Remark 1. When you are asked to hand in SCHEME code for a function, you may cut-and-paste the definition from your interpreter window into your solution set. **Always include with your code a number of illustrative examples.**

1. Recall from class the definition of number-sum, which computes the sum of the first *n* numbers:

(a) Adapt the function so that it computes the sum of the first n positive squares. (So your function, when evaluated at 4, should return the sum of the first 4 positive perfect squares: 1 + 4 + 9 + 16 = 30.) An immediate adaptation yields:

(b) Adapt the function so that it computes the sum of the first n even numbers. (So your function, when evaluated at 4, should return the sum of the first 4 even numbers: 2 + 4 + 6 + 8 = 20.)

Incidentally, evaluate your function at 1, 2, 3, 4, 5, 6, and 7. Does this sequence of numbers look familiar?

The kth even number is (2k). Thus we have:

You will notice that the sum of the first n even numbers is exactly n(n+1).

2. Write a recursive function that, given a positive integer k, computes the product

$$\underbrace{\left(1-\frac{1}{2}\right)\left(1-\frac{1}{3}\right)\cdots\left(1-\frac{1}{k}\right)}_{k-1}.$$

(Experiment with the results for some various values of k; this might suggest a simple non-recursive way to formulate this function.)

With a little experimentation, you'll discover that this product is always equal to 1/n.

3. Consider the problem of determining how many divisors a positive integer has. For example:

- The number 4 has three divisors: 1, 2, and 4;
- The number 5 has two divisors: 1 and 5;
- The number 10 has four divisors: 1, 2, 5, and 10.

In this problem you will write a SCHEME function (divisors n) that computes the number of divisors of a given number n.

The first tool you will need is a way to figure out if a given whole number ℓ divides another whole number n evenly. We provide the code for this, which you can just use as-is in your solution (it involves a function that we haven't talked about in class yet):

```
(define (divides a b) (= 0 (modulo b a)))
```

Once you have defined this function, (divides a b) will be #t if a divides b evenly, and #f if not. For example:

```
> (divides 2 4)
> (divides 3 5)
#f
> (divides 6 3)
```

At first glance, the problem of defining (divisors n) appears a little challenging, because it's not at all obvious how to express (divisors n) in terms of, for example, (divisors (- n 1)); in particular, it's not really clear how to express this function recursively.

To solve the problem, you need to introduce some new structure! Here's the idea. Focus, instead, on the function (divisors-upto n k) which computes the number of divisors n has between 1 and k (so it computes the number of divisors of *n* upto the value *k*). Now you will find that there is a straightforward way to compute (divisors-upto n k) in terms of (divisors-upto n (- k 1)). Specifically, notice that

```
(\text{divisors-upto n k}) = \begin{cases} 0 & \text{if } n = 0; & \text{(otherwise, n)} \\ 1 & \text{if } k = 1; \\ 1 + (\text{divisors-upto n (- k 1)}) & \text{if } k \text{ divides } n; \\ (\text{divisors-upto n (- k 1)}) & \text{if } k \text{ does not divide } n. \end{cases}
                                                                                                                                                                            if n = 0; (otherwise, n \ge 1, and)
```

Write the SCHEME code for the function divisors-upto; notice then that you can define

```
(define (divisors n) (divisors-upto n n))
```

In this case, we call divisors-upto a "helper" function. What did it do? It let us "re-structure" the problem we wish to solve in such a way that we can recursively decompose it.

```
(define (divides a b) (= 0 (modulo b a)))
(define (divisors-upto n k)
  (cond ((= k 1)
                       1)
        ((divides k n) (+ 1 (divisors-upto n (- k 1))))
                       (divisors-upto n (- k 1))))
(define (divisors n) (divisors-upto n n))
```

4. Write a function that, given a positive integer k, returns the sum of the first k terms of the infinite series:

$$\frac{4}{1} - \frac{4}{3} + \frac{4}{5} - \frac{4}{7} + \frac{4}{9} - \cdots$$

Use your function to sum the first 100 terms of this series.

Observe that signs alternate in this series; one easy way to implement an alternating sign is to use the function $(-1)^{\ell}$, which is 1 when ℓ is even, and -1 when ℓ is odd. You may wish to use the built-in SCHEME function (expt x t), which returns x^k (and so provides a straightforward way to compute $(-1)^{\ell}$).

Depending on how you wrote your code, SCHEME may have produced *exact* output of the form a/b. To coerce SCHEME to give you an approximation in decimal form, change the constant 4 in your code to 4.0.

Now compute the sum of the first 100,000 terms. Does this number look (roughly) familiar?

This series converges (rather slowly) to π .

- 5. Consider the definition of your last function.
 - (a) To compute the 300 terms, how many calls to expt were made? What are the actual values passed to each call?

The evaluation of the function required 300 evaluations of expt (one for each recursive call).

(b) Revise the function you just wrote to eliminate the repeated invocations of expt. (Your function should not use expt at all but somehow compute the signs on its own.) You can introduce a helper function with more arguments if you wish!

Another version, that doesn't even use even. This gives a slightly different recursive decomposition. The function pi-aux-b will return the sum of the terms numbered t through desired-length of the series, so long as you give it the sign of the first term (in the parameter s).

```
;; The workhorse, which computes the sum of terms t through ;; desired-depth of the series, so long as you give it the ;; sign (s) that it is supposed to start with.
```

6. (cf. SICP problem 1.4) Suppose we designed a new if function as follows:

Check that new-if works as you might expect by evaluating:

```
> (new-if (= 0 0) 4 5)
4
> (new-if (= 0 1) 4 5)
5
```

Recall now the factorial function that we defined and discussed in class:

```
(define (factorial n)
    (if (= n 0)
          1
          (* n (factorial (- n 1)))
    )
)
```

How does factorial behave if you replace the usage of if with new-if? Explain.

We discussed this at length in class: usage of new-if will result in a factorial function that never terminates.

7. Recall from high-school trigonometry the sin function: If T is a right triangle whose hypotenuese has length 1 and interior angles x and $\pi/2 - x$, $\sin(x)$ denotes the length of the edge opposite to the angle x (here x is measured in radians). You won't need any fancy trigonometry to solve this problem.

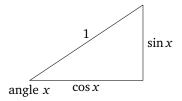


Figure 1: A right triangle.

It is a remarkable fact that for all real x,

$$\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \cdots$$

Write a SCHEME function new-sin so that (new-sin x n) returns the sum of the first (n + 1) terms of this power series evaluated at x. Specifically, (new-sin x 3), should return

$$x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!}$$

and, in general, (new-sin x n) should return

$$\sum_{k=0}^{n} (-1)^k \frac{x^{2k+1}}{(2k+1)!}.$$

You may use the built-in function (expt x k), which returns x^k . It might make sense, also, to define factorial as a separate function for use inside your new-sin function. (Aesthetic hint: Note that the value 2k is used several times in the definition of the kth term of this sum. Perhaps you can use a let statement to avoid computing this quantity more than once?)

One solution is the following: