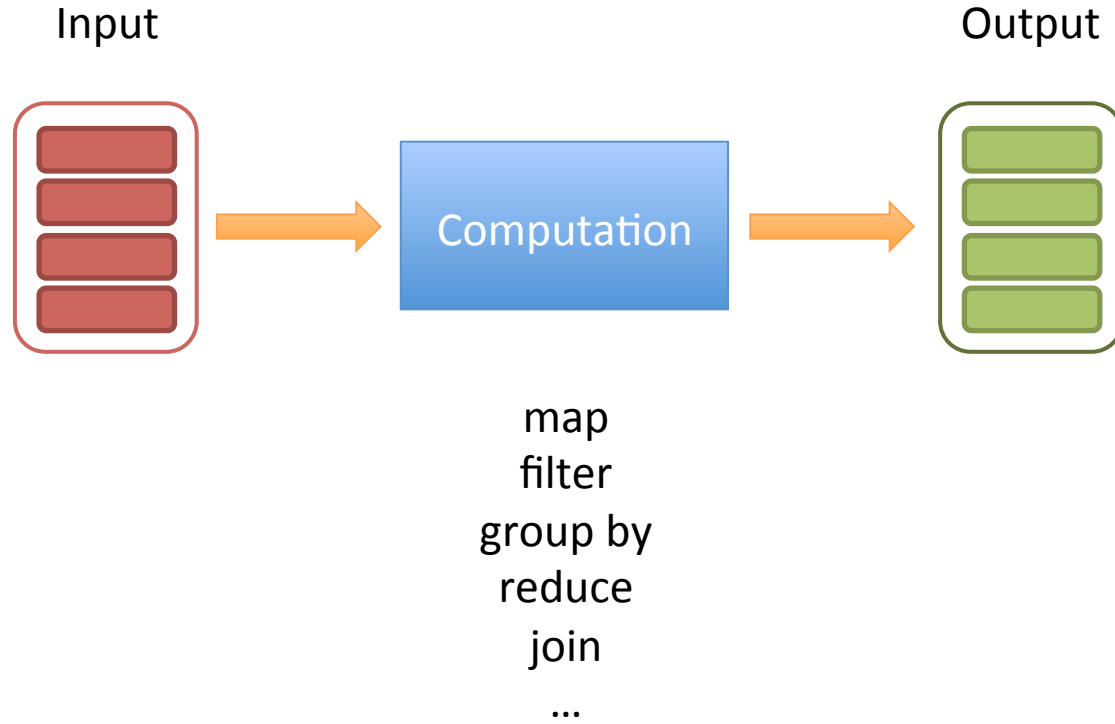


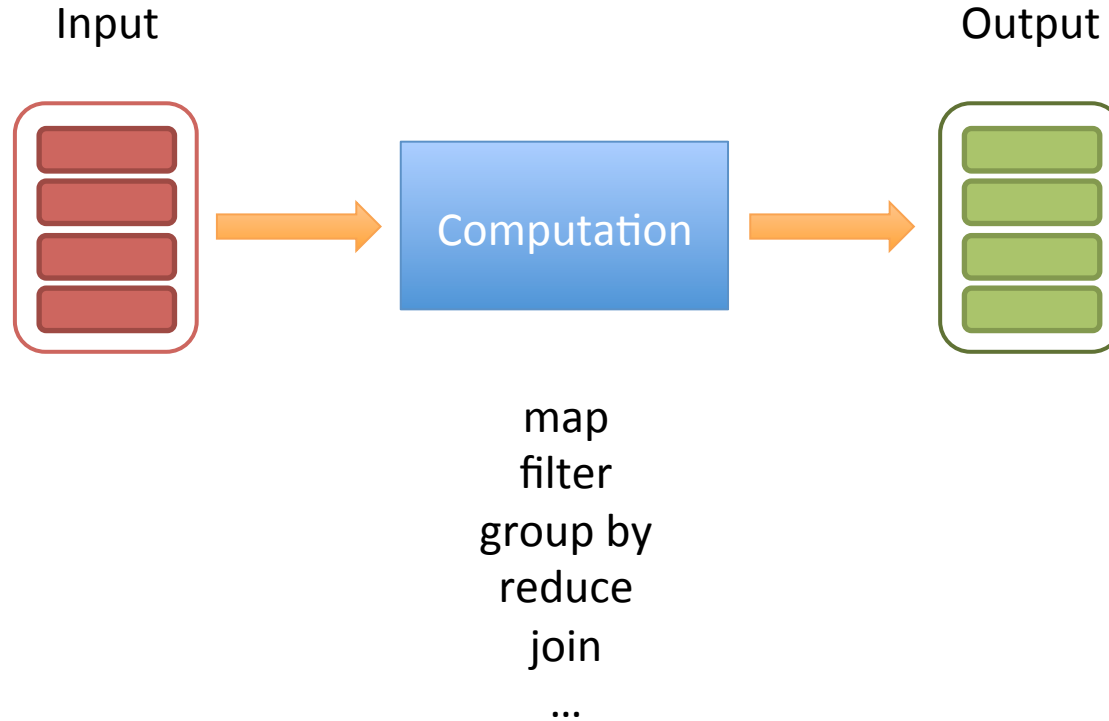
# Large-scale computation without sacrificing expressiveness

Sangjin Han    Sylvia Ratnasamy  
*UC Berkeley*

# Review: MapReduce and Friends



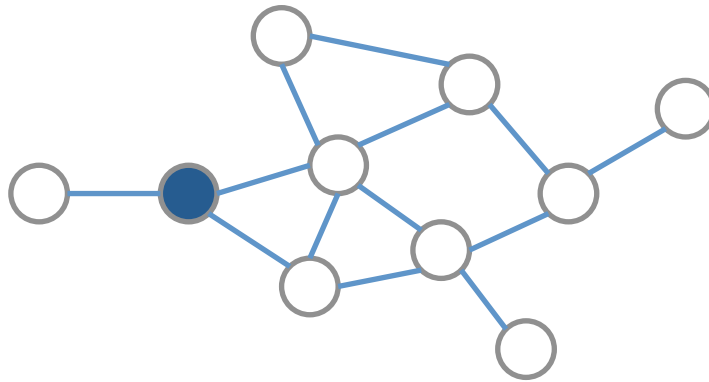
# Review: MapReduce and Friends



*Observation 1: **Bulk transformation of immutable data**  
(no fine-grained updates)*

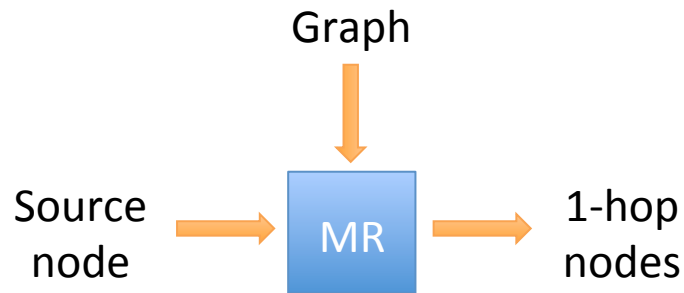
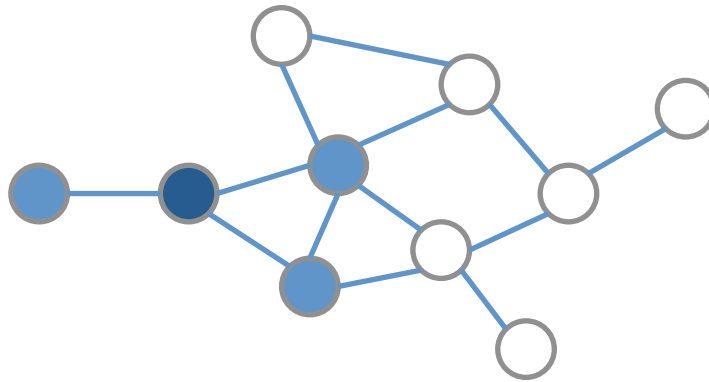
# Example 1: Sparse Operations

