

Linear Algebra and Matrix Theory  
MATH 3100

Marquette University

## **Google's Page Rank Algorithm**

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The Google PageRank algorithm is based on the following premise: a website is important if other important websites link to it. To create this thesis Sergey Brin and Larry Page, who later founded Google, made two assumptions. First, a hyperlink is a recommendation of a website. Second, a recommendation is more valuable if the recommender is selective. It follows that the more important a page  $P_i$ , the higher its PageRank,  $r(P_i)$ , which is defined as:

$$(1) \quad r(P_i) = \sum_{P_j \in B_{P_i}} \frac{r(P_j)}{|P_j|}$$

where  $P_i$  represents page  $i = 1, 2, \dots, n$ ,  $B_{P_i}$  is the set of pages that link to  $P_i$ , and  $|P_j|$  is the number of sites that  $P_j$  points into. These ranks fill the columns of the PageRank vector, defined as  $\pi^T = (r(P_1), r(P_2), \dots, r(P_n))$ , where  $n$  is the number of pages being ranked. To find the successive rank of a page, the calculations begin with an estimate of each page,  $r_o(P_i)$ , and the PageRank is updated iteratively with the following:

$$(2) \quad r_{k+1}(P_i) = \sum_{P_j \in B_{P_i}} \frac{r_k(P_j)}{|P_j|}$$

The hyperlink matrix  $\mathbf{H}$  has elements  $h_{ij} = \frac{1}{|P_i|}$  for each page  $P_j$  which is linked to by page  $P_i$  and  $h_{ij} = 0$  otherwise. If each row in  $\mathbf{H}$  sums to either 1 or 0, that matrix is called *substochastic*. This case results in page sinks, if a page has no link to other pages (when a row sums to 0). To avoid PageRank sink due to a row of zeros, replace each element in that row with  $1/n$  where  $n$  is the number of pages in the web.

With what is defined so far, there still exists a possibility for cycling between states. To avoid this, Google founders introduced the idea of the *random surfer*. According to this concept, there is some probability  $\alpha$  ( $0 \leq \alpha \leq 1$ ) that the web surfer moves following a hyperlink. There is also the probability that a web surfer moves by directly linking to one of the other  $n$  pages directly, represented by  $1-\alpha$ . With this concept and the hyperlink matrix  $\mathbf{H}$ , the Google PageRank update matrix was formulated.

$$(3) \quad \mathbf{G} = \alpha \mathbf{H} + \alpha \mathbf{a} \mathbf{1}^T + \frac{1-\alpha}{n} \mathbf{1} \mathbf{1}^T$$

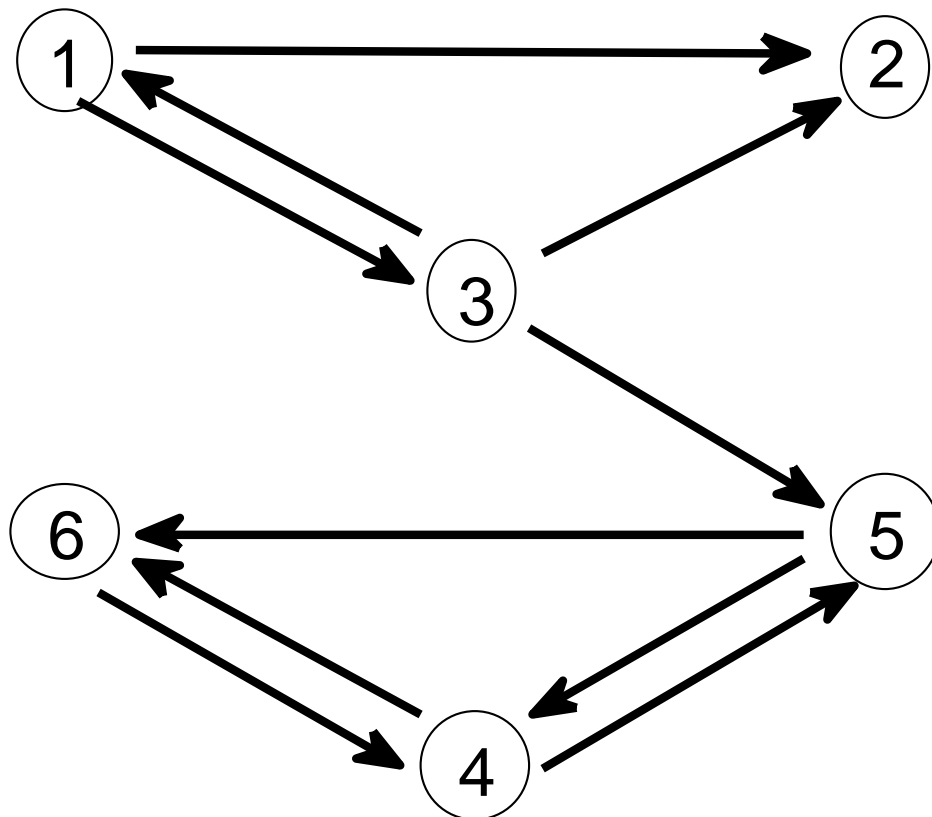
where  $\mathbf{a}$  is a column vector with elements  $a_i = \frac{1}{n}$  if  $P_i$  links to no pages or 0 otherwise and  $\mathbf{1}$  is a column vector of ones. Utilizing the PageRank update matrix, the PageRank vector can be iteratively updated as follows.

