

Homework 1

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1. Assume that a plaintext bit M is given with $\Pr[M = b] = p_b$, where $b \in \{0, 1\}$. Assume that random key K of the one-time pad encryption is chosen by $\Pr[K = 0] = 0.6$ and $\Pr[K = 1] = 0.4$. Consider the one-time pad encryption $C = M \oplus K$.
 - (a) Assume that an adversary A_1 guesses M randomly without even examining the ciphertext C . Show that the success probability of A_1 is exactly 0.5. (20 points)
 - (b) Suggest a good strategy A_2 of guessing M if p_0 and p_1 are known. (20 points)
2. The Euclidean algorithm computes $\gcd(a, b)$.
 - (a) Give the algorithm and show that its computation time is polynomial in the total length m of a and b , where $m = \text{len}(a) + \text{len}(b)$. (10 points)
 - (b) Solve the equation $r \times 128 + s \times 54 = 2$. (10 points)
3. In the SubBytes of AES, $f(x) = x^{-1} \bmod X^8 + X^4 + X^3 + X + 1$. Compute $f(11101001)$. (20 points)
4. Use the Chinese Remainder Theorem to compute $0 \leq x < 352$ for $x \bmod 3 = 1$, $x \bmod 11 = 3$, and $x \bmod 16 = 13$. (20 points)