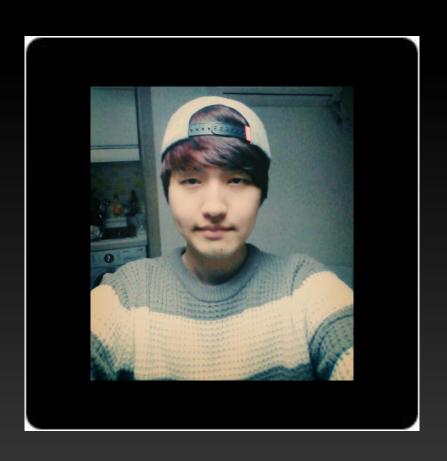


Hi, Futex! Hello, TowelRoot!

Lee JiHun wh_pesante@naver.com



Who am I?



- Name: Lee Ji Hun
- NickName: Pesante
- Type: Pwnable, Reversing
- 출몰지역: 대전, 서울
- Monsterz 잉여 회원
- Argos 잉여 회원

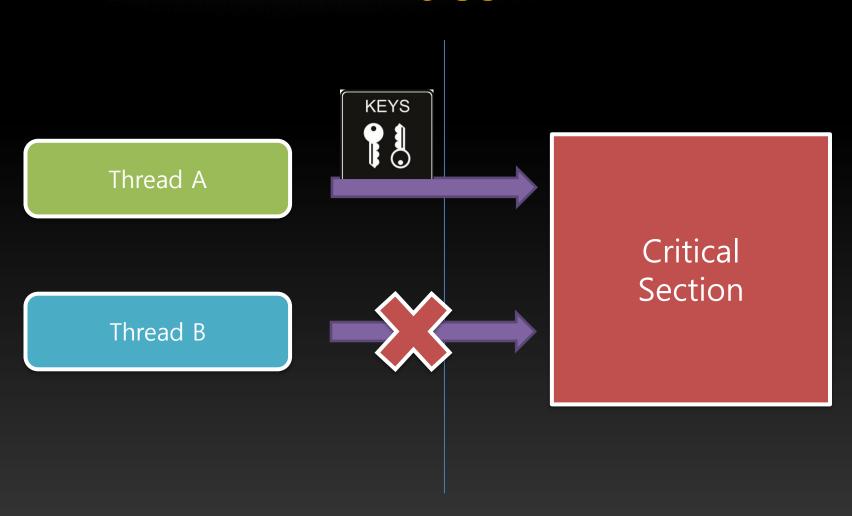


CONTENTS

- 1. Mutex & Condvar
- 2. Futex system call
- 3. Futex Requeue & PI
- 4. Simpler bug
- 5. Into the TowelRoot

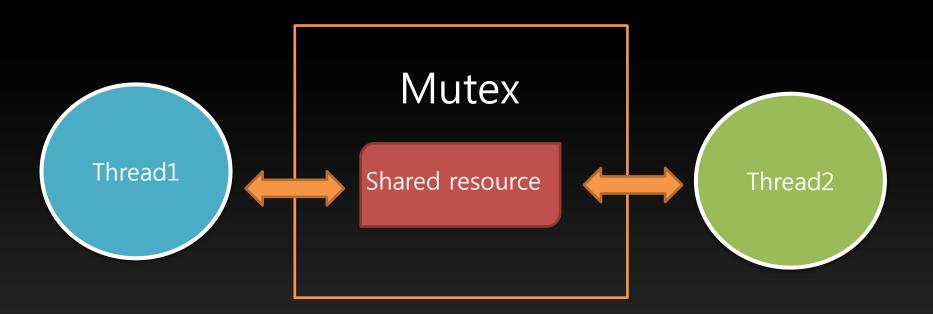


Mutex





Mutex



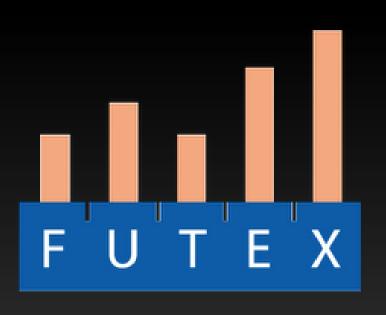
Wait/Release



Condvars

- "시그널 발생"과 "시그널에 대한 기다림" 을 통한 좀더 적극적인 동기화 방법을 제공
- 이러한 시그널을 기다리기 위해서 pthread_cond_wait() 와 pthread_cond_timedwait() 2개의 함수를 제공





- Linux kernel system call for locking
- Fast user-space mutexes
- 32-bit integer in user memory





Unlock=0

Lock=tid





Unlock State

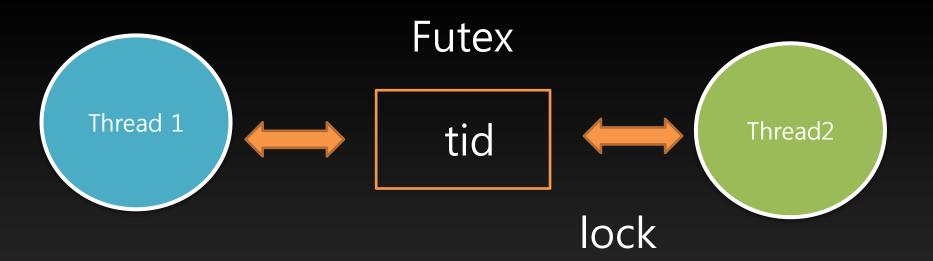
















Wait/Release



```
if (*uaddr == 0) {
    *uaddr = tid; /* Lock acquired */
}
```



장점

- 1. 빠르다.
- 2. User space에 있으므로 공유 메모리에 올려 쓰레드뿐만 아니라 프로세스끼리도 손쉽게 사용이 가능

단점

- 1. 멀티코어에서는 그다지 성능을 발휘하지 못함
- 2. Thread 수가 늘어나면 성능이 급격히 감소
- 3. 인터럽트를 지원하지 않음



Futex Syscall

int sys_futex(int *uaddr, int op, int val, const struct timespec *timeout, int *uaddr2, int val3);

- Action depends on the op arguments
- The arguments can be unused, or cast to differe nt types



FUTEX_WAIT & FUTEX_WAKE

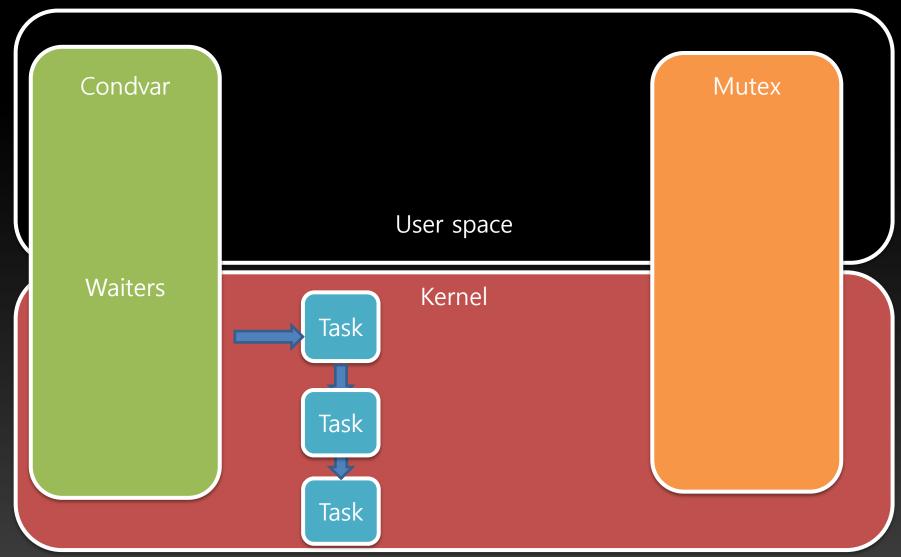
FUTEX_WAIT

This operation causes the thread to be suspended in the kernel until notified

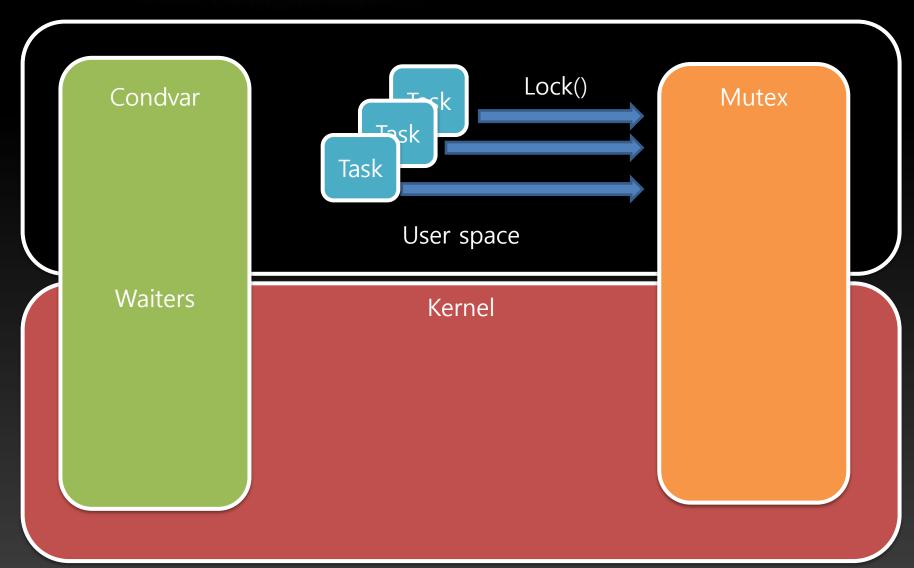
FUTEX_WAKE

To wake up one or more threads waiting on a futex this operation can be used.