$$(4) \quad \pi_{k+1}^T = \pi_k^T \mathbf{G}$$

Based on a given mini-web, different values for  $\alpha$ , and different values for  $k$ , we intend to determine what  $\alpha$  value Google should use. Using MatLab we will be able to easily compute the matrices and vectors necessary to determine an appropriate  $\alpha$  value.

Dr. Spiller has provided a mini-web consisting of six pages with links as shown below.

## Six page mini-web



Given this web, we start by determining the  $B_{P_i}$  and  $|P_j|$ , as follows.

Page	1	2	3	4	5	6
$B_{P_i}$	3	1, 3	1	5, 6	3, 4	4, 5
$ P_j $	2	0	3	2	2	1

From our values for  $B_{P_i}$  and  $|P_j|$ , we are able to calculate  $\mathbf{H}$ ,  $\mathbf{a}$ ,  $\mathbf{1}$ , and  $\mathbf{1}^T$ , where  $n = 6$ .

$$\mathbf{H} = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 1/3 & 1/3 & 0 & 0 & 1/3 & 0 \\ 0 & 0 & 0 & 0 & 1/2 & 1/2 \\ 0 & 0 & 0 & 1/2 & 1/2 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{bmatrix} \quad \mathbf{a} = \begin{bmatrix} 0 \\ 1/6 \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix} \quad \mathbf{1} = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{bmatrix} \quad \mathbf{1}^T = [1 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1]$$

Plugging the matrix and vectors determined above into equation 3, the transition matrix  $\mathbf{G}$  can be symbolically determined in terms of  $\alpha$ .

$$\mathbf{G} = \begin{bmatrix} 1/6 - \alpha/6 & \alpha/3 + 1/6 & \alpha/3 + 1/6 & 1/6 - \alpha/6 & 1/6 - \alpha/6 & 1/6 - \alpha/6 \\ 1/6 & 1/6 & 1/6 & 1/6 & 1/6 & 1/6 \\ \alpha/6 + 1/6 & \alpha/6 + 1/6 & 1/6 - \alpha/6 & 1/6 - \alpha/6 & \alpha/6 + 1/6 & 1/6 - \alpha/6 \\ 1/6 - \alpha/6 & 1/6 - \alpha/6 & 1/6 - \alpha/6 & 1/6 - \alpha/6 & \alpha/3 + 1/6 & \alpha/3 + 1/6 \\ 1/6 - \alpha/6 & 1/6 - \alpha/6 & 1/6 - \alpha/6 & \alpha/3 + 1/6 & 1/6 - \alpha/6 & \alpha/3 + 1/6 \\ 1/6 - \alpha/6 & 1/6 - \alpha/6 & 1/6 - \alpha/6 & 5\alpha/6 + 1/6 & 1/6 - \alpha/6 & 1/6 - \alpha/6 \end{bmatrix}$$

