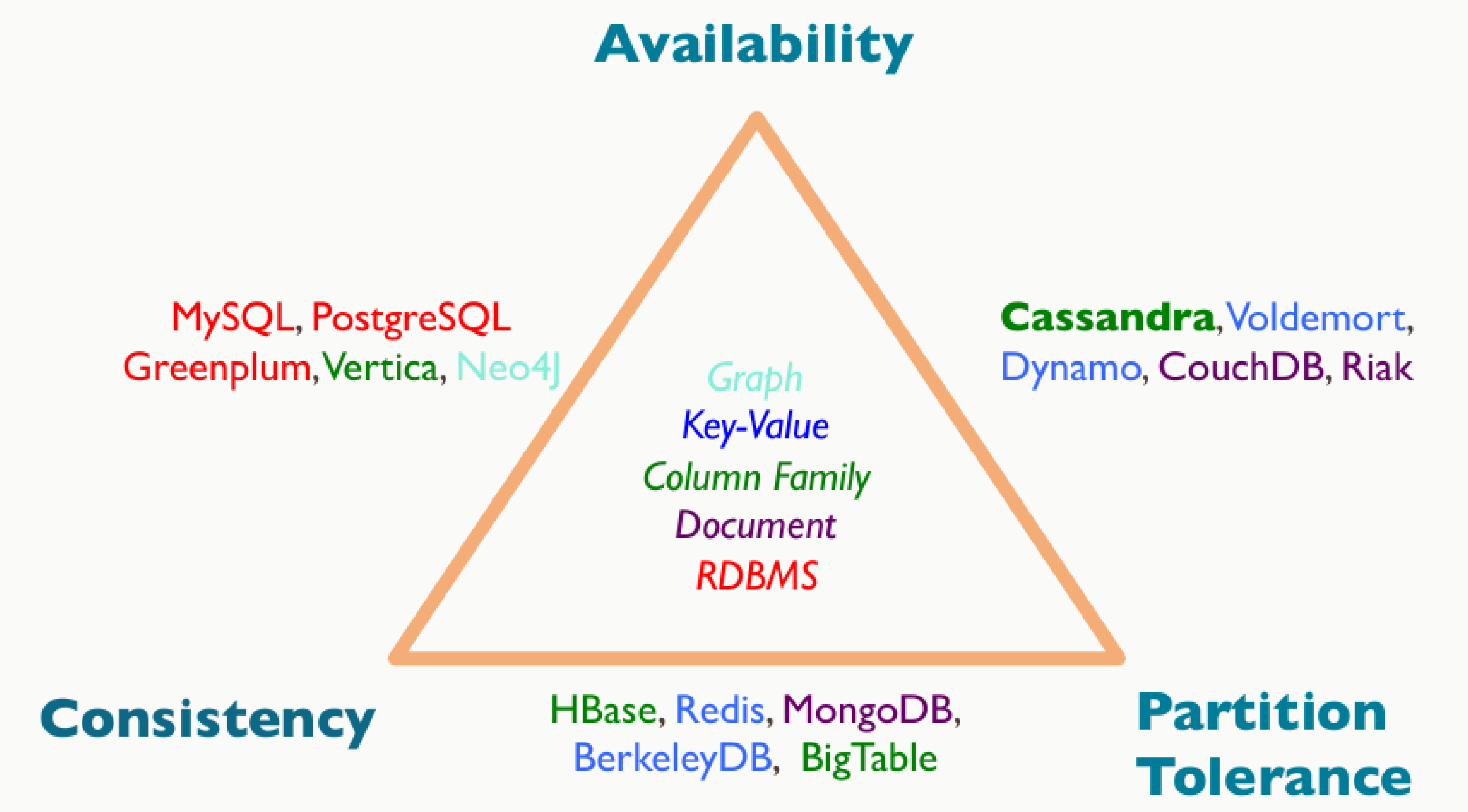


CAP theorem (I)

- Consistency Availability Partitioning trade-off
- Partitioning = Partition Tolerance = same network or not
- C* design choice:
  - A, P over C



CAP theorem (II)



Common Use Cases

- From:
  http://www.planetcassandra.org/apache-cassandra-use-cases/
- Product Catalog and Playlist
- Recommendation and Personalization
- Fraud Detection
- Messaging
- IOT and Sensor Data
- Marketing and Advertising
- Social Media and Networking

