

AVL – “self-balancing” binary tree

The heights h of the two child subtrees of any node x differ at most by one.

Balance Factor $bf(x) = (x \rightarrow \text{left} \rightarrow \text{height}) - (x \rightarrow \text{right} \rightarrow \text{height})$

$bf(x)$ must be -1, 0, or 1

After an insertion/deletion, if the heights of left and right subtrees of node x differ more than one, the tree must automatically rebalance.

```
class AVLTree
{
    public:
        class Node
        {
            public:
                Node(){}
                Node(int key) : key(key), height(1), right(nullptr), left(nullptr){}
                int key;
                int height;
                Node* right;
                Node* left;
        };
};
```

AVL

Insert 10 20 30

10

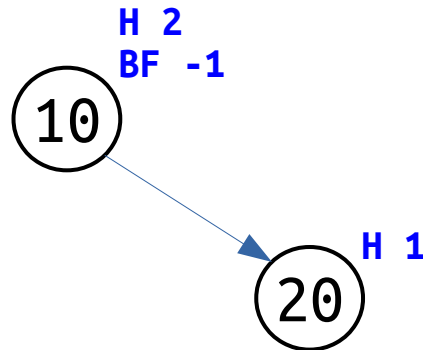
```
void insert(int key)
{
    _root = insert(_root, key);
}

// insert a key in the subtree rooted with "node"
// returns the new root of the subtree
Node* insert(Node* node, int key)
{
    // Step 1. Perform the insertion like in simple Binary Search Trees.
    if(node == nullptr)
    {
        return new Node(key);
    }

    if(key < node->key)
    {
        node->left = insert(node->left, key);
    }
    else if(key > node->key)
    {
        node->right = insert(node->right, key);
    }
    else
    {
        return node;
    }
}
```

AVL

Insert 10 20 30



```
int maxHeight(Node* left, Node* right)
{
    int l = left == nullptr ? 0 : left->height;
    int r = right == nullptr ? 0 : right->height;

    return l > r ? l : r;
}

int getBalanceFactor(Node* node)
{
    if(node == nullptr)
        return 0;

    int l = node->left == nullptr ? 0 : node->left->height;
    int r = node->right == nullptr ? 0 : node->right->height;

    return l - r;
}
```

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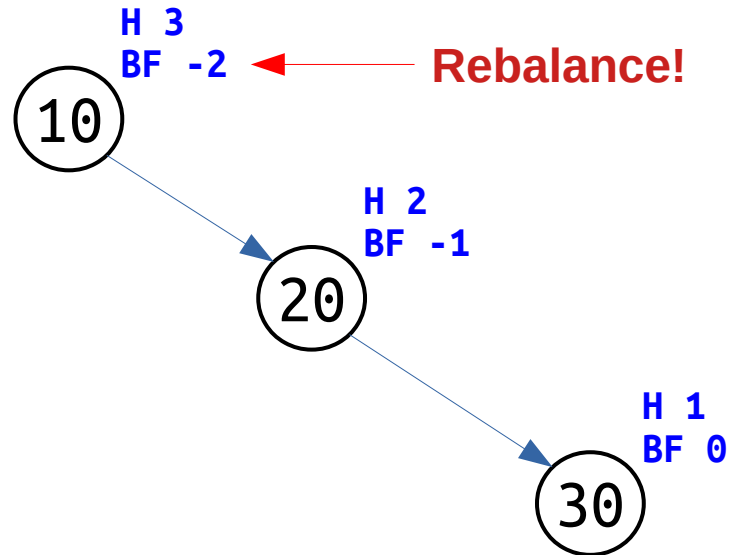
    if(key < node->key)
    {
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    }
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    {
        node->right = insert(node->right, key);
    }
    else
    {
        return node;
    }
}
```

```
// Step 2. Update height of this ancestor
node->height = 1 + maxHeight(node->left, node->right);
```

```
// Step 3. Check the balance factor of this ancestor
int bf = getBalanceFactor(node);
```

AVL

Insert 10 20 30



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    {
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    }
    else
    {
        return node;
    }
}
```

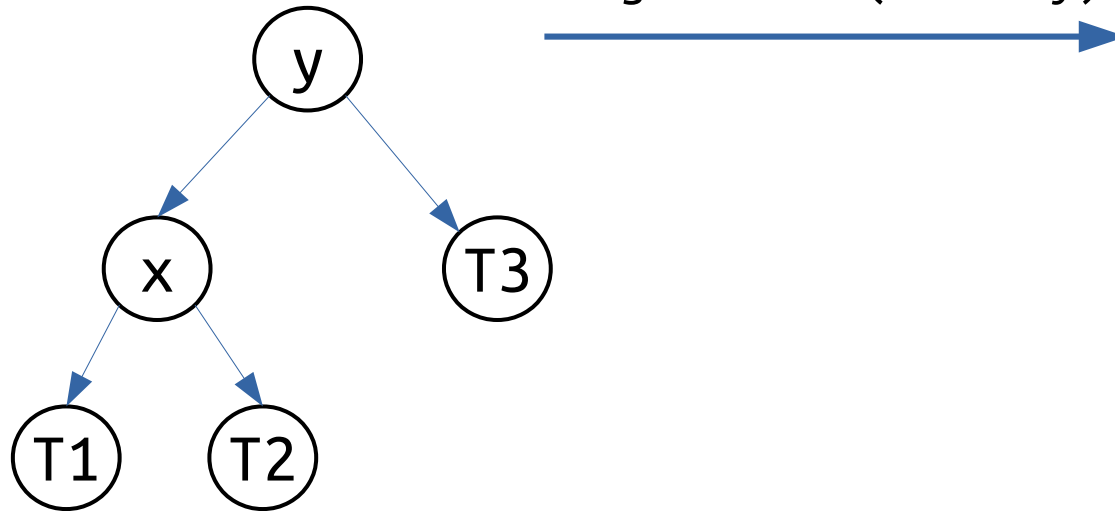
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// Step 4. If the subtree is unbalanced, check in which case we are, and balance!
```

AVL - Rotations

rightRotate(Node* y)



```
Node* rightRotate(Node *y)
```

```
//1. Let x be the left child of y
```

```
//2. Let t2 be the right child of x
```

```
--- rotation step ---
```

```
//3. The right child of x will be y
```

```
//4. The left child of y will be T2
```

```
--- update heights ---
```

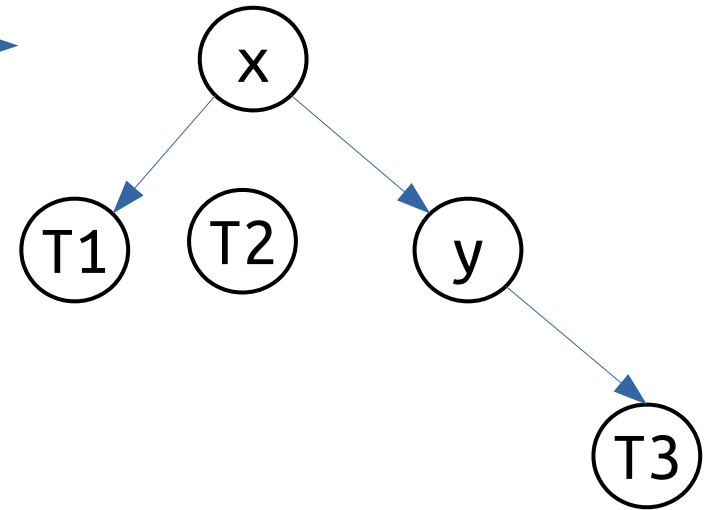
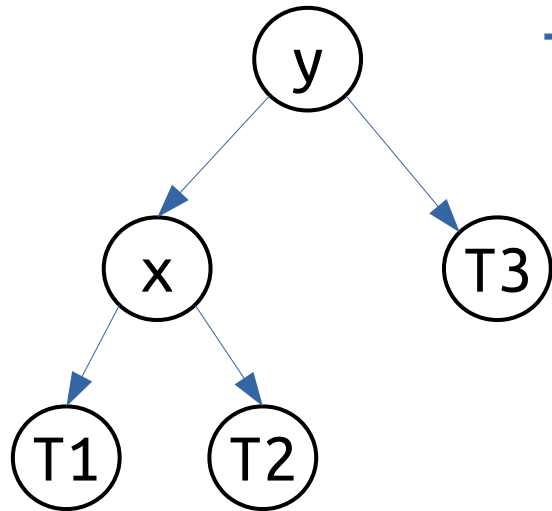
```
//5. height of y =(max height of his children) + 1
```

