

专题14-串口驱动程序设计

一、tty驱动架构

1.1、TTY概念解析

1.1.1、/dev/ttySAC0

在Linux系统中，终端是一类字符型设备，它包括多种类型，通常使用tty来简称各种类型的终端设备。

• 串口终端 (/dev/ttyS*)

串口终端是使用计算机串口连接的终端设备。Linux 把每个串行端口都看作是一个字符设备。这些串行端口所对应的设备名称是 /dev/ttySAC0; /dev/ttySAC1.....

1.1.2、/dev/tty1-n

虚拟终端 (/dev/tty*)

当用户登录时，使用的是虚拟终端。使用 Ctl+Alt+[F1—F6]组合键时，我们就可以切换到tty1、tty2、tty3等上面去。tty1-tty6等称为虚拟终端，而tty0则是当前所使用虚拟终端的一个别名。

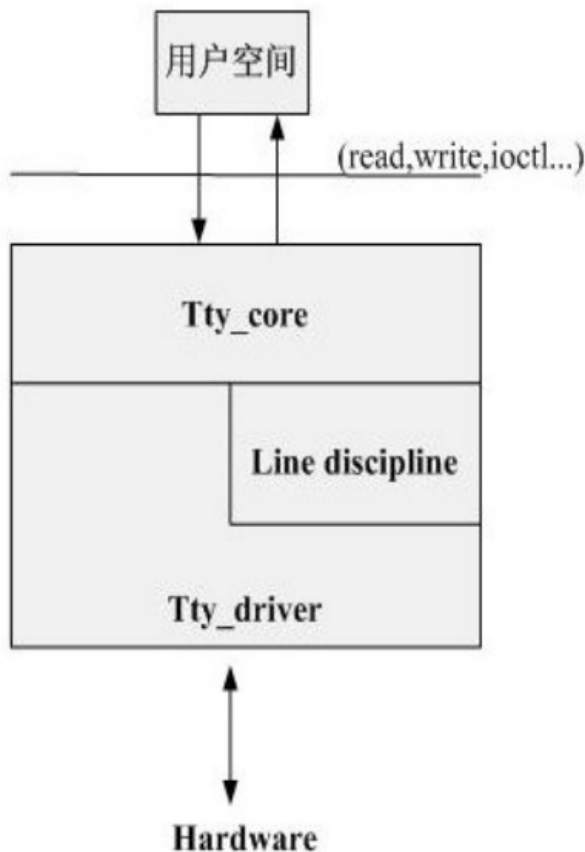
1.1.3、/dev/console

控制台终端 (/dev/console)