INTERNALS

Replication Factor

- Definition “how many copies of our data, do exist in a cluster” (RF)
- Data is always replicated
- RF is defined and configured for a KeySpace per Data Center
- A KeySpace is a “collection of Tables”

Multiple Data Centers

- DCs can be physical or logical
- Asynchronous replication to other Dcs

- CREATE KEYSPACE hospital
  WITH REPLICATION = {
      'class' : 'SimpleStrategy',
      'replication_factor' : 3
  };

Consistency Level

- Definition “How many replicas respond Properly to a query” in order to consider the query successful
  - A query can be a Read or a Write
- Examples: ALL, QUORUM, ONE
- Consistency Level (CL) affects performance and availability (fault-tolerance)
- CL is configured per query
  - This enables using C* even in **CA**P mode

Consistency Level Details

- Several are available
- Defined per request, by default ONE

Name	Description	Usage
ANY (writes only)	Write to any node, and store *hinted handoff* if all nodes are down.	Highest availability and lowest consistency (writes)
ALL	Check all nodes. Fail if any is down.	Highest consistency and lowest availability
ONE (TWO,THREE)	Check closest node to coordinator.	Highest availability and lowest consistency (reads)
QUORUM	Check quorum of available nodes.	Balanced consistency and availability

C*

Consistency Level Trade-Off

- Consistency Level ALL
  - **Consistent Read**,
    Highest latency, Lowest availability
- Consistency Level ONE
  - Maybe inconsistent Read,
    Lowest latency, **Highest availability**
- Consistency Level QUORUM
  - Consistent Read (if both Read/Write are QUORUM),
    Medium latency, Medium availability

Immediate Consistency (I)

- Immediate Consistency
  - Reads always return the most recent data
- We achieve this by configuring
  - CL per Read, Write
  - RF per KeySpace
- It must hold: $CL_{Read} + CL_{Write} > RF$

- Practically, does it worth it?
  - CL ONE is enough in most cases

Immediate Consistency (II)

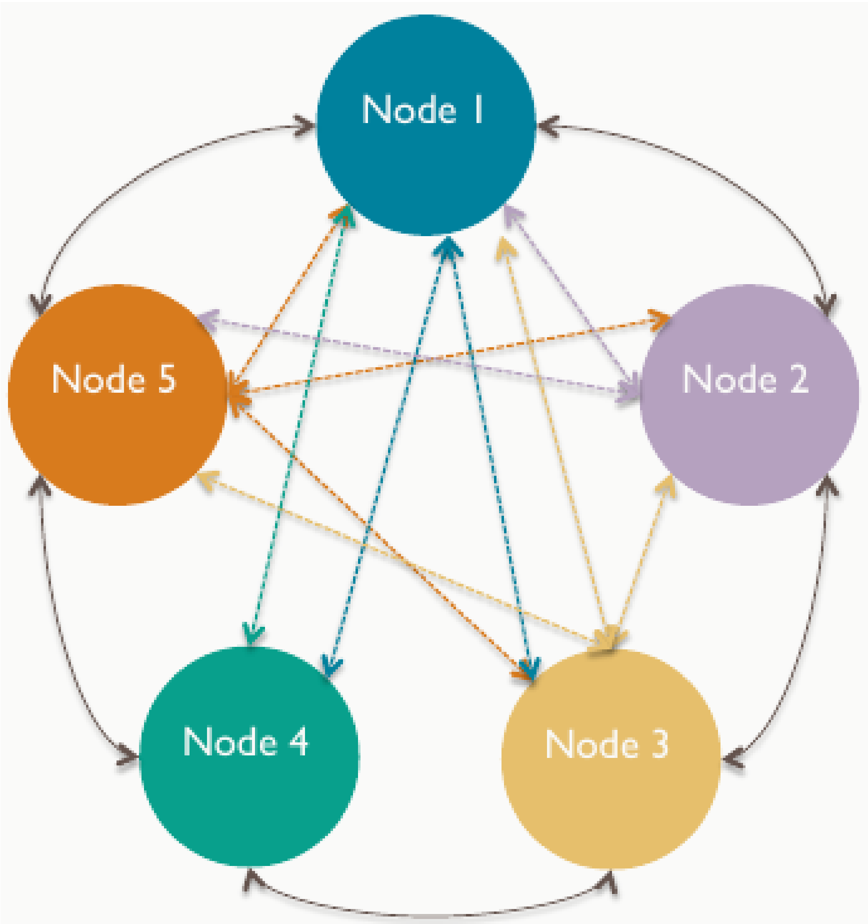
- Configuration examples for a Cluster with 4 Nodes
- Frequent Read operations:
  - RF = 3
  - CL_{Read} = QUORUM, CL_{Write} = QUORUM
- Frequent Write operations:
  - RF = 3
  - CL_{Read} = ALL, CL_{Write} = ONE

