Instead of grouping by identity of first addend, let's group by identity of all addends: {1,1,3} = {1,2,2} = 1 partitions of 5. 1+1+3 1+2+2 1+3+1 2+1+2 3+1+1. 2+2+1 Vet me still don't see a convenient way to count compositions. The trick is to break everything into 1: +9 1+1+3 = <1,1,1+1+1> 1+2+2= <1,1+1,1+1> Now of we squint at the nightcompositions of 5 wito 3 parts are obtained by filling 4 slots w/ 2 commons (putting 't's in the remainder): 1+1+3= 1.2121+1+1 1-2-12 121±121±1

Thus # compositions is the # ways commas it from a total of 4 stats You'll notice that this is essentially the "stars and bars" approach where the stars are "+" and the bars are Since the # slots is I less than 5 and # "bars" is 1 688 than # terms in composition, we have the Number of compositions of n into R parts is  $\binom{n-1}{k-1}$ . What about the number of compositions of n regardless of the # parts? Use the stars and board approach! Every composition takes the form n

where there are n-1 slots each of which is filled by either a "+" or a comma. - lead assignment of +s and commas - each composition determines a unique assignment of +5 and comma. Since there are n-1 vidependent binary choices to be made, there are 2<sup>n1</sup> compositions. Thus it must be the case that  $\sum_{k=1}^{\infty} \binom{n-1}{k-1} = 2^{n-1}$  $LHS = \frac{2^{-1}(n-1)}{m=0} = 2^{n-1}$ because: 2 (cf L3)

What if we now permit the integers composition to be 0? How there of n? n=7 R=5: Again breaking everything into 15, one encoding is: \*+\*\* Again, there are 4+0+1+2+0 = bal of 7+5-1 symbols: \* \* \* \* 1 | 1 | \* 1 | \* \* | Thus all we need to do is to bars into 7+5-1 slots, which is 7+5-1 in general, there are

weak compositions (those in which o is allowed) into k parts is an integer n

925 LVP Q. Probability that a word of length k has from an alphabet of size n. Assume all words are equally likely. A. How many words are there in total? #choices: nnn R times. => Hwonds = n How many of these have no repeated letter? #chuico: n n-1 n-2 n-(k-1) =) # words w/ no repeats = n(n-1) ... (n-(k-1) Since all words are equally likely, the prob. that a randonly c

word has no repeated letter is: # (words W/ no repeats) a What is the smallest word 3.t. words are more contain a respect the P(no repeat) 1 25 .89.79.66.53.41 P(repeat) P(nis repeat)

Birthday Paradox thow many people need to be in room before there is more than a 50% chance of a shared birthday? A To establish a connection with the previous counting problem, imagine asking for birthdays, one by one Jan 17 Feb 9 June 4

1st person 2rd person 46th person This is a word of length k (we don't specify the value of k yot) chosen from an alphabet of size n = 365. Thus, from the provious problem, P(k people don't share a birthday)  $=\frac{365!}{(365-k)!}$   $\frac{1}{365}$ For k=23, this is \$ 49%.

This problem is a "paradox" because the ## 00 people needed is surprisingly small compared to the # required to growantie a Shared birthday - 365 (by Pigeonhule principle) Now, how do things change of birthdays are not equally likely leg birthdates of professional hockey dayers). Turns out that this just increases the chance of birthday collisions. To see this, we need a new (probabilistic) approach: a Too people. Two possible birthdays
(A,B). What is the pro. they share A. Suppose P(A) = prob 4 being born on A P(B) = 1-p. birthdays: AA AB BA BB P(Shared b'day) = p2 + (1-p)2

Q. For what P is P(shoured b'day) smallest? dp P(shand) = 0. =) 2p+2(1-p)(-1)=0.  $\Rightarrow$  p=1/2Thus, when birthdays are not egnally likely las me initially assumed then prob. Ob alisions only become larger A conseguence is that really good bet that two people trave the semme birthday if there are only two possibilities for those pirthdays. Pleased 325 of LVP to see how a professor used this fact to make honey, over the course of his

-12-		
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		+
Summer 2 12 17		
Summary of L2-L5:		
Four kinds of counting.		
	Segnences (order matters)	SeAs
1	(order matters)	(order doesn't matter)
repetition allowed	n R	
no repetition	nl.	(n)
	(n-k)!	(k)
		*