```
//6. height of x =(max height of x's children) + 1
```

```
return x;
```

AVL - Rotations

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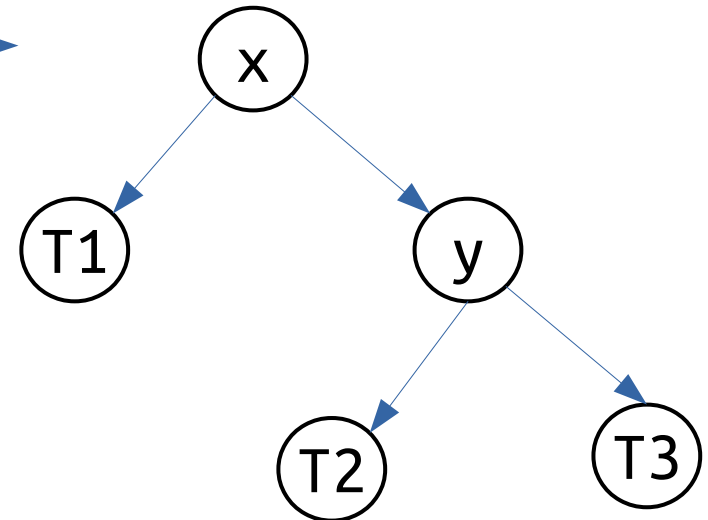
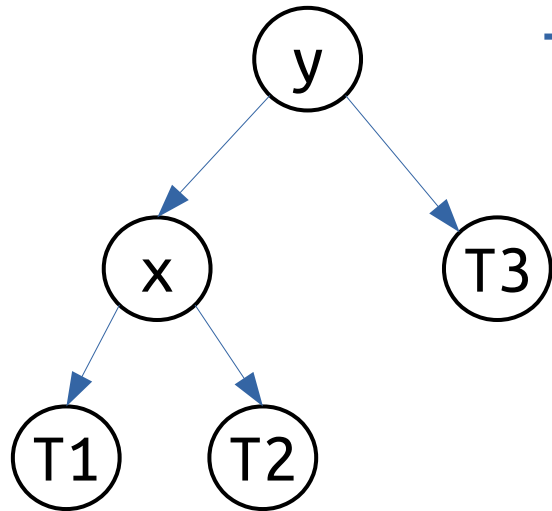
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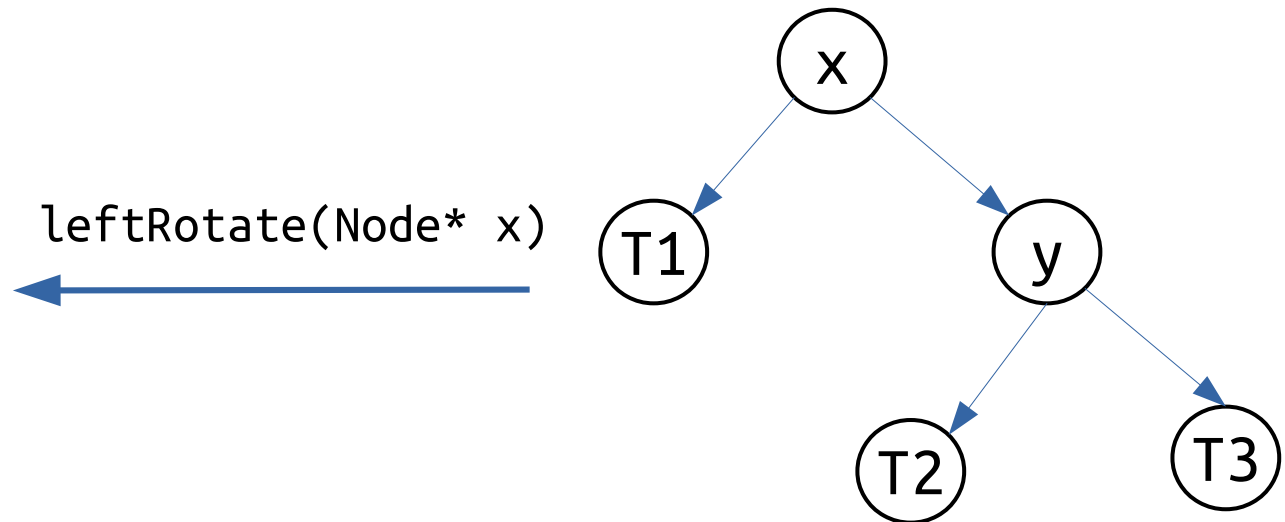
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AVL - Rotations



