编译原理第六章(三)

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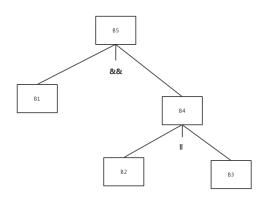
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- 1.(6.6.1)在图6-36的语法制导定义中添加处理下列控制流构造的规则
- 1) 一个repeat语句:repeat S while B
- 2)一个for循环语句: $for(S_1; B; S_2)S_3$

PRODUCTION	SEMANTIC RULES		
$S \rightarrow repeat \; S \; while \; B$	S = newlabel()		
	B.true = S		
	B.false = S.next		
	S.code = label(S) S.code B.code label(B.true) gen('goto'S) S.next		
	$S_3 = newlabel()$		
$S \rightarrow for (S_1; B; S_2)S_3$	B = newlabel()		
	$B.true = S_3$		
	B.false = S.next		
	$S_3.next = S_2$		
	$S_2.next = B$		
	$S.code = S_1.code B.code label(B.true) S_3.code S_2.code gen('goto'B) S.next$		

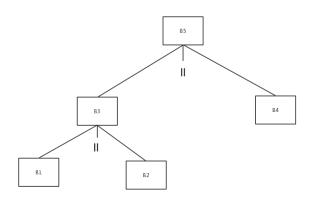
2.(6.7.1)使用图6-43的翻译方案翻译下列表达式。给出每个子表达式的truelist和falselist。你可以假设第一条被生成的指令地址是100

1)
$$a == b \&\& (c == d || e == f)$$



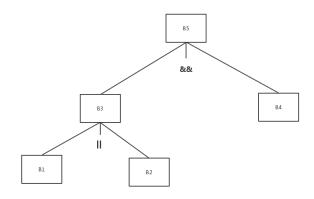
Block	T or F	
B_1	$B_1.truelist$	100
	$B_2.falselist$	101
B_2	$B_2.truelist$	102
	$B_2.falselist$	103
B3	$B_3.truelist$	104
	$B_3.falselist$	105
B4	$B_4.truelist$	102,104
	$B_4.falselist$	105
B5	$B_5.truelist$	102,104
D_0	$B_5.falselist$	101,105

$2) \ (a == b \mid\mid c == d) \mid\mid e == f$



Block	T or F	
B_1	$B_1.truelist$	100
	$B_2.falselist$	101
B_2	$B_2.truelist$	102
	$B_2.falselist$	103
B3	$B_3.truelist$	104
	$B_3.falselist$	105
B4	$B_4.truelist$	102,104
	$B_4.falselist$	103
B5	$B_5.truelist$	100, 102,104
	$B_5.falselist$	105

- $\begin{array}{lll} 100 & \text{if(a == b) goto true} \\ 101 & \text{goto } 102 \\ 102 & \text{if (c ==d) goto true} \\ 103 & \text{goto } 104 \\ 104 & \text{if(e == f) goto true} \\ 105 & \text{goto false} \\ \end{array}$
- 3) (a == b && c == d) && e == f



Block	T or F	
B_1	$B_1.truelist$	100
	$B_2.falselist$	101
B_2	$B_2.truelist$	102
	$B_2.falselist$	103
B3	$B_3.truelist$	104
	$B_3.falselist$	105
B4	$B_4.truelist$	102
	$B_4.falselist$	101,103
B5	$B_5.truelist$	104
	$B_5.falselist$	101,103,105

 $\begin{array}{lll} 100 & \text{if(a == b) goto } 102 \\ 101 & \text{goto false} \\ 102 & \text{if (c ==d) goto } 104 \\ 103 & \text{goto false} \\ 104 & \text{if(e == f) goto true} \\ 105 & \text{goto false} \\ \end{array}$