OSU Micro-Benchmark (OMB) Latency/Bandwidth Evaluation using ReFrame in UniLu HPC

Fabio, Shahmeer, Wei

MHPC

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Outline

- 1. Approach
- 2. Methodology
- 3. Results
- 4. Key takeaways

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According to the project requirements, it is clear that we need to compile and load the OSU Micro-Benchmark in three different ways, and then benchmark it on four system architectures across two different clusters.

Fetching the binary files is not particularly difficult, but using the ReFrame framework to handle everything from pulling source code, to compilation, to execution, and finally report generation is an interesting challenge.

Of course, we could split this into two steps: compile first, then hard-code the compiled binary paths into the ReFrame test. However, this approach lacks elegance.

Since benchmarking involves different code sources, different system topology parameters, different clusters, and different performance metrics, we explored ReFrame's documentation and eventually arrived at a clean and extensible solution for injecting these parameters:

Different code sources	Use fixtures to automate fetching and compiling from source	
Different ton alegar	Inject custom topology parameters using custom-defined	
Different topology	env_vars	
Different clusters	Configure cluster-specific logic via ReFrame's system	
	configuration	
Different metrics		

As a result, the final solution automatically handles the cluster-specific logic within the code. We only need to run a single command to execute the full benchmark on both Iris and Aion clusters.

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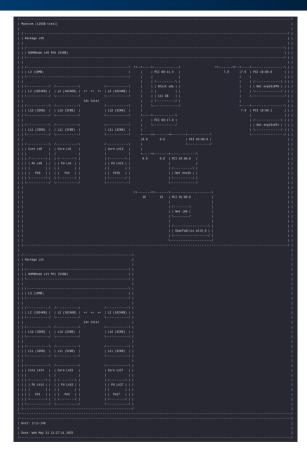


Figure 1: System topology of Iris cluster



Figure 2: System topology of Aion cluster

All tests are based on OSU Micro-Benchmarks, specifically:

- osu_latency for measuring latency, with a message size of 8192 Bytes
- osu_bw for measuring bandwidth, with a message size of 1MB (1048576 Bytes)

Metrics	Description
generic	Compiled from source
easybuild	Compiled with EasyBuild
eessi	Binaries loaded from the EESSI distribution

Metrics	Iris CPU binding strategy	Aion CPU binding strategy	Description
intra_numa	map_cpu:0,1	map_cpu:0,1	both processes are running on the same NUMA node
cross_numa	_	map_cpu:0,16	both processes are running on the same physical socket but different NUMA nodes
cross_socket	map_cpu:0,14	map_cpu:0,64	both processes are running on the same compute node but different sockets

			the 2 processes are
inter_node	cpu-bind=cores	cpu-bind=cores	running on different
			nodes

Each test uses 2 MPI processes, with the CPU binding strategy automatically configured based on the system architecture.

For example, map_cpu:0,1 corresponds to the intra_numa placement.

Specifically, we define the placement info in env_vars of config file for different systems(Aion and Iris).

```
// Aion's env_vars
                                                                           ison
   "env_vars": [
3
        ["intra numa", "map cpu:0,1"],
        ["cross numa", "map cpu:0,16"],
5
        ["cross_socket", "map_cpu:0,64"],
6
   ],
   "env_vars": [
        ["intra_numa", "map_cpu:0,1"],
        ["cross_socket", "map_cpu:0,14"],
10
```

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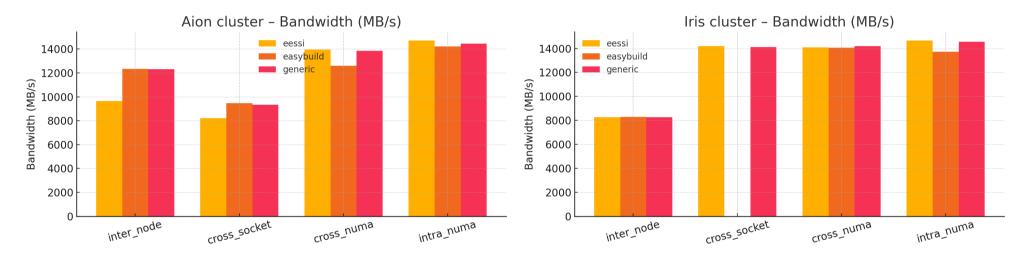


Figure 3: Bandwidth result in Aion & Iris cluster.