Cluster internal communication

- Nodes continuously communicate and exchange information
- Two central mechanisms
  - Gossip
  - Snitch

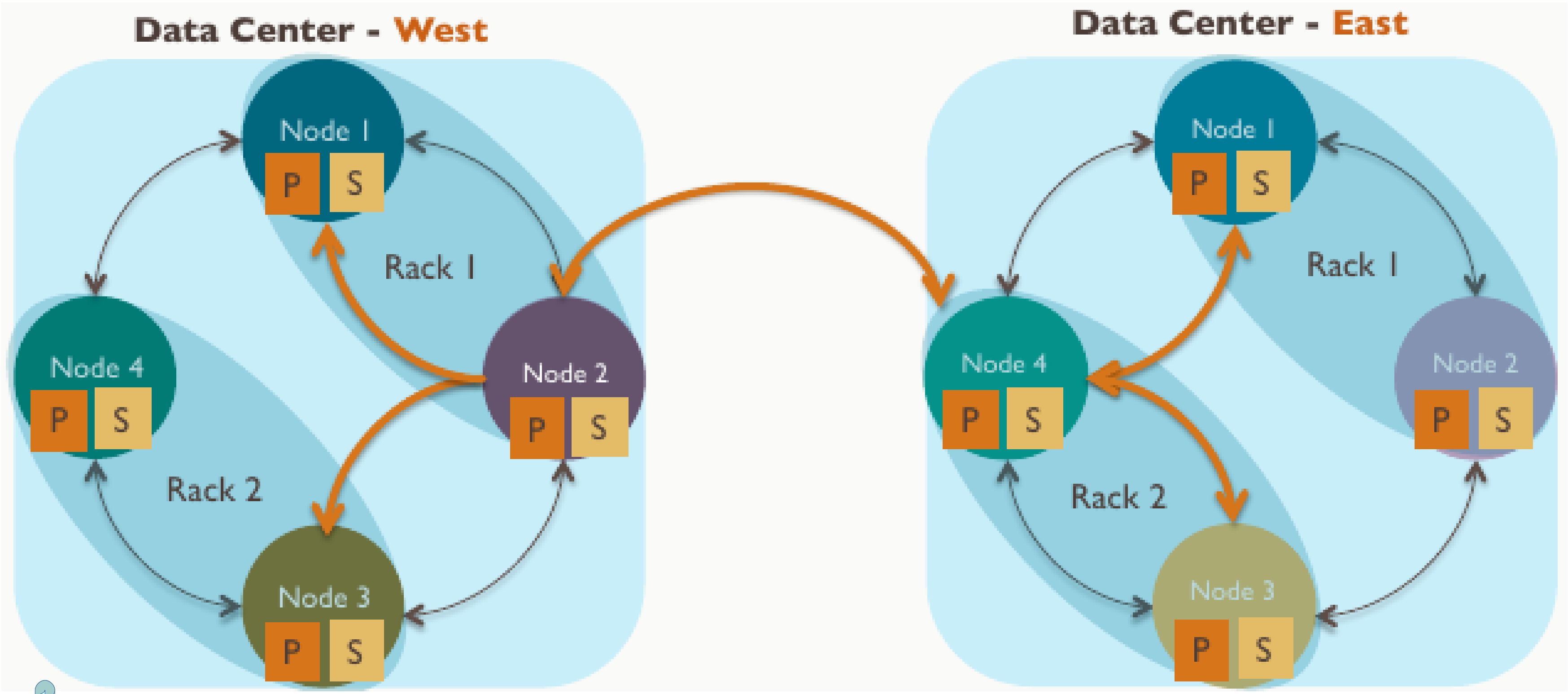
Gossip

- Every one second, each Node contacts 1 to 3 others, sending and requesting timestamped updates about known Nodes
  - states
  - locations