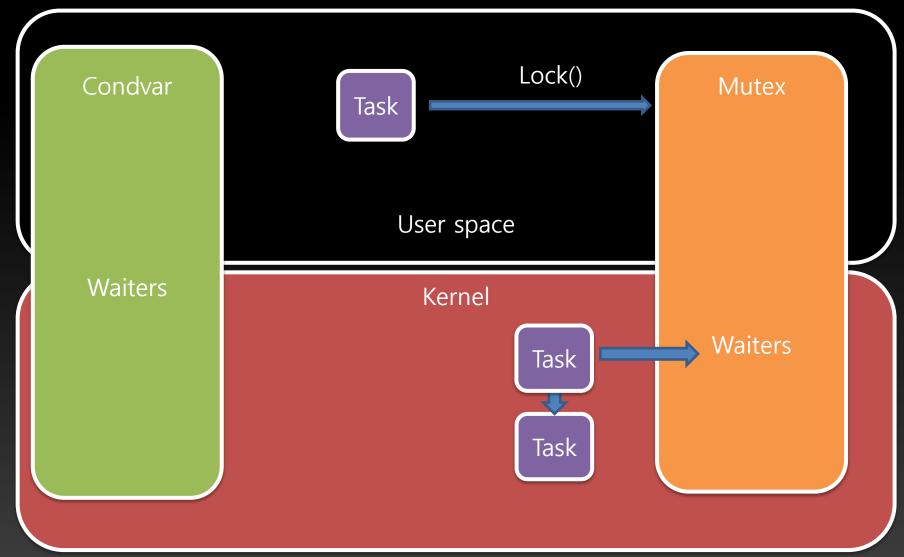




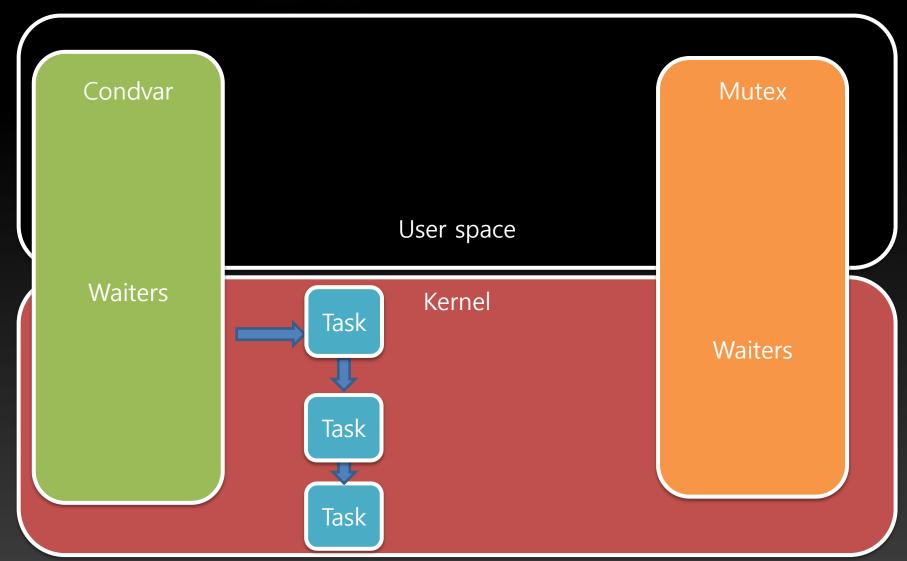




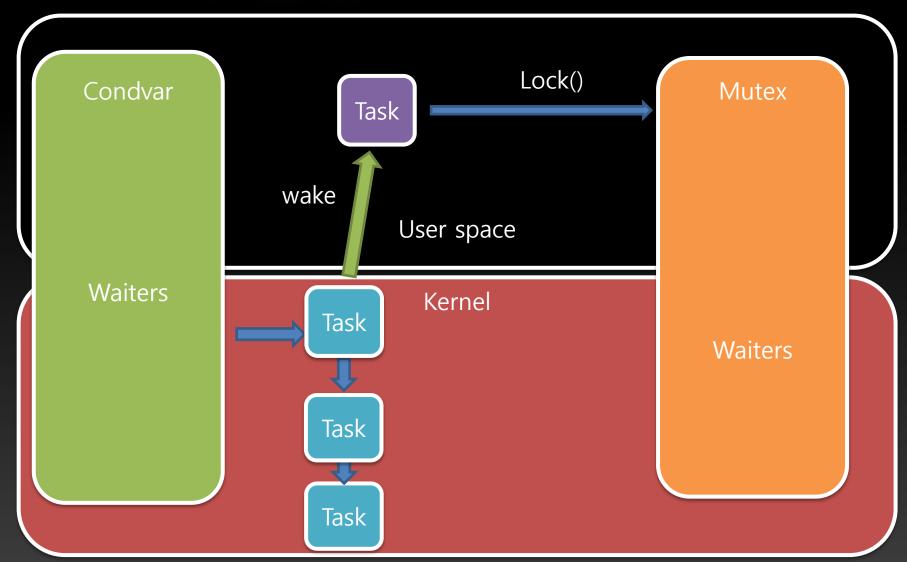




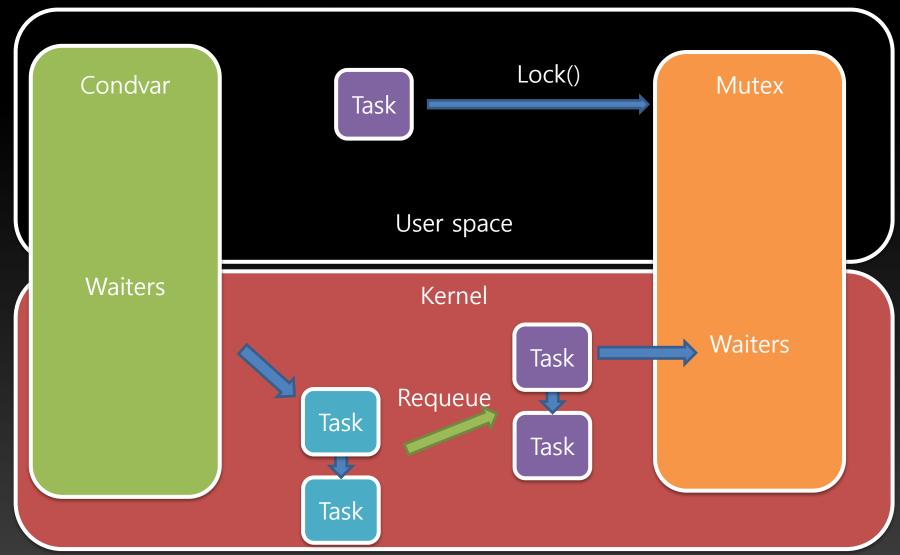




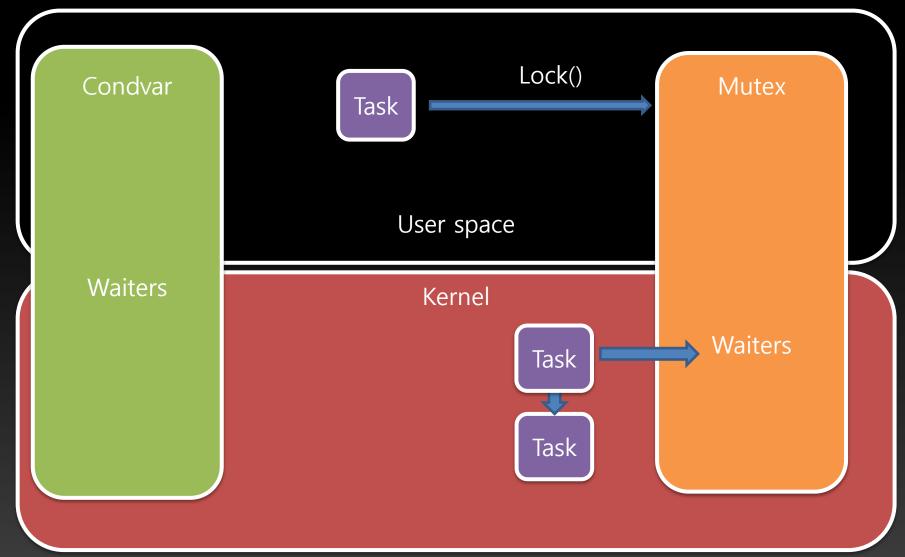












Priority



화장실



아, 화장실



Priority







먼저 쓰세요





Priority

Um...





Wating...



1등

2등

3등









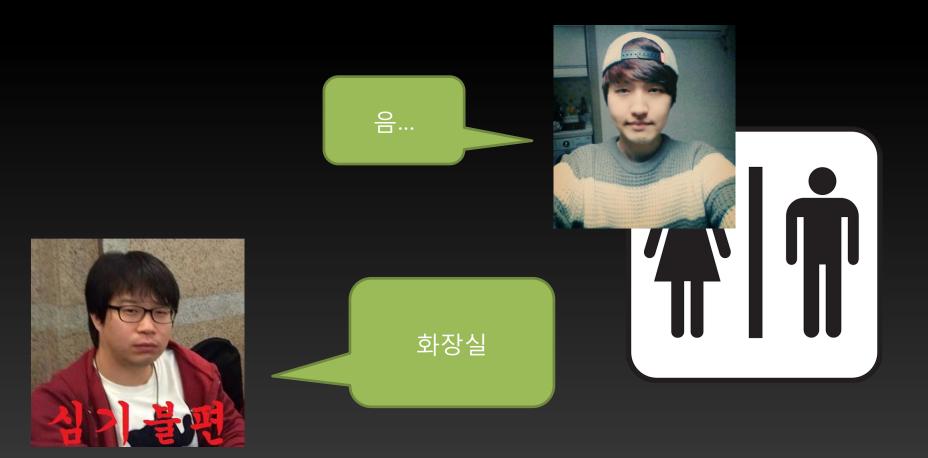




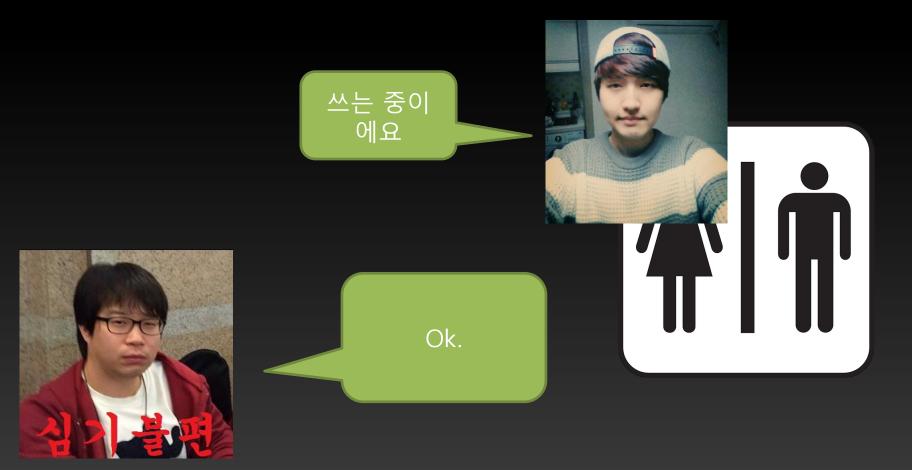
아, 화장실















CPU좀 줘봐. 다 른 거 할 거 있어.









일하는중..



Wating..















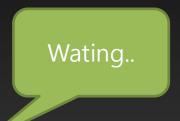


















Priority Inversion

1등



2등



3등



우선순위가 가장 높은 사람이 가장 늦게 처리됨



쓰는 중이 에요





그러면 나의 우 선순위를 받아 라!





CPU좀 줘봐. 다 른거 할 거 있어.



안돼요.























1등

2등

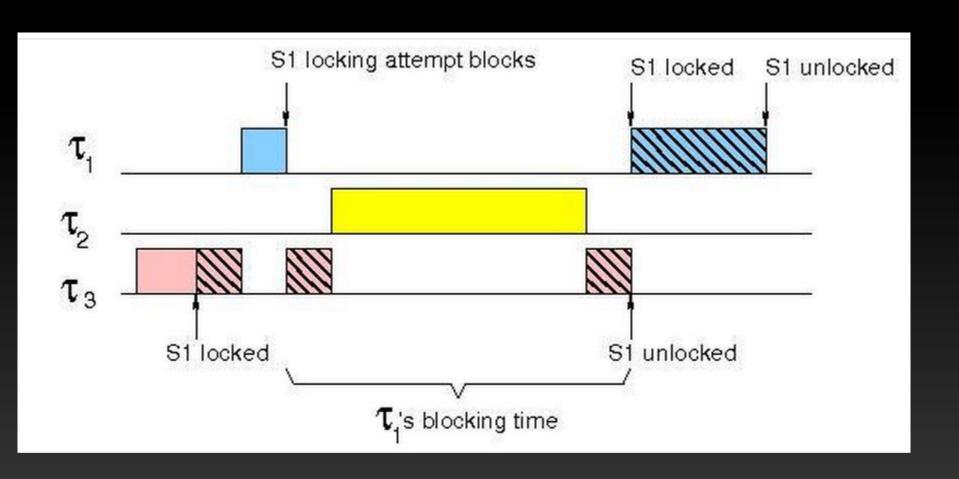
3등



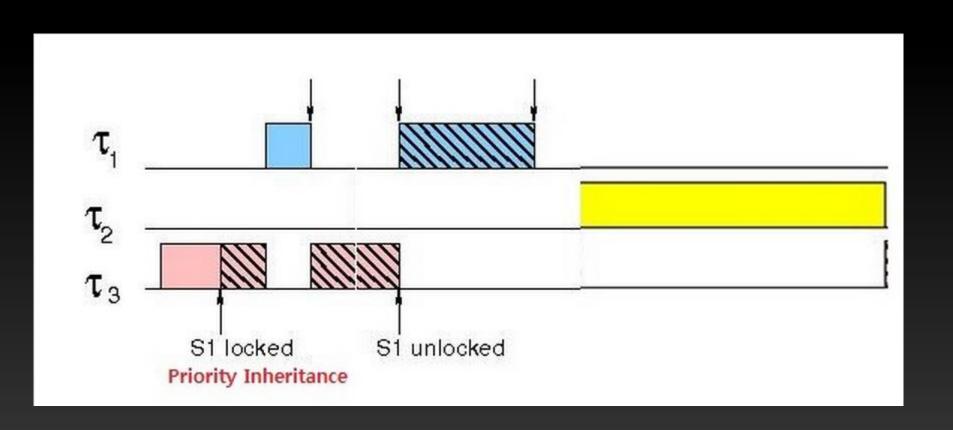














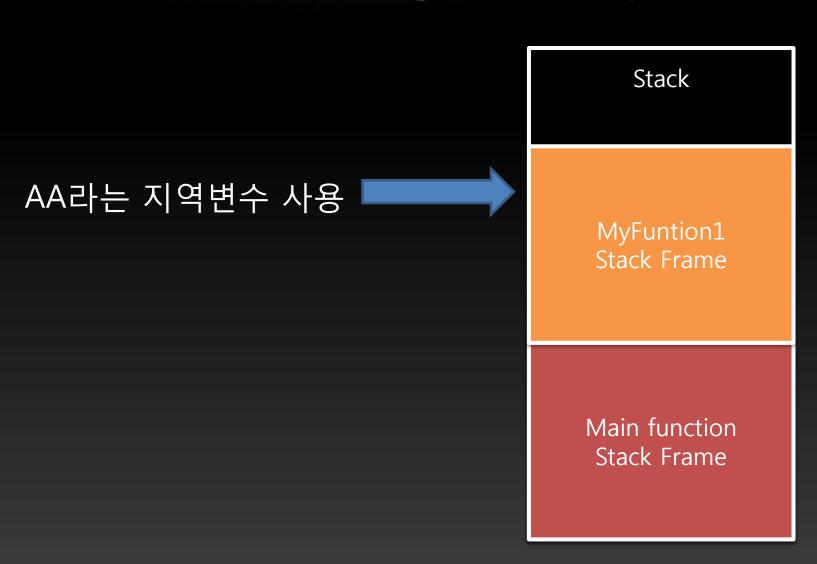
PI futexes

- The user-space futex value is zero for unlocked, or holds the thread ID of the owner
- The FUTEX_LOCK_PI and FUTEX_UNLOCK_PI calls are used instead of wait and wake
- Unlocking a PI-futex wakes only the highest priority waiter



```
#include <stdio.h>
int main()
      MyFunction()
      MyFunction2()
      • • • •
```







AA의 값은 그대로 남아있음 Stack

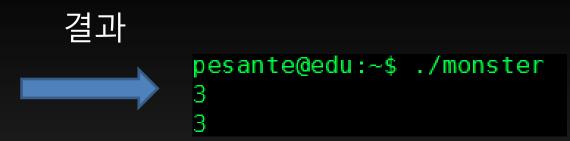


AA의 값은 그대로 남아있음 Stack

Myfuntion2 Stack Frame



```
#include <stdio.h>
int MyFunction1()
    int a=3;
    printf("%d\n", a);
}
int MyFunction2()
    int b;
    printf("%d\n", b);
int main()
    MyFunction1();
    MyFunction2();
```





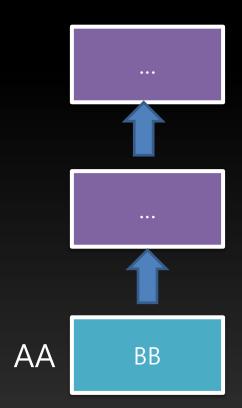
Stack

AA라는 지역변수 사용

MyFuntion1 Stack Frame



List



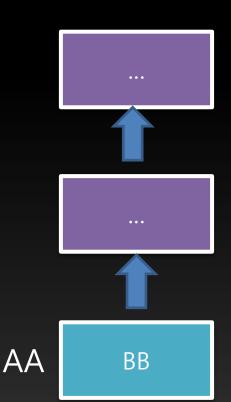
- 1. AA라는 지역변수 사용
- 2. 전역변수 List에 등록

Stack

MyFuntion1 Stack Frame



List

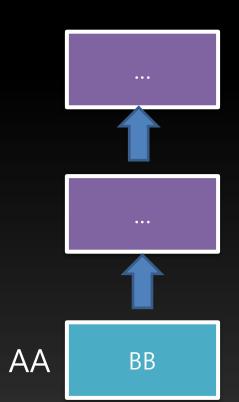


리스트를 해제하지 않고 MyFunction1 종료 Stack

MyFuntion1 Stack Frame



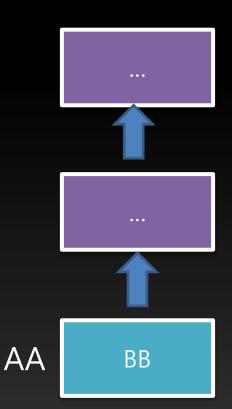
List



리스트를 해제하지 않고 MyFunction1 종료 Stack



List



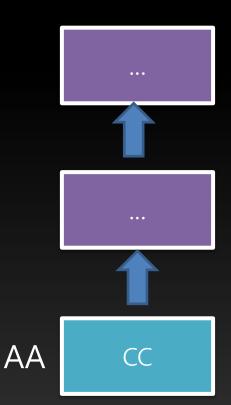
로컬 변수 사용

Stack

MyFuntion2 Stack Frame



List



로컬 변수 사용

Stack

MyFuntion2 Stack Frame



- 1. MyFunction1 내의 지역변수를 리스트에 추가
- 2. 지역변수(AA)를 리스트에서 해제하지 않고 함수 리턴
- 3. 다른 함수 내에서 지역변수(BB)를 사용
- 4. 리스트에 있던 AA의 값이 BB의 값으로 바뀜



Towelroot

because <u>activeroots</u> were <u>so last summer</u> <u>official XDA thread</u>

please donate on success #keepingrootfree you probably paid around \$500 for your phone think of the value added by this root these things are not easy to write and donations influence if i do the next one



root for at&t and verizon s5 but in seriousness should support all phones < jun 3 2014 click the lambda install the apk run towelroot ~ geohot ~

paypal: click the Donate button below bitcoin: 1DFhk6fs26AafaNiQvEgdGM7NSaLzVN2VE



- 안드로이드 루팅 어플
- Geohot에 의해 개발
- 커널 취약점을 이용하여 손쉽게 루팅
- CVE-2014-3153



- The vulnerability is based on two logical bugs in the Linux kernel futex mechanism
- one which allows releasing a pi-lock from userspace
 - another which allows requeuing a lock twice.

취약 버전: http://www.securityfocus.com/bid/67906



```
static int futex_wait_requeue_pi(.....)
   struct hrtimer_sleeper timeout, *to = NULL;
    struct rt mutex waiter rt waiter;
   struct rt_mutex *pi_mutex = NULL;
   struct futex_hash_bucket *hb;
   union futex_key key2 = FUTEX_KEY_INIT;
   struct futex_q q = futex_q_init;
   int res, ret;
   //...
   q.rt_waiter = &rt_waiter;
   //...
   //futex_wait_queue_me(hb, &q, to);
   if (!q.rt_waiter) {
   //...
```



rt_mutex_waiter



```
static int futex_wait_requeue_pi(.....)
    if (!q.rt_waiter) {
         if (q.pi_state && (q.pi_state->owner != current)) {
                   .........
    } else {
         ret = rt_mutex_finish_proxy_lock(pi_mutex, to, &rt_waiter, 1);
         .......
```



```
void requeue_pi_wake_futex(struct futex_q *q, union futex_key *key,
                     struct futex_hash_bucket *hb)
       get_futex_key_refs(key);
       q - key = key;
       __unqueue_futex(q);
      WARN_ON(!q->rt_waiter);
      q->rt_waiter = NULL;
      q - > lock_ptr = & hb - > lock;
      wake_up_state(q->task, TASK_NORMAL);
```



```
Thread B: futex_lock_pi( &futex2 )
Thread A: futex_wait_requeue_pi(&futex1,
&futex2)
Thread B: futex_cmp_requeue_pi( &futex1,
&futex2)
Thread B: futex2 = 0
Thread B: futex_cmp_requeue_pi( &futex2,
&futex2)
```



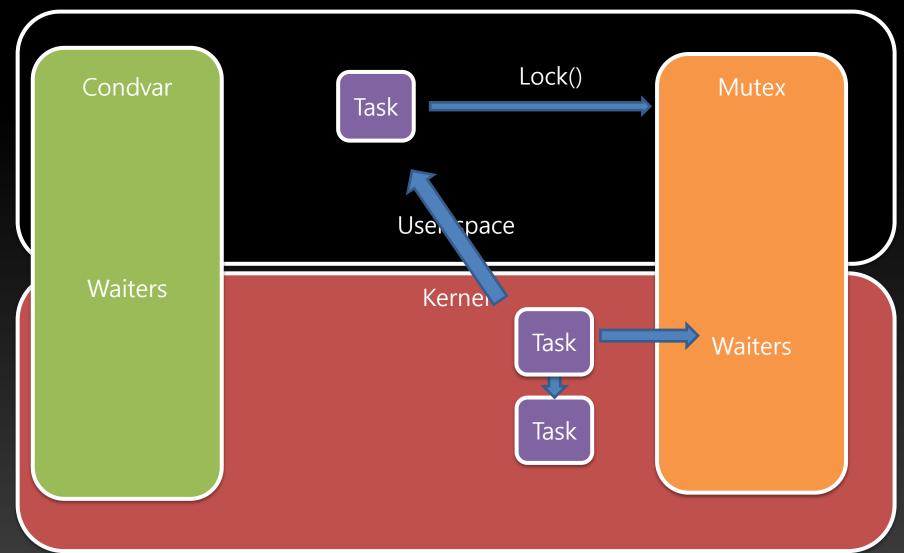
1)Futex_lock_pi(B) Create Thread 2)Futex_wait_requeue_pi(A, B) 3)futex_requeue(A,B) 4) B=0 5)futex_requeue(B,B) 6)overwrite!

Thread1

Thread2

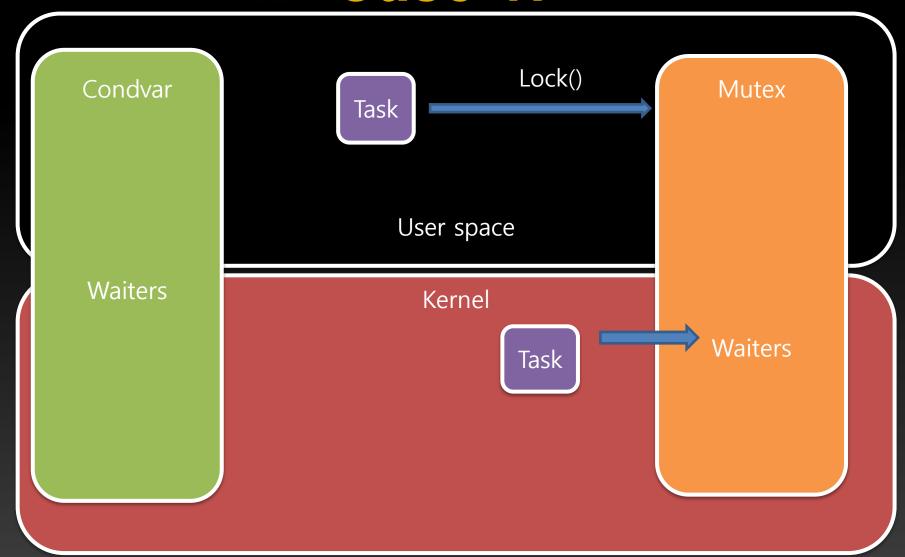


FUTEX_REQUEUE Case 1.



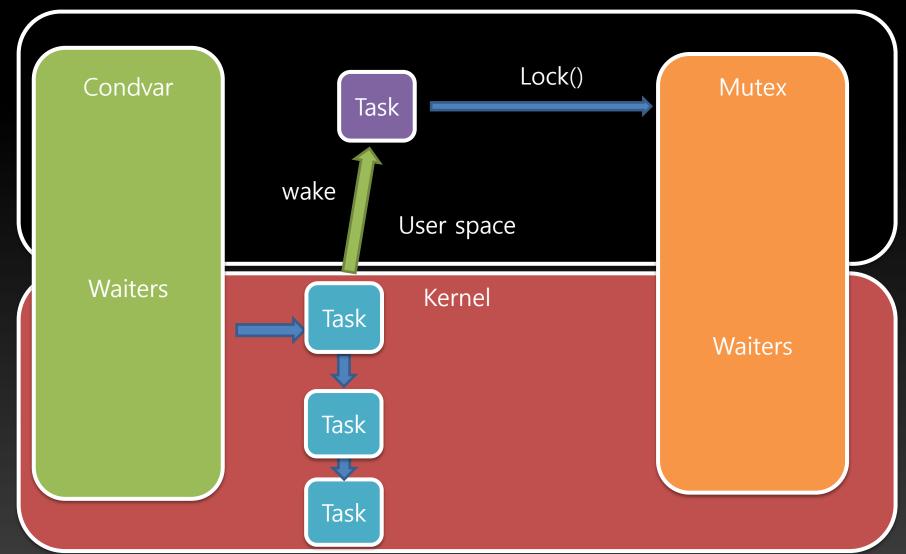


FUTEX_REQUEUE Case 1.