- k-hop reachability with iterative MapReduce



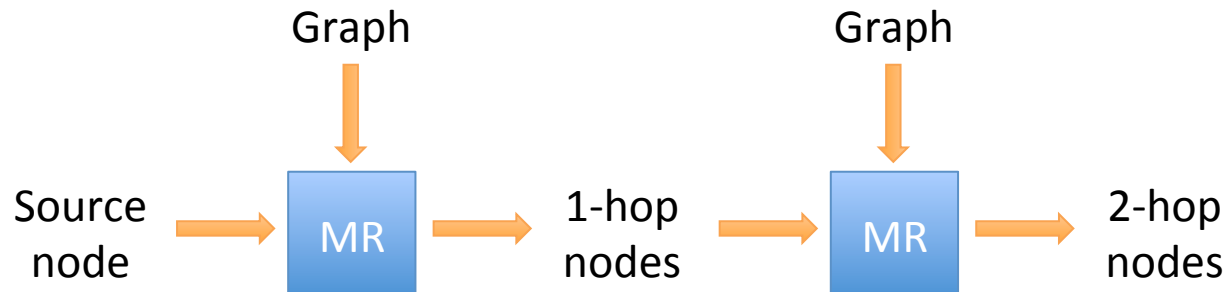
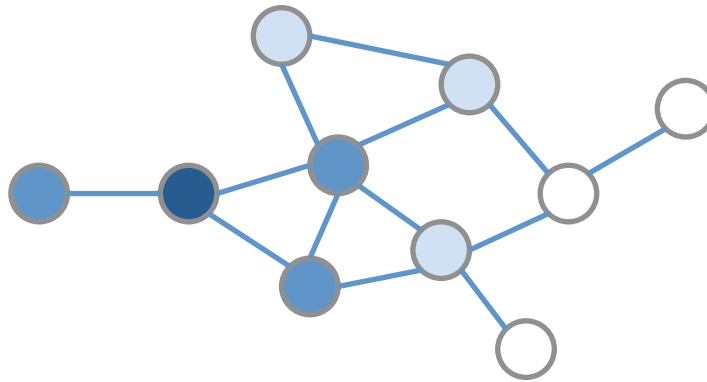
# Example 1: Sparse Operations