Snitch

- This is how Nodes know about the rack and data center topology



CONFIGURATION FILES & TOOLS

Installation

- Requirements for CPU, RAM, HDD
- Operating System
- NTP – C* requires synchronized clocks
- Disable memory swaps
- Java: Oracle JDK
- Network configuration
- C* installation
- C* configuration

Distributions

- Apache Cassandra
- DataStax Community Edition (DSC)
  - Additional tools for managing a Cluster
- DataStax Enterprise Edition (DSE)
  - More features than DSC, better for Analytics
  - Special program for start-ups
- www.planetcassandra.org/cassandra/

Configuration Files

- Located under **$CASSANDRA_HOME/conf/**
  - Example: dsc-cassandra-2.1.10/conf/
- Most important files:
  - cassandra.yaml
  - cassandra-env.sh
  - logback.xml
  - cassandra-rackdc.properties
  - cassandra-topology.properties

C* Tools

- Located under
  - **$CASSANDRA_HOME/bin/**
  - **$CASSANDRA_HOME/tools/**
- Tools
  - nodetool
  - cqlsh
  - cassandra-stress
  - sstable2json, json2sstable
  - Cassandra Cluster Management – CCM (DataStax)
  - DevCenter (DataStax)

C* DATA MODEL & CQL

C* Data Model (I)

- Data is stored and organized in a Column Family
- A Column Family is comprised of Rows
- A Row is the smallest unit that stores related data

C* Data Model (II)

- A Partition (old name: RowKey) uniquely identifies a Row in a Column Family
  - It stores data in Cells
  - Cell parts
    - column name
    - column value
    - data creation timestamp
  - Maximum cell size (column value)
    - 2 GB in theory
    - 100MB in praxis

C* Data Model (III)

- A Table is a 2D view of a column family
  - A table has Partitions
  - A Partition may be a single row or multiple rows
- A Partition Key uniquely identifies a Partition
  - Can be composite
  - It is hashed by the partitioner system to determine which Node will store it

C* Data Model (IV)

- A Primary Key uniquely identifies a row
  - Can be composite
  - It is comprised of two parts
    - the Partition Key
    - optionally, further columns
- Data Definition Language (DDL) describes Tables, Partition Keys, Primary Keys

Data in Clustering Columns (I)

- For table Videos below:

```
CREATE TABLE Videos (
   id INT,  name TEXT,  year int,  runtime int,
   PRIMARY KEY ((year), name)
  );
```

year	name	id	runtime
2015	Insurgent	1	119
2014	Interstellar	2	98
2014	Mockingjay	3	122

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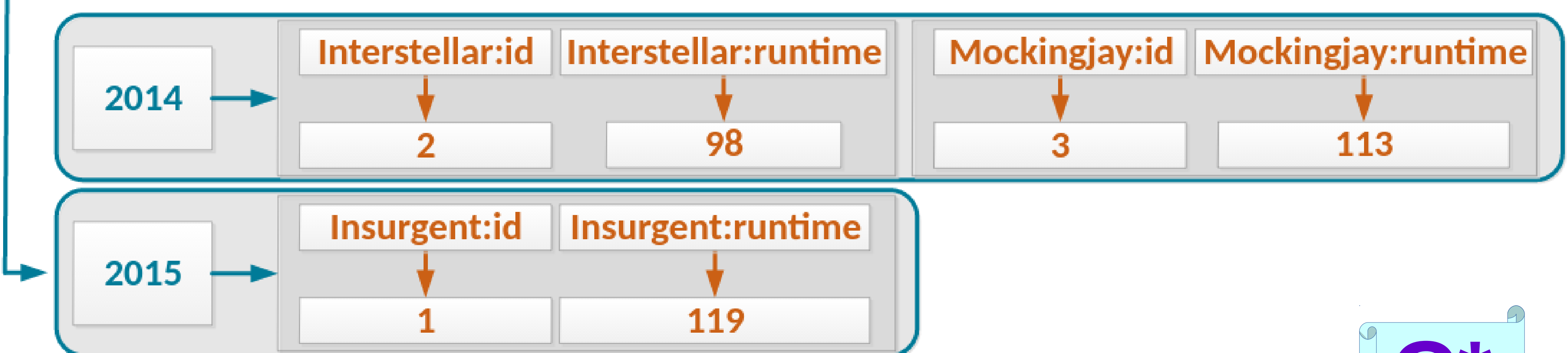
Data in Clustering Columns (II)