```
Node* leftRotate(Node *x)
```

```
//1. Let y be the right child of x
```

```
//2. Let t2 be the left child of x
```

```
--- rotation step ---
```

```
//3. The left child of y will be x
```

```
//4. The right child of x will be t2
```

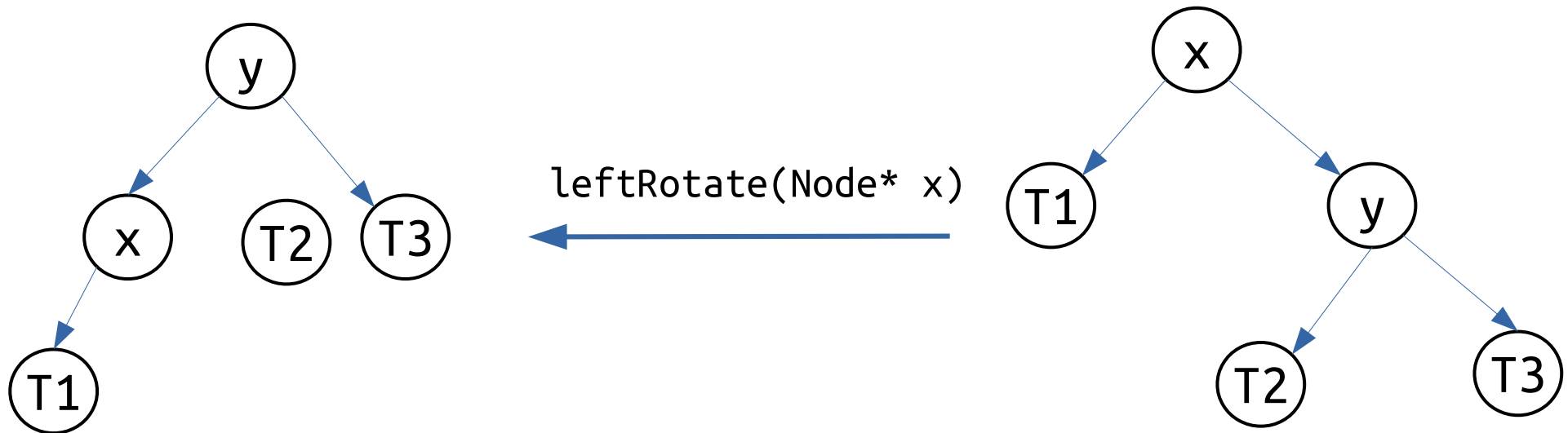
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AVL - Rotations



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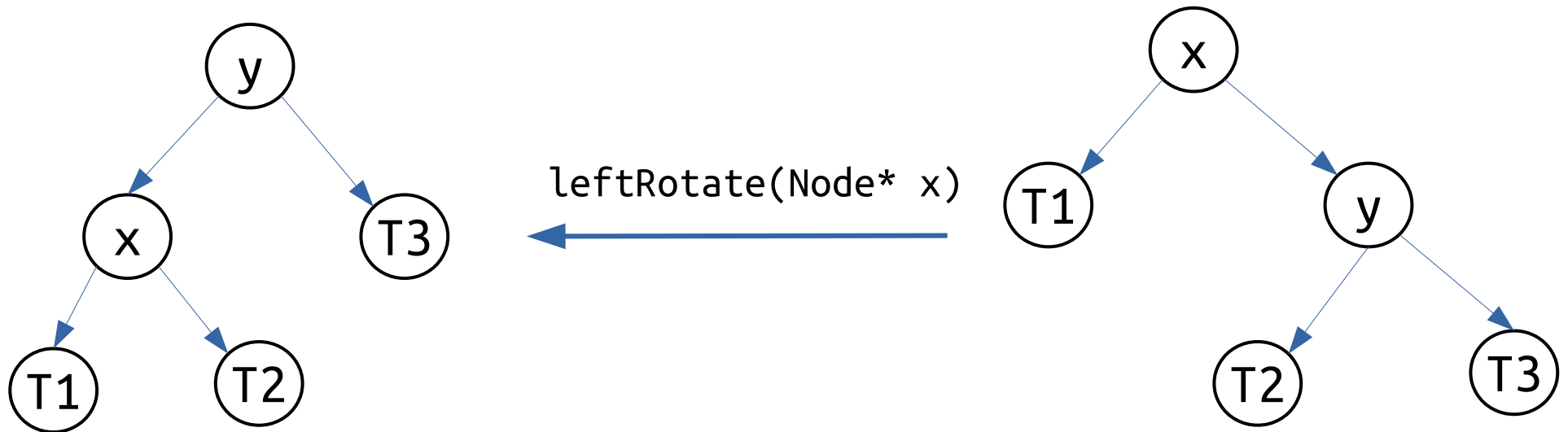
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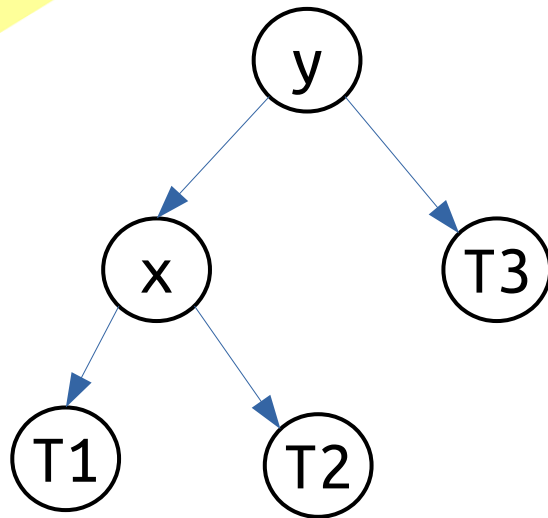
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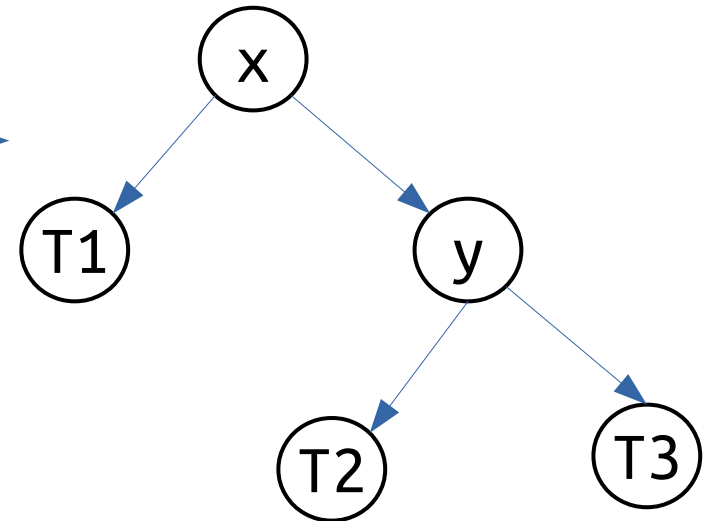
```
return y;
```



AVL - Rotations



rightRotate(Node* y)



leftRotate(Node* x)



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Node* rightRotate(Node *y)