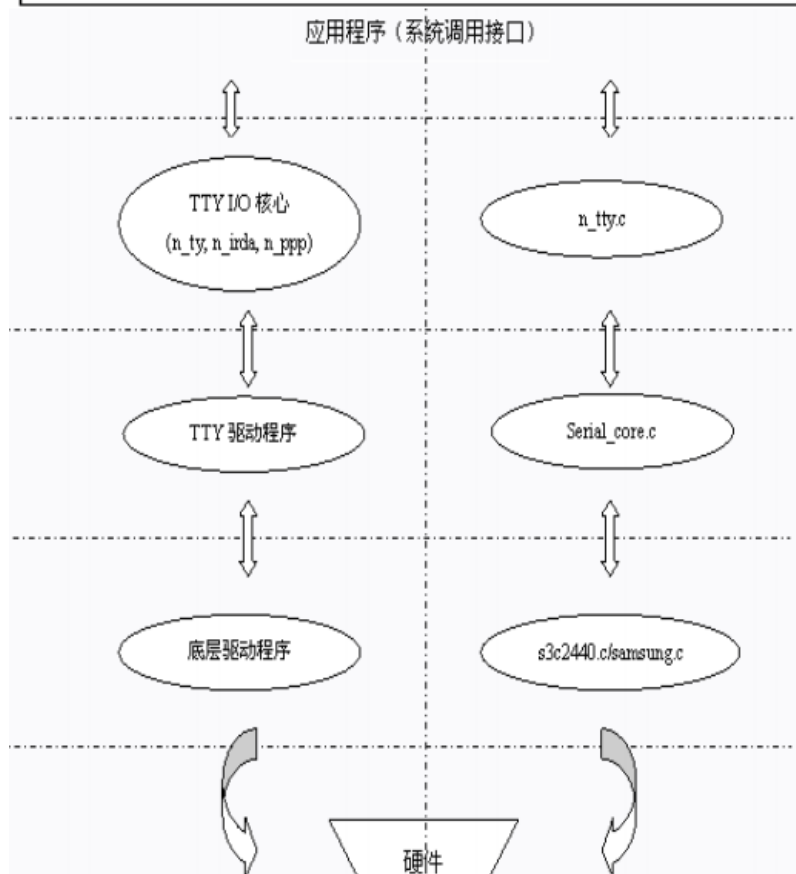
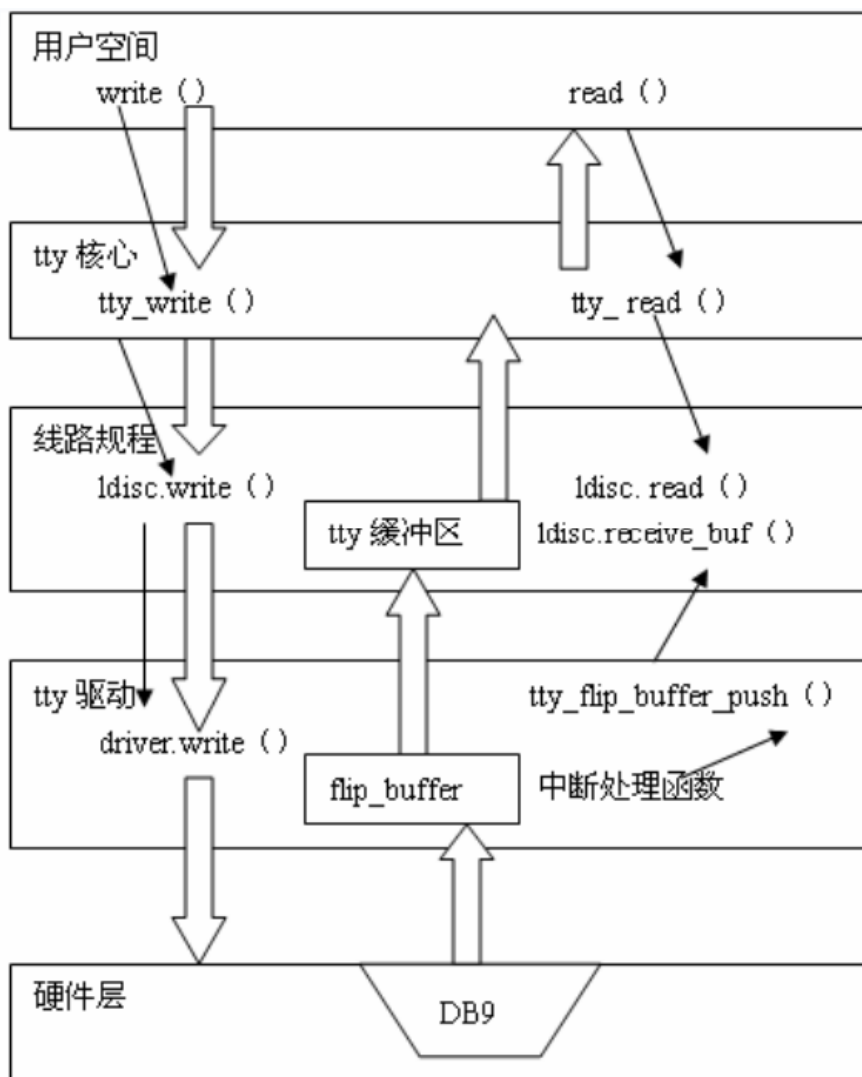
在Linux系统中，计算机的输出设备通常被称为控制台终端(Console),这里特指printf信息输出到的设备。

/dev/console是一个虚拟的设备，它需要映射到真正的tty上，比如通过内核启动参数“ console=ttySAC0” 就把console映射到了串口0

1.2、TTY结构分析



Linux tty子系统包含：tty核心，tty线路规程和tty驱动。tty核心是对整个tty设备的抽象，对用户提供统一的接口，tty线路规程是对传输数据的格式化，tty驱动则是面向tty设备的硬件驱动。