- Clustering columns divide Rows among partitions

year	name	id	runtime
2015	Insurgent	1	119
2014	Interstellar	2	98
2014	Mockingjay	3	122



Cassandra Query Language

- Language for communicating with the C* DB
- Abbreviated as CQL
- Similar to SQL
- It can create, modify, delete tables and data

CQL Basic Data Types

CQL Type	Constants	Description
ASCII	strings	US-ASCII character string
BIGINT	integers	64-bit signed long
BLOB	blobs	Arbitrary bytes (no validation), expressed as hexadecimal
BOOLEAN	booleans	true or false
COUNTER	integers	Distributed counter value (64-bit long)
DECIMAL	integers, floats	Variable-precision decimal
DOUBLE	integers	64-bit IEEE-754 floating point
FLOAT	integers, floats	32-bit IEEE-754 floating point
INET	strings	IP address string in IPv4 or IPv6 format*
INT	integers	32-bit signed integer
LIST	n/a	A collection of one or more ordered elements
MAP	n/a	A JSON-style array of literals: { literal : literal, literal : literal ... }
SET	n/a	A collection of one or more elements
TEXT	strings	UTF-8 encoded string
TIMESTAMP	integers, strings	Date plus time, encoded as 8 bytes since epoch
TUPLE	n/a	Up to 32k fields
UUID	uuids	A UUID in standard UUID format
TIMEUUID	uuids	Type 1 UUID only (CQL 3)
VARCHAR	strings	UTF-8 encoded string
VARINT	integers	Arbitrary-precision integer

C*

CQL: Create Table

```
CREATE TABLE cars_by_cost (
    brand TEXT,          # part of Partition Key
    model TEXT,          # part of Partition Key
    cost DECIMAL,    # Clustering Key
    merchant TEXT,
    PRIMARY KEY ((brand, model), cost)
) WITH CLUSTERING ORDER BY (cost ASC);
```

CQL: Modify Table

- ALTER TABLE cars_by_cost ADD cc INT;

- ALTER TABLE cars_by_cost
  ALTER cc TYPE BIGINT;

  - Types must be compatible

- ALTER TABLE cars_by_cost DROP cc;

CQL: Remove or Empty Table

- To fully remove a Table:

  DROP TABLE cars_by_cost;

- To clear all data (delete all Partitions) from a Table – but spare the Table:

  TRUNCATE cars_by_cost;

CQL: Read Data (I)

- General syntax:

  SELECT columns
  FROM table
  WHERE relations
  ORDER BY clustering_column ASC/DESC
  LIMIT number
  ALLOW FILTERING;

CQL: Read Data (II)

- Typical cases
  - Beware, these examples include Anti-Patterns!
- SELECT brand, merchant
  FROM cars_by_cost;
  - Avoid retrieving all partitions and rows unless absolutely necessary
- SELECT *
  FROM cars_by_cost;
  - Avoid retrieving all columns unless necessary

CQL: Read Data (III)

- SELECT * FROM cars_by_cost
  WHERE brand = “b” AND model = “m”;
  - To retrieve a partition, values for **all** partition columns are needed
- SELECT * FROM cars_by_cost
  WHERE brand = “b” AND
  model = “m” AND cost < 1000;
  - To retrieve a row, values for **all** partition and clustering columns (primary key) are needed

CQL: Read Data (IV)

- Secondary Indexes allow to query normal columns
  - Their usage is NOT a spontaneous decision, but a well thought one
- **CREATE INDEX** merchant_idx ON cars_by_cost (**merchant**);
  - SELECT * FROM cars_by_cost WHERE **merchant = “m”**;
  - SELECT * FROM cars_by_cost WHERE brand = “b” AND **merchant = “m”**;
- **DROP INDEX** merchant_idx;

CQL: Read Data (V)

- Trick of ALLOW FILTERING
  - Allows scanning over all partitions and the predicate needs not give values for all partition columns
  - May lead to slow queries with large result set

