



Accept / Reject Method



Suppose that we would like to generate random variates from a pmf $\{p_j, j \ge 0\}$ and we have an efficient way to generate variates from a pmf $\{q_i, j \ge 0\}$.

Let a constant c such that

$$\frac{p_j}{q_j} \le c$$
 for all j such that $p_j > 0$

In this case, use the following algorithm

- 1. Generate a random variate Y from pmf $\{q_i, j \ge 0\}$.
- 2. Generate u=U(0,1)
- 3. If u < pY/(cqY), set X=Y and return;
- 4. Else repeat from Step 1.



Show that the algorithm generates a random variate with the required pmf $\{p_i, j \ge 0\}$.

$$p_i = \Pr\{X = i\} = \sum_{k=1}^{\infty} \Pr\{\text{not stop for } 1, \dots, k-1\} \Pr\{X = i \text{ and stop at iteration } k\}$$

 $\Pr\{X = i \text{ and stop after } k\} =$

$$\Pr\{X = i \mid Y = i \text{ and is accepted}\} \Pr\{Y \text{ accepted} | Y = i\} \Pr\{Y = i\} = q_i \frac{p_i}{cq_i} = \frac{p_i}{c}$$

$$\Pr\{\text{Not stop up to } k-1\} = \left(\sum_{j} \Pr\{Y \text{ not accepted} | Y=j\} \Pr\{Y=j\}\right)^{k-1}$$
$$= \left(\sum_{j} \left(1 - \frac{p_{j}}{cq_{j}}\right) q_{j}\right)^{k-1} = \left(1 - \frac{1}{c}\right)^{k-1}$$

Therefore

$$p_i = \sum_{k=1}^{\infty} \left(1 - \frac{1}{c}\right)^{k-1} \frac{p_i}{c} = \qquad p_i$$



Determine an algorithm for generating random variates for a random variable that take values 1,2,..,10 with probabilities 0.11, 0.12, 0.09, 0.08, 0.12, 0.10, 0.09, 0.09, 0.10, 0.10 respectively.

 $c = \max_{i} \left\{ \frac{p_i}{q_i} \right\} = 1.2$

```
u1=U(0,1), u2=U(0,1)
Y=floor(10*u1 + 1);
while(u2 > p(Y)/c)
    u1= U(0,1); u2=U(0,1);
    Y=floor(10*rand + 1);
y(i)=Y;
```



Suppose that we would like to generate random variates from a pdf $f_x(x)$ and we have an efficient way to generate variates from a pdf $g_x(x)$.

Let a constant c such that

$$\frac{f_X(x)}{g_X(x)} \le c$$
 for all x

In this case, use the following algorithm

- 1. Generate a random variate Y from density $g_x(x)$.
- 2. Generate u=U(0,1)
- 3. If $u < f_x(Y)/(cg_x(Y))$, set X=Y and return;
- 4. Else repeat from Step 1.



The continue case is similar to the dicrete algorithm except the comparison step where rather than comparing the two probabilities we compare the values of the density functions.

Theorem

- ✓ The random variates generated by the Accept/Reject method have density $f_x(x)$.
- ✓ The number of iterations of the algorithm that are needed is a
 geometric random variable with mean c



Note: The constant *c* is important since is implies the number of iterations needed before a number is accepted, therefore it is required that it is selected so that it has its minimum value.



Use the method to generate random variates X that are normally distributed with mean 0 and variance 1, N(0,1).

 \checkmark First consider the pdf of the absolute value of |X|.

$$f_{|X|}(x) = \frac{2}{\sqrt{2\pi}} e^{-\frac{1}{2}x^2}$$

✓ We know how to generate exponentially distributed random variates Y with rate λ =1.

$$g_{y}(x)=e^{-x}, \quad x\geq 0$$



$$\frac{f_{|X|}(x)}{g_Y(x)} = \frac{2}{\sqrt{2\pi}} e^{x - \frac{1}{2}x^2} \implies c = \sqrt{\frac{2e}{\pi}}$$





Suppose we would like $Z\sim N(\mu, \sigma^2)$, then $Z:=\sigma X+\mu$

Example:



On average, how many random numbers must be simulated to generate 1000 variates from the Beta(α = 2, β = 2) distribution.



```
n=1000
       #conter for accepted
k=0
j=0
       #iterations
y=numeric(n)
while(k<n){
 u=runif(1)
 j=j+1
 x=runif(1)
 if (x^*(1-x)>u){
   #we accept x
  k=k+1
  y[k]=x
```

OBRIGADO!