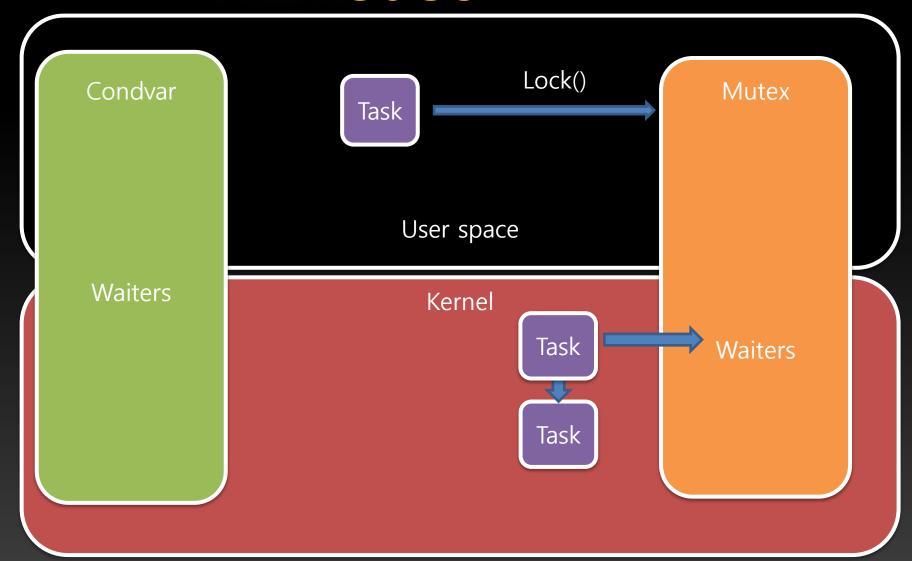


FUTEX_REQUEUE Case 2.





FUTEX_REQUEUE Case 2.





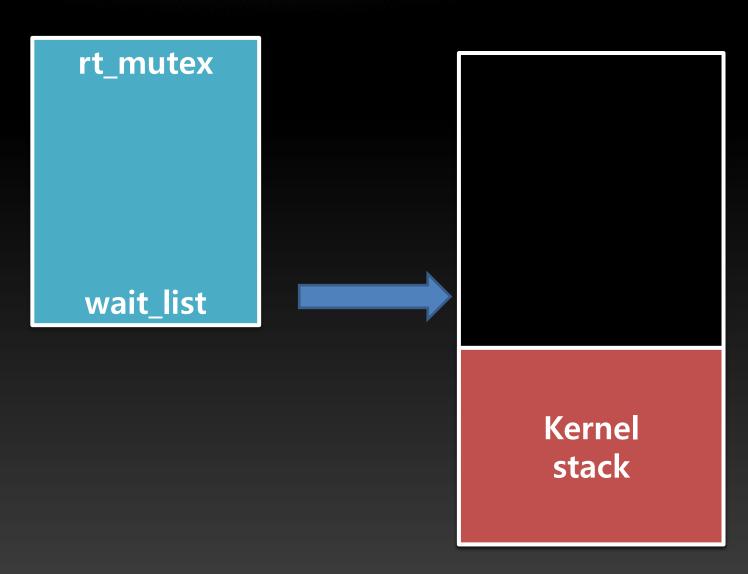


Futex_wait_requ eue_pi

rt_waiter

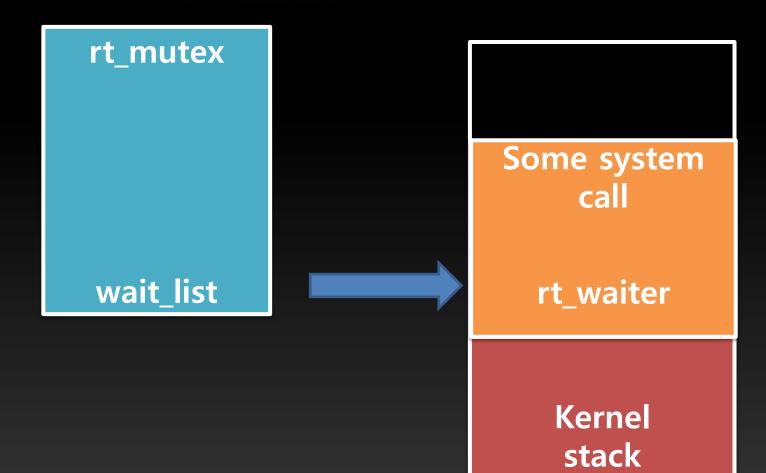
Kernel stack







CVE-2014-3153





Possible System call

systemcall

sendmmsg

recvmmsg

sendmsg

recvmsg

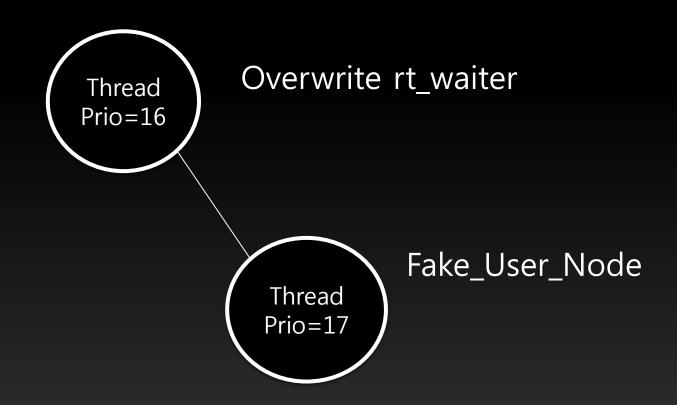
semctl



Exploitation

- 1. Kernel memory address leak
- 2. Arbitary memory corruption
- 3. Overwrite addr_limit
- 4. Overwrite cred struct
- 5. Get shell