As  $\mathbf{G}$  has been determined symbolically in terms of  $\alpha$ , we can plug in values for  $\alpha$ . With each increasing iteration for  $\pi_k$ , the values in the vector  $\pi_k$  converge to  $\pi$ . For  $\alpha = 0.9$ , a minimum of 15 iterations are required to obtain the steady state. The resulting PageRank vector after 15 iterations is

$$\pi_{15} = \pi = \begin{bmatrix} 0.0372 \\ 0.0540 \\ 0.0415 \\ 0.3750 \\ 0.2060 \\ 0.2862 \end{bmatrix}$$

$$\mathbf{D} = \begin{bmatrix} 1.0000 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0.6101 & 0 & 0 & 0 & 0 \\ 0 & 0 & -0.4500 & 0 & 0 & 0 \\ 0 & 0 & 0 & -0.4500 & 0 & 0 \\ 0 & 0 & 0 & 0 & -0.3705 & 0 \\ 0 & 0 & 0 & 1 & 0 & -0.0896 \end{bmatrix}$$

$$\mathbf{V} = \begin{bmatrix} 0.0715 & 0.2931 & 0.0000 & 0.0000 & -0.0647 & -0.2123 \\ 0.1036 & 0.5094 & 0.0000 & 0.0000 & 0.0139 & 0.8541 \\ 0.0797 & 0.3415 & 0.0000 & 0.0000 & 0.0730 & -0.3636 \\ 0.7204 & -0.5891 & -0.7071 & 0.7071 & -0.6606 & 0.0184 \\ 0.3957 & -0.1414 & 0.7071 & -0.7071 & 0.7376 & -0.3047 \\ 0.5498 & -0.4135 & 0.0000 & 0.0000 & -0.0992 & 0.0082 \end{bmatrix}$$