- k-hop reachability with iterative MapReduce



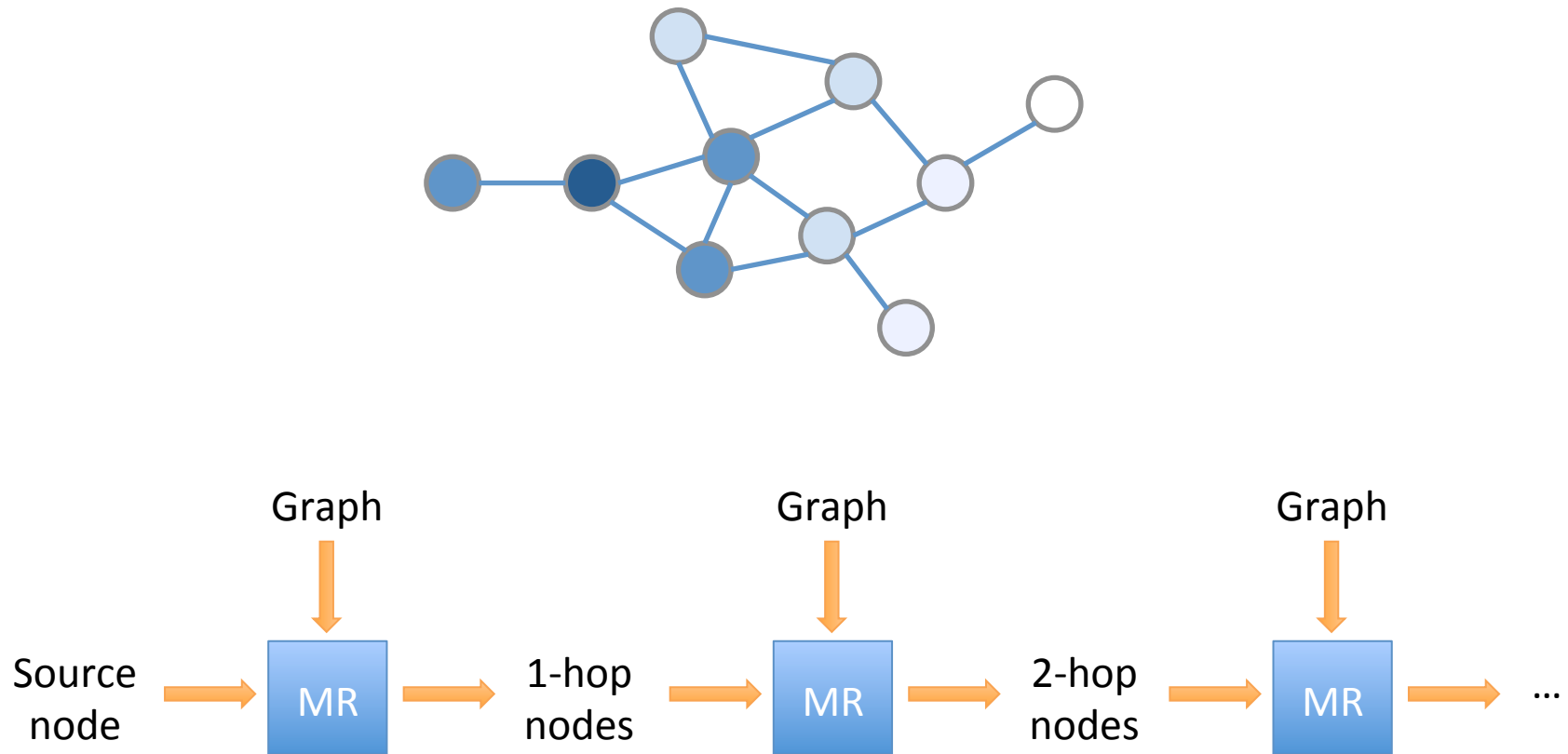
# Example 1: Sparse Operations

- k-hop reachability with iterative MapReduce



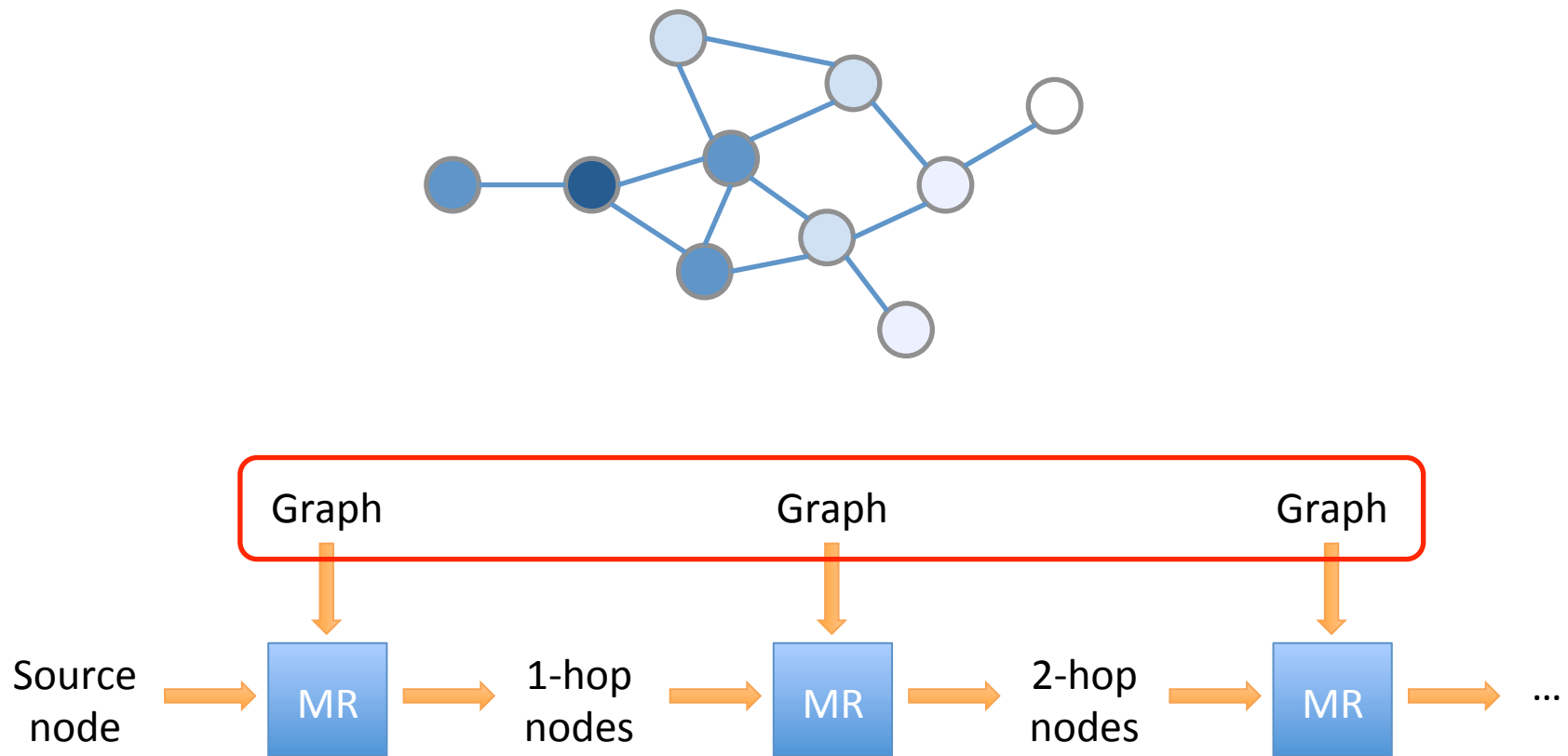
# Example 1: Sparse Operations

- k-hop reachability with iterative MapReduce



# Example 1: Sparse Operations

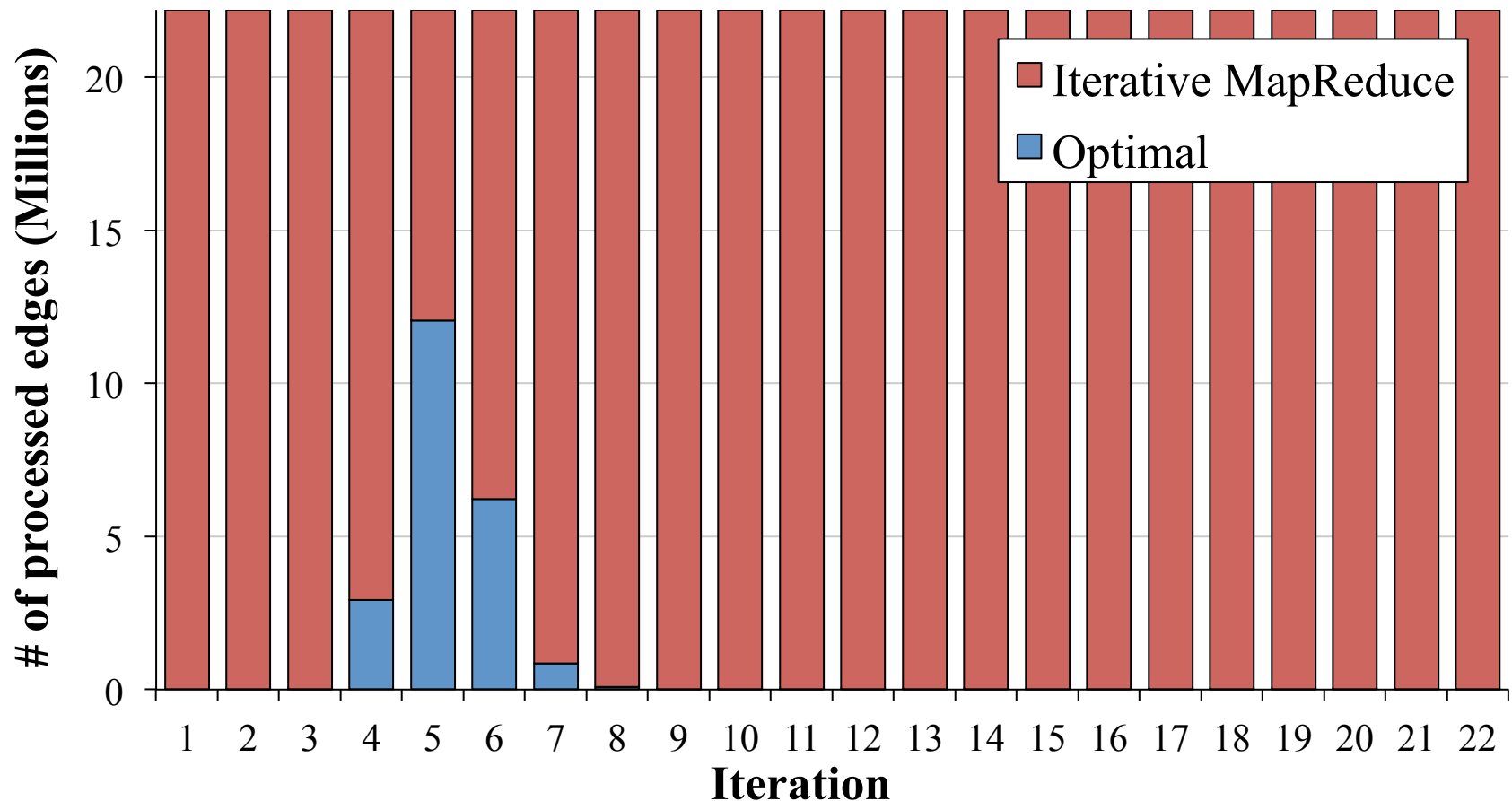
- k-hop reachability with iterative MapReduce