1.3、回溯串口发送流程

`dump_stack()`;用于显示调用信息。

```

[<0049048>] (dump_backtrace+0x0/0x10c) from [<037429c>] (dump_stack+0x18/0x1c)
r7:c3942002 r6:00000000 r5:c38bda00 r4:c04caf20
[<0374284>] (dump_stack+0x0/0x1c) from [<018ea60>] (s3c24xx_serial_start_tx+0x14/0x64)
[<018ea4c>] (s3c24xx_serial_start_tx+0x0/0x64) from [<018ad78>] (uart_start+0x68/0x6c)
r5:c38bda00 r4:60000013
[<018ad10>] (uart_start+0x0/0x6c) from [<018c5f8>] (uart_write+0xc0/0xe4)
r5:c38bda00 r4:00000000
[<018c538>] (uart_write+0x0/0xe4) from [<0178034>] (n_tty_write+0x1d8/0x448)
[<0177e5c>] (n_tty_write+0x0/0x448) from [<0175790>] (tty_write+0x14c/0x244)
[<0175644>] (tty_write+0x0/0x244) from [<0175910>] (redirected_tty_write+0x88/0x98)
[<0175888>] (redirected_tty_write+0x0/0x98) from [<00a6e68>] (vfs_write+0xb4/0xe8)
r9:c397e000 r8:c0045008 r7:00000004 r6:c397ff78 r5:40000000
r4:c3953700
[<00a6db4>] (vfs_write+0x0/0xe8) from [<00a6f80>] (sys_write+0x4c/0x84)
r7:00000004 r6:c3953700 r5:00000000 r4:00000000
[<00a6f34>] (sys_write+0x0/0x84) from [<0044e60>] (ret_fast_syscall+0x0/0x2c)
r6:001d27f8 r5:40000000 r4:00000002

```

tty核心：

```

ssize_t redirected_tty_write(struct file *file, const char __user *buf,
    size_t count, loff_t *ppos)

```

```

{
    struct file *p = NULL;

    spin_lock(&redirect_lock);
    if (redirect) {
        get_file(redirect);
        p = redirect;
    }
    spin_unlock(&redirect_lock);

    if (p) {
        ssize_t res;
        res = vfs_write(p, buf, count, &p->f_pos);
        fput(p);
        return res;
    }
    return tty_write(file, buf, count, ppos);
}

```

```

static ssize_t tty_write(struct file *file, const char __user *buf,
    size_t count, loff_t *ppos)

```

```

{
    struct inode *inode = file->f_path.dentry->d_inode;
    struct tty_struct *tty = file_tty(file);
    struct tty_ldisc *ld;
    ssize_t ret;

    if (tty_paranoia_check(tty, inode, "tty_write"))
        return -EIO;
    if (!tty || !tty->ops->write ||
        (test_bit(TTY_IO_ERROR, &tty->flags)))
        return -EIO;
    /* Short term debug to catch buggy drivers */
    if (tty->ops->write_room == NULL)
        printk(KERN_ERR "tty driver %s lacks a write_room method.\n",
            tty->driver->name);
    ld = tty_ldisc_ref_wait(tty);
    if (!ld->ops->write)
        ret = -EIO;
    else
        ret = do_tty_write(ld->ops->write, tty, file, buf, count);
    tty_ldisc_deref(ld);
    return ret;
}

```

线路规程：

```

static ssize_t n_tty_write(struct tty_struct *tty, struct file *file,
    const unsigned char *buf, size_t nr)

