

## **A GPU-based Survey for Millisecond Radio Transients Using ARTEMIS**

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**Abstract.** Astrophysical radio transients are excellent probes of extreme physical processes originating from compact sources within our Galaxy and beyond. Radio frequency signals emitted from these objects provide a means to study the intervening medium through which they travel. Next generation radio telescopes are designed to explore the vast unexplored parameter space of high time resolution astronomy, but require High Performance Computing (HPC) solutions to process the enormous volumes of data that are produced by these telescopes. We have developed a combined software/hardware solution (code named ARTEMIS) for real-time searches for millisecond radio transients, which uses GPU technology to remove interstellar dispersion and detect millisecond radio bursts from astronomical sources in real-time. Here we present an introduction to ARTEMIS. We give a brief overview of the software pipeline, then focus specifically on the intricacies of performing incoherent de-dispersion. We present results from two brute-force algorithms. The first is a GPU based algorithm, designed to exploit the L1 cache of the NVIDIA Fermi GPU. Our second algorithm is CPU based and exploits the new AVX units in Intel Sandy Bridge CPUs.

### **1. Introduction**

**ARTEMIS** stands for *Advanced Radio Transient Event Monitor and Identification System*. It is a project being carried out at Oxford (Astrophysics, OeRC, Engineering Science) aimed at the real-time processing of high-time resolution data from radio astronomy (Karastergiou et al. in preparation). The project aims to develop the software and piece together the hardware for surveys of fast transients and pulsars using next generation radio telescopes such as LOFAR and MeerKAT. Real-time processing is essential to ensure that broadband data streams are reduced to manageable rates both for storage and further processing. The ARTEMIS servers perform (in real-time) all the operations necessary to discover short duration radio pulses from pulsars and fast transients (Wayth et al. 2011), thanks to a modular software structure operating in a C++

scalable framework (PELICAN, developed at the OeRC). AMPP (ARTEMIS Modular PELICAN Pipelines) is the software that we have developed for receiving the data, for further channelisation in finer frequency, generation of Stokes parameters, excision of radio frequency interference, integration, real-time dispersion searches and detection of interesting signals across multiple telescopes, in high-throughput CPU and GPU code. This article describes the results of the GPU implementation of the real-time incoherent de-dispersion aspect and gives a comparison to a CPU implementation. Recently incoherent de-dispersion (MDSM). Here we present a new, optimised kernel, which we have used with the MDSM wrapper.

## 2. Searching for Radio Transients

The quadratic cold plasma dispersion law of the interstellar medium results in radio pulses at lower frequencies arriving at Earth later than their high frequency counterparts (Lorimer & Kramer 2005). To take advantage of the fact that astrophysical radio bursts are typically broadband, integration over frequency is essential to increase signal to noise. Incoherent de-dispersion is the process of shifting the (power) data inside each individual frequency channel to counter the effect of interstellar dispersion before frequency integration. Given the quadratic relationship between time delay and frequency, the phenomenon is governed by a single free parameter, known as the dispersion measure (DM), which is the integrated electron number density along the line of sight to the source. Figure 1 shows simulated filterbank data (using the SIGPROC package), with a dispersed radio signal sitting in noise. In a blind search for dispersed radio bursts, the DM is unknown. A large range of DM values is typically searched, by shifting each frequency channel by the appropriate amount of time for each DM being searched. This results in each data point (in the frequency-time domain) being used many times for all the dispersion curves that it contributes to, a useful quality for GPU acceleration.

## 3. Acceleration via GPU Computing

In order to produce a GPU kernel that can achieve a significant proportion of the peak performance of the GPU we need to ensure **three things**. The first is that the accumulator that stores the integrated value of the intensity (along the trial dispersion curve) sits in the fastest area of memory. The second is that the correct data from the  $(f, t)$  domain is always available to the streaming multiprocessors. The third is that the shifting value is calculated using as few operations as possible. The GPU algorithm presented is designed to exploit the new fast L1 cache present on the NVIDIA Fermi hardware. The algorithm is designed to reuse cache-lines that are present in the L1 cache, vastly

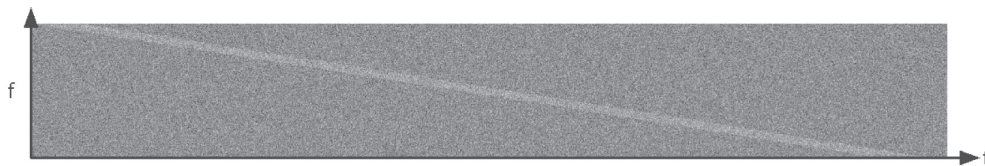


Figure 1. An example of a test data-set (total intensity as a function of  $(f, t)$ ) for 496 channels. A weak signal can be seen moving from top left to bottom right.