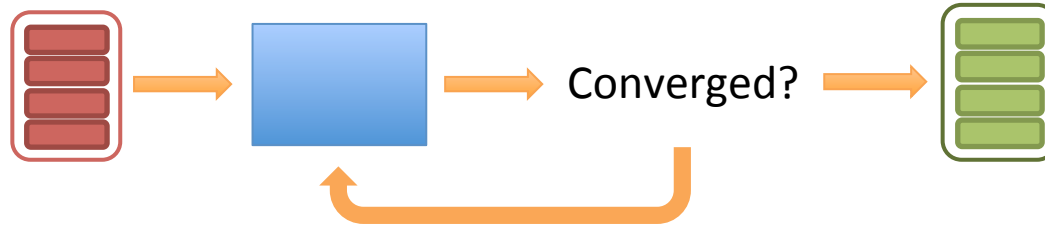
# Example 1: Sparse operations

- k-hop reachability with iterative MapReduce

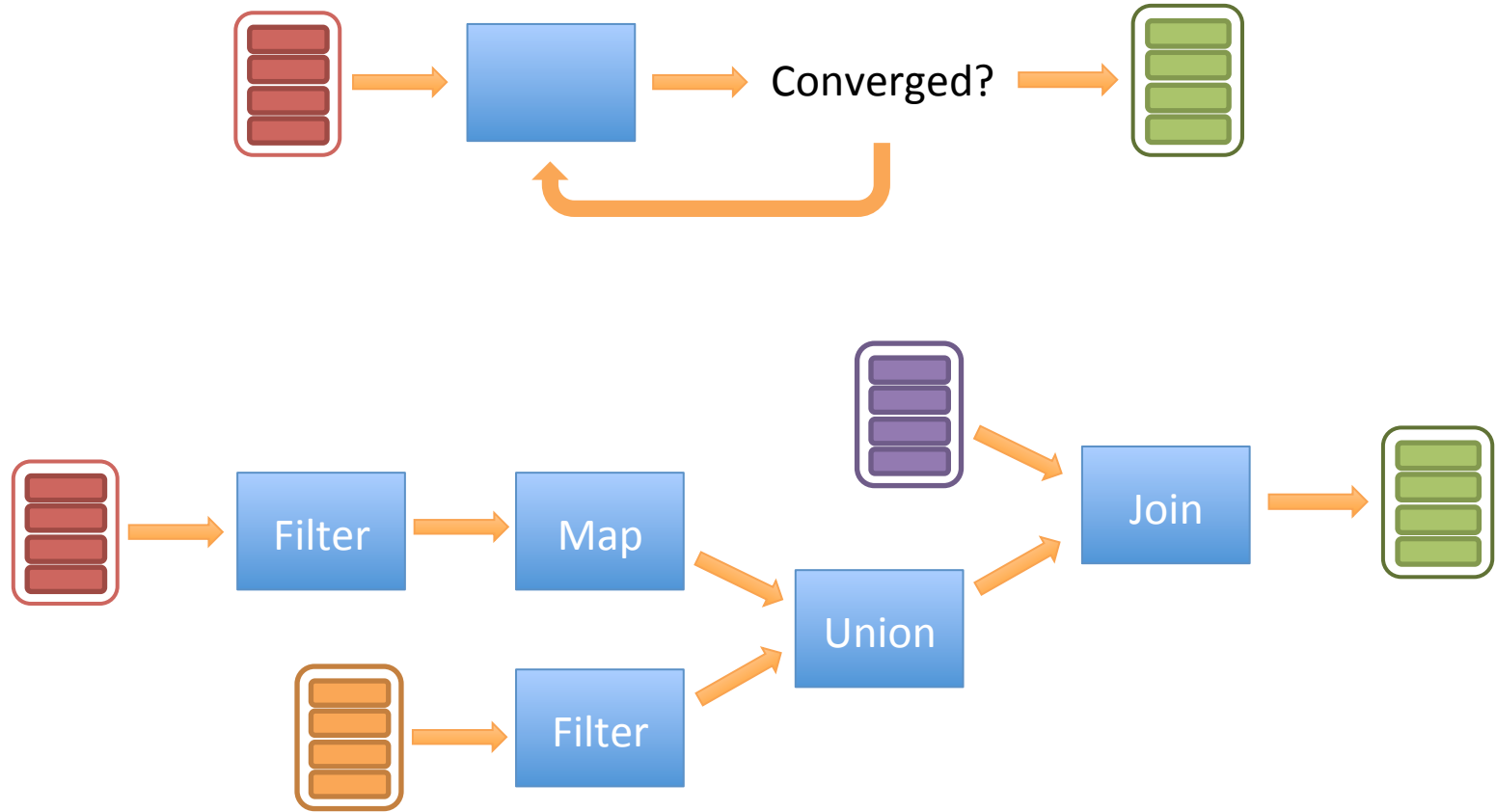


Internet router topology graph (1.7M nodes, 22.2M edges)

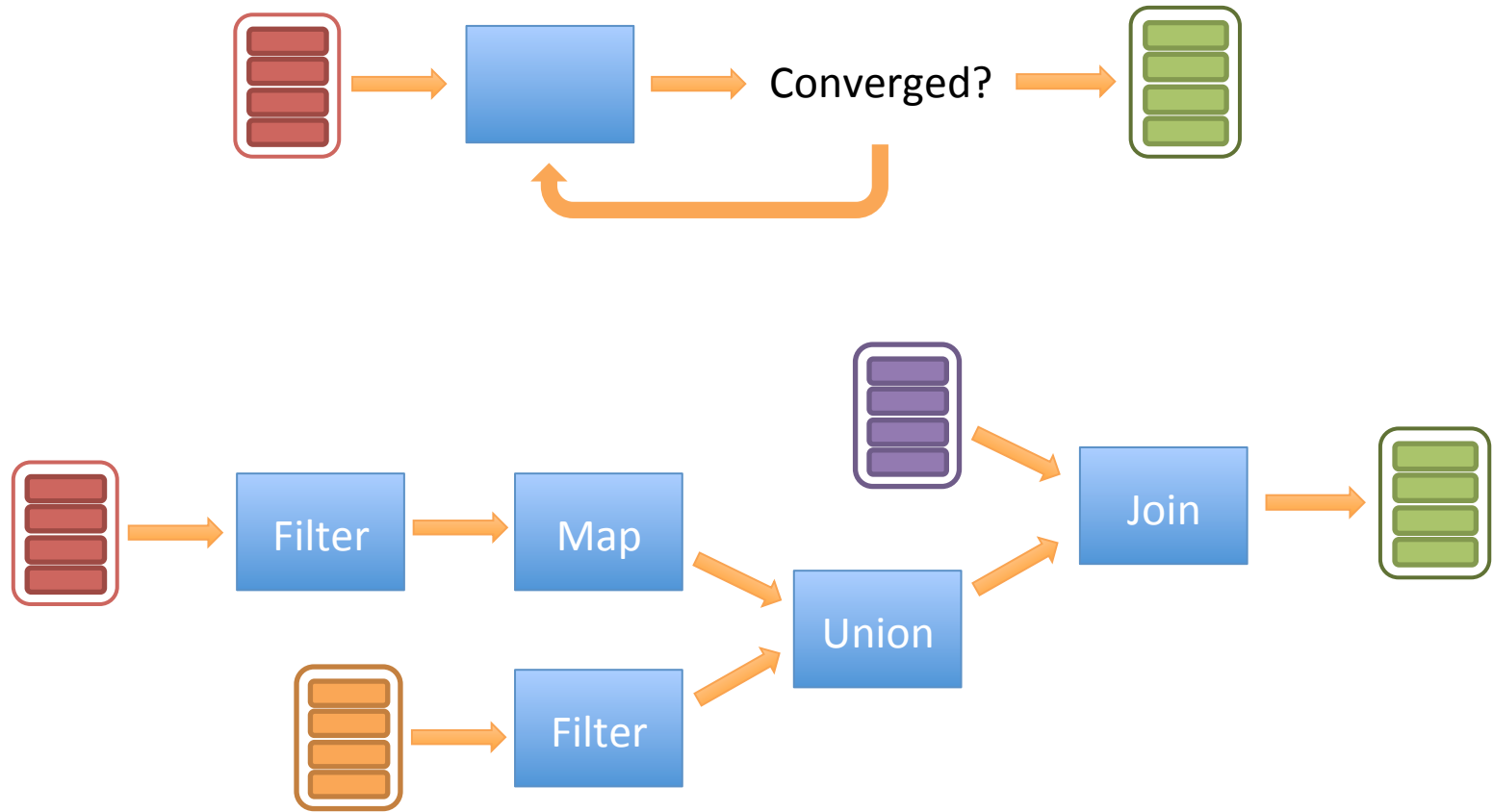
# Review: MapReduce and Friends (cont'd)



# Review: MapReduce and Friends (cont'd)



# Review: MapReduce and Friends (cont'd)



*Observation 2: **Static dataflow**  
(no data-dependent control flow)*

## Example 2: Irregular parallelism

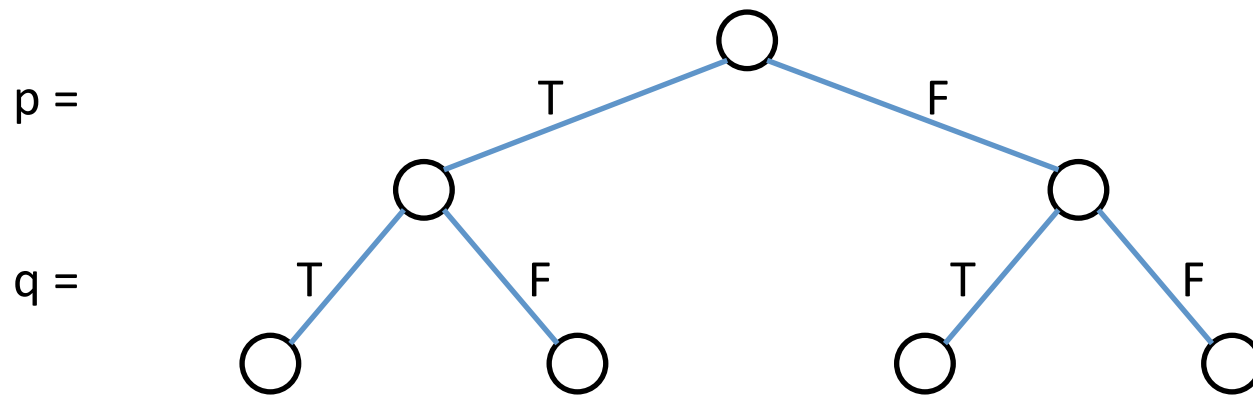
- Parallel SAT solver

$$E = (p \vee !q) \wedge (!p \vee r \vee s) \wedge (q \vee !s \vee !t) \wedge (!p \vee s) \wedge \dots$$

