### Structures and All That

October 8, 2019

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We can get you prepared for the 31st century
With advanced programming and quad rendering
And Java plus plus plus scripting language
We offer advanced job placement assistance"

from Upgrade by Deltron 3030

We have taken a weird approach by fixating on data structured in the form of "or" first.

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- Let's say that in Java that you have some person class with a first and last name represented as strings.
- It is easy to define a method that returns the person's full name by concatenating the first and last name.

So, why do we need compound data?

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- We can represent a grid with one number in the same sense that we can simulate a 10x10 2D array with a 100 element array.

## Structs Make Things Easier

Personally, I like doing things the easy way.



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```
def first_name(tup):
   return tup[0]
```

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We can actually define other data structures in terms of things like lists.

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- So Python gives classes (or named tuples) as a way to more easily define such structured data.

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- We will return to discussing lists in more detail later, since they are extremely important.
- But for now, remember that we wanted to avoid the inconveniences given by using other existing data types to represent some piece of compound data!

We said that we didn't want to represent all of our compound data with existing structures like lists are tuples, so let's *finally* talk about structs.

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- To get the x-coordinate: (point-x one-two)

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Let's add functional examples as tests:

```
(check-expect (distance-to-0 (point 0 5)) 5)
(check-expect (distance-to-0 (point 7 0)) 7)
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previous version of this function that took in two parameters.
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Testing that function is simple, so let's just move on to talking about structs in general.

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  - 2. One selector per field, which extracts the value of the field from a structure instance; and
  - 3. One structure predicate, which, like ordinary predicates, distinguishes instances from all other kinds of values.

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- The selectors per field are (point-x point-val) and (point-y point-val). The general form of a selector for a specific field is (struct-name-field-name val)
- 3. A predicate for checking types is automatically created, for example: (point? point-val) and in general a predicate struct? is created.

Here are some basic examples of structs:

• (struct movie [title producer year] \#:transparent)

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- Sample Problem Develop a structure type definition for a
  program that deals with bouncing balls,. The balls location is
  a single number, namely the distance of pixels from the top.
  Its constant speed is the number of pixels it moves per clock
  tick. Its velocity is the speed plus the direction in which it
  moves.

# Designing Our Ball Struct

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- Let's first consider defining a 2D vector struct as follows:
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- Now, we can represent a ball as a point (which only has positive components) and a vector (which can have negative components):

```
(struct 2D-ball [position vec] \#:transparent)
```

Our 2D Ball struct has nested occurrences of other structs. This is a natural thing, and even recursive descriptions of data are natural, i.e. linked lists and binary trees. But we can also consider using a *flat representation* for our 2D Ball, which doesn't nest structs.

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(struct point [x y] \#:transparent)
; A Point is a structure:
; (point Number Number)
; interpretation a point x pixels from left, y from top
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- Let us consider how to relate a struct description to a diagram that illustrates its "structure".
- (struct centry [name home office cell] \#:transparent)

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• (define pl (centry "Al Abe" "666-7771" "lee@x.me"))
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So, for cell phone entry structs we had three fields that can contain possible values. Consider the following concrete one:

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```
name phone email "lee@x.me"
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 In a sense, calling (centry-name p1) is unlocking a box in the struct, with a specific key that allows you to retrieve the underlying value. In general we can think of field access as a kind of "unboxing".

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- Using a "key" to unlock the wrong box raises a runtime error (centry-name (point 1 2)) → entry-name:expects a centry, given (point 42 5)

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- **Denotation** the object or concept to which a term refers, or the set of objects of which a predicate is true.

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- My definition A formal language that has a computable semantics specifying the denotation of a "sentence" written in the language. This denotation can be given by compilation into some other formal language, or by a direct interpretation.

