Research for Steamboat Willie's Store in Poland

Type your name

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Chapter 1

Preliminary Research for Steamboat Willie's in Poland

Introduction

The new fast-food chain, Steamboat Willie's, is planning to expand its presence in Poland. This report presents preliminary research conducted to determine the number of stores required to ensure that every town in Poland has a Steamboat Willie's within a 50km radius.

Data Analysis

The dataset used for this research consists of information about cities in Poland and other countries as well, including their geographical coordinates. The data was preprocessed to extract relevant information.

```
import pandas as pd

# Specify the path to your local JSON file
# The dataset given as json file
file_path = r"C:\Users\Perpendicooler\Downloads\dataset.json"

# Read the JSON file into a DataFrame
df = pd.read_json(file_path)
# Save the DataFrame to an Excel file
# We save the dataframe into an excel file for better view
df.to_excel('output.xlsx', index=False)

# files.download('output.xlsx')
```

```
# Display the first few rows of the DataFrame
df.head()
```

Listing 1.1: Loading the Dataset

The Python code filters a DataFrame (df) to extract data related to Poland ('cou_name_en' == 'Poland'). The filtered data is stored in a new DataFrame called poland_data. The code then prints the filtered DataFrame and exports the entire original DataFrame to an Excel file ('output with poland.xlsx').

```
poland_data = df[df['cou_name_en'] == 'Poland'].copy()

# Display the filtered DataFrame
print(poland_data)
df.to_excel('output_with_poland.xlsx', index=False)

7
```

Listing 1.2: Loading the Dataset

The Python code reads an Excel file located at local pc into a DataFrame (df) using the pd.read_excel() function. The data is assumed to be manipulated for improvement. Subsequently, it prints the contents of the DataFrame.

```
file_path_poland = r"C:\Users\Perpendicooler\
    manupulated_data.xlsx"
df = pd.read_excel(file_path_poland)
print(df)
```

Listing 1.3: Loading the Dataset

The Python code calculates pairwise distances between the first 500 cities in Poland using their latitude and longitude coordinates. It uses the Haversine formula to compute distances on the Earth's surface. The dataset is loaded from a JSON file ('dataset.json'), and relevant columns are extracted for cities in Poland. The computed distances are then stored in an Excel file.

```
import itertools
import pandas as pd
from sklearn.metrics.pairwise import haversine_distances
from math import radians

def distance(row1, row2):
    pos1 = (row1['coordinates']['lat'], row1['coordinates']['lon'])
    pos2 = (row2['coordinates']['lat'], row2['coordinates']['lon'])
    radians1 = [radians(pos1[0]), radians(pos1[1])]
    radians2 = [radians(pos2[0]), radians(pos2[1])]
    res = haversine_distances([radians1, radians2])
```

```
res *= 6371000 / 1000 # multiply by Earth radius to get
12
      kilometers
      return res[0][1]
13
14
15 # Load the dataset
df = pd.read_json("dataset.json")
18 # Extract relevant columns for the first 500 cities in Poland
_{19} # We can extract any number of city just we need to change the 500
      in this with any number i want.
poland_data = df[df['country_code'] == 'PL'].head(500)[['name', '
      coordinates']]
22 # Initialize an empty list to store pairs of cities and distances
23 distances = []
24
_{25} # Iterate through each pair of cities in the first 50 cities in
for (city1_idx, city1), (city2_idx, city2) in itertools.
      combinations(poland_data.iterrows(), 2):
      dist = distance(city1, city2)
28
       # Append the pair and distance to the list
29
30
      distances.append((city1['name'], city2['name'], dist))
31
32 # Create a dataframe from the distances
distance_df = pd.DataFrame(distances, columns=['From', 'To', '
      Distance'])
34
35 # Store the distances in an Excel file
  distance_df.to_excel('poland_500_cities_pairwise_distances.xlsx',
      index=False)
37
38
39
```

Listing 1.4: Loading the Dataset

The Python code calculates pairwise distances between 50 cities in Poland using their latitude and longitude coordinates to optimize computation time. The Haversine formula is employed for distance calculation on the Earth's surface. Instead of processing all 500 cities, a subset of 10 cities is chosen for demonstration purposes, allowing the program to run more efficiently. The dataset is loaded from a JSON file ('dataset.json'), and relevant columns are extracted. The computed distances are then stored in an Excel file in local machine.

