

# CS/ECE 374 P14

Jiawei Tang, Junquan Chen, Pengxu Zheng

TOTAL POINTS

**50 / 100**

QUESTION 1

1 Problem 14.A. **15 / 20**

✓ - **5 pts** Missing explanation or largely incorrect justification of algorithm correctness

QUESTION 2

2 Problem 14.B. **25 / 40**

✓ - **10 pts** Missing explanation or largely incorrect justification of algorithm correctness

✓ - **5 pts** Small error in algorithm, justification of algorithm, or number of calls calculations

💬 "Placeholders" should likely be lists that are appended to rather than strings that are assigned to.

QUESTION 3

3 Problem 14.C. **10 / 40**

✓ - **30 pts** IDK

Submitted by:

- <<Pengxu Zheng>>: <<pzheng5>>
- <<Junquan Chen>>: <<junquan2>>
- <<Jiawei Tang>>: <<jiaweit2>>

## 14

### Solution:

14.A.

```
FindMissing (A[0, ..., n - 1]):
result = "" // initialize result to an empty string
for j in l:
    num0 <- 0 // set the counter for the number of 0's to 0
    num1 <- 0 // set the counter for the number of 1's to 0
    for i in n:
        if FetchBit(i, j) == 0:
            num0++
        else:
            num1++
    if num0 < num1:
        result += '0'
    else:
        result += '1'
return result
```

Since  $n = 2^l - 1$ ,  $l = \log_2(n + 1)$ . The total number of FetchBit calls is  $n * l$ , which is  $n \log(n + 1)$ . Thus, we conclude that our algorithm runs in  $O(n \log n)$  time.

14.B.

```
FindMissing (A[0, ..., n - 1]):
result = "" // initialize result to an empty string
partial0, partial1 = "" // declare two partial placeholders for A.
for j = l - 1; j >= 0; j--:
    partial0, partial1 = "" // reset the partial placeholders
    for i = 0; i < strlen(A); i++:
        if FetchBit(i, j) == 0:
            partial0 = A[i]
        else:
            partial1 = A[i]
    if strlen(partial1) < strlen(partial0):
        result += '0'
        A = partial1
    else:
        result += '1'
        A = partial0
```

1 Problem 14.A. 15 / 20

✓ - 5 pts Missing explanation or largely incorrect justification of algorithm correctness

Submitted by:

- <<Pengxu Zheng>>: <<pzheng5>>
- <<Junquan Chen>>: <<junquan2>>
- <<Jiawei Tang>>: <<jiaweit2>>

## 14

### Solution:

14.A.

```
FindMissing (A[0, ..., n - 1]):
result = "" // initialize result to an empty string
for j in l:
    num0 <- 0 // set the counter for the number of 0's to 0
    num1 <- 0 // set the counter for the number of 1's to 0
    for i in n:
        if FetchBit(i, j) == 0:
            num0++
        else:
            num1++
    if num0 < num1:
        result += '0'
    else:
        result += '1'
return result
```

Since  $n = 2^l - 1$ ,  $l = \log_2(n + 1)$ . The total number of FetchBit calls is  $n * l$ , which is  $n \log(n + 1)$ . Thus, we conclude that our algorithm runs in  $O(n \log n)$  time.

14.B.

```
FindMissing (A[0, ..., n - 1]):
result = "" // initialize result to an empty string
partial0, partial1 = "" // declare two partial placeholders for A.
for j = l - 1; j >= 0; j--:
    partial0, partial1 = "" // reset the partial placeholders
    for i = 0; i < strlen(A); i++:
        if FetchBit(i, j) == 0:
            partial0 = A[i]
        else:
            partial1 = A[i]
    if strlen(partial1) < strlen(partial0):
        result += '0'
        A = partial1
    else:
        result += '1'
        A = partial0
```

## 2 Problem 14.B. 25 / 40

✓ - **10 pts** Missing explanation or largely incorrect justification of algorithm correctness

✓ - **5 pts** Small error in algorithm, justification of algorithm, or number of calls calculations

💬 "Placeholders" should likely be lists that are appended to rather than strings that are assigned to.

return result

Since each time we cut the string in half, the total number of FetchBit calls is  $2^l - 1 + 2^{l-1} - 1 + \dots + 2^1 - 1$ , which is less than  $2^{l+1}$  calls and that is also less than  $2(n + 1)$  calls. Hence, we conclude that is an  $O(n)$  algorithm.

14.C.

IDK. sorry

3 Problem 14.C. 10 / 40

✓ - 30 pts IDK