

Integer Linear Programming and its Application in Carleton's CS Match

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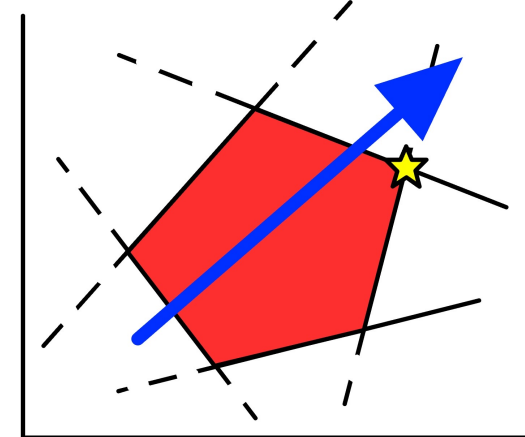
Carleton

1. Introduction

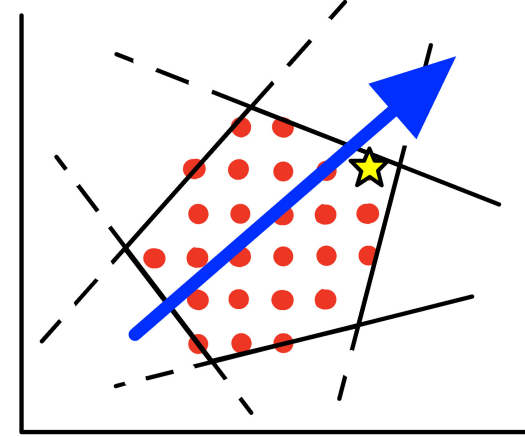
1. Understanding Integer Linear Programming

Solve for x_1, x_2, \dots, x_N

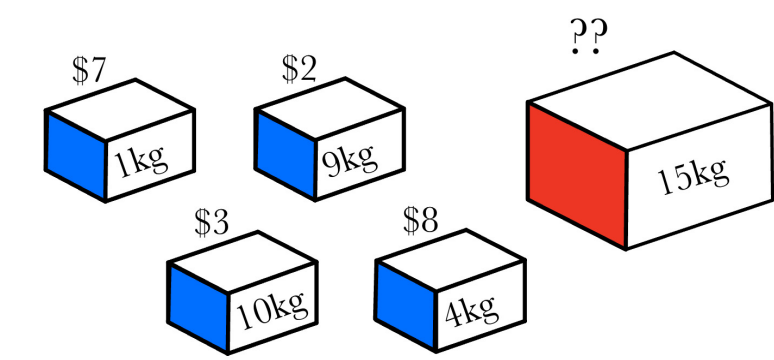
Objective Function: $\sum_{i=1}^N c_i x_i$
Constraints: $\sum_{i=1}^N a_{j,i} x_i \leq b_j$



Integer Variables



2. Example: Knapsack Problem



Maximize the total value
without exceeding the
weight limit

Let $x_i = 1$ if the i th item is taken and $x_i = 0$ if otherwise.

Maximize $7x_1 + 2x_2 + 3x_3 + 8x_4$
While respecting $1x_1 + 9x_2 + 10x_3 + 4x_4 \leq 15$
Where $0 \leq x_i \leq 1$

3. Carleton's CS Match

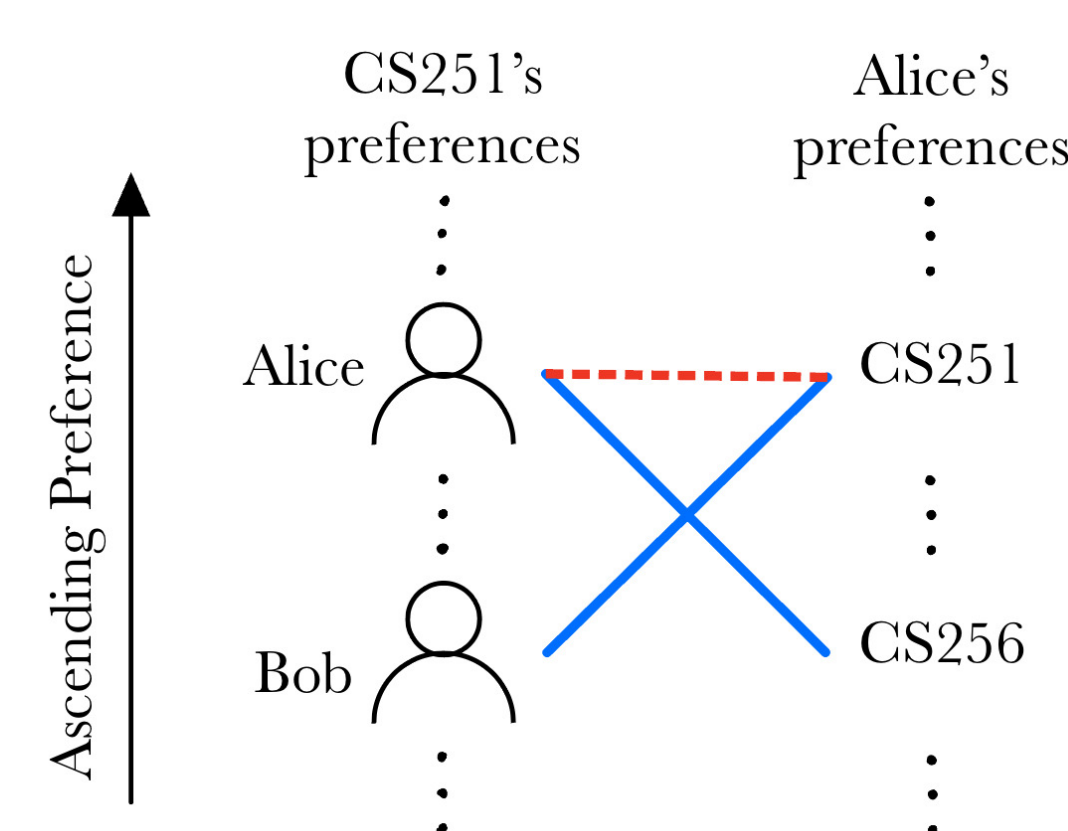
≥ 300 students
 ~ 10 courses
Only **74%** first-choice matches
Currently uses the Gale Shapley Algorithm