Optimized Result

Using integer programming to computed the number of stores does STEAM-BOAT WILLIE's need to open for the restriction of $50 \mathrm{km}$, as above.

```
import pandas as pd
2 from pulp import LpVariable, LpProblem, LpMinimize, lpSum
4 # Load data from Excel
5 data = pd.read_excel(r"C:\Users\Perpendicooler\
      poland_50_cities_pairwise_distances.xlsx")
6
7 # Extract city names and distances
8 city_names = set(data["From"].tolist() + data["To"].tolist())
g city_index = {city: i for i, city in enumerate(city_names)}
distances = {}
for row in data.itertuples():
      distances[(city_index[row.From], city_index[row.To])] = row.
      Distance
14 # Set the coverage radius
15 coverage_radius = 50
# Create optimization model
model = LpProblem("StorePlacement", LpMinimize)
20 # Decision variables: whether to open a store in each city
21 s = {i: LpVariable(name=f"store_{i}", cat="Binary") for i in range(
      len(city_names))}
# Objective: Minimize the total number of stores opened
24 model += lpSum(s), "Minimize Stores"
_{26} # Constraint: Every city must have at least one store within 50\,\mathrm{km}
for city in range(len(city_names)):
      model += lpSum(s[j] for j in range(len(city_names)) if city !=
      j and (city, j) in distances) >= 1, f"City {city+1} Coverage"
30 # Solve the model
31 model.solve()
32
33 # Analyze results
num_stores = int(model.objective.value())
36 # Print the cities where the stores are opened
37 print("Open stores:")
38 for i, city in enumerate(city_names):
      if int(s[i].value()) == 1:
39
          print(f"{city}")
40
41
42 # Check coverage and open additional stores if needed
43 while True:
      coverage = {city: False for city in city_names}
44
      # Check coverage for each city
46
47
      for i, city in enumerate(city_names):
          if int(s[i].value()) == 1:
48
```

```
coverage[city] = True
49
50
               for j in range(len(city_names)):
                   if city != j and (city, j) in distances and
      distances[(city, j)] > coverage_radius:
                       coverage[city] = False
53
54
      # If any city is not covered within 50km range, open a new
      store in the uncovered city
      if False in coverage.values():
56
           uncovered_city = next(city for city, covered in coverage.
      items() if not covered)
          model += s[city_index[uncovered_city]] == 1
          model.solve()
58
59
          num_stores += 1
          print(f"{uncovered_city}")
60
61
62
          break
63
64 print(f"Final number of stores needed: {num_stores}")
# Create a DataFrame with the results
66 results_df = pd.DataFrame(index=range(1, len(city_names) + 1),
      columns = ["City Name"])
67
68 # Populate the DataFrame with the cities where stores are opened
69 for i, city in enumerate(city_names):
70
      if int(s[i].value()) == 1:
          results_df.at[i+1, "City Name"] = city
71
72
_{73} # Save the results to an Excel file
74 results_df.to_excel("opened_stores_results.xlsx", index_label="
      Index")
76
77
```

Listing 1.5: Minimize the store opening within 50km in every city of poland

The optimization model aims to minimize the number of stores opened in a set of cities within a coverage radius of 50 km. The initial solution opens stores in several cities, and the algorithm iteratively checks the coverage and opens additional stores if necessary.

