Fast quantification of splice junctions from RNA-seq data by *sjcount* v3.0

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1 Synopsis

The purpose of sjcount is to provide a fast utility for counting splice junctions in BAM files. It is the annotation-agnostic version of bam2ssj. This document describes the version $\mathbf{v3.0}$ of sjcount. The older versions of sjcount (v1.0, v2.0) is also included in the package inder the name deprecated.

2 Changes since v2.0

The has been a substrantial change between v2.0 and v3.0.

- 1. The utility now counts and reports reads with multisplits
- 2. Accordingly, the output format has changed to account for multisplits
- 3. A simplier and more efficient data structure is now used to store and parse multisplits
- 4. Test rountines are now added to check the quality and integrity of the output as compared to the output of a perl script with easily controlled syntax

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3 Installation and usage

See README.md file for installation instructions. The program *sjcount* is used from the command line with the following keys

```
sjcount -bam bam_file [-ssj junctions_output] [-ssc boundary_output]
        [-read1 0|1] [-read2 0|1] [-unstranded] [-nbins number_of_bins]
        [-lim number_of_lines] [-quiet]
```

where

- bam_file is a sorted input BAM file with a header
- junctions_output is the output file with junction counts
- boundary_output is the output file with boundary counts
- read1 0/1, reverse complement read1 no/yes (default=no)
- read2 0/1, reverse complement read2 no/yes (default=no)
- unstranded, force strand=0
- **nbins** number of offset bins, (default=1)
- maxnh the max value of the NH tag, (default=none)
- lim stop after reading these many lines, (default=no limit)
- quiet suppress verbose output NOTE: use -quiet if you redirect stderr to a file!

The output consists of two files. First, a tab-delimited file containing multisplit counts is produced as follows

chr1_100_200_+	1	34	1
chr1_100_200_+	1	36	1
chr1_100_200_+	1	37	6
chr1_100_200_+	1	38	3
chr1_100_200_300_400_+	2	49	1
chr1_100_400_+	1	33	1

where the first column contains the coordinates of the split alignment (including multi-splits, see below). The second column contains the numer of splits. The third column contains the offset sefined as the distance within the short read sequence of the latest split (defined below). The last column is the respective count, i.e., the number of split-mapped reads with the given combination of alignment coordinates and offset.

For instance, chr1_100_200_+ denotes an alignment that was split once between positions 100 and 200 on the '+' strand, while chr1_100_200_300_400_+ denotes an alignment that was split twice, first between positions 100 and 200, and then between positions 300 and 400. The coordinates are 1-based and always refer to terminal *exonic* nucleotides. The strand is denoted by '+' and '-' for stranded data or by '.' for unstranded data.

The second output is also a tab-delimited file which contains the counts of read alignments that *overlap* exon boundries (exon boundries are defined by splice junctions in the previous file). In this version all alignments that overlap an exon boundary by at least one nucleotide are counted (in older versions only continuous alignments were counted. This second file is optional and is needed to compute the completness of splicing index [2, 3].

4 Method

By definition, we say that we observe a *splice junction* each time we see an 'N' symbol in the CIGAR attribute of the alignment. If the CIGAR attribute contains several N's, then we have a *multi-split* or n-split, where n is the number of N's in CIGAR. In this terms, each 1-split defines one splice junction while each n-split defines n splice junctions.

Each multi-split is counted according to the number of splits so that, for example, the alignment chr1_100_200_300_400_500_600_+ is counted once as a 3-split, two times as a 2-split (chr1_100_200_300_400_+ and chr1_300_400_500_600_+), and three times as single-split (chr1_100_200_+, chr_300_400_+, and chr1_500_600_+). The positions of splits are decided entirely by the mapper which produced the alignment.