# Some Important People







Kurt Gödel

Alan Turing

Alonzo Church

# Some Important People



Kurt Gödel



Alan Turing

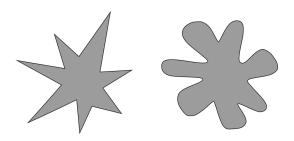


Alonzo Church



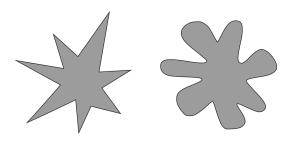
### Natural Interpretations?

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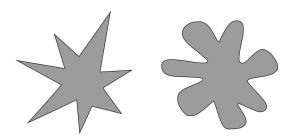
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### Natural Interpretations?

Here's a fun little experiment. Look at these two shapes:



- Which of these shapes is named Bouba and which Kiki?
- This is likely due to some physical attributes of sound, but our strong preference for naming the round one Bouba and the sharp one Kiki shows that humans have innate preferences about the naming and representation of things.

### Interpreting Structure

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; A Ball-1d is a structure:
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; interpretation 1 distance to top and velocity
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• Interestingly, we can give 2 interpretations, depending on if the ball is moving vertically (interpretation 1) or horizontally (interpretation 2)

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```
; A Ball-2d is a structure:
; (ball Point Vel)
; interpretation a 2-dimensional position and velocity

(struct vel [deltax deltay] \#:transparent)
; A Vel is a structure:
; (vel Number Number)
; interpretation (vel dx dy) means a velocity of
; dx pixels [per tick] along the horizontal and
; dy pixels [per tick] along the vertical direction
```

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- The T just communicates that we are defining linked lists that work over any type (parametric polymorphism otherwise known as generics in Java)

### Structs and Program Design

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```
(define MTS (empty-scene 100 100))
(define DOT (circle 3 "solid" "red"))
; A Point represents the state of the world.
: Point -> Point
(define (main p0)
  (big-bang p0
    [on-tick x+]
    [on-mouse reset-dot]
    [to-draw scene+dot]))
```

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 Testing this is uninteresting, so let's consider if we were asked to define the x+ function, which takes in a Point and returns a new Point with an x-coordinate that is 3 units further to the right of the old point.

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(check-expect (x+ (point 0 0)) (point 3 0))
(check-expect (x+ (point 10 10)) (point 13 10))
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 To take inventory we project out the x and y fields from our point, as usual:

```
(define (x+p) (... (point-x p) ... (point-y p) ...)
```

```
(define (x+ p)
  (point (+ (point-x p) 3) (point-y p)))
```

To finish coding, we need to first add 3 to the x-coordinate and then pack the result back into a new point structure.

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(define (x+ p)
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- If we had more fields than y, we simply are projecting out the old field as an argument when creating the new struct value.
- This adds a lot of boilerplate code...

#### Struct Boilerplate

We might want to define a function, point-set-x which takes in a point and a value and produces a new point where the x-coordinate is the given value and the y-coordinate is taken from the old point.

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- point-set-x is known as a functional setter, similar to a more traditional setter in languages like Java.
- However, defining an update operation on a complicated structure can get very complicated, and we can get around this uses *lenses* (we may discuss this later in the course).

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(check-expect
  (reset-dot (point 10 20) 29 31 "button-down")
  (point 29 31))
(check-expect
  (reset-dot (point 10 20) 29 31 "button-up")
  (point 10 20))
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- For real this time:

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(define (reset-dot p x y me)
  (cond
    [(mouse=? "button-down" me) (... p ... x y ...)]
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Finally, though the skeleton looks complicated the final version of the function is relatively simple.

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- The takeaway here is that our skeletons can end up being more complicated than the actual final version.
- This is especially true when we consider skeletoning out code for which a parameter is a struct.