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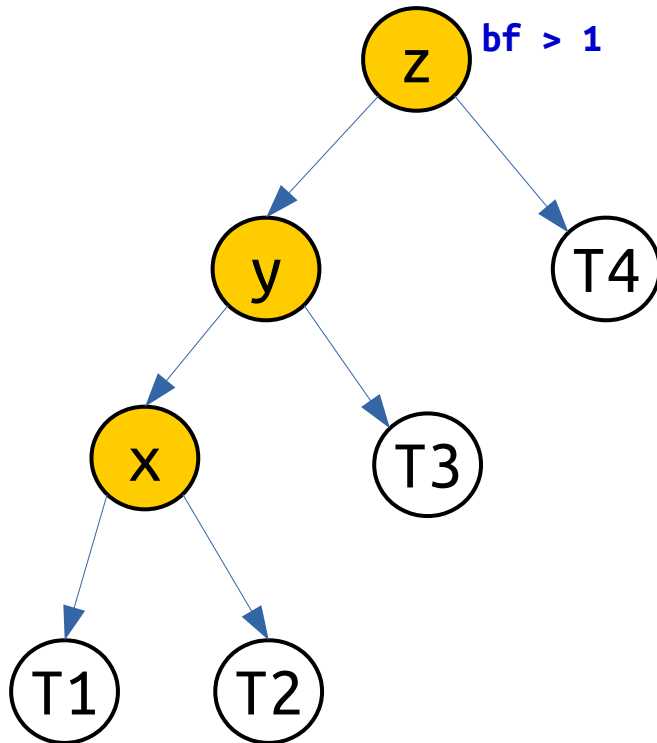
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return y;
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AVL - Cases

Case 1. LEFT LEFT

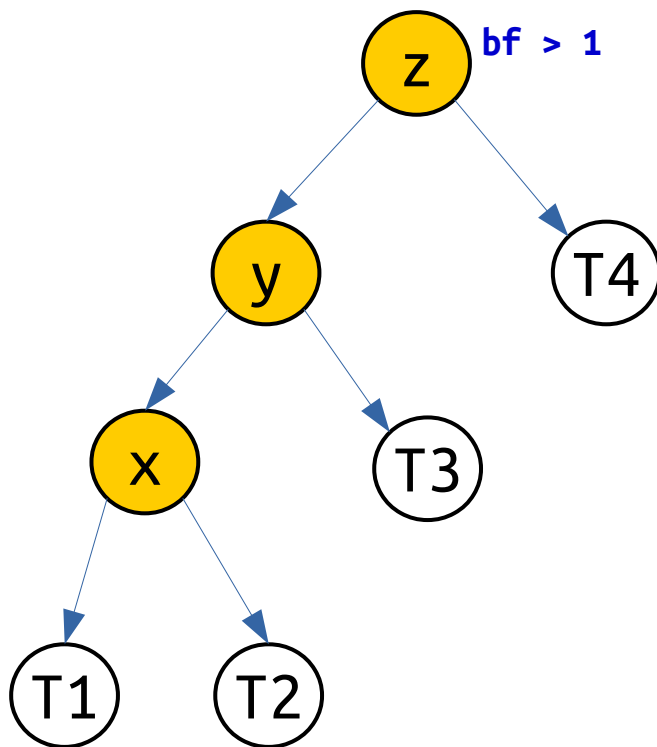
$bf > 1$ and $key < node \rightarrow left \rightarrow key$



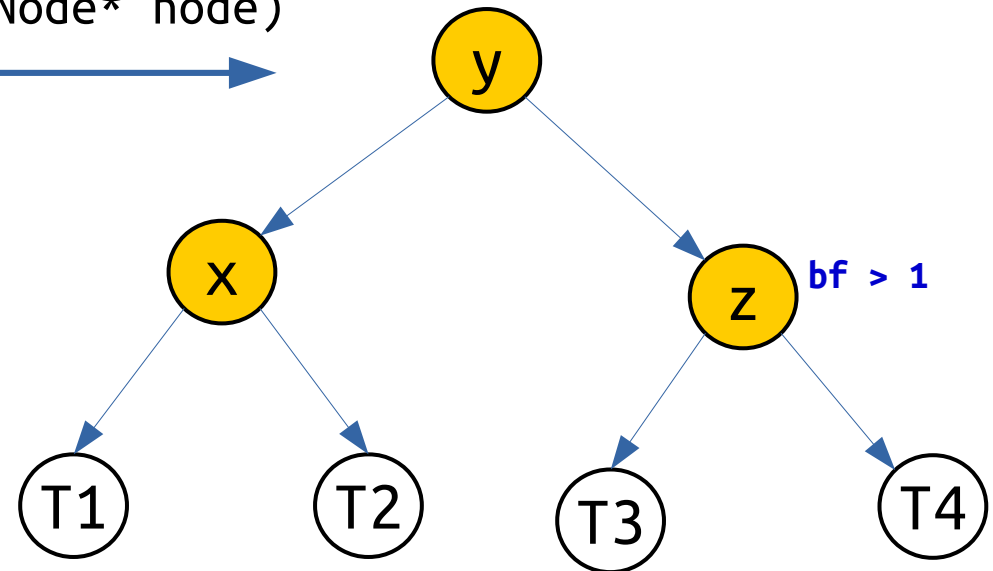
AVL - Cases

Case 1. LEFT LEFT

$bf > 1$ and $key < node \rightarrow left \rightarrow key$



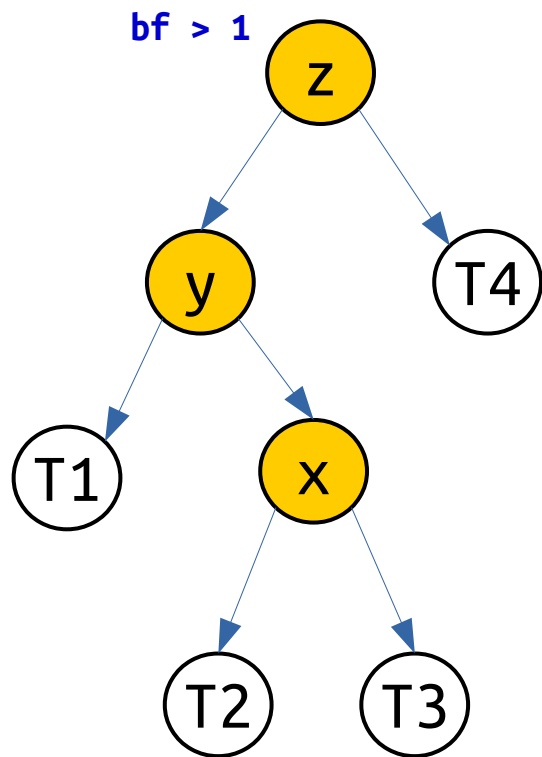
`rightRotate(Node* node)`



AVL - Cases

Case 2. LEFT RIGHT

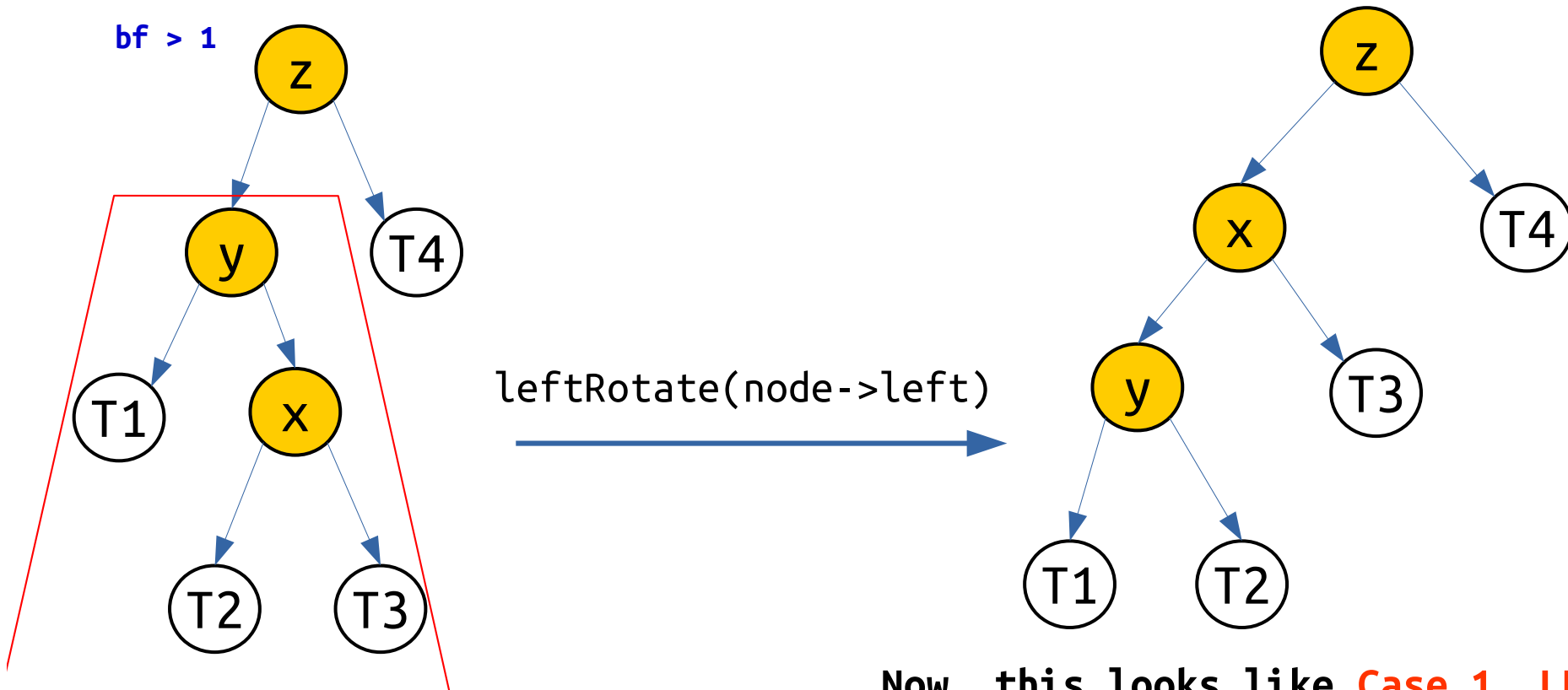
$bf > 1$ and $key > node \rightarrow left \rightarrow key$



AVL - Cases

Case 2. LEFT RIGHT

$bf > 1$ and $key > node \rightarrow left \rightarrow key$

