

Path-Coloring Algorithms for Plane Graphs

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June 1, 2017

2010 Mathematics Subject Classification. Primary 05C38; Secondary 05C10, 05C15.

Key words and phrases. Path coloring, list coloring, algorithm.

Abstract

A path coloring of a graph G is a vertex coloring of G such that each color class induces a disjoint union of paths. We present two efficient algorithms to construct a path coloring of a plane graph.

The first algorithm, based on a proof of Poh, is given a plane graph; it produces a path coloring of the given graph using three colors.

The second algorithm, based on similar proofs by Hartman and Škrekovski, performs a list-coloring generalization of the above. The algorithm is given a plane graph and an assignment of lists of three colors to each vertex; it produces a path coloring of the given graph in which each vertex receives a color from its list.

Implementations of both algorithms are available.

1 Introduction

All graphs will be finite, simple, and undirected. See West [9] for graph theoretic terms.

A *path coloring* of a graph G is a vertex coloring (not necessarily proper) of G such that each color class induces a disjoint union of paths. A graph G is *path k -colorable* if G admits a path coloring using k colors.

Broere & Mynhardt conjectured [1, Conj. 16] that every planar graph is path 3-colorable. This was proven independently by Poh [7, Thm. 2] and by Goddard [5, Thm. 1].

Theorem 1.1 (Poh 1990, Goddard 1991). *If G is a planar graph, then G is path 3-colorable.* \square

It is easily shown that the “3” in Theorem 1.1 is best possible. In particular, Chartrand & Kronk [4, Section 3] gave an example of a planar graph whose vertex set cannot be partitioned into two subsets, each inducing a forest.

Hartman [6, Thm. 4.1] proved a list-coloring generalization of Theorem 1.1 (see also Chappell & Hartman [3, Thm. 2.1]). A graph G is *path k -choosable* if, whenever each vertex of G is assigned a list of k colors, there exists a path coloring of G in which each vertex receives a color from its list.

Theorem 1.2 (Hartman 1997). *If G is a planar graph, then G is path 3-choosable.* \square

Essentially the same technique was used by Škrekovski [8, Thm. 2.2b] to prove a result slightly weaker than Theorem 1.2.

We discuss two efficient path-coloring algorithms based on proofs of the above theorems. We distinguish between a *planar* graph—one that can be drawn in the plane without crossing edges—and a *plane* graph—a graph with a given embedding in the plane.

In Section 2 we outline our graph representations and the basis for our computations of time complexity.

Section 3 covers an algorithm based on Poh’s proof of Theorem 1.1. The algorithm is given a plane graph; it produces a path coloring of the given graph using three colors.

Section 4 covers an algorithm based Hartman’s proof of Theorem 1.2, along with the proof of Škrekovski mentioned above. The algorithm is given a plane graph and an assignment of a list of three colors to each vertex; it produces a path coloring of the given graph in which each vertex receives a color from its list.

Implementations of both algorithms are available; see Bross [2].

2 Graph Representations and Time Complexity

We will represent a graph via *adjacency lists*: a list, for each vertex v , of the neighbors of v . A vertex can be represented by an integer $0 \dots n - 1$, where n is the order of the graph.

A plane graph will be specified via a *rotation scheme*: a circular ordering, for each vertex v , of the edges incident with v , in the order they appear around v in the plane embedding; this completely specifies the combinatorial embedding of the graph. Rotation schemes are convenient when we represent a graph using adjacency lists; we simply order the adjacency list for each vertex v in clockwise order around v ; no additional data structures are required.

We will assume an integer RAM model of computation. The input for each algorithm will be a connected planar graph with n vertices and m edges, represented via adjacency lists. The input size will be n , the number of vertices. Note that $\mathcal{O}(m) = \mathcal{O}(n)$.

In Section 4, given an edge uv , we will need an efficient operation to find v 's entry in u 's adjacency list from u 's entry in v 's list. An *augmented adjacency list* is an adjacency list such that for any edge uv , a reference to v 's entry in u 's list is stored in u 's entry in v 's list, and vice versa. Given an adjacency list representation of a graph, an augmented adjacency list representation may be constructed in $\mathcal{O}(m)$ time via the following algorithm.