CQL: Additional Functions

- Aggregation related
  - count(), min(), max(), sum(), avg(), ...
- Time related
  - now(), dateof(), …
- Blob conversion related
  - bigintAsBlob, blobAsBigint, …
- User Defined Functions are also possible!
  - To be executed within C*, thus written in Java

CQL: Create (Insert) Data

- INSERT INTO cars_by_cost (brand, model, cost, merchant) VALUES (“volvo”, “xc90”, 9999, “daves”);

- What does it do?
  - Creates non-existing partitions
  - **But also updates existing partitions**

CQL: Update Data

- UPDATE cars_by_cost
  SET merchant = “pauls”
  WHERE brand = “volvo”
  AND model = “xc90”
  AND cost = 9999;

- What does it do?
  - Updates existing partitions
  - **But also creates non-existing partitions**

CQL: Upsert Data

- Insert and Update have the notion of Upsert
  - Update or Insert

- Why?
  - Because of the way data is organized into Clustering columns

CQL: Delete Data (I)

- Deleting a Partition

DELETE FROM cars_by_cost
WHERE brand = “b” AND model = “m”;

- Deleting a Row

DELETE FROM cars_by_cost
WHERE brand = “b”
AND model = “m”
AND cost = 1000;

CQL: Delete Data (II)

- Deleting (setting to NULL) a cell from a Row

  DELETE merchant FROM cars_by_cost
  WHERE brand = “b”
  AND model = “m”
  AND cost = 1000;

- To clear all data (delete all Partitions) from a Table – but spare the Table:

  TRUNCATE cars_by_cost;

ACID & TRANSACTIONS

ACID

- Not in the usual RDBMS sense
- Atomicity
  - Per Partition
- Consistency
  - Configurable via CL
- Isolation
  - Per Partition
- Durability
  - Write operations are indeed persisted

Lightweight Transactions

- Two ways to accomplish, as
  - Compare-And-Set (CAS) operations
  - Batch Statements
- Both affect performance
  - Negative impact

Compare-And-Set Ops. (I)

- It performs a Read operation, checks a Condition, and if that one holds, proceeds with the Write operation
  - All atomically

Compare-And-Set Ops. (II)

- INSERT INTO
  cars_by_cost (brand, model, cost, merchant)
  VALUES ("volvo", "xc90", 9999, "daves")
  IF NOT EXISTS;

- UPDATE cars_by_cost
  SET merchant = "pauls"
  WHERE brand = "volvo"
  AND model = "xc90"
  AND cost = 9999
  IF EXISTS;

Batch Statements (I)

- BATCH statement
- Offers Atomicity for a series of operations
  - Write-Operations
    - INSERT, UPDATE, DELETE
  - All these operations receive the same timestamp
  - Order of operations is NOT guaranteed
- Does NOT offer isolation
  - Other statements can read/write data affected by the batch

Batch Statements (II)

```
BEGIN BATCH
    UPDATE user SET lock = true IF lock = false;
    WHERE performer = 'The Beatles' AND
                      year = 1966 AND title = 'Revolver';
    INSERT INTO albums_by_performer (performer, year, title,
                                              genre)
    VALUES ('The Beatles', 1966, 'Revolver', 'Rock');
    UPDATE user SET lock = false;
APPLY BATCH;
```