# Example 2: Irregular parallelism

- Parallel SAT solver

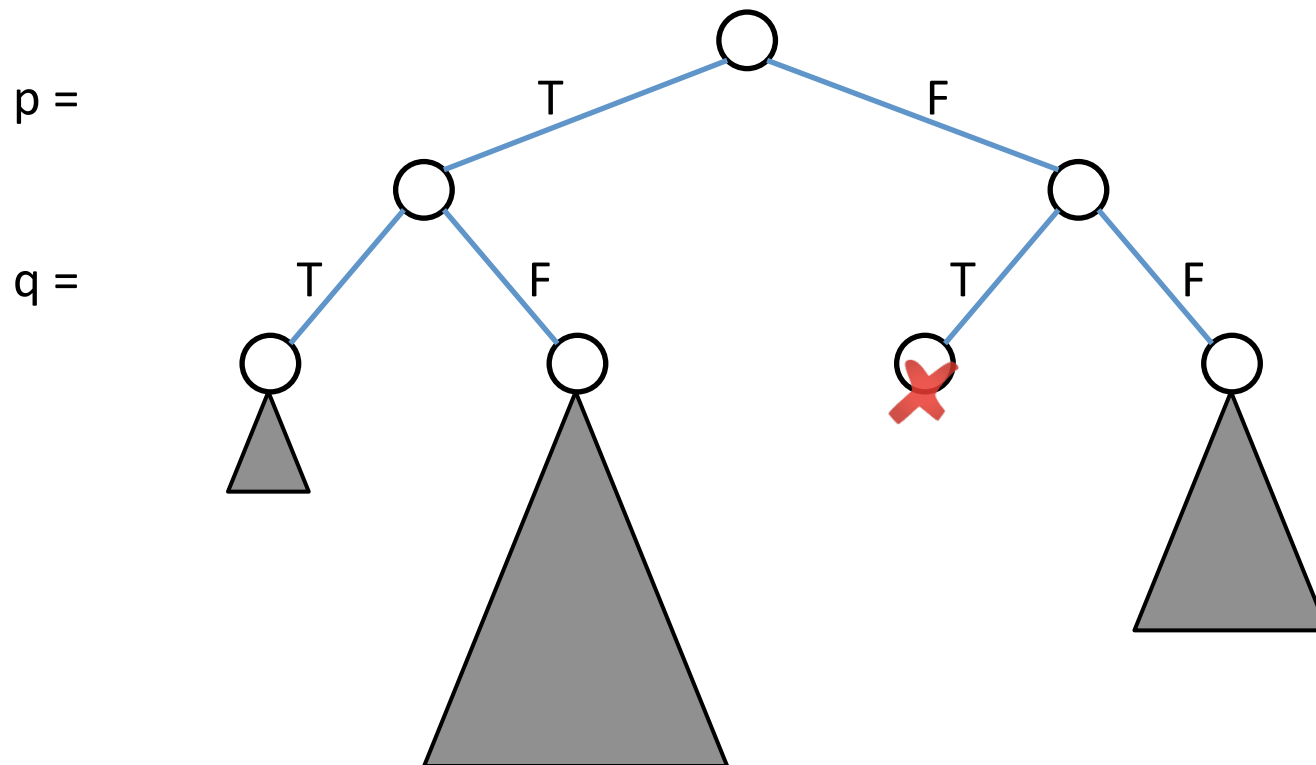
$$E = (p \vee !q) \wedge (!p \vee r \vee s) \wedge (q \vee !s \vee !t) \wedge (!p \vee s) \wedge \dots$$



# Example 2: Irregular parallelism

- Parallel SAT solver

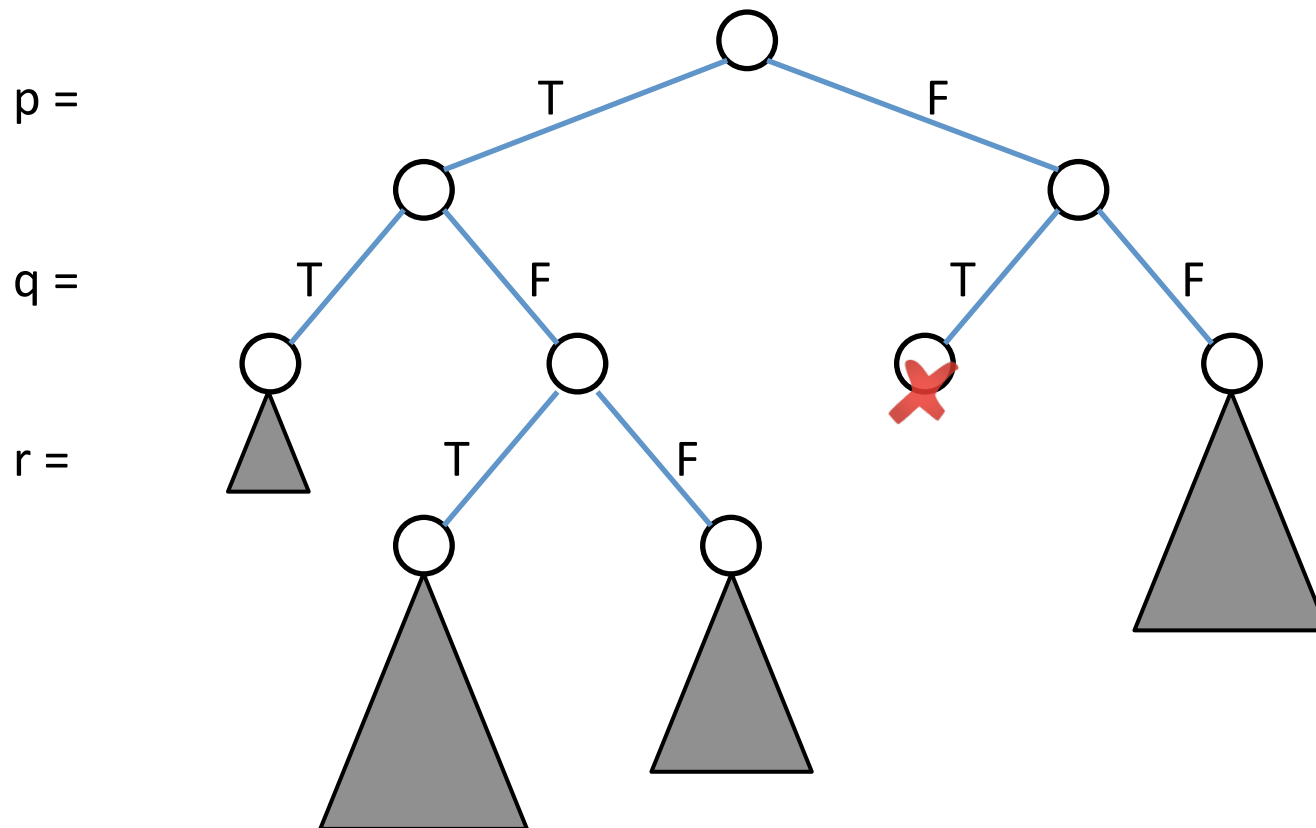
$$E = (p \vee !q) \wedge (!p \vee r \vee s) \wedge (q \vee !s \vee !t) \wedge (!p \vee s) \wedge \dots$$



# Example 2: Irregular parallelism

- Parallel SAT solver

$$E = (p \vee !q) \wedge (!p \vee r \vee s) \wedge (q \vee !s \vee !t) \wedge (!p \vee s) \wedge \dots$$





MapReduce-like frameworks assume:

1. Bulk transformation of immutable data
2. Static dataflow

~~Existing frameworks assume:~~

Our work:

1. ~~Bulk transformation of immutable data~~

Fine-grained operations on mutable data

2. ~~Static dataflow~~

Dynamic, data-dependent control flow

Yet we still want elastic scalability and fault tolerance

Spinning a small twist to Linda

# **CELIAS PROGRAMMING MODEL**

Programming model =  
data model + computation model

# Data Models for Mutable Shared Memory

# Data Models for Mutable Shared Memory

Global address space: UPC, X10, Fortress...



Too low level

# Data Models for Mutable Shared Memory

Global address space: UPC, X10, Fortress...



Too low level

Key-value tables: RAMCloud, Dynamo, Piccolo...

Key	Value
...	...

Limited lookup ability

Consistency concerns

# Data Models for Mutable Shared Memory

Global address space: UPC, X10, Fortress...



Too low level

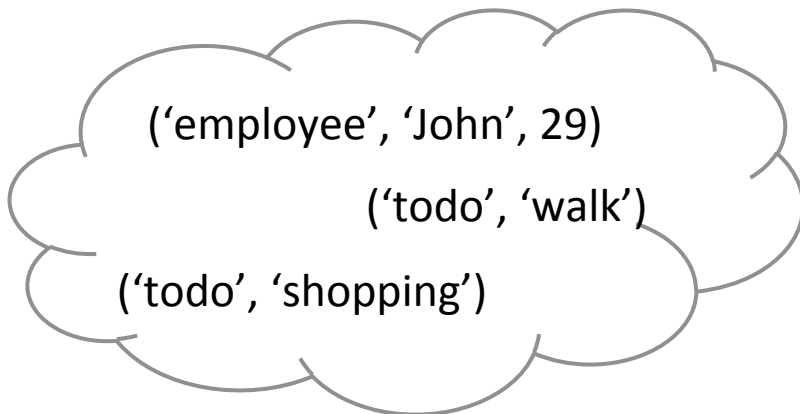
Key-value tables: RAMCloud, Dynamo, Piccolo...