As an example, consider the multi-split alignment shown in Figure 1 below. In the output file it will be counted in three lines: in chr1_31_52_+ as having 1 split, in chr1_64_78_+ as having 1 split, and in chr1_31_52_64_78_+ as having 2 splits. One may want to subset the output to regular splice junctions by requiring the second column be equal to one.

Artifacts may arise from combining counts that come from different starting

```
10
         20
              30
                   40
                         50
                              60
                                    70
                                         80
                              1
              П
   chr1
Query
     CTAGGAGACGG**TAGGAG......ATCTA*AAAACAT......GATa
               |<---->|
                                 |<--- SJ2 --->|
The corresponding SAM line is:
Query
    123
       chr1 14
              255
                 5M1I5M2D6M20N5M1D7M13N3M1S 1234
```

Figure 1: An example alignment and its CIGAR attribute

positions of the alignment. We define the offset to be the distance (in the query sequence!) from the first alignment position to the corresponding 'N'. For instance, the junction SJ_1 in Figure 1 has offset 17, while the junction SJ_1 has offset 29. The offset of the multi-split is defined to be the offset of it's last N, i.e., 29 in this case. Since the offset is defined as a position in the query sequence, its value cannot exceed the read length.

Some offsets may give artifactually large read counts corresponding to PCR artefacts [1]. In Figure 2 we show six split reads supporting the same splice junction with offsets 14 (Q1), 12 (Q2–Q4), and 8 (Q5–Q6). Note that offsets appear decreasing when sequentially processing lines a sorted BAM file.

	10 	20 	30 	40 	50 	60 I	70 	80 I				
	123456789012345678901234567890123456789012345678901234567890123456789012											
Ref	Ref AGTCTAGGGACGCATAGGAGGTGAGCATTTGTGTACGCAGATCTACAAAACATGTGTCACGGATAGGATCG											
Q1	GGACGGCATAGGAGATCT											
Q2	ACGGCATAGGAGATCTAC											
QЗ	ACGGCATAGGAGATCTAC											
Q4	ACGGCATAGGAGATCTAC											
Q5	CATAGGAGATCTACAAAA											
Q6	CATAGGAGATCTACAAAA											

Figure 2: Split-mapped reads support the same splice junction with different offsets

The quantification of abundance is done as follows. We initialize and keep nbins separate counters for each n-split. For each instance of n-split, we incre-

ment the counter corresponding to its offset. If the offset is larger than or equal to nbins then it is set to be equal to nbins - 1.

For example, in the default settings we have nbins = 1. This means that the bin number will be 1 - 1 = 0 for all supporting reads, regardless of their offset (t = 14 for Q1, t = 12 for Q2–Q4, and t = 8 for Q5–Q6 in Figure 2). Therefore, there is only one counter to increment, and the result will be so called "collapsed" counts. The output corresponding to Figure 2 will then be

By contrast, if we set *nbins* equal to read length, there will be a separate counter for each offset and the output corresponding to Figure 2 will be

One single number is usually reported for each splice junction as an endpoint. Normally, the user wants to know how many reads aligned to a certain split regardless of the offset. This number is equal to the sum of counts for the given alignment over all values of offset. In other words, the total number of counts is obtained from offset-specific counts by aggregation using the function $f(x_1, \ldots, x_n) = x_1 + \cdots + x_n$.

We also report two other important feature that are obtained from offsetspecific counts by aggregation using different functions. The number of staggered counts is the result of aggregation using $f(x_1, ..., x_n) = \theta(x_1) + \cdots + \theta(x_n)$, where $\theta(x) = 1$ for x > 0 and $\theta(x) = 0$ for $x \le 0$. The entropy of the distribution of counts by offsets is obtained by using

$$f(x_1, \dots, x_n) = \log_2(\sum_{i=1}^n x_i) - \frac{\sum_{i=1}^n x_i \log_2(x_i)}{\sum_{i=1}^n x_i}.$$

The entropy and the number of staggered reads are two important characteristics that can be used to filter out non-uniform distibution of read counts.

References

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