C*

DATA MODELING FOR C*

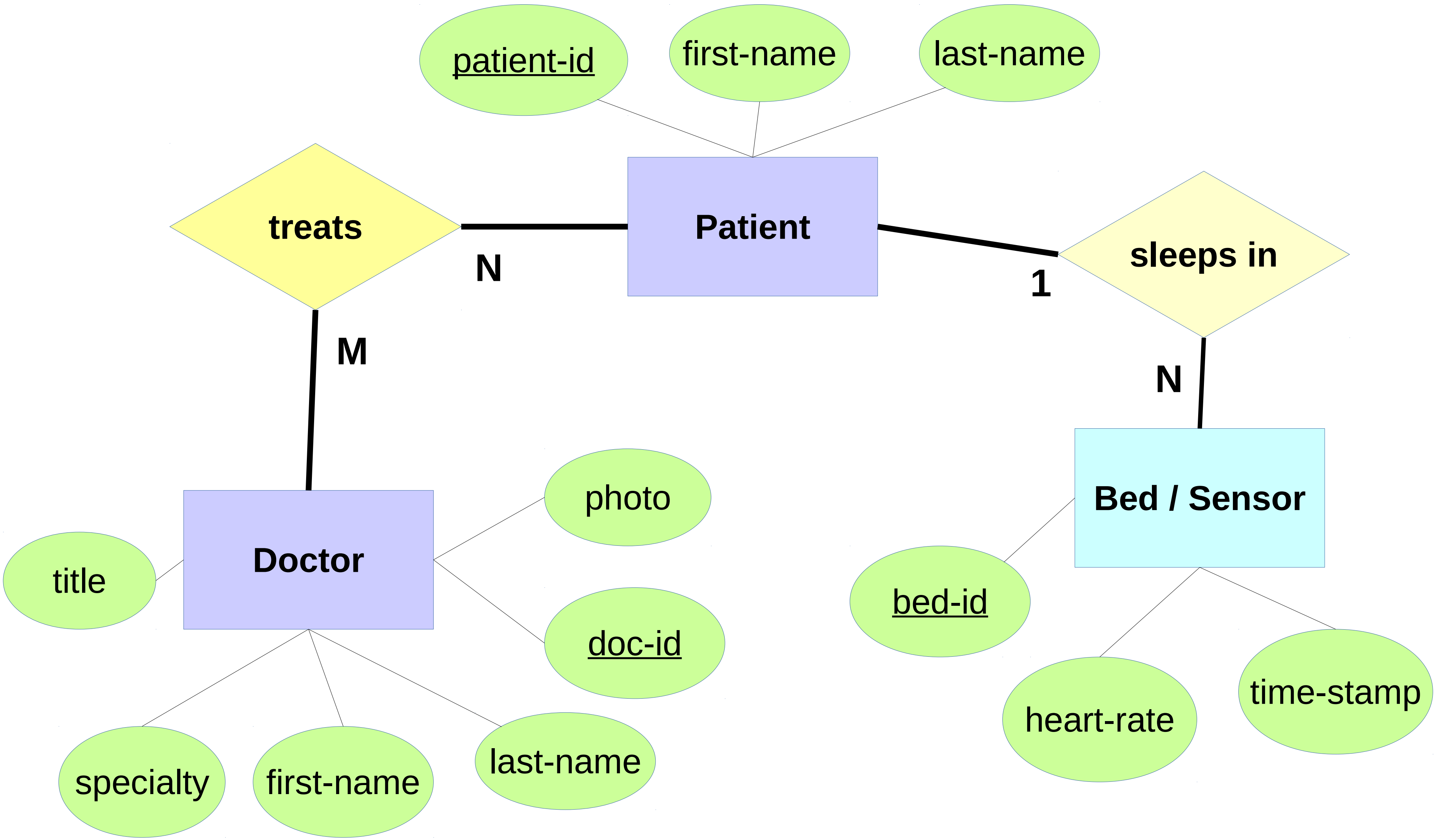
Data Modeling (DM)

1. Conceptual Data Model
2. Query-Driven Schema Design
    - Access Patterns
3. Logical Data Model
4. Analysis for Partition Size and Data Duplication
5. Physical Data Model
    - CQL

DM: Example

- Toy Application related to heart rate measurements of patients in a hospital
- We start with a Conceptual Model and some Queries in the form of Requirements

DM: Conceptual Data Model



DM: Queries – Reqs. (I)

Q1) Retrieve all information for a Doctor given his/her full name

Q2) Retrieve all information excluding the photograph for a Doctor given his/her full name

Q3) Retrieve the names and ids of the Doctors treating a given Patient, who is known by his/her patient-id. Sort them alphabetically.

DM: Queries – Reqs. (II)

Q4) Find the average heart rate for a given Patient – known by his/her patient-id – on a single given date and given hour range

Q5) Find the average heart rate for a given Patient – known by his/her patient-id – on a given date range.
Assume that the date range will have at most ten days.

