

Homework 4

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Problem 9.10

The riddle can be expressed in first-order logic as:

$$\begin{aligned} & \text{Male}(m) \wedge \\ & (\neg \exists s : \text{Brother}(s, I) \vee \text{Sister}(s, I)) \wedge (\exists f : \text{Son}(f, \text{Father}(I)) \wedge f = \\ & \text{Father}(m)) \end{aligned}$$

Where I represents the speaker in the riddle, and m represents “that man”. From this, we can use the axioms of the kinship domain to deduce the identity of “that man”:

$$\begin{aligned} & \text{Male}(m) \wedge \\ & (\neg \exists s : \text{Brother}(s, I) \vee \text{Sister}(s, I)) \wedge \\ & (\exists f : \text{Son}(f, \text{Father}(I)) \wedge f = \text{Father}(m)) \end{aligned} \tag{1}$$

$$\begin{aligned} & \text{Male}(m) \wedge \\ & (\neg \exists s : (\text{Sibling}(s, I) \wedge \text{Male}(s)) \vee (\text{Sibling}(s, I) \wedge \text{Female}(s))) \wedge \\ & (\exists f : \text{Son}(f, \text{Father}(I)) \wedge f = \text{Father}(m)) \end{aligned} \tag{2}$$

$$\begin{aligned} & \text{Male}(m) \wedge \\ & (\neg \exists s : \text{Sibling}(s, I) \wedge (\text{Male}(s) \vee \text{Female}(s))) \wedge \\ & (\exists f : \text{Son}(f, \text{Father}(I)) \wedge f = \text{Father}(m)) \end{aligned} \tag{3}$$

$$\begin{aligned} & \text{Male}(m) \wedge \\ & (\neg \exists s : \text{Sibling}(s, I)) \wedge \\ & (\exists f : \text{Son}(f, \text{Father}(I)) \wedge f = \text{Father}(m)) \end{aligned} \tag{4}$$

$$\begin{aligned} & \text{Male}(m) \wedge \\ & (\neg \exists s : \text{Sibling}(s, I)) \wedge \\ & (\exists f : \text{Son}(f, \text{Father}(I)) \wedge \text{Parent}(f, m) \wedge \text{Male}(f)) \end{aligned} \tag{5}$$

$$\begin{aligned}
& \text{Male}(m) \wedge \\
& (\neg \exists s : \text{Sibling}(s, I)) \wedge \quad (6) \\
(\exists f : \text{Child}(f, \text{Father}(I)) \wedge \text{Parent}(f, m) \wedge \text{Male}(f))
\end{aligned}$$

$$\begin{aligned}
& \text{Male}(m) \wedge \\
& (\neg \exists s : s \neq I \wedge (\exists p : \text{Parent}(p, s) \wedge \text{Parent}(p, I))) \wedge \quad (7) \\
(\exists f : \text{Child}(f, \text{Father}(I)) \wedge \text{Parent}(f, m) \wedge \text{Male}(f))
\end{aligned}$$

$$\begin{aligned}
& \text{Male}(m) \wedge \\
& (\forall s : s = I \vee \neg(\exists p : \text{Parent}(p, s) \wedge \text{Parent}(p, I))) \wedge \quad (8) \\
(\exists f : \text{Child}(f, \text{Father}(I)) \wedge \text{Parent}(f, m) \wedge \text{Male}(f))
\end{aligned}$$

$$\begin{aligned}
& \text{Male}(m) \wedge \\
& (\forall s : s = I \vee \neg(\exists p : \text{Parent}(p, s) \wedge \text{Parent}(p, I))) \wedge \quad (9) \\
(\exists f, g : g = \text{Father}(I) \wedge \text{Child}(f, g) \wedge \text{Parent}(f, m) \wedge \text{Male}(f))
\end{aligned}$$

$$\begin{aligned}
& \text{Male}(m) \wedge \\
& (\forall s : s = I \vee \neg(\exists p : \text{Parent}(p, s) \wedge \text{Parent}(p, I))) \wedge \\
(\exists f, g : \text{Parent}(g, I) \wedge \text{Child}(f, g) \wedge \text{Parent}(f, m) \wedge \text{Male}(f) \wedge \text{Male}(g)) \quad (10)
\end{aligned}$$

$$\begin{aligned}
& \text{Male}(m) \wedge \\
& (\forall s : s = I \vee \neg(\exists p : \text{Parent}(p, s) \wedge \text{Parent}(p, I))) \wedge \\
(\exists f, g : \text{Parent}(g, I) \wedge \text{Parent}(g, f) \wedge \text{Parent}(f, m) \wedge \text{Male}(f) \wedge \text{Male}(g)) \quad (11)
\end{aligned}$$

$$\begin{aligned}
& \text{Male}(m) \wedge \\
& (\forall s : (\exists p : \text{Parent}(p, s) \wedge \text{Parent}(p, I)) \rightarrow s = I) \wedge \\
(\exists f, g : \text{Parent}(g, I) \wedge \text{Parent}(g, f) \wedge \text{Parent}(f, m) \wedge \text{Male}(f) \wedge \text{Male}(g)) \quad (12)
\end{aligned}$$

$$\begin{aligned}
& \text{Male}(m) \wedge \\
(\forall s : (\exists p : \text{Parent}(p, s) \wedge \text{Parent}(p, I)) \rightarrow s = I) \wedge & (13) \\
(\exists f : f = I \wedge \text{Parent}(f, m) \wedge \text{Male}(f))
\end{aligned}$$

$$\text{Male}(m) \wedge \text{Parent}(I, m) \wedge \text{Male}(I) \quad (14)$$

$$\text{Male}(m) \wedge \text{Child}(m, I) \quad (15)$$

$$\text{Son}(m, I) \quad (16)$$

From this, we can conclude that “that man” is, in fact, the speaker’s son.

In prenex normal form, the original riddle can be expressed as follows:

$$\begin{aligned}
\exists f : \forall s : (\neg \text{Brother}(s, I) \vee \neg \text{Son}(f, \text{Father}(I)) \vee \neg f = \text{Father}(m) \vee \\
\neg \text{Male}(m)) \wedge (\neg \text{Sister}(s, I) \vee \neg \text{Son}(f, \text{Father}(I)) \vee \neg f = \text{Father}(m) \vee \\
\neg \text{Male}(m))
\end{aligned}$$

Problem 9.23

- a. Premise: $\forall h : \text{Horse}(h) \rightarrow \text{Animal}(h)$
Conclusion: $\forall c, q : (\text{HeadOf}(c, q) \wedge \text{Horse}(q)) \rightarrow \exists a : (\text{HeadOf}(c, a) \wedge \text{Animal}(a))$
- b. Premise: $\neg \text{Horse}(h) \vee \text{Animal}(h)$
Conclusion: $\text{HeadOf}(C, Q) \wedge \text{Horse}(Q) \wedge (\neg \text{HeadOf}(C, a) \vee \neg \text{Animal}(a))$
- c.
$$\frac{(\neg \text{HeadOf}(C, Q) \vee \neg \text{Animal}(Q)), \text{HeadOf}(C, Q)}{\neg \text{Animal}(Q)}$$

$$\frac{\neg \text{Animal}(Q), (\neg \text{Horse}(h) \vee \text{Animal}(h))}{\neg \text{Horse}(Q)}$$

$$\frac{\neg \text{Horse}(Q), \text{Horse}(Q)}{\perp}$$