编译原理第三章(二)

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- 1. 为下面语言设计一个DFA或NFA
- 1) 包含5个元音的所有小写字母串,这些串中元音按照顺序出现

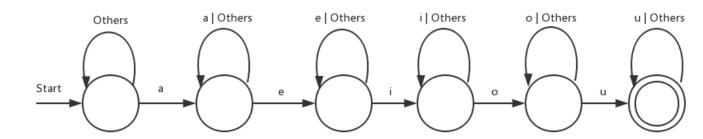


Figure 1: $Others = \{b|c|d|f|g|h|j|k|l|m|n|p|q|r|s|t|v|w|x|y|z\}$

2) 所有由a,b组成得且不含子串abb的串

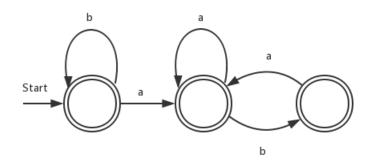


Figure 2: 由a,b组成得且不含子串abb的串

2. 用算法3.22模拟图3-29中得NFA在处理输入aabb时得过程.

$$1)S=0; c=nextChar()=a;\\$$

$$2)c! = EOF \rightarrow S = \varepsilon - closure(move(S, a)) = \{0, 1\}, c = nextChar() = a$$

3)
$$c! = EOF \rightarrow S = \varepsilon - closure(move(S, a)) = \{0, 1, 2\}, c = nextChar() = b$$

4)
$$c! = EOF \rightarrow S = \varepsilon - closure(move(s, b)) = \{0, 1, 2, 3\}, c = nextChar() = b$$

5)
$$c! = EOF \rightarrow S = \varepsilon - closure(move(s, b)) = \{0, 1, 2, 3\}, c = nextChar() = EOF$$

6)
$$c == EOF \rightarrow S \cap F! = \emptyset \ return \ yes$$

- 3. 使用算法3.23和3.20将下述正则表达式转换为DFA,并尝试化简该DFA
- 1) $((\varepsilon|a)b^*)^*$

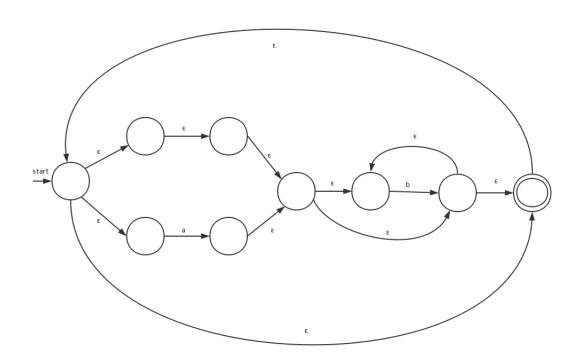


Figure 3: NFA $((\varepsilon|a)b^*)^*$

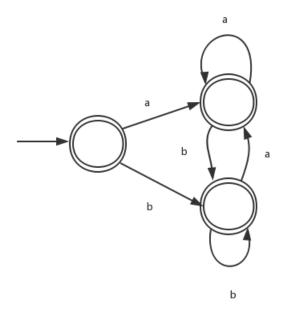


Figure 4: DFA $((\varepsilon|a)b^*)^*$

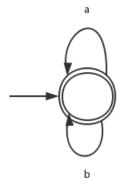


Figure 5: DFA化筒 $((\varepsilon|a)b^*)^*$

 $2) \ (a|b)^*abb(a|b)^*$

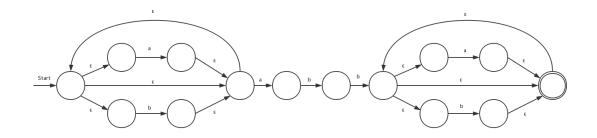


Figure 6: NFA $(a|b)^*abb(a|b)^*$

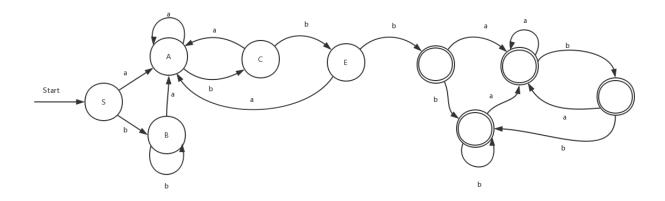


Figure 7: DFA $(a|b)^*abb(a|b)^*$

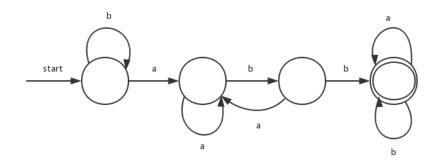


Figure 8: DFA化简(a|b)*abb(a|b)*