Algorithm 2.1.

Input: An adjacency list representation Adj of a graph G .

Output: An augmented adjacency list representation Adj' of G with the same rotation scheme as Adj .

Step 1: Construct an augmented adjacency list copy Adj' of Adj with the reference portion of each entry uninitialized.

Step 2: For each vertex v construct an array $\text{Wrk}[v]$ storing vertex-reference pairs. For each v from 0 to $n - 1$ iterate through $\text{Adj}'[v]$. For each neighbor u in $\text{Adj}'[v]$ let $r(v, u)$ be a reference to u 's entry in $\text{Adj}'[v]$ and append the pair $(v, r(v, u))$ to $\text{Wrk}[u]$.

Step 3: For each v from $n - 1$ to 0 iterate through $\text{Wrk}[v]$. At the pair $(u, r(u, v))$ in $\text{Wrk}[v]$ note that the last element of $\text{Wrk}[u]$ is $(v, r(v, u))$. Look up and assign references for the edge uv in Adj' . Remove $(v, r(v, u))$ from the back of $\text{Wrk}[u]$.

After completing Step 2 in Algorithm 2.1 the array $\text{Wrk}[v]$ contains a pair $(u, r(u, v))$ for each neighbor of v , sorted in increasing order of u .

At a given iteration of Step 3 in Algorithm 2.1 let v be the current vertex. For each edge $uw \in E(G)$ such that $u < w$ and $v < w$, prior iterations of Step 3 will have initialized the references for uw in $\text{Adj}'[u]$ and $\text{Adj}'[w]$, and also removed the pair $(w, r(w, u))$ from $\text{Wrk}[u]$. The current iteration will handle all edges $uv \in E(G)$ with $v > u$.

3 Path Coloring: the Poh Algorithm

We will first describe Poh's algorithm for path 3-coloring plane graphs. If C is a cycle in a plane graph G we will use $\text{Int}(C)$ to denote the subgraph of G consisting of C and all interior vertices and edges.

Algorithm 3.1 (Poh 1990).

Input: A weakly triangulated plane graph G with outer face a cycle $C = v_1, v_2, \dots, v_k$. Additionally, a 2-coloring of C such that each color class induces a path, $P_1 = v_1, v_2, \dots, v_l$ and $P_2 = v_{l+1}, v_{l+2}, \dots, v_k$ respectively.

Output: An extension of the path 2-coloring of C to a path 3-coloring of G such that no vertex in $G - C$ receives the same color as a neighbor in C .

Step 1: If $G = C$ then G is already path 3-colored and we are done. Otherwise there are two cases to consider.

Case 1.1: Suppose C is an induced cycle in G . Let $u, w \in V(G) - V(C)$ such that uv_1v_k and wv_lv_{l+1} are faces; note that u and w are unique, but may not be distinct. Since C is induced and $G \neq C$, $G - C$ is connected. Let $P_3 = u_1u_2 \dots u_r$ be a shortest u, w -path

in $G - C$. Color each vertex of P_3 with the third color not used in the 2-coloring of C . Let $C_1 = v_1v_2 \dots v_lu_ru_{r-1} \dots u_1$ and $C_2 = u_1u_2 \dots u_rv_{l+1}v_{l+2} \dots v_k$.

Case 1.2: Suppose C is not an induced cycle. Then there exists an edge $v_iv_j \in E(G) - E(C)$ such that $i \leq l < j$. Let $C_1 = v_1v_2 \dots v_iv_jv_{j+1} \dots v_k$ and $C_2 = v_iv_{i+1} \dots v_j$.

Step 2: Apply Algorithm 3.1 to $\text{Int}(C_1)$ and $\text{Int}(C_2)$.

Note that in Algorithm 3.1 the graph G is finite and the recursive step applies the algorithm to two proper subgraphs of G . Therefore Algorithm 3.1 must terminate.

4 Path List Coloring: the Hartman-Škrekovski Algorithm

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