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- It would look like the following:

```
(define (complex-skeleton ball)
  (... (point-x (ball-position ball)) ...
     ... (point-y (ball-position ball)) ...
     ... (vec-delta-x (ball-vector ball)) ...
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 For some of the classes you see in Java, skeletoning out code in such a manner is infeasible...

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- Of course, things are easier if we picked the write functions to add in our skeleton, but this keeps the structure of the skeleton much simpler.
- We will return to this point later in the course.

Signatures can tell us a lot about a function. I'll put some signatures up, and what kind of functions could we write for these signatures?

Number Number → Number

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- The idea that a function signature can tell you about the behavior of the function is an extremely powerful idea.

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- For example, look at this data definition:

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• If this is the domain of some function, we restricted ourselves to needing to handle three values.

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With itemizations we specified subsets of the existing data universe.



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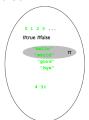
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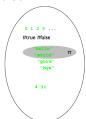
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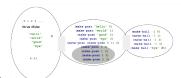
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Notice that we specify our point struct as containing two numbers.

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- Without that, it is up to us programmers to validate that points are only constructed with valid arguments.
- This might mean creating a function
   (define (point x y) ...) that checks that both x and y
   actually receive integers before calling the point constructor.
   More on this later.

Data definitions are a critical part of programming.

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- For itemizations, deal with each part separately
- For data definitions for structures, follow the natural language description; that is, use the constructor and pick an example from the data collection named for each field.

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- **Sample Problem** Design a function that computes the distance of objects in a 3-dimensional space to the origin.
- The first step to our design process has a bigger change than was typical for previous revisions to our design process.

 (1) When a problem calls for the representation of pieces of information that belong together or describe a natural whole, you need a structure type definition. It requires as many fields as there are relevant properties. An instance of this structure type corresponds to the whole, and the values in the fields correspond to its attributes.

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- (1) In the end, we (and others) must be able to use the data definition to create sample structure instances. Otherwise, something is wrong with our data definition. To ensure that we can create instances, our data definitions should come with data examples.

Here's how we apply our step 1 changes to our sample problem: (struct r3 [x y z] \#:transparent); ; An R3 is a structure: ; (r3 Number Number)

```
(define ex1 (r3 1 2 13))
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- (2) Consider:

```
;;R3 -> Number
;;Computes the distance of a R3 triple to the origin
(define (r3-distance triple) 0)
```

• (3) Your functional examples should use the examples generated by step (1):

```
(check-expect (r3-distance ex1) (sqrt 174))
(check-expect (r3-distance ex2) (sqrt 10))
```

We've talked about changes to taking inventory previously:

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- (5) Use the selector expressions from the template when you define the function. Delete unneeded selections.
- (6) As usual, except that your tests are based on examples from step (1) and other examples.

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- Do you find that plausible? What are the implications?

## Language and Cognition

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- The strong version says that language determines thought and that linguistic categories limit and determine cognitive categories.
- The weak version says that linguistic categories and usage only influence thought and decisions.
- Do you find that plausible? What are the implications?
- The strong version has been rejected, but the weaker version has been supported. Could something similar apply to the first programming language you learn?

Noam Chomsky was a big opponent of linguistic determinism in the strong sense.

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- Our first programming language affects the questions about the ones we learn later.
- We tend to back translate new concepts into languages that we are familiar with.
- Now returning to Racket, let's think about how we design big-bang applications around structs.

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- Let's come up with a data interpretation. Let's assume the ufo descends 2-dimensionally with random jumps to the left or right.

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- What would we need if we wanted to store lots of enemies and not just a single UFO?
- Also think about how we would randomly add enemies to a scene in a Geometry Wars style game.
- Continue to think about other kinds of games and what kind of state we need.

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- Let's imagine a struct for the state of a text editor that looks like the following:
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- If we press space it updates to the following:
- hello world
- Here is a start to our data interpretation:

```
(struct editor [pre post] \#:transparent)
; An Editor is a structure:
; (editor String String)
; interpretation (editor s t) describes an editor
; whose visible text is (string-append s t) with
; the cursor displayed between s and t
```

So, we are creating an editor structure that tracks the text before and after the cursor.

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- 3. Need to isolate a specific character form a string? Thou shalt use string-ref to extract this character?

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Consider representing our Editor like this:

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- How does this relate to Sapir-Whorf?

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 Consider a space invaders like game where we can shoot a single bullet at a time at a single enemy.