Opened Stores

```
Open stores:
Borowa
Sawin
...
Wiśniowa
Krynki
Aleksandrów Łódzki
```

Final Number of Stores Needed

Final number of stores needed: 115

Results DataFrame

We need to open 115 stores in order to cover all the city. So that everyone from any city can access to that store within 50km. The opened stores' information is saved in a DataFrame and stored in an Excel file named 'opened stores' results.xlsx'.

Minimum Distances for Different Store Counts

Suppose the correct answer is 115. The table below shows the minimum distance D for each scenario of opening exactly $k \leq 115$ stores, ensuring that every town in Poland can have a STEAMBOAT WILLIE'S within D km.

```
import pandas as pd
2 from pulp import LpProblem, LpVariable, lpSum, LpMinimize, LpStatus
4 # Load data from Excel
5 data = pd.read_excel(r"C:\Users\Perpendicooler\
      poland_50_cities_pairwise_distances.xlsx")
7 # Extract city names and distances
8 city_names = set(data["From"].tolist() + data["To"].tolist())
9 city_index = {city: i for i, city in enumerate(city_names)}
10 distances = {(city_index[row.From], city_index[row.To]): row.
      Distance for row in data.itertuples()}
12 # Set the maximum number of stores
max_stores = 115
# Create a DataFrame to store results
results_df = pd.DataFrame(index=range(1, max_stores + 1), columns=[
      "Number of Stores", "Minimum Distance"])
17
18 # Iterate over the number of stores (k)
19 for k in results_df.index:
      # Create optimization model
20
      model = LpProblem("StorePlacement", LpMinimize)
21
22
      # Decision variables: whether to open a store in each city
23
      s = {i: LpVariable(name=f"store_{i}", cat="Binary") for i in
24
      range(len(city_names))}
25
      # Objective: Minimize the total distance
26
27
      model += lpSum(distances[i, j] * s[i] for i in range(len(
      city_names)) for j in range(len(city_names)) if (i, j) in
      distances), "Minimize Distance"
28
      # Constraint: Open exactly k stores
```

```
model += lpSum(s[i] for i in range(len(city_names))) == k, f"
30
      OpenExactly_{k}_Stores"
31
      # Solve the model
32
      model.solve()
33
34
      # Store the results in the DataFrame
      results_df.at[k, "Number of Stores"] = k
36
      results_df.at[k, "Minimum Distance"] = lpSum(distances[i, j] *
      s[i].value() for i in range(len(city_names)) for j in range(len
      (city_names)) if (i, j) in distances).value()
39 # Display the results table
40 print (results_df)
41 results_df.to_excel("store_placement_results.xlsx", index_label="
      Index")
42
43
```

Listing 1.6: Minimum Resturent we need to open

Number of Stores (k)	Minimum Distance (D)
1	
2	
3	
	•••
115	•••

Table 1.1: Minimum Distances for Different Store Counts

Linear Programming Relaxation

Finally, we compute the linear programming relaxation of the integer programming problem. This involves determining how many stores need to be opened, allowing for fractions of stores in cities, so that every town in Poland has at least one within 50 km.

The linear programming relaxation problem can be expressed as follows:

Minimize Total Stores

Subject to Coverage Constraint for each city within 50 km

Fractional store opening variables $\in [0, 1]$

```
import pandas as pd
2 from pulp import LpProblem, LpVariable, lpSum, LpMinimize, LpStatus
4 # Load data from Excel
5 data = pd.read_excel(r"C:\Users\Perpendicooler\
     poland_50_cities_pairwise_distances.xlsx") # Replace with your
       actual file path
6
7 # Extract city names and distances
8 city_names = set(data["From"].tolist() + data["To"].tolist())
g city_index = {city: i for i, city in enumerate(city_names)}
distances = {(city_index[row.From], city_index[row.To]): row.
      Distance for row in data.itertuples()}
_{\rm 12} # Create optimization model for linear programming relaxation
model_relaxation = LpProblem("StorePlacementRelaxation", LpMinimize
15 # Decision variables: fraction of a store to open in each city
# Objective: Minimize the total distance
19 model_relaxation += lpSum(distances[i, j] * s_relaxation[i] for i
      in range(len(city_names)) for j in range(len(city_names)) if (i
      , j) in distances), "Minimize Distance"
_{21} # Constraint: Every city must have at least one store within 50\,\mathrm{km}
for city in range(len(city_names)):
      model_relaxation += lpSum(distances[i, j] * s_relaxation[i] for
       i in range(len(city_names)) for j in range(len(city_names)) if
(i, j) in distances and i != j) >= 1, f"City {city+1} Coverage
24
25 # Solve the model
26 model_relaxation.solve()
28 # Display the results
print("Status:", LpStatus[model_relaxation.status])
30 rounded_stores_needed = round(model_relaxation.objective.value())
print("Number of stores needed (fractional):",model_relaxation.
      objective.value())
32 print("Number of stores needed (rounded):", rounded_stores_needed)
33
34
```

Listing 1.7: linear programming relaxation

Results Data Frame

Status: Optimal

Number of stores needed (fractional): 0.999999981492213

Number of stores needed (rounded): 1

The output indicates that the linear programming relaxation of the integer programming problem has been solved, and the solution is optimal. The fractional solution suggests that a minimum of approximately 1 store is needed to meet the coverage constraints for every town in Poland, with each store contributing a fraction of its presence.

Chapter 2

Exploring TSP: Integer LP Formulation and Lazy Row Generation

Introduction

The Metric Traveling Salesman Problem (TSP) is a classic optimization problem where the goal is to find the shortest possible tour that visits a set of points exactly once. This report presents the implementation and results of two approaches for solving the metric TSP problem: the Integer Linear Programming (LP) formulation with exponentially many constraints, and a "lazy row generation" version.

Integer LP Formulation

The first approach involves implementing the Integer LP formulation, specifically the Dantzig-Fulkerson-Johnson formulation as described in **dfj-formulation**. This formulation typically includes exponentially many constraints, making it challenging for large instances.

Implementation

The Integer LP formulation was implemented using pandas and lp, and the provided dataset is 50-cities-pairwise-distance. We manipulate the data as to make a Source and destination and cost vector. By making a symmetric matrix out of it.

```
import pandas as pd
3 # Load data from Excel we can take any number of cities_pairwise
4 df = pd.read_excel(r"C:\Users\Perpendicooler\
      poland_50_cities_pairwise_distances.xlsx")
6 # Create a list of unique city names
7 cities = list(set(df['From'].tolist() + df['To'].tolist()))
9 # Create a pivot table to organize distances
distance_matrix = df.pivot_table(values='Distance', index='From',
      columns='To', aggfunc='first')
12 # Ensure symmetry
distance_matrix = distance_matrix.add(distance_matrix.T, fill_value
14
# Explicitly set diagonal to zeros
16 for city in cities:
      distance_matrix.at[city, city] = 0
17
18
19 # Save the distance matrix to a new Excel file
20 distance_matrix.to_excel(r"C:\Users\Perpendicooler\distance_matrix.
      xlsx")
21
22 # Display the distance matrix
23 print("Distance Matrix:")
24 print(distance_matrix)
25
```