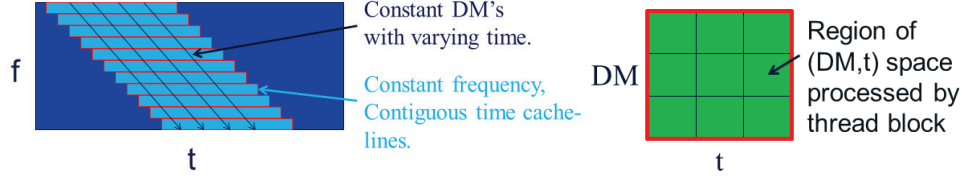


Figure 2. Left: A single thread loads data from cache-lines of constant frequency, contiguous time and increments accumulators for multiple points in  $(DM, t)$  space. Right: Each thread-block processes an area of  $(DM, t)$  space.

reducing the need to transfer the same data from main graphics memory multiple times. This is achieved by each thread processing several time elements for its given value of dispersion, holding these values in local registers (Figure 2, left). This gives rise to each thread-block processing a rectangular area of the dispersion-time  $(DM, t)$  space, ensuring cache-lines of  $(f, t)$  data are reused multiple times (Figure 2, right).

#### 4. Comparisons of GPU and CPU Algorithms

In this section we present results from our GPU kernel and compare these results to a vector-parallel CPU code that exploits the SSE registers on a multiprocessor Intel Xeon machine, or the AVX registers on a new Intel i7 Sandy Bridge based machine (Overclocked from 3.2GHz to 4.2GHz, employing 1600MHz DDR3 SD-RAM). The CPU code has been designed with maximum cache-line usage in mind and use the Intel intrinsics in the vector parts of the code. Results from a vectorised code using the Intel auto-vectorizer have not been presented because they are consistently slower (approximately 3x, in our region of interest) compared with our vectorised code. Figure 3 shows results from the CPU code, demonstrating results that are in exact agreement with our simulated data. In Figure 4 (left) we present the proportion of real-time taken by the CPU/GPU codes (including different platforms) against a varying number of frequency channels. Importantly we hold the maximum dispersion measure at 200. However to ensure that we do not sub-sample the data we set the total number of dispersion mea-

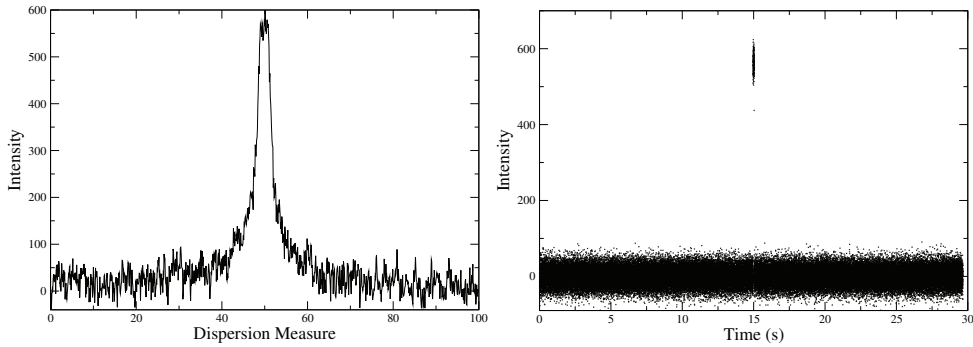


Figure 3. Left: The result of a dispersion search on simulated data with a DM of  $50 \text{ pc cm}^{-3}$ . Right: The square pulse signal recovered at the identified DM has all the characteristics of the simulated signal.

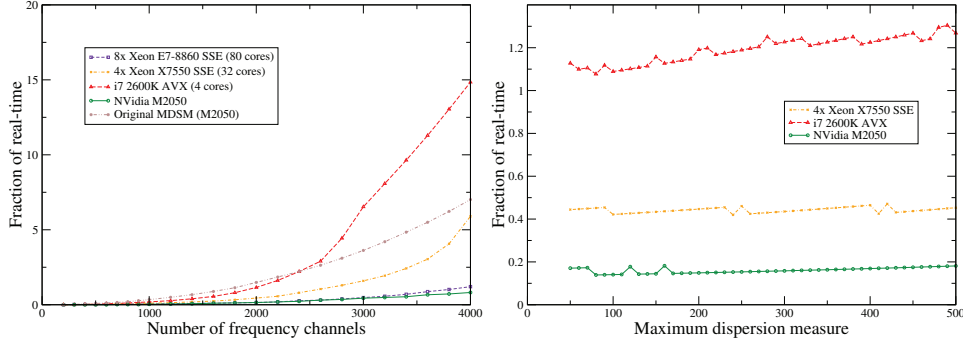


Figure 4. Left: Plot varying the number of channels showing execution time as percentage of real-time. Here we hold the total number of trial dispersion searches equal to the number of channels. Right: Plot varying the maximum dispersion measure, whilst holding the total number of trial dispersions constant (number of channels = number of trial dispersion searches = 2000).

asures equal to the total number of channels. Figure 4 (right) shows how the different codes perform as we increase the maximum value of the dispersion measure with a fixed number of frequency channels. Here we hold the number of dispersion measures equal to the number of channels but we change the value of the equally spaced dispersion curves. In both cases we observe better performance from the GPU code.

## 5. Conclusions and Future Work

With typical parameters of a dispersion search ( $\sim 2000$  frequency channels and dispersion measures to search), we estimate our kernel achieves approximately 40% – 50% of peak GPU performance. This leaves little margin for improvement for GPU based, brute-force, incoherent dedispersion algorithms and makes real-time dispersion searches a possibility under many different observing situations. Our kernel has been tested successfully in a real environment within ARTEMIS. To try and achieve a balance between the CPU and GPU computing powers our future work will focus on implementing vectorisation in the poly-phase filter using AVX registers and the Intel intrinsics.

## References

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