4. Preferences and Stable Matching

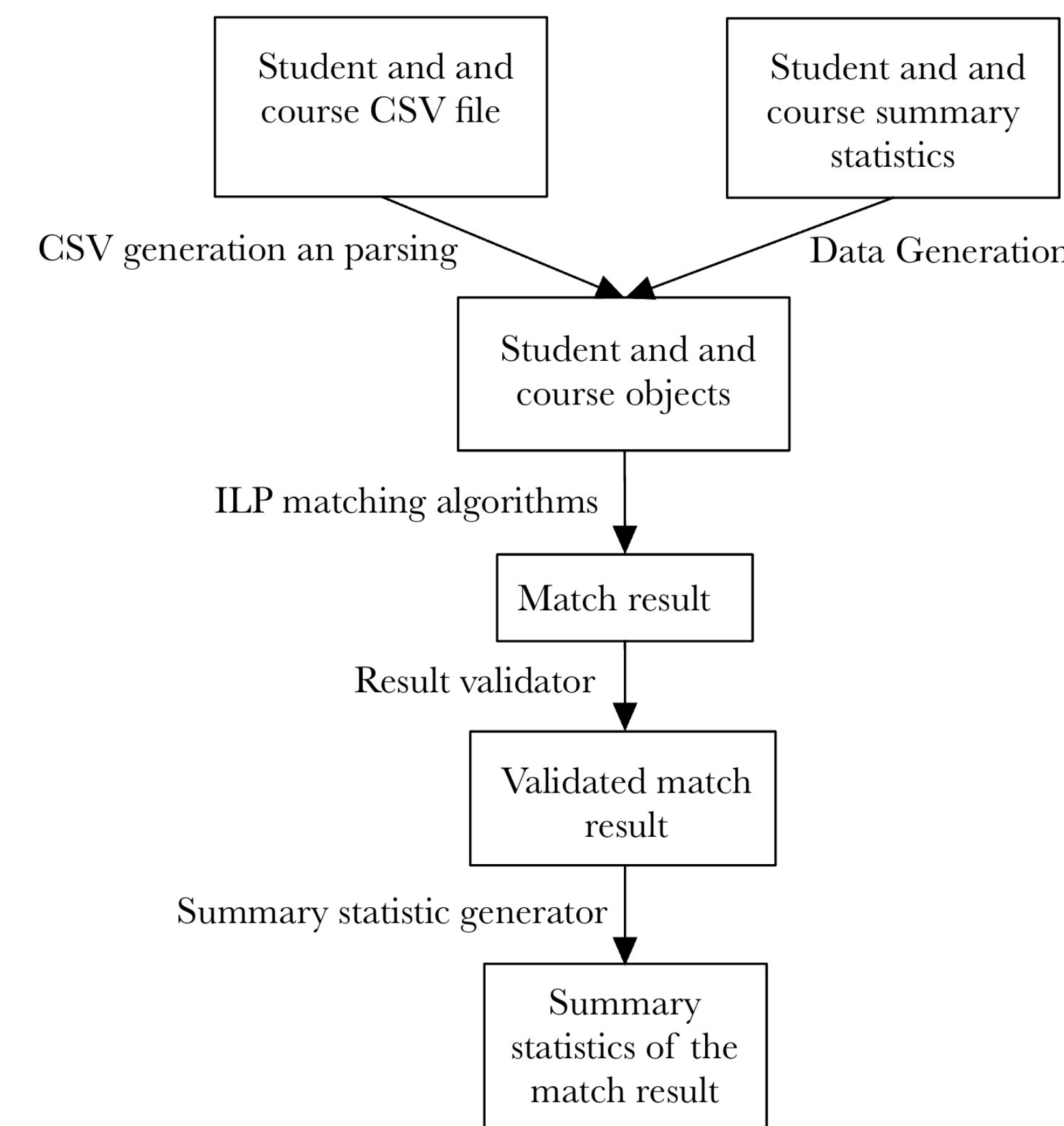
Student and Course preference Lists: Each student submits a preference list of courses

Courses rank students based on seniority, number of courses taken, etc.

