# 15.4 Connectivity and Recursive Graph Traversal

Now that we have covered the basics of a Python representation of graphs, we are ready to perform some interesting graph computations! First, recall this graph method from Section 15.3:

```
class Graph:
    def adjacent(self, item1: Any, item2: Any) -> bool:
        """Return whether item1 and item2 are adjacent vertices in this
        graph.

    Return False if item1 or item2 do not appear as vertices in this
        graph.
        """

    if item1 in self._vertices and item2 in self._vertices:
        v1 = self._vertices[item1]
        return any(v2.item == item2 for v2 in v1.neighbours)
    else:
        # We didn't find an existing vertex for both items.
        return False
```

Our goal in this section will be to implement a generalization of this method that goes from checking whether two vertices are adjacent to whether two vertices are connected:<sup>1</sup>

Recall that two vertices are connected when there exists a path between them.

```
class Graph:
    def connected(self, item1: Any, item2: Any) -> bool:
        """Return whether item1 and item2 are connected vertices in this graph.

    Return False if item1 or item2 do not appear as vertices in this graph.

>>> g = Graph()
>>> g.add_vertex(1)
>>> g.add_vertex(2)
>>> g.add_vertex(3)
>>> g.add_vertex(4)
>>> g.add_edge(1, 2)
>>> g.add_edge(2, 3)
>>> g.connected(1, 3)
True
```

```
>>> g.connected(1, 4)
False
"""
```

We can start implementing this method in the same fashion as Graph.adjacent, by finding the vertex corresponding to item1:

```
class Graph:
    def connected(self, item1: Any, item2: Any) -> bool:
        """..."""
        if item1 in self._vertices and item2 in self._vertices:
            v1 = self._vertices[item1]
            ...
        else:
            return False
```

The problem, of course, is in the final else branch. We can't simply check whether item2 is an immediate neighbour of item1, since there are many other possible cases for two vertices being connected: item2 could equal item1 itself, or it could be a neighbour, or the neighbour of a neighbour, or the neighbour of a neighbour, etc.

So you might be tempted to start writing code that tries to cover all these cases:

```
class Graph:
    def connected(self, item1: Any, item2: Any) -> bool:
        """..."""
    if item1 in self._vertices and item2 in self._vertices:
        v1 = self._vertices[item1]
        if v1.item == item2:
            return True
        elif any(u.item == item2 for u in v1.neighbours):
            return True
        elif any(any(u1.item == item2 for u1 in u.neighbours)
            for u in v1.neighbours):
            return True
        ...
        else:
        return False
```

You can see why this gets unmanageable very quickly, and fact cannot possibly cover all possibilities! Because we need to handle an arbitrary number of "neighbour" links, we can't write a fixed amount of code that enumerates the possibilities; instead, we either need to use a loop or recursion. We'll develop a recursive approach in this section, although it is possible to translate this approach into one that uses loops as well.

# Recursively determining connectivity

Here is our "complete" implementation of Graph.connected. As we did in earlier chapters, we're going to move the trickiest part (filling in that else branch) into a helper method.

```
class Graph:
    def connected(self, item1: Any, item2: Any) -> bool:
        """..."""
        if item1 in self._vertices and item2 in self._vertices:
            v1 = self._vertices[item1]
            return v1.check_connected(item2, set()) # Pass in an empty
        "visited" set
        else:
            return False
```

Note that our helper method, check\_connected is called with vertex to the left of the dot—that means that it's a method of the \_Vertex class, not Graph. Here is its initial specification:

```
class _Vertex:
    def check_connected(self, target_item: Any) -> bool:
        """Return whether this vertex is connected to a vertex corresponding
            to the target_item.
            """
```

Now let's recall our definition of connectedness: two vertices are connected when there exists a path between them. The challenge with this definition is that it only requires that a path exists, but we don't know how long it is. But there's another, equivalent way of defining connectedness *recursively*. Given two vertices  $v_1$  and  $v_2$ , they are connected when:

- $v_1 = v_2$ , or
- there exists a neighbour u of  $v_1$  such that u and  $v_2$  are connected.

This definition maps onto the recursive definition structure we've seen so far, with the first point being the base case and the second being the recursive part. However, this definition is different from all the previous recursive definitions in one important respect: it isn't exactly *structural*, since it doesn't break down the data type—a graph—into a smaller instance with the same structure. We'll come back to this point further down below.

