



UNIVERSITY OF  
BIRMINGHAM

# ARTIFICIAL INTELLIGENCE 1 INFORMED SEARCH

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SCHOOL OF COMPUTER SCIENCE

2023/2024 - Week 9

# AIMS OF THE SESSION

This session aims to help you:

- Understand the concept of a heuristic function
- Explain the steps involved in  $A^*$
- Analyse the performance of  $A^*$  and apply the algorithm to solve search problems

# OUTLINE

1 Recap: Uninformed Search

2 Informed Search

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# RECAP: QUIZ

Q1. What does it mean for an environment to be **known**?

- ☐ There are only finitely many actions at any state
- ☐ Possible to determine which states are reached by which action
- ☐ The agent is able to know the current state
- ☐ Each action has exactly one outcome

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- Q2. If we relax the assumption for an environment to be discrete, BFS is still complete.
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- ☐ True
  - ☒ False, as  $b$  would be infinite



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- ☐ Initial state
- ☐ Goal test
- ☐ Path cost function
- ☐ Action set

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- ☐ False

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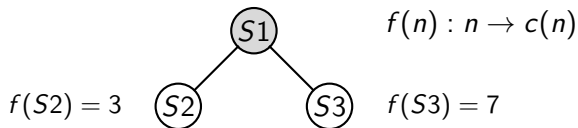
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# INFORMED SEARCH

- **Informed search** strategies use problem-specific knowledge beyond the definition of the problem itself
- In general, informed search strategies can find solutions more efficiently compared to uninformed search
- The general approach, called best-first search, is to determine which node to expand based on an evaluation function  $f(n)$



- This function acts as a cost estimate: the node with the lowest cost is the one that is expanded next

# HEURISTICS

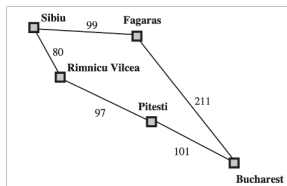
- The evaluation function  $f(n)$  for most best-first search algorithms includes a heuristic function as a component:  
 $h(n)$ : estimated cost of the cheapest path from node  $n$  to a goal node
- Heuristic functions are the most common form in which new knowledge is given to the search algorithm. If  $n$  is a goal node, then

$$h(n) = 0$$

- A heuristic can be a rule of thumb, common knowledge; it is quick to compute, but not guaranteed to work (nor to yield optimal solutions)

# HEURISTICS

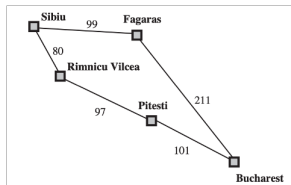
- Consider the following problem, where we want to find the shortest path to Bucharest (from, e.g., Sibiu) in Romania
- A possible heuristic is the straight-line distance heuristic,  $h_{SLD}$
- This is a useful heuristic as it is correlated with actual road distances





# HEURISTICS

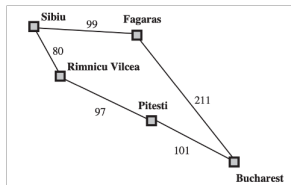
- Consider the following problem, where we want to find the shortest path to Bucharest (from, e.g., Sibiu) in Romania
- The straight-line distances  $h_{SLD}(n)$  are shown in the table below
- For example, the SLD from Sibiu would be  $h_{SLD}(\text{Sibiu}) = 253$



Arad	366	Mehadia	241
Bucharest	0	Neamt	234
Craiova	160	Oradea	380
Drobeta	242	Pitesti	100
Eforie	161	Rimnicu Vilcea	193
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Lugoj	244	Zerind	374

# GREEDY BEST-FIRST SEARCH

- Consider the following problem, where we want to find the shortest path to Bucharest (from, e.g., Sibiu) in Romania
- If we use  $f(n) = h_{SLD}(n)$ , then from Sibiu we expand Fagaras
- This is because  $h_{SLD}(Fagaras) = 176 < h_{SLD}(Rimnicu\ Vilcea) = 193$
- When  $f(n) = h(n)$ , we call this strategy **greedy best-first search**



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# A<sup>\*</sup> SEARCH

- The most widely known informed search strategy is **A<sup>\*</sup>**
- This search strategy evaluates nodes using the following evaluation function

$$f(n) = g(n) + h(n)$$

where  $g(n)$  is the cost to reach node  $n$  and  $h(n)$  is the heuristic from node  $n$  to the goal

- This is equivalent to the cost of the cheapest solution through node  $n$

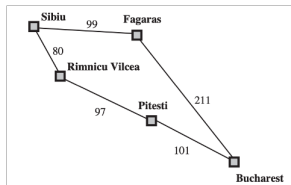
# A\* SEARCH

- Steps of the A\* algorithm:
  - **Expand** the node in the frontier with smallest  $f(n) = g(n) + h(n)$
  - **Repeated states and loopy paths.** If a node is in the list of visited nodes, **do not add** it to the frontier.  
If the state of a given child is in the frontier:
    - If the frontier node has a larger  $g(n)$ , place the child into the frontier and remove the node with larger  $g(n)$  from the frontier
  - **Stop** when a goal node is visited

# A\* SEARCH

**Activity.** Consider the following problem, where we want to find the shortest path to Bucharest from Sibiu in Romania. In pairs or small groups, calculate  $f(n)$  at each step of the A\* algorithm.