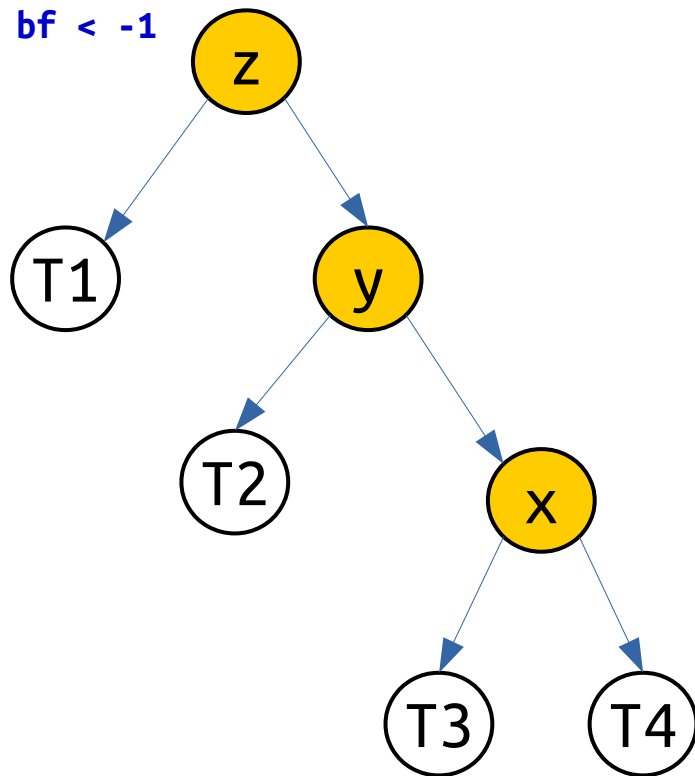


Now, this looks like **Case 1. LEFT LEFT**
Do right rotation on node z.

Case 3. RIGHT RIGHT

$bf < -1$ and $key > node \rightarrow right \rightarrow key$

$bf < -1$

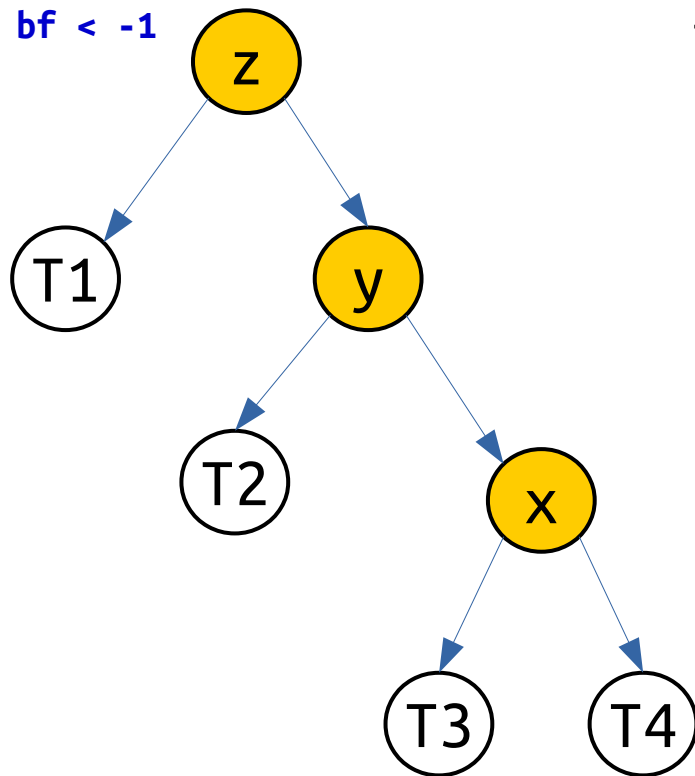


AVL - Cases

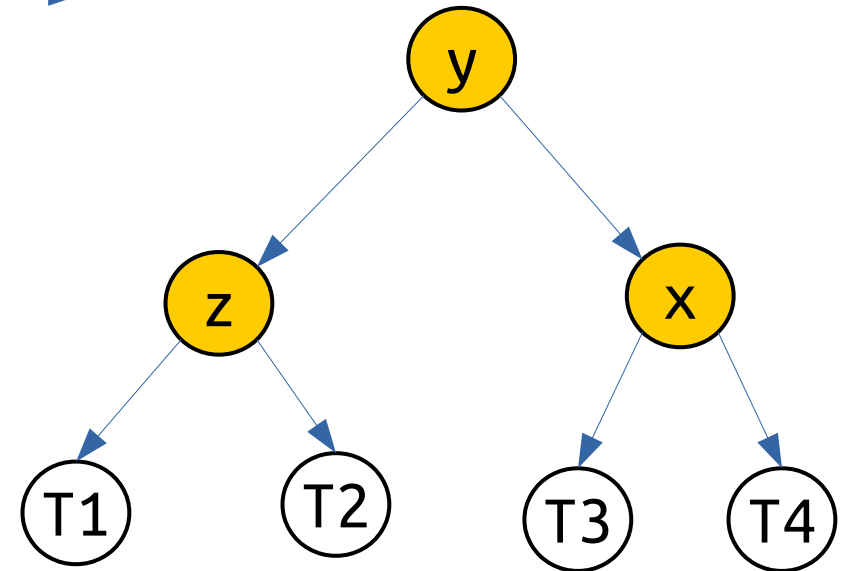
Case 3. RIGHT RIGHT

$bf < -1$ and $key > node \rightarrow right \rightarrow key$

$bf < -1$



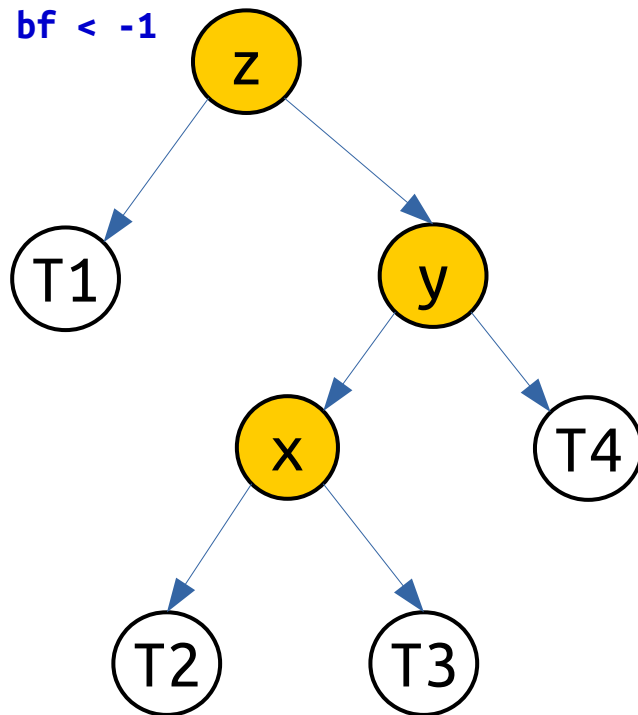
leftRotate(Node* node)



Case 4. RIGHT LEFT

$bf < -1$ and $key < node \rightarrow right \rightarrow key$

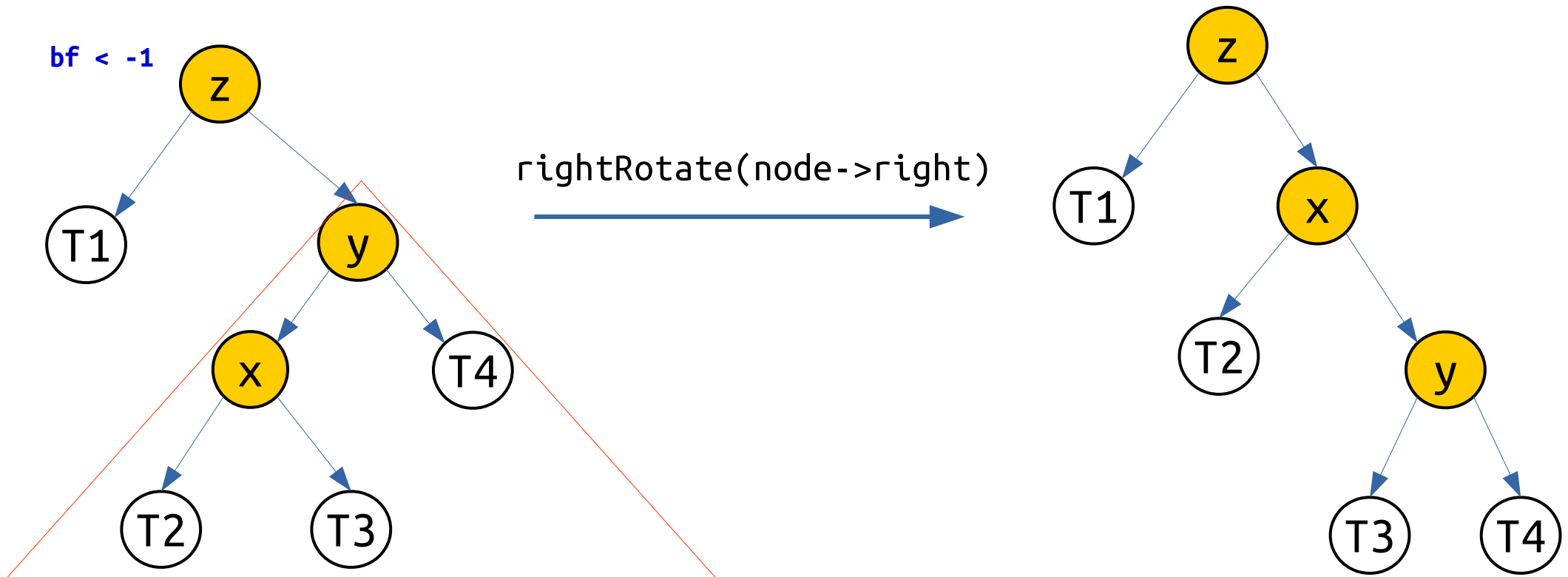
$bf < -1$



AVL - Cases

Case 4. RIGHT LEFT

$bf < -1$ and $key < node \rightarrow right \rightarrow key$



Now, this looks like **Case 3. RIGHT RIGHT**
Do left rotation on node z .

AVL - Cases