But first, well, let's take this recursive definition and run with it! Even though we're still new to graphs, everything we've learned about translating recursive definitions into Python still applies, and our code actually turns out to be surprisingly simple:<sup>2</sup>

<sup>&</sup>lt;sup>2</sup> This version uses a loop, but you could certainly write a version with the built-in any function as well.

Pretty simple! But unfortunately, our implementation has a flaw that leads to a new kind of error:

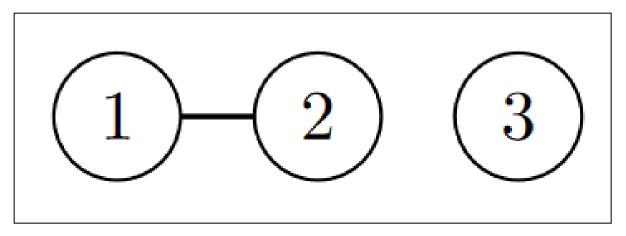
```
>>> g = Graph()
>>> g.add_vertex(1)
>>> g.add vertex(2)
>>> g.add edge(1, 2)
>>> g.connected(1, 3) # Should return False!
Traceback (most recent call last):
  File "<input>", line 1, in <module>
 File "path/to/graph.py", line 163, in connected
    return vertex.check_connected(item2)
 File "path/to/graph.py", line 41, in check_connected
    if u.check_connected(target_item):
  File "path/to/graph.py", line 41, in check connected
    if u.check_connected(target_item):
 File "path/to/graph.py", line 41, in check_connected
    if u.check_connected(target_item):
  [Previous line repeated 985 more times]
 File "path/to/graph.py", line 36, in check_connected
    if self.item == target item:
RecursionError: maximum recursion depth exceeded in comparison
```

What's going on?

# RecursionError and the danger of infinite recursion

Our implementation for \_Vertex.check\_connected is simple, and exactly follows the recursive definition of connectedness we gave above. So what we're running into here is not an error in our logic, but a limitation of how we've expressed that logic in our code. To understand what's gone wrong, let's do a bit of manual tracing of our example above.

First, we have a graph g with three vertices, containing the items 1, 2, and 3. In our initial call to g.connected(1, 3), we first find the vertex corresponding to 1, which we'll name  $v_1$ . Then our code calls \_Vertex.check\_connected on  $v_1$  and 3. This is where the problem begins.



- 1. For  $v_1$ , self.item == 1, which is not 3. We enter the else branch, where the loop (for u in self.neighbours) executes. The only neighbour is the vertex containing 2, which we'll name  $v_2$ . We make a recursive call to \_Vertex.check\_connected on  $v_2$  and 3.
- 2. For  $v_2$ , self.item == 2, which is not 3. We enter the else branch, where the loop (for u in self.neighbours) executes. The only neighbour is  $v_1$ . We make a recursive call to \_Vertex.check\_connected on  $v_1$  and 3.
- 3. For  $v_1$ , self.item == 1, which is not 3. We enter the else branch...

Hey, wait a second! We're back to the first step, repeating a call to \_Vertex.check\_connected with arguments  $v_1$  and 3. That's not good, because we know what happens in this case: we recurse on  $v_2$ .

This is an instance of **infinite recursion**, meaning we've expressed a recursive computation that does not stop by reaching a base case. In our example, the Python interpreter alternates between calling <code>\_Vertex.check\_connected</code> with  $v_1$  and 3 and with  $v_2$  and 3, and never stops. From a technical point of view, each of these function calls adds a new *stack frame* onto the function call stack, which uses up more and more computer memory. To prevent your computer from running out of memory, the Python interpreter enforces a limit on the total number of stack frames that can be on the call stack at any one point in time—and when that limit is reached, the Python interpreter raises a RecursionError.

## Fixing \_Vertex.check\_connected

Okay, so now that we understand what the problem is, how do we fix it? Intuitively, we need to prevent repeated recursive calls to the same node. Before changing our code, let's first revisit our recursive definition of connectedness.

Given two vertices  $v_1$  and  $v_2$ , they are connected when:

• there exists a neighbour u of  $v_1$  such that u and  $v_2$  are connected.

Our problem is in the recursive step. When we say that "u and  $v_2$  are connected", we're allowing for the possibility that these two vertices are connected by a path going through  $v_1$ . But that's not necessary: if  $v_1$  and  $v_2$  are connected, then we should be able to find a path between a neighbour u and  $v_2$  that doesn't use  $v_1$ , and then add  $v_1$  to the start of that path.

So we can modify this definition as follows. Given two vertices  $v_1$  and  $v_2$ , they are connected when:

- $v_1 = v_2$ , or
- there exists a neighbour u of  $v_1$  such that u and  $v_2$  are connected by a path that does not use  $v_1$ .