```

```

{
const unsigned char *b = buf;
DECLARE_WAITQUEUE(wait, current);
int c;
ssize_t retval = 0;

/* Job control check -- must be done at start (POSIX.1 7.1.1.4). */
if (L_TOSTOP(tty) && file->f_op->write != redirected_tty_write) {
    retval = tty_check_change(tty);
    if (retval)
        return retval;
}

/* Write out any echoed characters that are still pending */
process_echoes(tty);

add_wait_queue(&tty->write_wait, &wait);
while (1) {
    set_current_state(TASK_INTERRUPTIBLE);
    if (signal_pending(current)) {
        retval = -ERESTARTSYS;
        break;
    }
    if (tty_hung_up_p(file) || (tty->link && !tty->link->count)) {
        retval = -EIO;
        break;
    }
    if (O_OPOST(tty) && !(test_bit(TTY_HW_COOK_OUT, &tty->flags))) {
        while (nr > 0) {
            ssize_t num = process_output_block(tty, b, nr);
            if (num < 0) {
                if (num == -EAGAIN)
                    break;
                retval = num;
                goto break_out;
            }
            b += num;
            nr -= num;
            if (nr == 0)
                break;
            c = *b;
            if (process_output(c, tty) < 0)
                break;
            b++; nr--;
        }
        if (tty->ops->flush_chars)
            tty->ops->flush_chars(tty);
    } else {
        while (nr > 0) {
            c = tty->ops->write(tty, b, nr);
            if (c < 0) {
                retval = c;
                goto break_out;
            }
            if (!c)
                break;
            b += c;
            nr -= c;
        }
    }
    if (!nr)
        break;
    if (file->f_flags & O_NONBLOCK) {
        retval = -EAGAIN;
        break;
    }
    schedule();
}
}

```

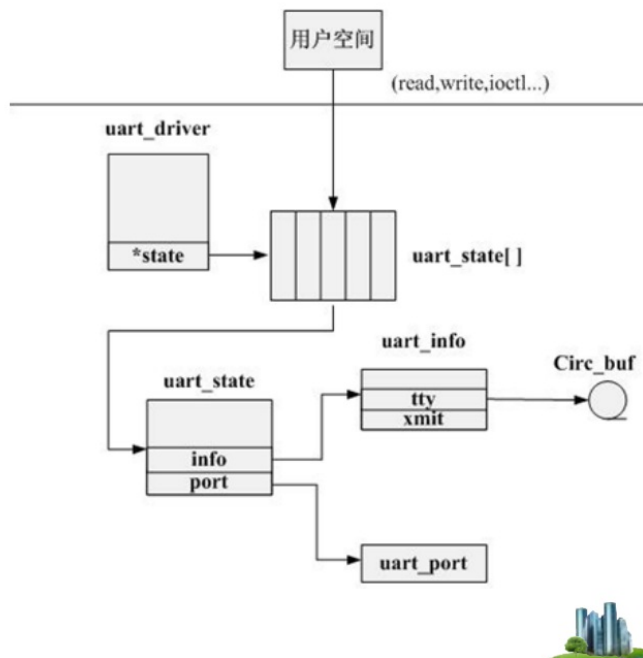
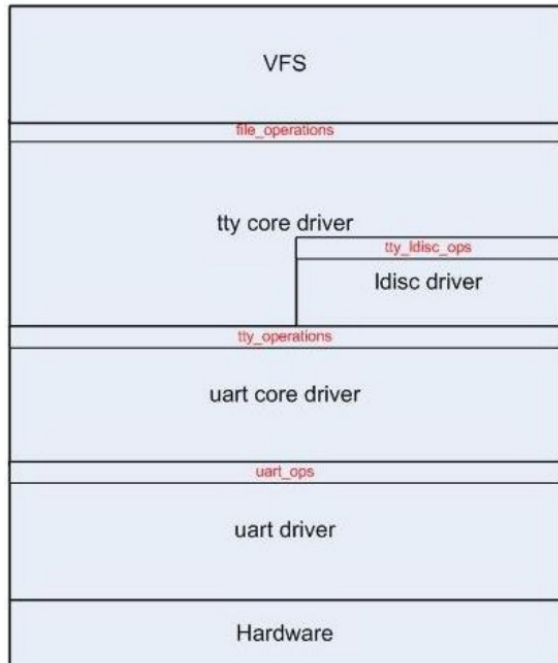
```

break_out;
__set_current_state(TASK_RUNNING);
remove_wait_queue(&tty->write_wait, &wait);
if (b - buf != nr && tty->fasync)
    set_bit(TTY_DO_WRITE_WAKEUP, &tty->flags);
return (b - buf) ? b - buf : retval;
}