$$f(\text{Fagaras}) = g(\text{Fagaras}) + h(\text{Fagaras})$$



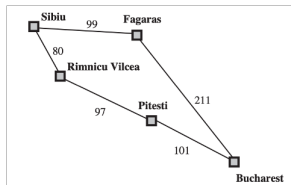
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# A\* SEARCH

**Activity.** Consider the following problem, where we want to find the shortest path to Bucharest from Sibiu in Romania. In pairs or small groups, calculate  $f(n)$  at each step of the A\* algorithm.

$$f(\text{Fagaras}) = 99 + 176 = 275$$

$$f(\text{Rimnicu Vilcea}) = g(\text{Rimnicu Vilcea}) + h(\text{Rimnicu Vilcea})$$



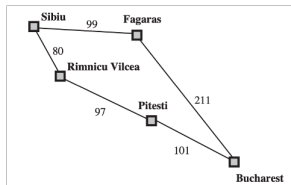
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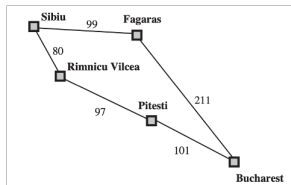
$$f(\text{Fagaras}) = 99 + 176 = 275$$

$$f(\text{Rimnicu Vilcea}) = 80 + 193 = 273$$

$$f(\text{Bucharest}) = (99 + 211) + 0 = 310$$

$$f(\text{Pitesti}) = (80 + 97) + 100 = 277$$

$$f(\text{Bucharest}) = (80 + 97 + 101) + 0 = 278$$

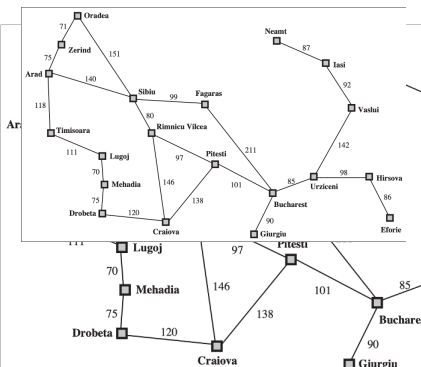


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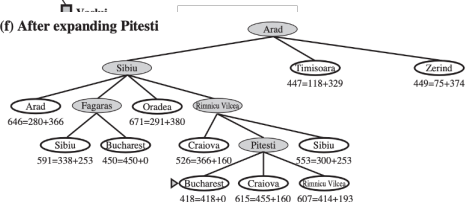
# A\* SEARCH

**Activity.** Consider the following problem, where we want to find the shortest path to Bucharest from Arad in Romania. In pairs or small groups, determine the solution retrieved by A\*.



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(f) After expanding Pitesti



# PERFORMANCE OF $A^*$

- $A^*$  is **complete** and **optimal** if  $h(n)$  is **consistent**

**Definition:** A heuristic is said to be **consistent** (or **monotone**), if the estimate is always no greater than the estimated distance from any neighbouring node  $n'$  to the goal, plus the cost of reaching that neighbour:

$$h(n) \leq \text{cost}(n, n') + h(n')$$

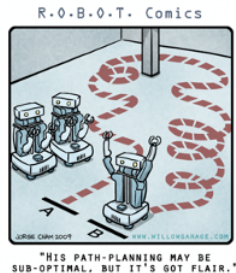
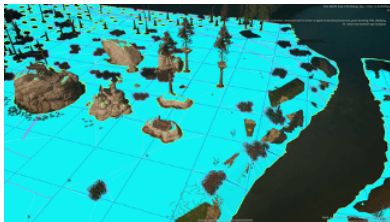
- The number of states generated by  $A^*$  is exponential in the length of the solution, namely for constant step costs:  $\mathcal{O}(b^{\epsilon d})$
- If  $h^*$  is the actual cost from root node to goal node,  $\epsilon = \frac{h^* - h}{h^*}$  is the relative error
- Space is the main issue with  $A^*$ , as it keeps all generated nodes in memory, not suitable for large-scale problems

# PERFORMANCE OF $A^*$

- Let us summarise the performance of  $A^*$ .
  - **Completeness:** if the heuristic  $h(n)$  is consistent, then  $A^*$  is complete
  - **Optimality:** if the heuristic  $h(n)$  is consistent, then  $A^*$  is optimal
  - **Time complexity:**  $\mathcal{O}(b^{\epsilon d})$ , where  $\epsilon$  is the relative error of the heuristic
  - **Space complexity:**  $\mathcal{O}(b^d)$ , since we keep in memory all expanded nodes and all nodes in the frontier

# APPLICATIONS OF $A^*$

- $A^*$  is one of the most widely adopted search algorithms
- The most common settings are in games and in robotics



# References



Russell, A. S., and Norvig, P., *Artificial Intelligence A Modern Approach*, 4<sup>th</sup> Edition. Prentice Hall.



Chapter 3 – “Solving Problems by Searching”, Section 3.5 “Informed (Heuristic) Search Strategies” up to and including Section 3.6.2 “Generating Admissible Heuristics From Relaxed Problems”.

# AIMS OF THE SESSION

You should now be able to:

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Thank you!