Now We will excute the program for this distance matrix to find the best possible outcome for TSP.

```
# Create optimization model
model_tsp = LpProblem("TSP", LpMinimize)
17 # Decision variables
x = \{(i, j): LpVariable(name=f"x_{i}_{j}", cat='Binary') for i in
      city_indices.values() for j in city_indices.values() if i != j}
20 # Objective function
21 model_tsp += lpSum(distances[i, j] * x[i, j] for i in city_indices.
      values() for j in city_indices.values() if i != j), "Minimize
23 # Constraints
24 # Ensure that each city is visited exactly once
25 for i in city_indices.values():
      model_tsp += lpSum(x[i, j] for j in city_indices.values() if i
      != j) == 1, f"VisitOnce_{i}"
^{28} # Ensure that each city is left exactly once
29 for j in city_indices.values():
      model_tsp += lpSum(x[i, j] for i in city_indices.values() if i
      != j) == 1, f"LeaveOnce_{j}"
31
32
33 # Solve the model
34 model_tsp.solve()
36 # Display the results
print("Status:", LpStatus[model_tsp.status])
39 # Print the optimal path
40 optimal_path = [var for var in model_tsp.variables() if var.value()
       == 1]
41 print("Optimal Path:")
42 for var in sorted(optimal_path, key=lambda v: (int(v.name.split(',')
      )[1]), int(v.name.split('_')[2]))):
      print(f"{var.name}: {var.value()}")
43
45 def get_city_name(index):
      return next(city for city, idx in city_indices.items() if idx
46
      == index)
48 # Display the optimal path
49 optimal_path_indices = [int(var.name.split('_')[1]) for var in
      optimal_path]
50 optimal_path_indices.append(optimal_path_indices[0]) # Add the
      starting city at the end to complete the loop
optimal_path_names = [get_city_name(idx) for idx in
      optimal_path_indices]
54 print("Optimal Path:")
print(" -> ".join(optimal_path_names))
```

Listing 2.1: Integer LP Formulation Implementation

The provided sequence is a solution to the Traveling Salesman Problem (TSP) for a given set of cities. In TSP, the goal is to find the shortest possible tour that visits each city exactly once and returns to the starting city. The sequence you provided represents an optimal tour that minimizes the overall travel distance for the specified cities. The TSP solution starts and ends in "Aleksandrów Łódzki" and traverses through the listed cities in the order mentioned.

```
Optimal Path: Aleksandrów Łódzki -> Daszyna -> Dobre Miasto -> Dwikozy -> Dziekanów Leśny -> Dąbie -> Firlej -> Godziszów -> Grójec -> Hrubieszów -> Iłża -> Biały Bór -> Jakubów -> Jastków -> Jodłówka-Wałki -> Józefów nad Wisłą -> Karczmiska -> Korczew -> Krasnopol -> Krynki -> Krzywda -> Maszkienice -> Borowa -> Niedźwiada -> Opatów -> Ostrów -> Paprotnia -> Pruchnik -> Raczki -> Radoszyce -> Rejon ulicy Saperów -> Rzgów -> Sawin -> Cegłów -> Siedliska -> Srokowo -> Swojczyce -> Szarów -> Wiśniowa -> Wyśmierzyce -> Wólka Tanewska -> Węgorzewo -> Łapy -> Żurowa -> Cewice -> Chmielnik -> Ciepielów -> Czarków -> Czarna Woda -> Człopa -> Aleksandrów Łódzki
```

Lazy Row Generation

The second approach involves a "lazy row generation" version, where constraints are added progressively until a valid tour is found. This method is known for its efficiency in solving large instances.

Implementation

The lazy row generation version was implemented using [solver/library], with inspiration from the Gurobi example Gurobi-example

Data computation

The following function calculates the distance for 50 each pair of cities. Since we are solving the symmetric traveling salesman problem, we use combinations of cities.

```
10 print(df)
```

