

Letters after Replacement (hard)

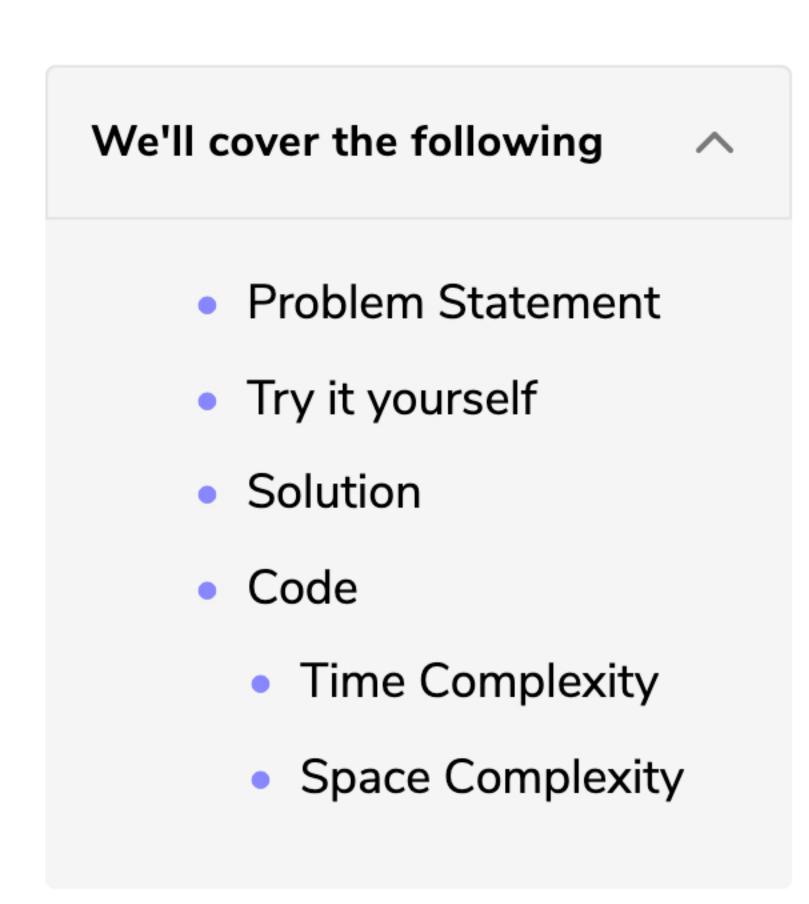
Longest Subarray with Ones after

Replacement (hard)

Problem Challenge 1

educative

Longest Substring with Distinct Characters (hard)



Problem Statement

Given a string, find the length of the longest substring, which has all distinct characters.

Example 1:

```
Input: String="aabccbb"
Output: 3
Explanation: The longest substring with distinct characters is "abc".
```

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? Ask a Question

Example 2:

```
Input: String="abbbb"
Output: 2
Explanation: The longest substring with distinct characters is "ab".
```

Example 3:

```
Input: String="abccde"
Output: 3
Explanation: Longest substrings with distinct characters are "abc" & "cde".
```

Try it yourself

Try solving this question here:

Solution

This problem follows the **Sliding Window** pattern, and we can use a similar dynamic sliding window strategy as discussed in Longest Substring with K Distinct Characters. We can use a **HashMap** to remember the last index of each character we have processed. Whenever we get a duplicate character, we will shrink our sliding window to ensure that we always have distinct characters in the sliding window.

Code

Here is what our algorithm will look like:

```
Python3
                        ⓒ C++
                                     Js JS
👙 Java
    import java.util.*;
     class NoRepeatSubstring {
       public static int findLength(String str) {
        int windowStart = 0, maxLength = 0;
         Map<Character, Integer> charIndexMap = new HashMap<>();
         // try to extend the range [windowStart, windowEnd]
         for (int windowEnd = 0; windowEnd < str.length(); windowEnd++) {</pre>
          char rightChar = str.charAt(windowEnd);
          // if the map already contains the 'rightChar', shrink the window from the beginning so that
          // we have only one occurrence of 'rightChar'
           if (charIndexMap.containsKey(rightChar)) {
 12
             // this is tricky; in the current window, we will not have any 'rightChar' after its previous in
 13
 14
             // and if 'windowStart' is already ahead of the last index of 'rightChar', we'll keep 'windowStar
15
             windowStart = Math.max(windowStart, charIndexMap.get(rightChar) + 1);
 16
17
          charIndexMap.put(rightChar, windowEnd); // insert the 'rightChar' into the map
          maxLength = Math.max(maxLength, windowEnd - windowStart + 1); // remember the maximum length so far
 18
 20
         return maxLength;
21
 22
23
       public static void main(String[] args) {
24
25
        System.out.println("Length of the longest substring: " + NoRepeatSubstring.findLength("aabccbb"));
        System.out.println("Length of the longest substring: " + NoRepeatSubstring.findLength("abbbb"));
26
        System.out.println("Length of the longest substring: " + NoRepeatSubstring.findLength("abccde"));
27
 Run
                                                                                                  Reset
                                                                                         Save
```

Time Complexity

The above algorithm's time complexity will be O(N), where 'N' is the number of characters in the input string.

Space Complexity

The algorithm's space complexity will be O(K), where K is the number of distinct characters in the input string. This also means K <= N, because in the worst case, the whole string might not have any duplicate character, so the entire string will be added to the **HashMap**. Having said that, since we can expect a fixed set of characters in the input string (e.g., 26 for English letters), we can say that the algorithm runs in fixed space O(1); in this case, we can use a fixed-size array instead of the **HashMap**.