• At its most complicated our program only has three objects.

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- At its most complicated our program only has three objects.
- It has the player vehicle, an enemy space ship, and a bullet.
- When the player doesn't shoot there are only two things, the enemy ship and the player ship.
- We can consider the state where a bullet destroys the enemy ship as the exit condition and so our world state doesn't have to worry about a state where the enemy ship isn't present.

Our state of the world became an itemization of different structs. As we consider adding a score, levels, etc. how must our data definition for the state of the world change?

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- What happens to our code?
- We have to cond on different massive structs.
- Is having multiple god-structs really a nice way to structure our code?
- Even in object oriented programs, it's typical to have god objects and it's better than having hidden state across many different modules.
- Just ensure that the god object defers functionality to other modules. And similarly, our god struct is separated from event handling code

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 Now we need to design structs representing our individual objects and also the alternatives of our itemization. Let's start with aim and fired

The data interpretation for aim is a bit simple, after looking at our itemization.

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    Similarly, for our design for our fired state

  (struct fired [ufo tank missile] #:transparent)
  ; fired is a structure
  : (aim UFO Tank Missile)
  ; Interpretation, UFO represents the UFO position
  ; Tank represents the Tank position.
  ; Missile represents the missile's position
  (define ex2 (fired ...))
```

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Is there something a bit strange about our fired state?

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- Think about how this affects our on-tick event.
- Let's return to designing the rest of our structs for data representation.

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• Is our player vehicle (tank) more complicated than the UFO, what does its data description look like?

```
(struct tank [loc vel] #:transparent)
; A Tank is a structure:
; (tank Number Number).
; interpretation (tank x dx) specifies the position:
; (x, HEIGHT) and the tank's speed; dx pixels/tick
```

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- Our tank had a movement vector (which can contain negative values) associated with it, while the other two just have a position that is updated.
- This returns to our discussion of making sure our data lives in the correct domain with a clear interpretation.
- Now that we are done with specifying data interpretations, what kind data examples can we give for our SIGS itemization?

(tank 100 3) (point 22 103)))

 We only had two states for our itemization, why did we need three examples?

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- We had an *implicit* state arising from the fired state.
- Can we think of another implicit state?
- Yes, when the bullet goes off screen.
- Certain states shouldn't require new data definitions but must be handled!

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- (1) An itemization of different forms of dataincluding collections of structuresis required when your problem statement distinguishes different kinds of information and when at least some of these pieces of information consist of several different pieces.
- Steps (2) and (3) remain the same.
- (4) Write a cond for dealing with itemization input and have a
  predicate for each data from in the itemization. If one of the
  items in the itemization is a struct, add selectors. With a
  separate data definition, for items, instead of writing selectors,
  write a separate function.

# Design Process 3.0 (cont.)

(5) Fill out the code, starting with the easy parts. If dealing
with an itemization and you're stuck, analyze the functional
examples dealing with this item. If your template calls another
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- (6) Same as before.

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; computes the area of a disk with radius v,
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(define (checked-area-of-disk v)
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The world can get wrong too quickly.

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- To help with this kind of problem, big-bang comes with an optional check-with clause that accepts a predicate for world states.
- If, for example, we chose to represent all world states with Number, we could express this fact easily like this:

```
(define (main s0)
  (big-bang s0 ... [check-with number?] ...))
```

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 Now we can add a check-with clause that uses a predicate to check the output of all world producing events:

Let's talk about equality again...at least in programming languages.

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- We can also define our own equality predicates like light=?
   so that an error is thrown if a string is given that isn't a traffic light.
- Let me design something similar

We have now finished the first part of this book. To summarize what was covered:

 Basic Racket syntax and constructs. Defining functions and using conditional expressions.

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- Designing basic functions in Racket with a design process.
- Using different design processes for scripts and event driven programs.
- Using itemizations to specify different kind of values to handle from other existing data types
- Using structs to group related data and altering our design process for world programs with itemizations and structs