Key	Value
...	...

Limited lookup ability

Consistency concerns

Tuplespace: Linda



Flexible lookup with any attributes

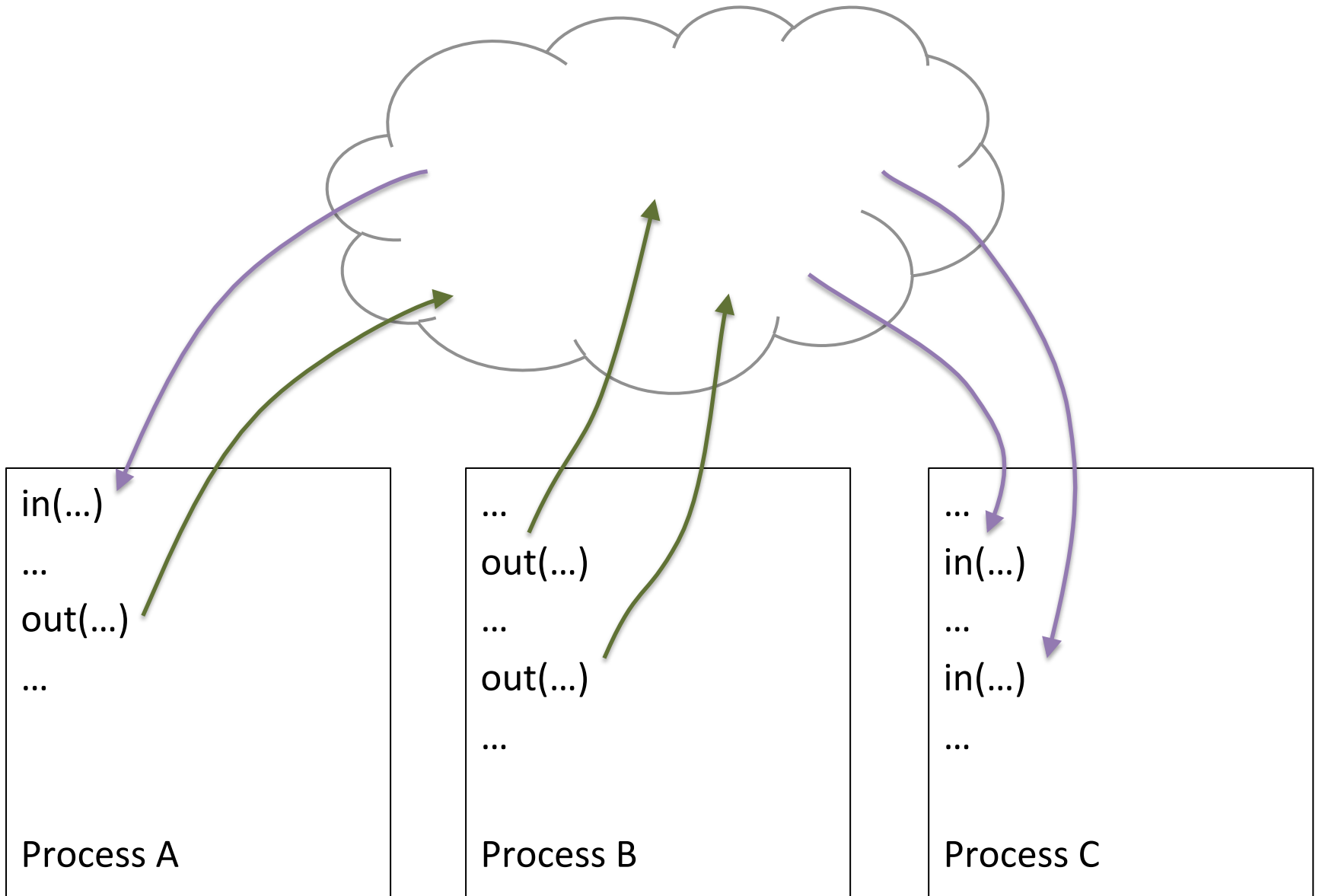
Individual tuples are immutable



Programming model =  
data model + computation model

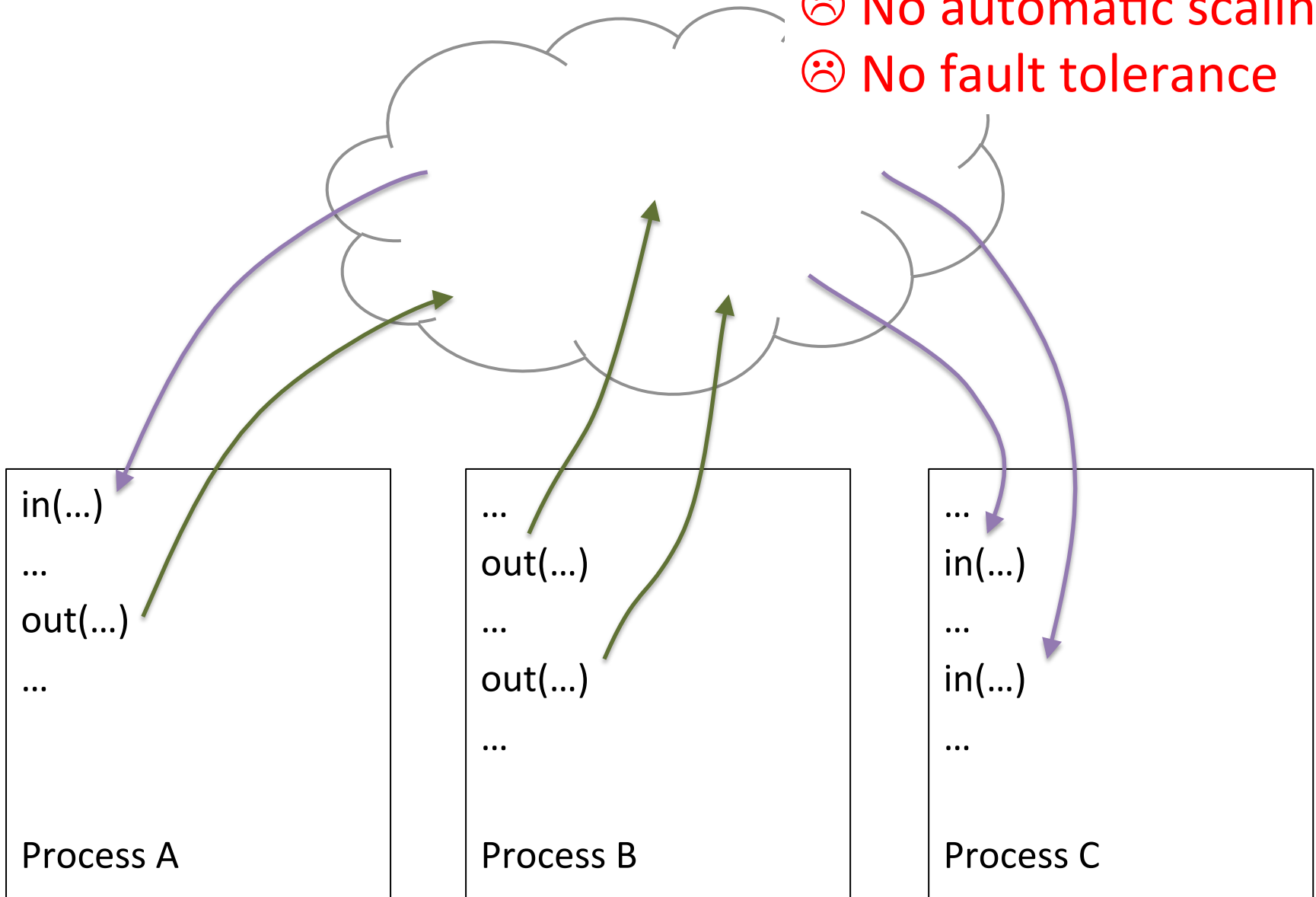
Linda =  
Tuplespace + Linda processes

# Linda Processes




# Linda Processes

- ☹ No automatic scaling
- ☹ No fault tolerance



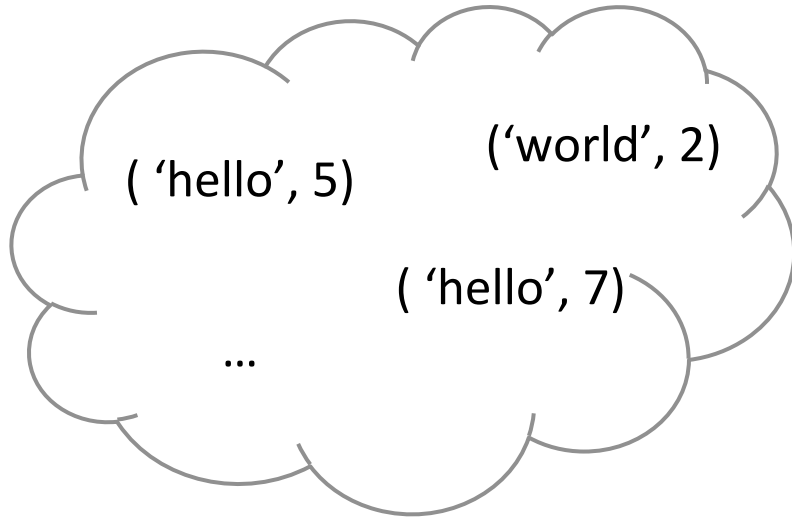
Programming model =  
data model + computation model

Linda =  
Tuplespace + Linda processes

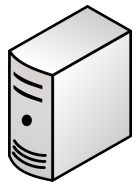


Celias =  
Tuplespace + microtasks

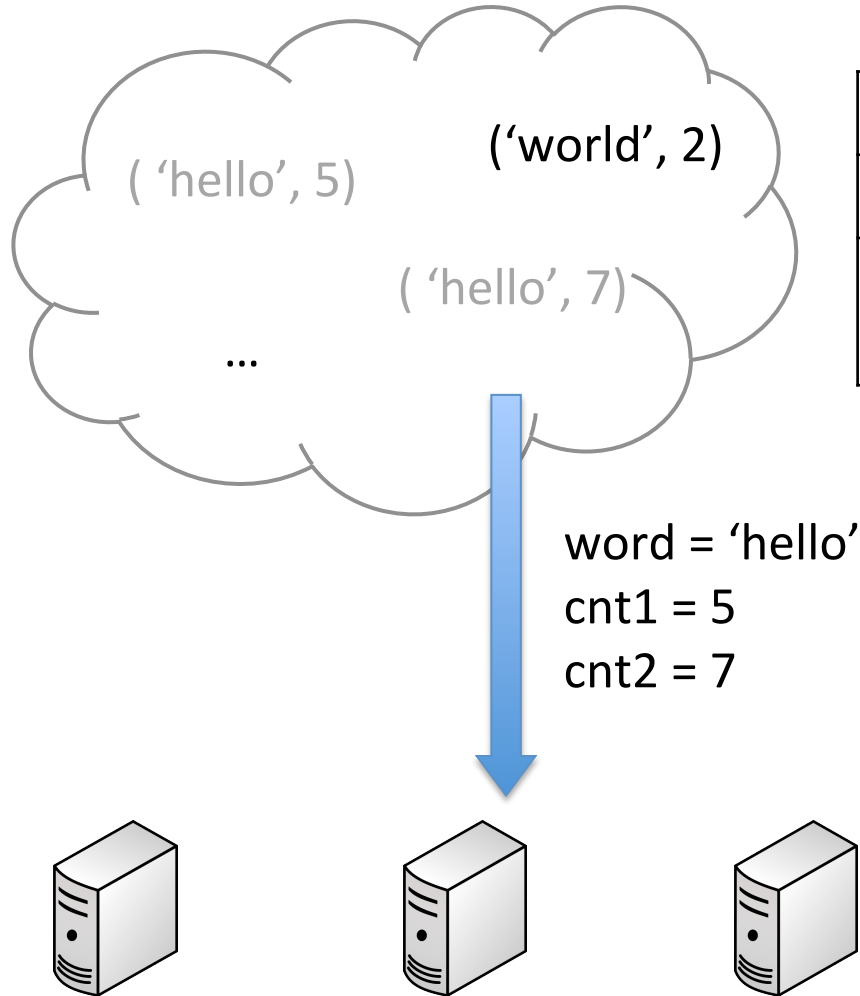