```
// Step 4. If the subtree is unbalanced, check in which case we are, and balance!

// Case 1. LEFT LEFT CASE
if(bf > 1 && key < node->left->key)
{
    return rightRotate(node);
}

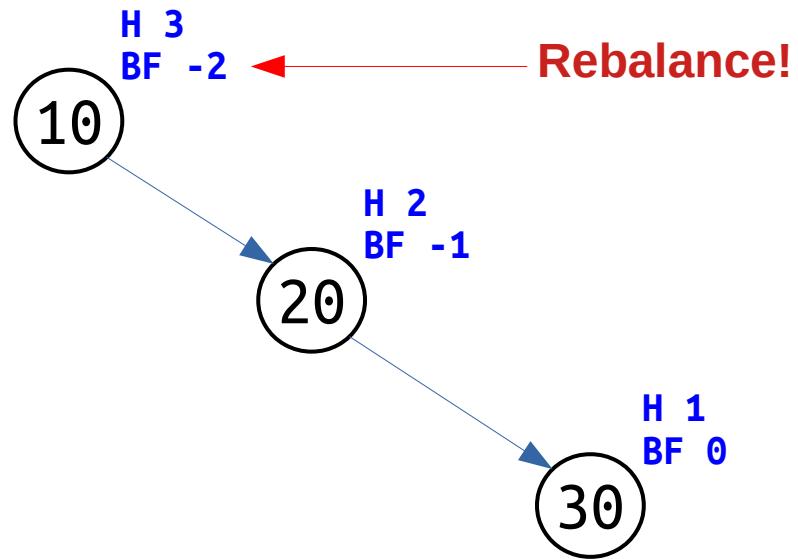
// Case 2. LEFT RIGHT CASE
if(bf > 1 && key > node->left->key)
{
    node->left = leftRotate(node->left);
    return rightRotate(node);
}

// Case 3. RIGHT RIGHT CASE
if(bf < -1 && key > node->right->key)
{
    return leftRotate(node);
}