DM: Logical Data Model (I)

- Q1, Q2:
  - last_name TEXT → Partition Key column (*)
  - first_name TEXT → Partition Key column (*)
  - doc_id INT → Clustering column (**)
  - specialty TEXT
  - title TEXT
  - photo BLOB
- (*) We search by it
- (**) Needed for uniqueness across a row - case of two different doctors with same names

DM: Logical Data Model (II)

- Q3:
  - patient_id TEXT → Partition Key column (*)
  - doc_last_name TEXT → Clustering column (**)
  - doc_first_name TEXT → Clustering column (**)
  - doc_id INT → Clustering column (***)

- (*) We search by it
- (**) We order the result by them
- (***) Needed for uniqueness across a row - Corner case of two different doctors with same names treating one patient

DM: Logical Data Model (III)

- Q4, Q5:
  - patient_id TEXT → Partition Key column (*)
  - patient_last_name TEXT
  - patient_first_name TEXT
  - bed_id TEXT
  - when TIMESTAMP → Clustering column (**)
  - heart_rate INT
- (*) We search by it
- (**) We perform range-search by it

DM: Analysis (I)

- Q1, Q2:
  - A photo is only needed for Q1 but can slow down Q2
  - We can duplicate data to make Q2 faster
- Q1 =

  [last_name **K**, first_name **K**, doc-id **C↑**, specialty, title, **photo**]
- Q2 =
  [last_name **K**, first_name **K**, doc_id **C↑**, specialty, title]

DM: Analysis (II)

- Q3: No surprises
- Q3 =
  [patient_id **K** ,
  doc_last_name **C↑**,
  doc_first_name **C↑**,
  doc_id **C↑**]

DM: Analysis (III)

- Q4, Q5:
  - We have at most 1 measurement per minute, that is 1440 per day, 14 400 for ten days
  - The range is defined by the timestamp is quite large
  - Although we keep a full timestamp, we only query for days or hours
  - Information about the patient's names is repeated

DM: Analysis (IV)

- Q4, Q5 =

[patient_id **K**,
patient_last_name **S**,
patient_first_name **S**,
bed_id,
when_date **TIMESTAMP** **C↓**,
when_day_minutes **INT** **C↓**,
heart_rate]

- This is good enough for bounded Partition size

DM: Physical and CQL

- Q1) TABLE docs_w_photos_by_name
- Q2) TABLE docs_by_name
- Q3) TABLE docs_by_patient
- Q4, Q5)
  TABLE heart_rate_by_patient_and_time

DM: CQL and Application

Show me some Code

MORE ABOUT C*

Get Started (I)

- Apache C* - http://cassandra.apache.org/
- DataStax C* - http://www.datastax.com/
- DataStax C* Drivers and Tools - https://academy.datastax.com/downloads/welcome
- DataStax and Apache C* Drivers -
  - DataStax
  - Apache

Get Started (II)

- DataStax and Apache CQL Documentation -
  - http://docs.datastax.com/en/cql/3.3/cql/cqlIntro.html
  - https://cassandra.apache.org/doc/cql/CQL.html

- **DataStax Startup Program** - http://www.datastax.com/datastax-enterprise-for-startups

Learn More

- DataStax Community Service
  - http://www.planetcassandra.org/
- DataStax Cassandra Academy
  - https://academy.datastax.com/
- Stack Overflow Tags:
  - cassandra, datastax, datastax-enerprise
- Books
  - visit Amazon

Get Involved

- Events about Cassandra - http://www.datastax.com/company/events
  - Conferences
  - Webinars
- Meetup Groups - http://www.meetup.com/
  - Join a group, or
  - **CREATE ONE!**

References – Sources (I)

- These (already mentioned) web pages:
  - http://cassandra.apache.org/
  - http://www.datastax.com/
  - http://docs.datastax.com/en/cql/3.3/cql/cqlIntro.html
  - https://cassandra.apache.org/doc/cql/CQL.html
  - https://academy.datastax.com/
  - http://www.planetcassandra.org/

References – Sources (II)

- These additional articles/web pages:
  - Leslie Lamport, "Time, clocks, and the ordering of events in a distributed system", http://research.microsoft.com/en-us/um/people/lamport/pubs/time-clocks.pdf
  - Mark Burgess, "Deconstructing the 'CAP theorem' for CM and DevOps", http://markburgess.org/blog_cap.html

References - Acknowledgments

- Tables and Graphics in slides marked with

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- We really thank them for that.

Let's Talk!

THANK YOU