# Microtasks



Function wordcount()	
Signature	(?word, ?cnt1), (?word, ?cnt2)
Code	sum := cnt1 + cnt2 emit (word, sum)



# Microtasks

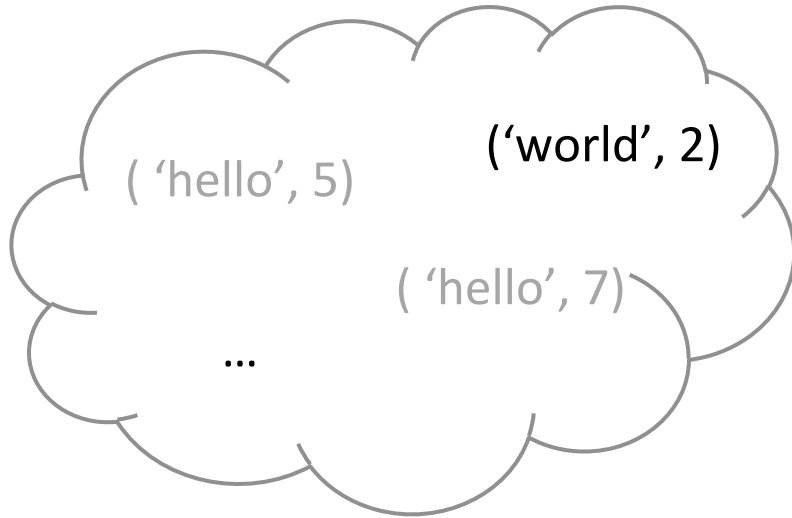


Function wordcount()	
Signature	(?word, ?cnt1), (?word, ?cnt2)
Code	sum := cnt1 + cnt2 emit (word, sum)

When a signature matches:

1. microtask launch

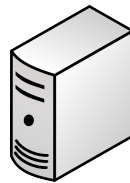
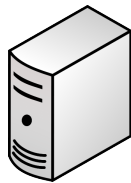
# Microtasks



Function wordcount()	
Signature	<code>(?word, ?cnt1), (?word, ?cnt2)</code>
Code	<code>sum := cnt1 + cnt2</code> <code>emit (word, sum)</code>

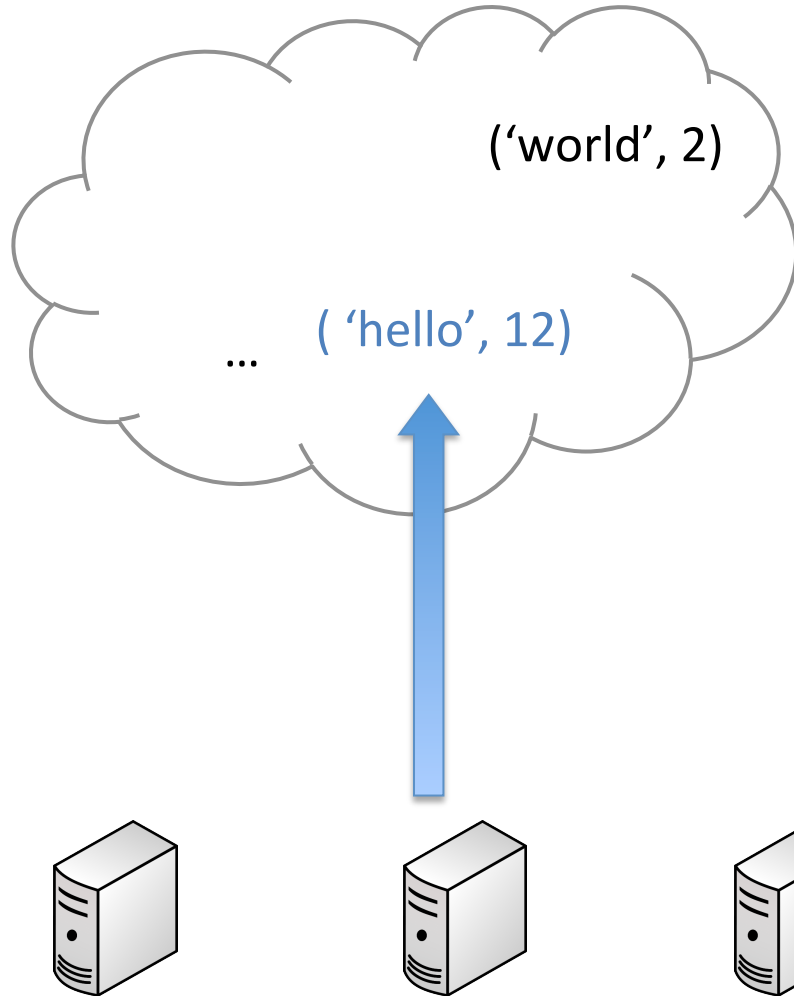
When a signature matches:

1. microtask launch
2. code execution



`5 + 7 = ??`

# Microtasks



Function wordcount()	
Signature	(?word, ?cnt1), (?word, ?cnt2)
Code	sum := cnt1 + cnt2 emit (word, sum)

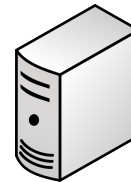
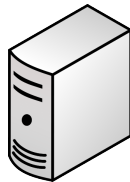
When a signature matches:

1. microtask launch
2. code execution
3. atomic replacement



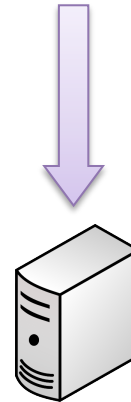
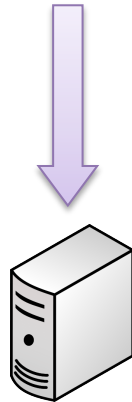
Two functions: add() and multiply()

$$(A + B) \times (C + D)$$

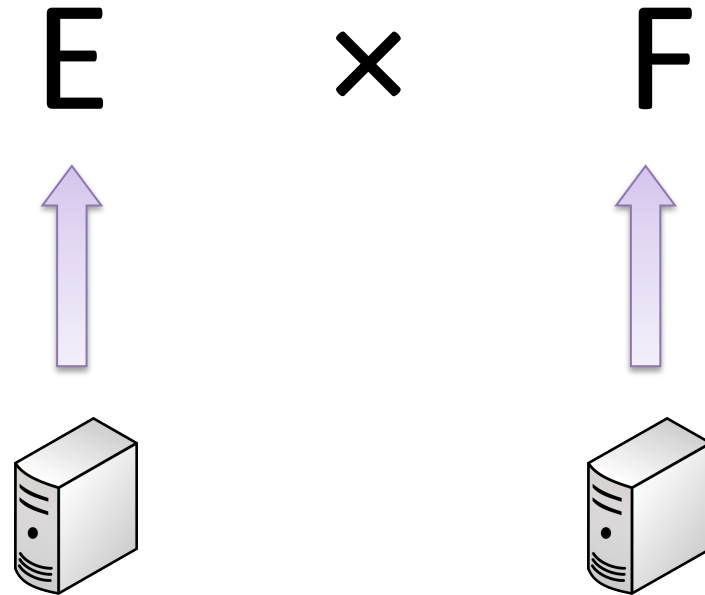


Two functions: add() and multiply()

$$(A + B) \times (C + D)$$

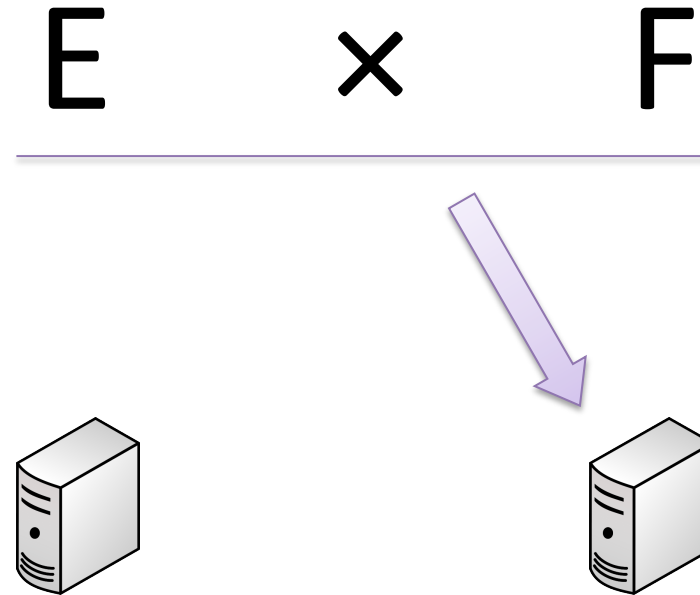


Two functions: add() and multiply()



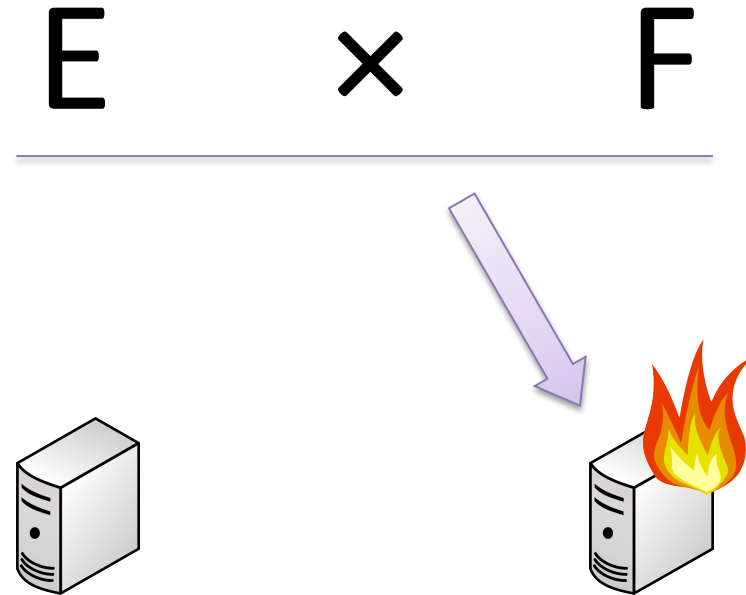
😊 Automatic scaling

Two functions: add() and multiply()



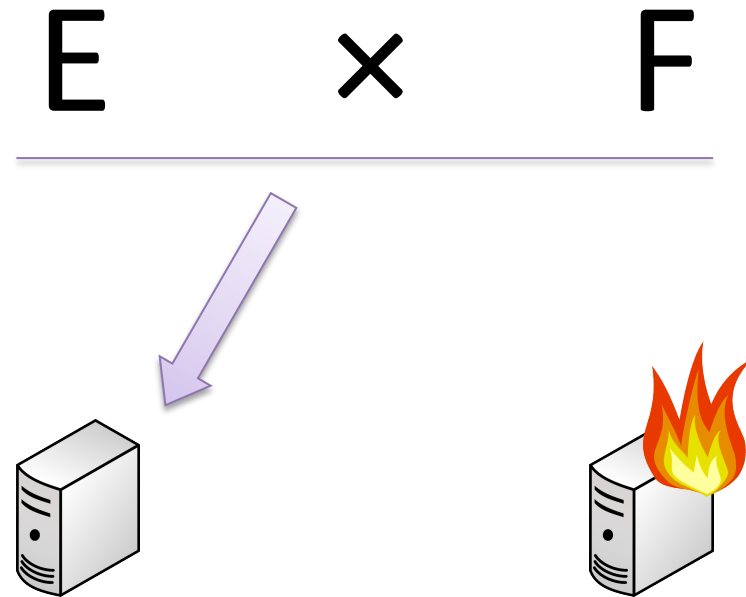
😊 Automatic scaling

Two functions: add() and multiply()



😊 Automatic scaling

Two functions: add() and multiply()



😊 Automatic scaling

😊 Fault tolerance

# More Examples in the Paper...

- MapReduce
  - Celas is ~~Turing-complete~~ MapReduce-complete!
  - without any artificial sync. barriers
- Single-source shortest path
  - Pregel-style graph processing
- Quicksort
  - Recursive control flow

# Summary

- MapReduce-like frameworks are not suitable for algorithms with:
  - Sparse/incremental/fine-grained computation
  - Dynamic dataflow
- Celas comes to our rescue, yet it is also
  - automatically scalable
  - fault tolerant



# Open Questions

- Microtask abstraction: good enough? went too far?
- Feasibility of an efficient implementation
  - Reliable tuplespace
  - Signature matching
  - Microtask transactions
- ... what is a killer app of Celas?
- <Your questions here>