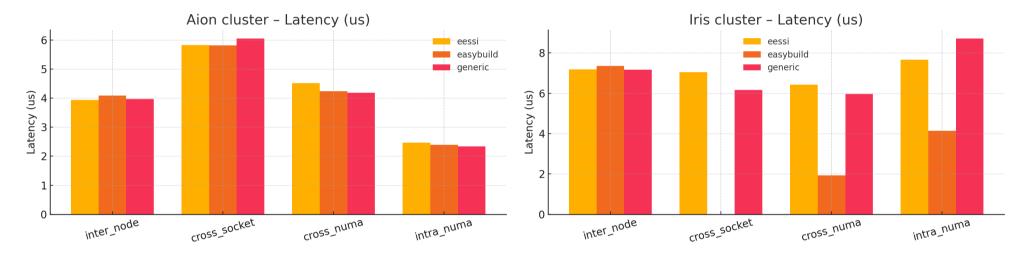
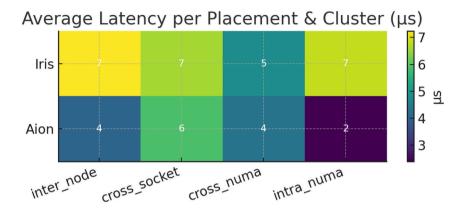


Figure 4: Latency result in Aion & Iris cluster.



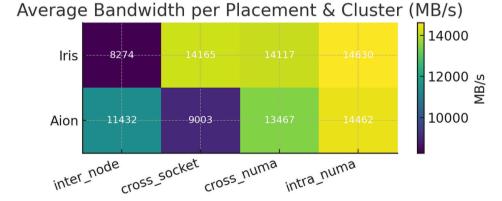


Figure 5: Average Bandwidth & Latency

From Figure 3 and Figure 4 we can observation:

- 1. Iris excels in local bandwidth (especially for cross_socket/cross_numa).
- 2. Aion shows lower latency, likely due to faster CPU scheduling or better NUMA balancing.

Figure 5 shows that for latency:

- 1. Iris: Averages 7 μs across most placements.
- 2. Aion: Significantly faster—down to 2.3 μs in intra_numa.

Also for bandwidth:

- 1. Iris dominates across all placements, with cross_numa/cross_socket 14.1 GB/s+
- 2. Aion achieves similar intra_numa bandwidth, but underperforms in inter_node (11.4 GB/s)

All tests ran smoothly on the Aion cluster. However, on the Iris cluster, the EasyBuild version of the binary occasionally exhibits instability:

```
--- rfm_job.out ----- rfm_job.err (last 10 lines) ---
                                                                        sh
    4 0x0000000000040440e start() ???:0
   ==== backtrace (tid:1240770) ====
5
    0 0x0000000000012d10 funlockfile() :0
6
    1 0x0000000000040cf3b omb ddt assign()
                                          ???:0
    2 0x00000000004037ac main() ???:0
    3 0x000000000003a7e5 libc start main() ???:0
9
    4 0x0000000000040440e start() ???:0
10
```

Some Clues:

- The error only occurs occasionally.
- The problem only affects the EasyBuild binary.
- omb_ddt_assign() is where the crash occurs: this function may use data-dependent logic or SIMD instructions.

Reasonable Explanation:

If the binary was compiled on a node with newer instruction sets, but later executed on an older Iris node without such capabilities, this would result in a SIGILL (illegal instruction) error like the one shown.

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4.1 Platform Recommendations:

- Aion is better suited for low-latency computing tasks (lower latency in intra- and cross-NUMA scenarios).
- Iris is better suited for **high-bandwidth communication** workloads (higher internode bandwidth observed).

4.2 Binary Selection Recommendations:

- Generic and EasyBuild binaries show balanced and stable performance across scenarios.
- EESSI binaries may present compatibility issues, although they perform well in certain placement cases.

4.3 Insights for Benchmark Design:

- Benchmarks should thoroughly account for system **topology** (NUMA domains, socket layout, and inter-node communication).
- Compatibility of high-performance libraries must not be overlooked; for cross-platform deployment, locally built or EasyBuild-based binaries are preferred.