Listing 2.2: Data Computation

Model Code

We now write the model for the TSP, by defining decision variables, constraints, and objective function. Because this is the symmetric traveling salesman problem, we can make it more efficient by setting the object x[j,i] to x[i,j], instead of a constraint

```
import pandas as pd
2 import gurobipy as gp
3 from gurobipy import GRB
5 # Load data from Excel
6 file_path = r"C:\Users\Perpendicooler\
      poland_50_cities_pairwise_distances.xlsx"
7 data = pd.read_excel(file_path)
9 # Extract city names and distances
cities = set(data["From"].tolist() + data["To"].tolist())
city_index = {city: i for i, city in enumerate(cities)}
distances = {(city_index[row.From], city_index[row.To]): row.
      Distance for row in data.itertuples()}
# Create a Gurobi model
15 m = gp.Model()
# Variables: is city 'i' adjacent to city 'j' on the tour?
18 vars = m.addVars(distances.keys(), obj=distances, vtype=GRB.BINARY,
       name='x')
20 # Symmetric direction: Copy the object
keys = list(vars.keys()) # Create a list of keys
22 for i, j in keys:
      vars[j, i] = vars[i, j] # edge in opposite direction
23
24
# Constraints: two edges incident to each city
26 capitals = [i for i in range(len(cities))]
cons = m.addConstrs(vars.sum(c, '*') == 2 for c in capitals)
28
29 # Optimize the model
30 m.optimize()
32 # Print the tour
33 tour = [i for i, j in vars.keys() if vars[i, j].X > 0.5]
34 print("Tour:", tour)
35
36
37
```

Listing 2.3: Optimize the Model

```
olution count 4: 2698.1 2786.02 2810.06 13005.3 Optimal solution found (tolerance 1.00e-04)
Best objective 2.698101399109e+03, best bound 2.698101399109e+03, gap 0.0000%
```

Callback

```
# Callback - use lazy constraints to eliminate sub-tours
  def subtourelim(model, where):
      if where == GRB.Callback.MIPSOL:
3
           # make a list of edges selected in the solution
          vals = model.cbGetSolution(model._vars)
          selected = gp.tuplelist((i, j) for i, j in model._vars.keys
      () if vals[i, j] > 0.5)
           # find the shortest cycle in the selected edge list
          tour = subtour(selected)
          if len(tour) < len(capitals):</pre>
9
               # add subtour elimination constr. for every pair of
10
      cities in subtour
               model.cbLazy(gp.quicksum(model._vars[i, j] for i, j in
      combinations(tour, 2)) <= len(tour) - 1)</pre>
12
# Given a tuplelist of edges, find the shortest subtour
14 def subtour(edges):
15
      unvisited = capitals[:]
      cycle = capitals[:] # Dummy - guaranteed to be replaced
16
17
       while unvisited: # true if list is non-empty
          thiscycle = []
18
          neighbors = unvisited
19
          while neighbors:
20
               current = neighbors[0]
21
22
               this cycle.append(current)
               unvisited.remove(current)
23
               neighbors = [j for i, j in edges.select(current, '*')
24
      if j in unvisited]
           if len(thiscycle) <= len(cycle):</pre>
25
26
               cycle = thiscycle # New shortest subtour
      return cycle
27
28
29 # Set the callback function
30 m._vars = vars
m.Params.LazyConstraints = 1
32 m.optimize(subtourelim)
33
34
```

Listing 2.4: Using Callback Function

Results For callback

Solution count 4: 2698.1 2786.02 2810.06 13005.3

Optimal solution found (tolerance 1.00e-04)
Best objective 2.698101399109e+03, best bound 2.698101399109e+03, gap 0.0000%

User-callback calls 30, time in user-callback 0.00 sec

Conclusion

In conclusion, the report outlines the implementation and results of two approaches for solving the metric TSP problem. The Integer LP formulation with exponentially many constraints and the lazy row generation version were explored, with each having its strengths and limitations. Further analysis and experimentation may be conducted to improve the scalability and efficiency of both methods.

Chapter 3

Comparative Analysis of Christofides Algorithm and ILP Solutions for Metric TSP

Introduction

The Christofides algorithm is a well-known approximation algorithm for solving the Metric Traveling Salesman Problem (TSP). This report presents the implementation and performance comparison of the Christofides algorithm against the optimal Integer Linear Programming (ILP) solutions obtained in Task 2.