// Case 4. RIGHT LEFT CASE
if(bf < -1 && key < node->right->key)
{
    node->right = rightRotate(node->right);
    return leftRotate(node);
}
```

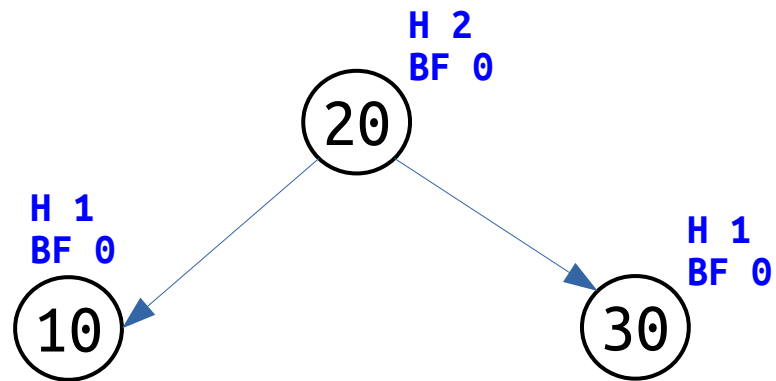
AVL

Insert 30



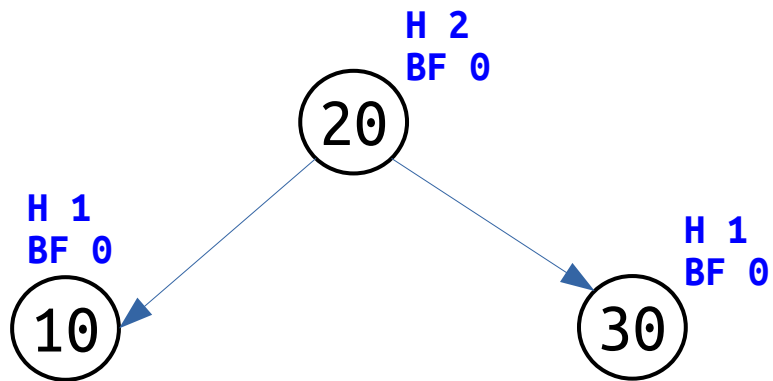
AVL

Insert 30



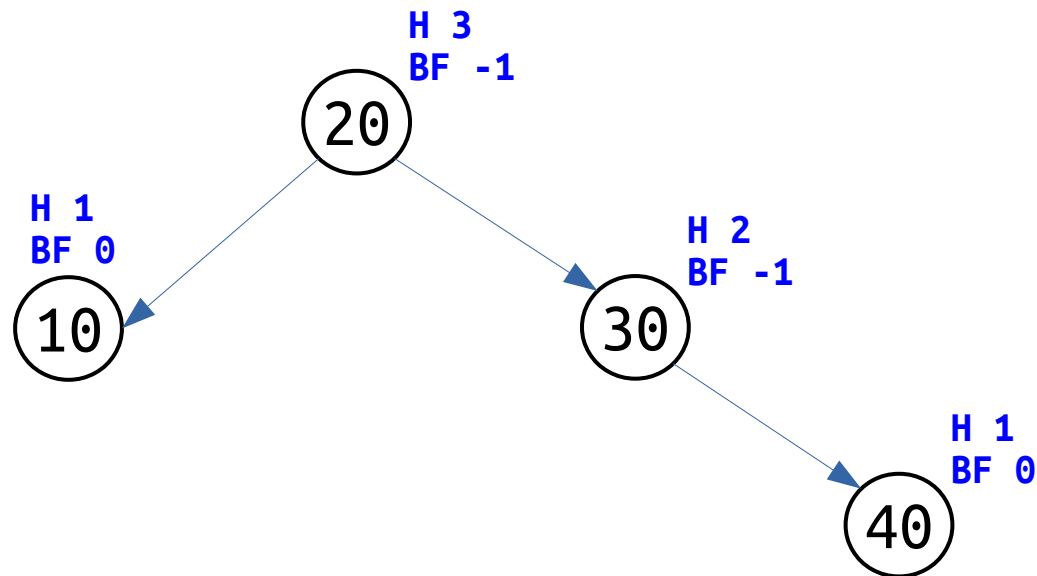
AVL

Insert 40



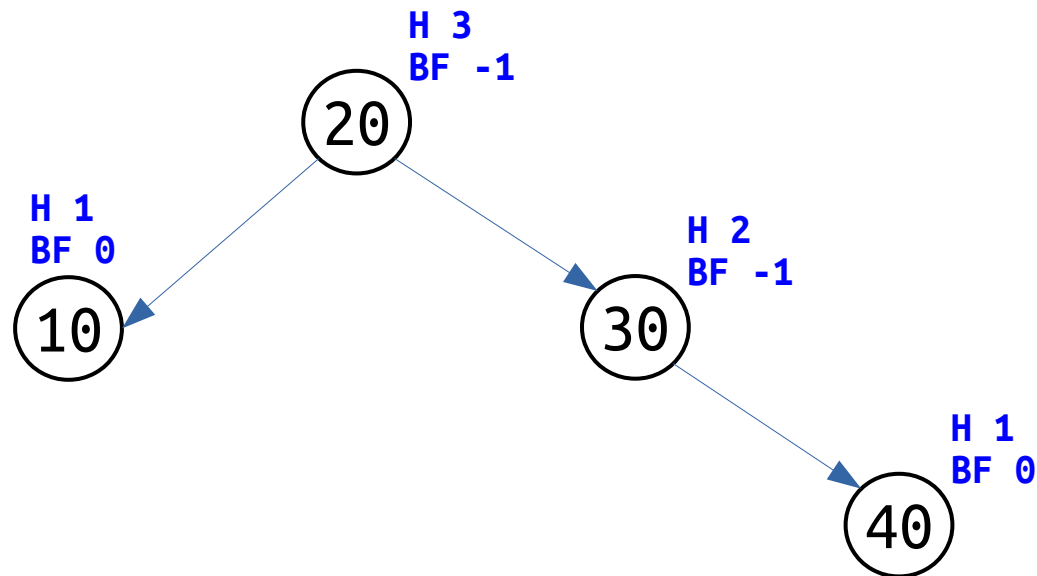