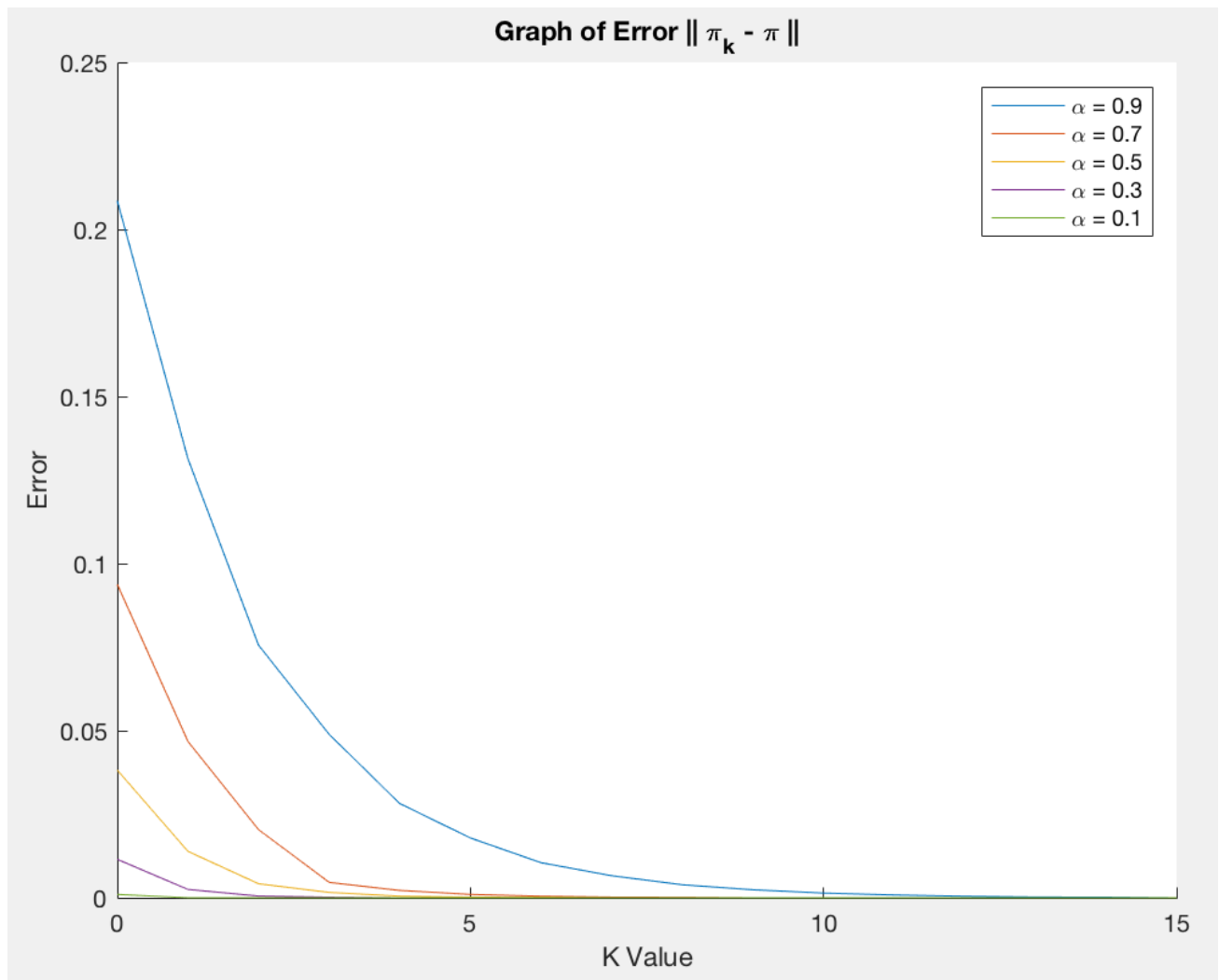
$\mathbf{D}$  shows the eigenvalues and  $\mathbf{V}$  shows the corresponding eigenvectors. The PageRank vector is updated iteratively until it converges. The largest eigenvalue and corresponding eigenvector,  $\lambda$  and  $\mathbf{v}_{max}$ , are important for the update relationship of the Google Matrix,  $\mathbf{G}$ . Because this is designed as a stochastic matrix, as  $K$  goes to infinity, the steady state is obtained with the  $\pi$  vector. The  $\lambda$  values are raised to the power of  $K$  and multiplied by the corresponding coefficient and eigenvector, so that the page vector does not become infinitely large, as would happen if this were a *substochastic* matrix design. The eigenvalue and eigenvector  $\lambda$  and  $\mathbf{v}_{max}$  are the last iteration values that contribute to the Page Rank vector, and are the last changes that cause the matrix to converge. The initial vector  $\pi_0$  is an arbitrary vector in  $\mathbf{R}^6$ , and that can be written as a linear combination of eigenvectors of  $\mathbf{G}$ . This is due to the fact that the eigenvectors of  $\mathbf{G}$  form a basis of  $\mathbf{R}^6$ . While iterating through the  $\pi_{k+1}$

Based on the vector,  $\pi_{15}$ , the ranks of the websites would be ordered  $P_4, P_6, P_5, P_2, P_3, P_1$ .

In addition, we can calculate  $\pi_k$  for  $\alpha = 0.7, 0.5, 0.3$ , and  $0.1$ . For each value of  $\alpha$ , there is a different number of iterations needed in order for the values of  $\pi_k$  to converge to  $\pi$ . The table below shows the number of iterations necessary and the value of  $\pi$  for each value of  $\alpha$ .

$\alpha$	0.9	0.7	0.5	0.3	0.1
$k$	15	10	7	4	1
$\pi$	[0.0372]	[0.0852]	[0.1162]	[0.1392]	[0.1581]
	0.0540	0.1150	0.1452	0.1601	0.1661
	0.0415	0.0932	0.1245	0.1456	0.1607
	0.3750	0.2899	0.2390	0.2044	0.1781
	0.2060	0.1866	0.1759	0.1699	0.1670
	[0.2862]	[0.2302]	[0.1992]	[0.1808]	[0.1700]

Using this information, we obtain the error that occurs as the iteration increases for each value of  $\alpha$  shown in the following graph.



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As shown, while the value of  $\alpha$  decreases, the number of iterations needed for  $\boldsymbol{\pi}_k$  to converge also decreases. Based on the gathered information the most appropriate value that Google should use is 0.1 because it requires only 1 iteration for it to be the most accurate.