```

二、串口驱动分析-初始化

2.1、串口驱动程序结构



2.2、串口驱动中的重要数据结构

- UART驱动程序结构: `struct uart_driver`
- UART端口结构: `struct uart_port`
- UART相关操作函数结构: `struct uart_ops`
- UART状态结构: `struct uart_state`
- UART信息结构: `struct uart_info`

2.3、初始化分析

samsung.c是2440,S3C6410,210通用的文件，s3c6400.c是s3c6410处理器的文件，这两个文件合起来就是串口文件。

```

static int uart_write(struct tty_struct *tty,
    const unsigned char *buf, int count)
{
    struct uart_state *state = tty->driver_data;
    struct uart_port *port;
    struct circ_buf *circ;
    unsigned long flags;
    int c, ret = 0;

    /*
     * This means you called this function _after_ the port was
     * closed. No cookie for you.
     */
    if (!state) {
        WARN_ON(1);
        return -EL3HLT;
    }

    port = state->uart_port;
    circ = &state->xmit;

    if (!circ->buf)
        return 0;

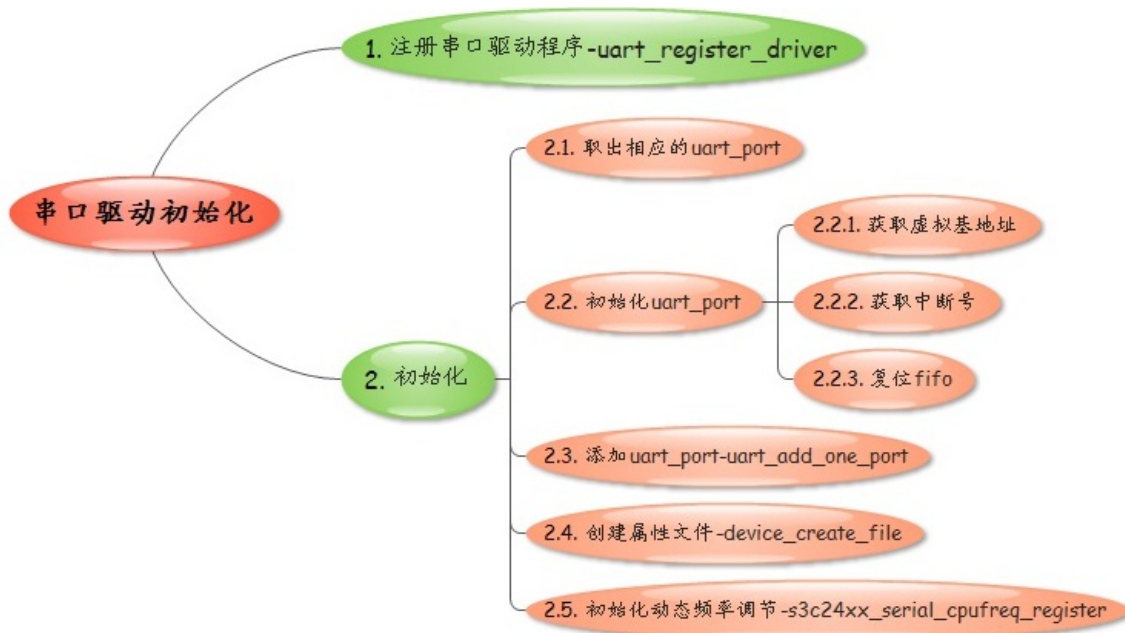