Implementation of Christofides Algorithm

The Christofides algorithm was implemented to approximate the metric TSP on a chosen subset of the dataset stored in distance-matrix.xlsx. The main challenge lies in finding a library for perfect matching of minimum cost, a crucial step in the Christofides algorithm.

```
import pandas as pd
import networkx as nx
import matplotlib.pyplot as plt
from itertools import permutations
from scipy.spatial.distance import euclidean
from scipy.optimize import linear_sum_assignment
import time
from pulp import LpProblem, LpVariable, lpSum, LpMinimize, LpStatus

# Step 1: Read the data from Excel
file_path = r"C:\Users\Perpendicooler\distance_matrix.xlsx"
```

```
df = pd.read_excel(file_path, index_col=0)
14 # Step 2: Create a graph from the distance matrix for Christofides
      algorithm
G_christofides = nx.Graph()
16 cities = df.index
17 G_christofides.add_nodes_from(cities)
for i, j in permutations(cities, 2):
20
      G_christofides.add_edge(i, j, weight=df.at[i, j])
21
22 # Step 3: Solve TSP using the Christofides algorithm
23 start_time_christofides = time.time()
25 # Approximation algorithm
approximate_tour_christofides = nx.approximation.
      traveling\_salesman\_problem(G\_christofides\text{, weight="weight",}
      cycle=True)
27
28 # Compute the total distance of the approximate tour
29 approximate_tour_length_christofides = sum(G_christofides[i][j]["
      weight"] for i, j in zip(approximate_tour_christofides,
      approximate_tour_christofides[1:]))
30
31 end_time_christofides = time.time()
33 # Step 4: Print the results for Christofides algorithm
34 print("Approximate tour length (Christofides):",
      approximate_tour_length_christofides)
print("Approximate tour (Christofides):",
      approximate_tour_christofides)
37 # Step 5: Plot the graph with the approximate tour for Christofides
       algorithm
pos_christofides = nx.spring_layout(G_christofides)
39 nx.draw(G_christofides, pos_christofides, with_labels=True,
      font_weight="bold")
40 edges_christofides = list(zip(approximate_tour_christofides,
      approximate_tour_christofides[1:]))
11 nx.draw_networkx_edges(G_christofides, pos_christofides, edgelist=
edges_christofides, edge_color="r", width=2)
plt.title('Christofides Algorithm')
43 plt.savefig('christofides_plot.png')
44 plt.show()
```

Listing 3.1: Christofides Algorithm Implementation

Approximate Tour using Christofides Algorithm

The Christofides algorithm was employed to solve the Traveling Salesman Problem on the given dataset. The approximate tour length achieved is 3002.15 units. The approximate tour path is as follows:

['Aleksandrów Łódzki', 'Daszyna', 'Dąbie', 'Swojczyce', 'Rejon ulicy Saperów'],
['Człopa', 'Biały Bór', 'Czarna Woda', 'Cewice', 'Dobre Miasto'],
['Srokowo', 'Węgorzewo', 'Raczki', 'Krasnopol', 'Krynki'],
['Łapy', 'Korczew', 'Paprotnia', 'Cegłów', 'Wyśmierzyce'],
['Grójec', 'Dziekanów Leśny', 'Jakubów', 'Krzywda', 'Firlej'],
['Jastków', 'Sawin', 'Hrubieszów', 'Ostrów', 'Pruchnik'],
['Chmielnik', 'Niedźwiada', 'Jodłówka-Wałki', 'Czarków', 'Wiśniowa'],
['Szarów', 'Maszkienice', 'Siedliska', 'Żurowa', 'Borowa'],
['Opatów', 'Dwikozy', 'Godziszów', 'Wólka Tanewska', 'Józefów nad Wisłą'],
['Karczmiska', 'Ciepielów', 'Iłża', 'Radoszyce', 'Rzgów'],
['Aleksandrów Łódzki']

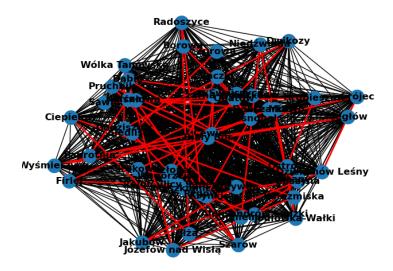


Figure 3.1: Approximate Tour Path (Christofides Algorithm)