One way to visualize is that in the recursive step, we should be able to *remove*  $v_1$  *from the graph* and still find a path between u and  $v_2$ . And this is how we can take our original recursive definition and impose some structure on it: we take the original graph and make it smaller by "removing" the vertex  $v_1$ , and then recursively checking for the connectivity of each neighbour u and  $v_2$ .

#### Adding a visited parameter

Now, let's take this idea and apply it to fix our code. It might be tempting to actually mutate our Graph object to remove the " $v_1$ " vertex at each recursive call, but we don't actually want to mutate the original graph when we call Graph.connected. So instead, we'll use a common technique when traversing graphs: keep track of the items that have been already visited by our algorithm, so that we don't visit the same vertex more than once.

Here is our new \_Vertex.check\_connected specification:

```
class _Vertex:
    def check_connected(self, target_item: Any, visited: set[_Vertex]) ->
        bool:
        """Return whether this vertex is connected to a vertex corresponding
        to the target_item,
        WITHOUT using any of the vertices in visited.

        Preconditions:
            - self not in visited
        """
```

And before getting to our updated implementation, let's update how we call this helper method in Graph.connected. When we make the initial call to \_Vertex.check\_connected, we haven't yet visited any vertices:

```
class Graph:
    def connected(self, item1: Any, item2: Any) -> bool:
        """..."""
        if item1 in self._vertices and item2 in self._vertices:
            v1 = self._vertices[item1]
            return v1.check_connected(item2, set()) # Pass in an empty
        "visited" set
        else:
            return False
```

And finally, let's modify our implementation of \_Vertex.check\_connected to use this new parameter. We need to make two changes: first, we add self to visited before making any recursive calls (to indicate that the current \_Vertex has been visited by our algorithm); second, when looping over self.neighbours, we only make recursive calls into nodes that have not yet been visited.

```
class _Vertex:
    def check_connected(self, target_item: Any, visited: set[_Vertex]) ->
        """Return whether this vertex is connected to a vertex corresponding
         to the target item,
        WITHOUT using any of the vertices in visited.
        Preconditions:
            - self not in visited
        if self.item == target_item:
            # Our base case: the target item is the current vertex
            return True
        else:
            new_visited = visited.union({self}) # Add self to the set of
         visited vertices
            for u in self.neighbours:
                if u not in new visited: # Only recurse on vertices that
         haven't been visited
                    if u.check_connected(target_item, new_visited):
                        return True
            return False
```

And with this version, we've eliminated our infinite recursion error:

```
>>> g = Graph()
>>> g.add_vertex(1)
>>> g.add_vertex(2)
>>> g.add_edge(1, 2)
>>> g.connected(1, 3) # Should return False!
False
```

### Bonus optimization

One last thing: you might have noticed that in our final implementation of \_Vertex.check\_connected, we created a new set of visited vertices (new\_visited) rather than mutating visited, which would have been simpler and more efficient. Consider this alternate version:

```
class _Vertex:
    def check connected(self, target item: Any, visited: set[ Vertex]) ->
        """Return whether this vertex is connected to a vertex corresponding
        to the target item,
        WITHOUT using any of the vertices in visited.
        Preconditions:
           - self not in visited
        if self.item == target_item:
            # Our base case: the target item is the current vertex
            return True
        else:
            visited.add(self) # Add self to the set of visited
         vertices
           for u in self.neighbours:
                if u not in visited: # Only recurse on vertices that haven't
         been visited
                    if u.check_connected(target_item, visited):
                        return True
            return False
```

This call uses our familiar set.add method, which is indeed constant time. However, there is an important difference with this version: now there is only one visited set object that gets shared and mutated across all recursive calls, including ones that return False.

So for example, suppose self has two neighbours u0 and u1, and we first recursive on u0. If u0 is not connected to the target vertex, then the check\_connected call will return False, as expected. But u0 would still be in visited after that recursive call ends, meaning the subsequent check\_connected call with u1 will never visit u0. This is somewhat surprising, but technically still correct: once we've established that u0 is not connected to the target item, we *shouldn't* recurse on it ever again.

But even though this implementation is correct, and much more efficient than the previous one, it does come with an important warning we want you to remember for the future. Whenever you use recursion with a mutable argument, be very careful when choosing

whether to mutate that argument or create a modified copy—if you choose to mutate the argument, know that all recursive calls will mutate it as well!

CSC110/111 Course Notes Home