```

```

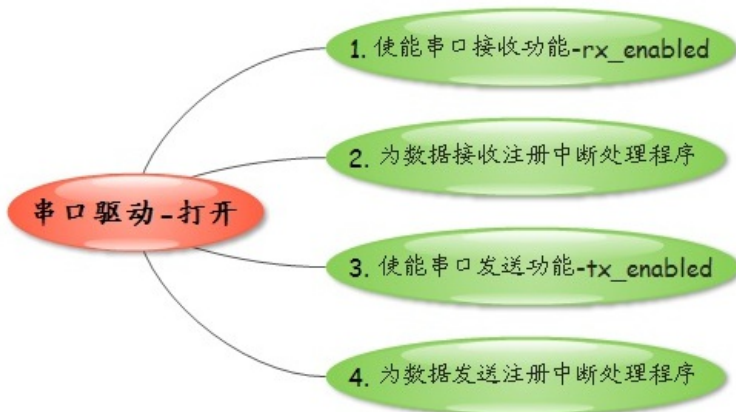
spin_lock_irqsave(&port->lock, flags);
while (1) {
    c = CIRC_SPACE_TO_END(circ->head, circ->tail, UART_XMIT_SIZE);
    if (count < c)
        c = count;
    if (c <= 0)
        break;
    memcpy(circ->buf + circ->head, buf, c);
    circ->head = (circ->head + c) & (UART_XMIT_SIZE - 1);
    buf += c;
    count -= c;
    ret += c;
}
spin_unlock_irqrestore(&port->lock, flags);

uart_start(tty);
return ret;
}

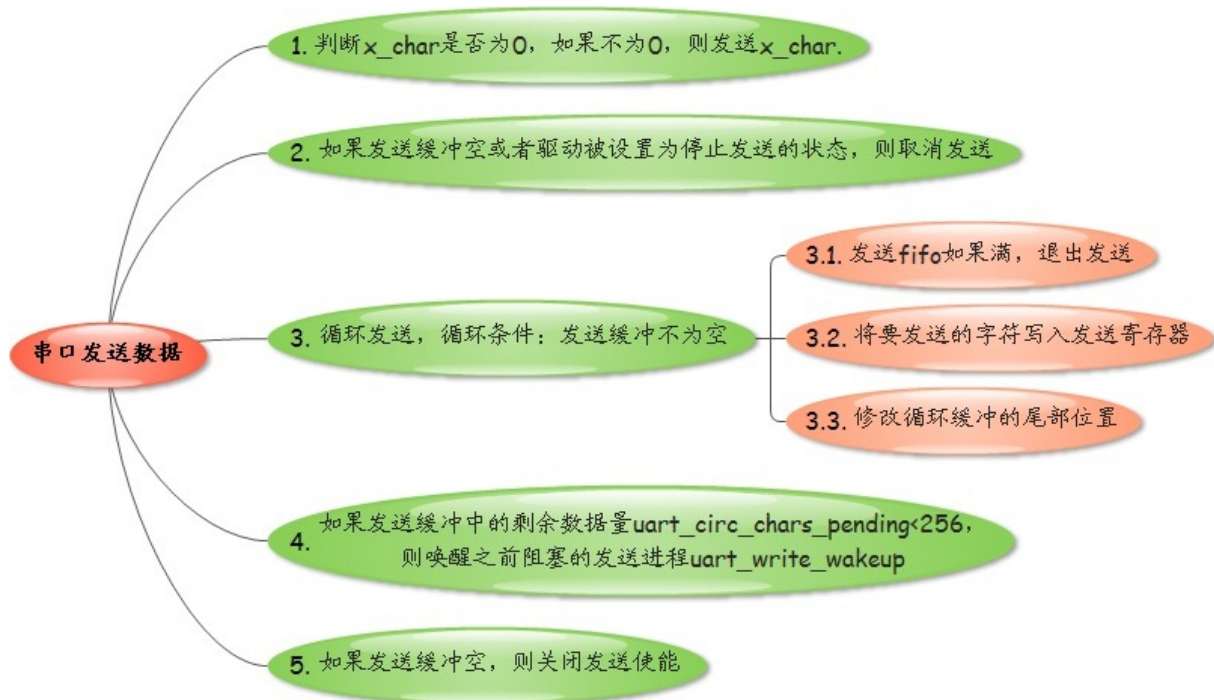
```



三、串口驱动分析-打开设备



四、串口驱动分析-数据发送



4.1、数据传递-循环缓冲

五、串口驱动分析-数据接收



5.1、流控

硬流控的RTS、CTS：

(现在做串口使用RTS/CTS必看内容，因为MTK/)

RTS (Require To Send, 发送请求) 为输出信号，用于指示本设备准备好可接收数据，低电平有效，低电平说明本设备可以接收数据。
CTS (Clear To Send, 发送允许) 为输入信号，用于判断是否可以向对方发送数据，低电平有效，低电平说明本设备可以向对方发送数据。

六、串口驱动编程实现

6.1、分析架构-子系统

6.2、分析驱动-流程图

6.3、实现代码