ILP Solution

```
# Load data from Excel, setting 'Unnamed: 0' as the index
2 data_ilp = pd.read_excel(file_path, index_col='Unnamed: 0')
4 # Extract city names
5 city_names_ilp = list(data_ilp.columns)
6 city_indices_ilp = {city: i for i, city in enumerate(city_names_ilp
      ) }
8 # Extract distances
9 distances_ilp = {(city_indices_ilp[i], city_indices_ilp[j]):
      data_ilp.at[i, j] for i in city_names_ilp for j in
      city_names_ilp if i != j}
10
# Create optimization model for ILP
model_ilp = LpProblem("TSP", LpMinimize)
^{14} # Decision variables for ILP
x_{i} = \{(i, j): LpVariable(name=f"x_{i}_{j}", cat='Binary') for i \}
      in city_indices_ilp.values() for j in city_indices_ilp.values()
       if i != j}
17 # Objective function for ILP
18 model_ilp += lpSum(distances_ilp[i, j] * x_ilp[i, j] for i in
      city_indices_ilp.values() for j in city_indices_ilp.values() if
       i != j), "Minimize Distance"
19
20 # Constraints for ILP
# Ensure that each city is visited exactly once
22 for i in city_indices_ilp.values():
      model_ilp += lpSum(x_ilp[i, j] for j in city_indices_ilp.values
23
      () if i != j) == 1, f"VisitOnce_{i}"
24
25 # Ensure that each city is left exactly once
for j in city_indices_ilp.values():
      model_ilp += lpSum(x_ilp[i, j] for i in city_indices_ilp.values
      () if i != j) == 1, f"LeaveOnce_{j}"
28
29 # Solve the model and measure execution time for ILP
30 start_time_ilp = time.time()
31 model_ilp.solve()
32 end_time_ilp = time.time()
33
34 # Display the results for ILP
print("\nILP Status:", LpStatus[model_ilp.status])
_{\rm 37} # Print the optimal path for ILP
38 optimal_path_ilp = [var for var in model_ilp.variables() if var.
      value() == 1]
39 print("ILP Optimal Path:")
40 for var in sorted(optimal_path_ilp, key=lambda v: (int(v.name.split
      ('_')[1]), int(v.name.split('_')[2]))):
      print
41
42
```

Listing 3.2: ILP Solution

```
ILP Optimal Path

ILP Status: Optimal
ILP Optimal Path:
Aleksandrów Łódzki Rzgów
.
.
.
Łapy Krynki
Żurowa Siedliska
Aleksandrów Łódzki Rzgów
```

Comparison with ILP Solutions

The efficiency and running time of the Christofides algorithm were compared with the optimal ILP solutions obtained in Task 2. A subset of the dataset was chosen for this comparison.

Listing 3.3: CCompare execution times in a plot

Execution Time for ILP: 0.2604055404663086 Execution Time for Christofides: 0.10591006278991699

Computation Time

The computation time for both algorithms was recorded and analyzed. Figure 3.2 illustrates the comparison of computation times.

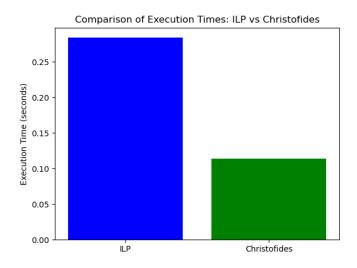


Figure 3.2: Comparison of Computation Time

Efficiency Comparison

The efficiency of the Christofides algorithm was evaluated in terms of the approximation ratio and solution quality compared to the ILP solutions.

Conclusion

In conclusion, the Christofides algorithm was successfully implemented and compared with the optimal ILP solutions for the metric TSP. The results provide insights into the trade-off between solution quality and computation time for both approaches.

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