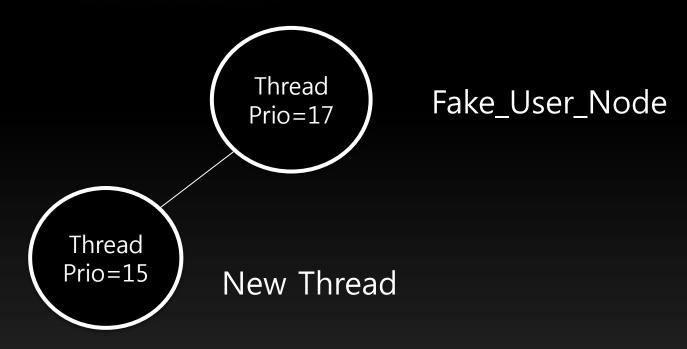






Fake_User_Node

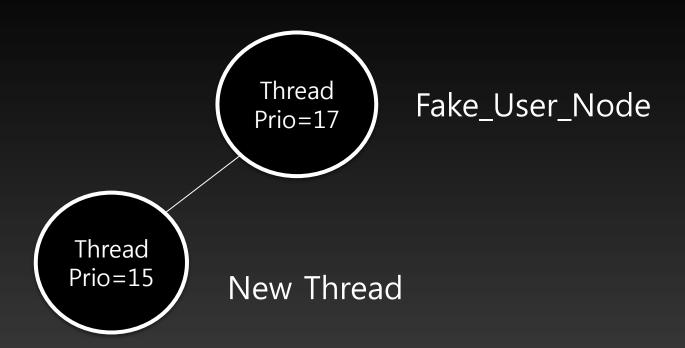




```
futex_wait_requeue_pi(&srcfutex,&destfutex) and futex_cmp_requeue_pi(&srcfutex,&destfutex) or Futex_lock_pi
```



- 1. Fake_User_Node는 유저영역
- 2. 현재 Fake_User_Node의 left가 커널 메모리를 소유





Thread Info

```
struct thread_info {
    ......
    mm_segment_t addr_limit; /* address limit */
    struct task_struct *task; /* main task structure */
    .......
};
```

thread_info # 프로세스에 관한 저수준 정보와 커널용 스택을 관리한다.



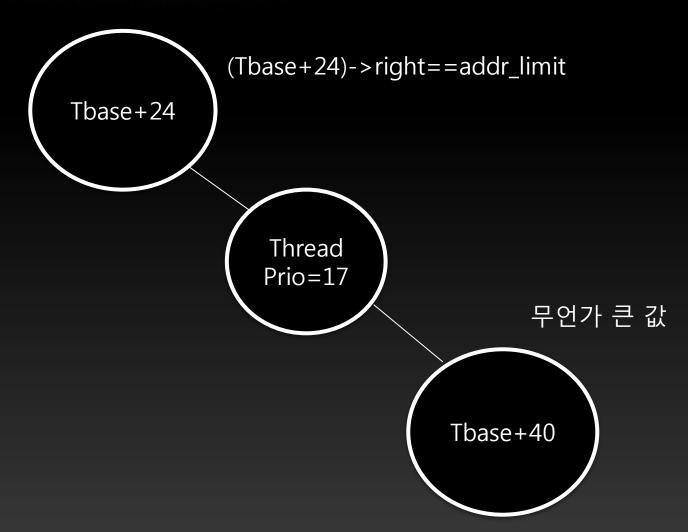
Thread Info

```
static inline struct thread_info *current_thread_info(void)
{
    register unsigned long sp asm ("sp");
    return (struct thread_info *)(sp & ~(THREAD_SIZE - 1));
}
```

tbase = (unsigned long)kernel_waiter & 0xfffffffffffe000;

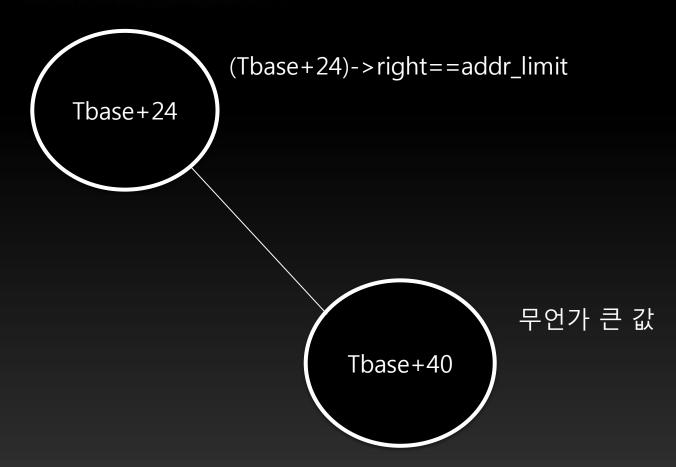


Arbitary memory corruption





Arbitary memory corruption





Cred Struct

- The security context of a task



Cred Struct

```
struct cred {
                           /* effective UID of the task */
    uid_t
               euid;
    gid_t
                           /* effective GID of the task */
                egid;
          • • • • • • • • •
};
                                                    Exploit!
             /bin/sh
                                                      root!
```

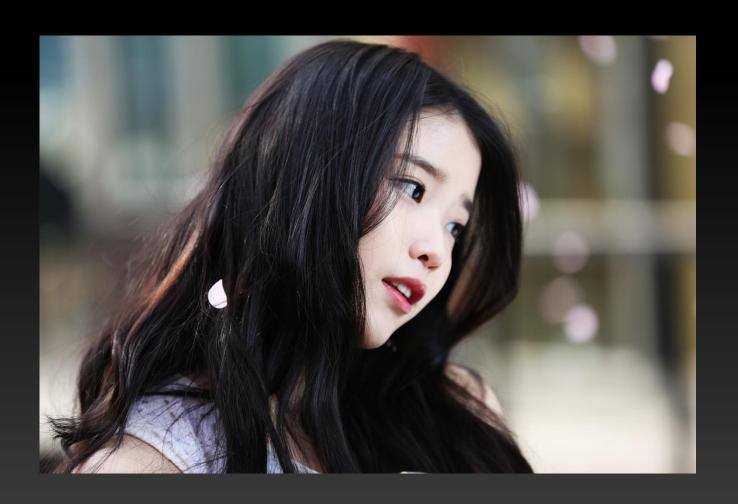


시연





QnA



THANK YOU!