AVL

Insert **40**



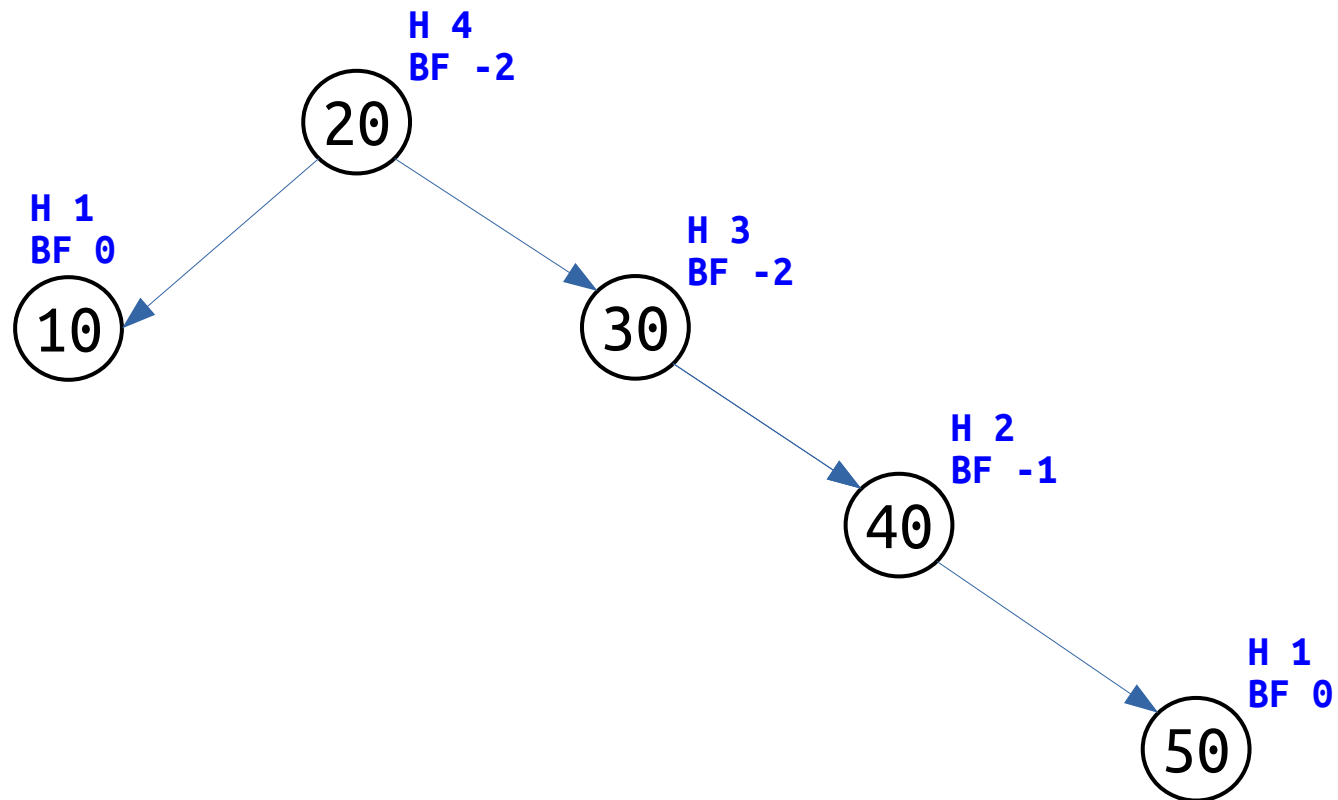
AVL

Insert **50**



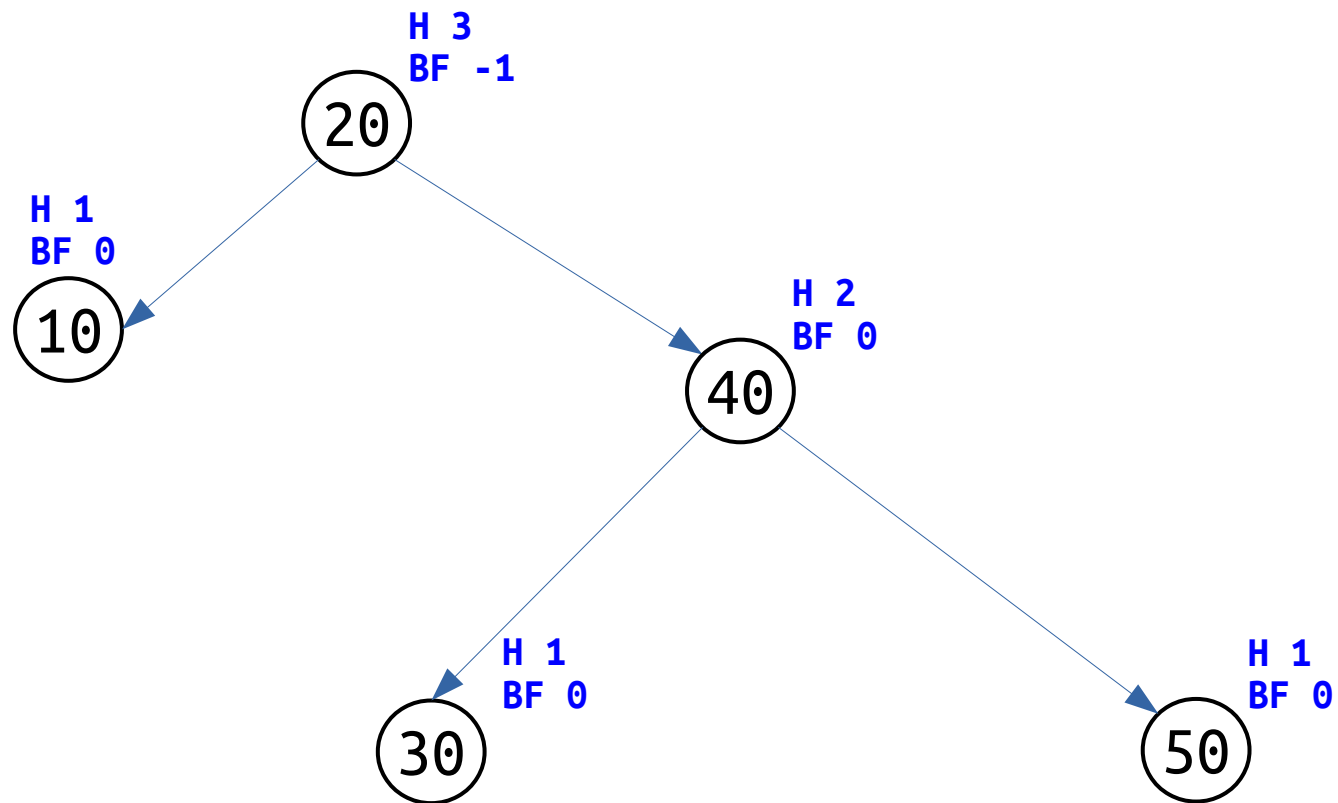
AVL

Insert **50**



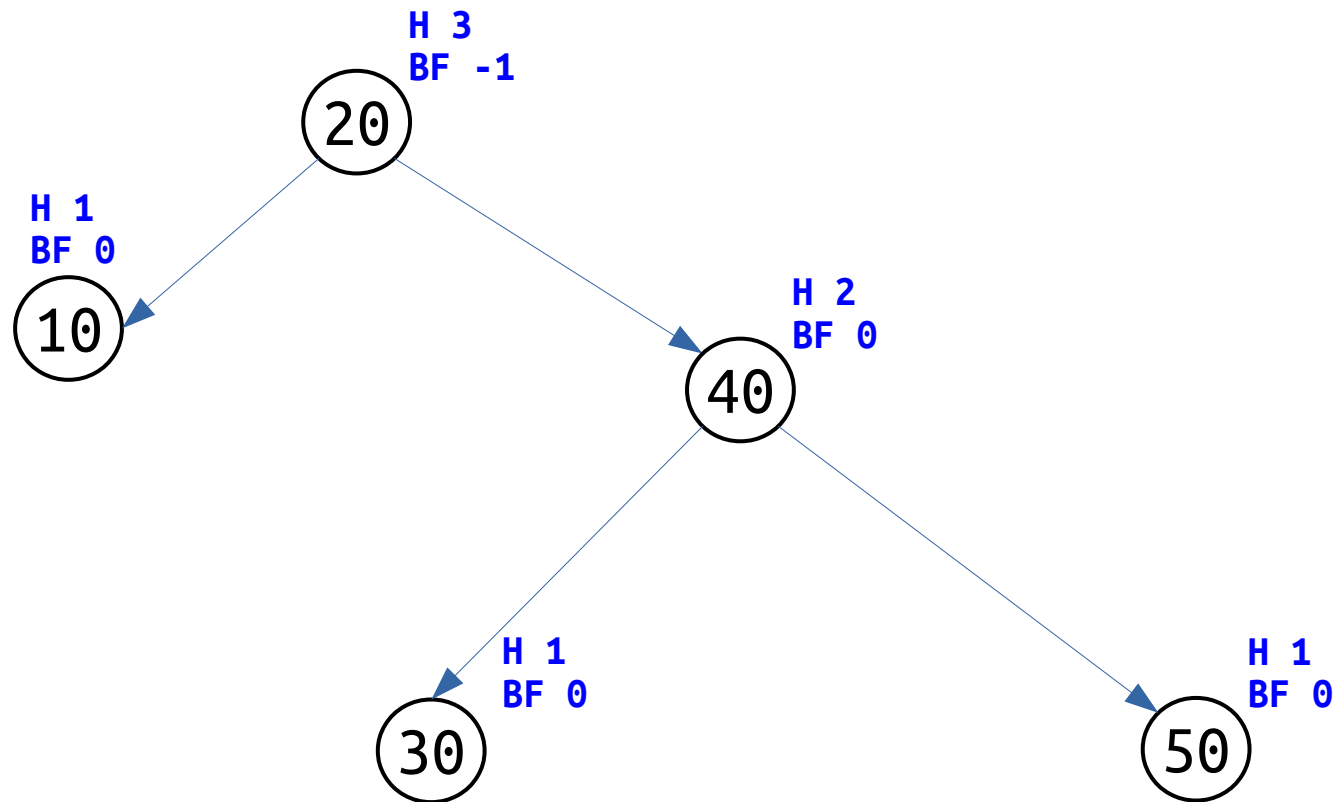
AVL

Insert **50**



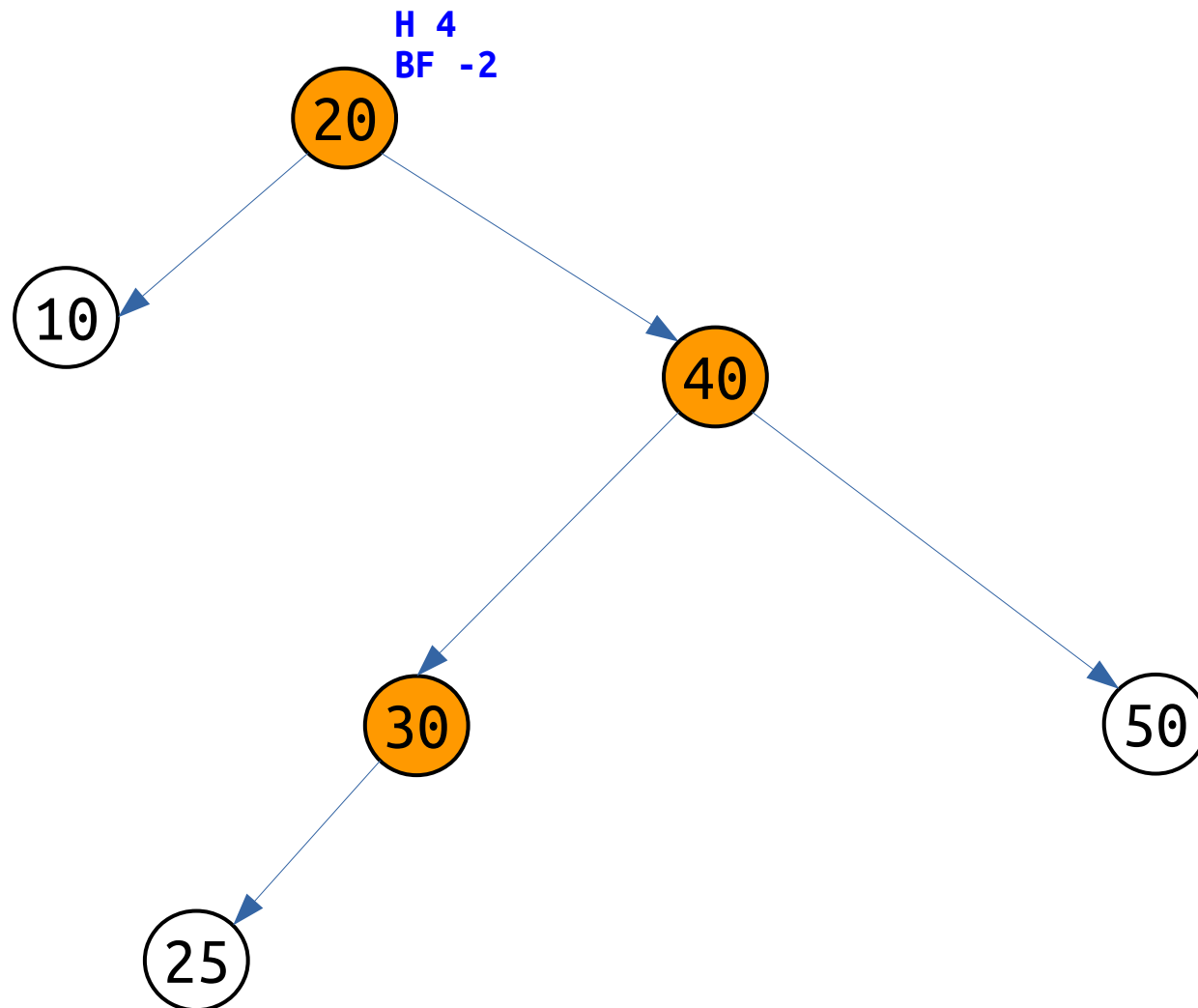
AVL

Insert 25



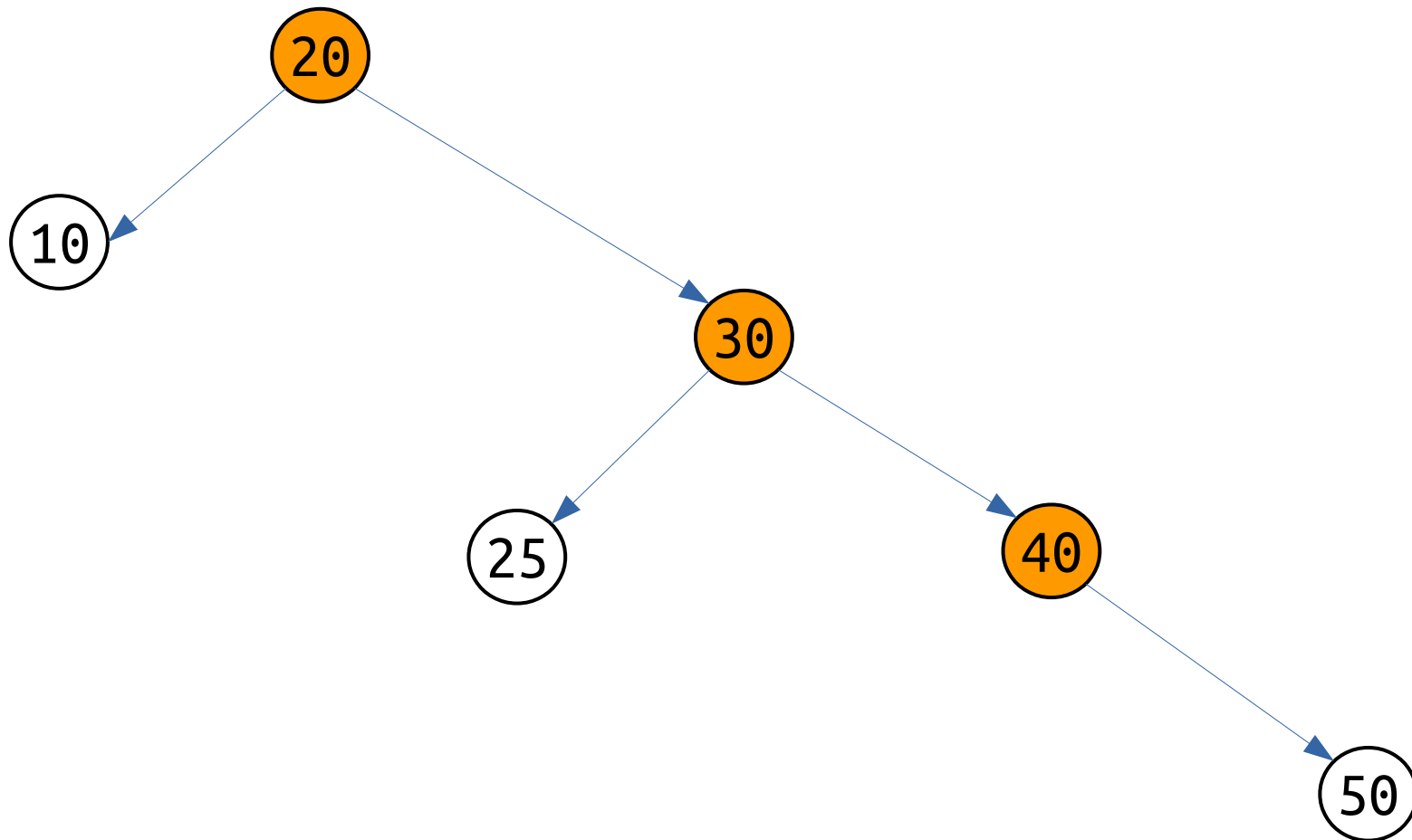
AVL

Insert 25



AVL

Insert **25**



AVL

Insert 25