Stable Matching: After a match, no unmatched pair wants each other more than what they are matched to



2. Work Structure



3. ILP formulations of the CS Match

For i th student and j th course, let $x_{ij} = \begin{cases} 0, & \text{if } i \text{ and } j \text{ didn't match} \\ 1, & \text{if } i \text{ and } j \text{ matched} \end{cases}$

1. Students get matched to at most one course

$$\sum_{j=1}^C x_{ij} \leq 1$$

2. Courses don't exceed their capacity

$$\sum_{i=1}^S x_{ij} \leq \text{Capacity}_j$$

3. "Stable Matching" constraint

For i th student and j th course
let $S = \{i' | i' \text{th student is not worse than } i \text{th student for } j \text{th course}\}$
let $C = \{j' | j' \text{th course is not worse than } j \text{th course for } i \text{th student}\}$

$$\text{Capacity}_j (1 - \sum_{j' \in C} x_{ij'}) \leq \sum_{i' \in S} x_{i'j}$$

4. Objective functions

Hospital Resident Matching with Ties.

$$\text{Maximize } \sum_{i=1}^S \sum_{j=1}^C x_{ij}$$

Gale Shapley Emulation

$$\text{Maximize } \sum_{i=1}^S \sum_{j=1}^C x_{ij} (10000 - \text{index}(i, j))$$

Where $\text{index}(i, j)$ is the **order** of j th course for i th student

Weighted Matching

$$\text{Maximize } \sum_{i=1}^S \sum_{j=1}^C x_{ij} (10000 + \text{weight}(i, j))$$

Where $\text{weight}(i, j)$ is the **weight** of j th course for the i th student

4. Specialization: Simplex Algorithm

Standard Form

Minimize $\sum_{i=1}^N c_i x_i$
With respect to $\sum_{i=1}^N a_{j,i} x_i = b_j$
Where $x_i \geq 0$

Tableau Form

x_1	x_2	x_3	x_4	x_5	f	RHS
2	2	0	1	2	0	7
1	1	1	3	0	0	6
1	2	3	1	0	0	6
2	2	-1	2	1	-1	0

Canonical Form

x_1	x_2	x_3	x_4	x_5	f	RHS
1	0	0	15/2	-1	0	15/2
0	1	0	-7	2	0	-4
0	0	1	5/2	-1	0	5/2
0	0	0	7/2	-2	-1	-9/2

Pivot

x_1	x_2	x_3	x_4	x_5	f	RHS
1	1/2	0	4	0	0	11/2
0	1/2	0	-7/2	1	0	-2
0	1/2	1	-1	0	0	1/2
0	1	0	-7/2	0	-1	-17/2

Repeat

x_1	x_2	x_3	x_4	x_5	f	RHS
1/4	1/8	0	1	0	0	11/8
7/8	15/16	0	0	1	0	45/16
1/4	5/8	1	0	0	0	15/8
7/8	23/16	0	0	0	-1	-59/16

Results $x_1 = x_2 = 0, x_3 = \frac{15}{8}, x_4 = \frac{11}{8}, x_5 = \frac{45}{16}, f = \frac{59}{16}$

Finiteness of the Algorithm

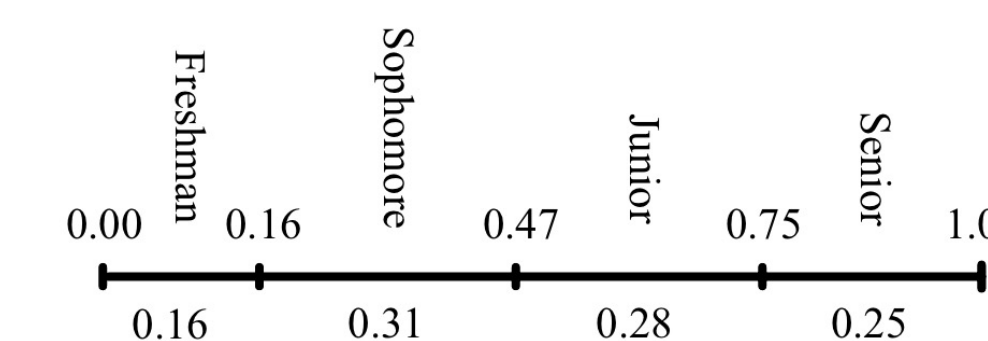
Bland Rule: Out of equivalent choices (rows or columns), pick one that has the minimum index

5. Specialization: Data Generation

Accessed only summary statistics to respect the privacy of students

	# of total students	Mean # of classes wanted
Freshman	57	4.15
Sophomore	96	4.51
Junior	97	3.85
Senior	61	2.43

1. Generate Class Year
 $c \sim U(0, 1)$



2. Generate course weights based on class year
 $w_i \sim U(0, 1)$

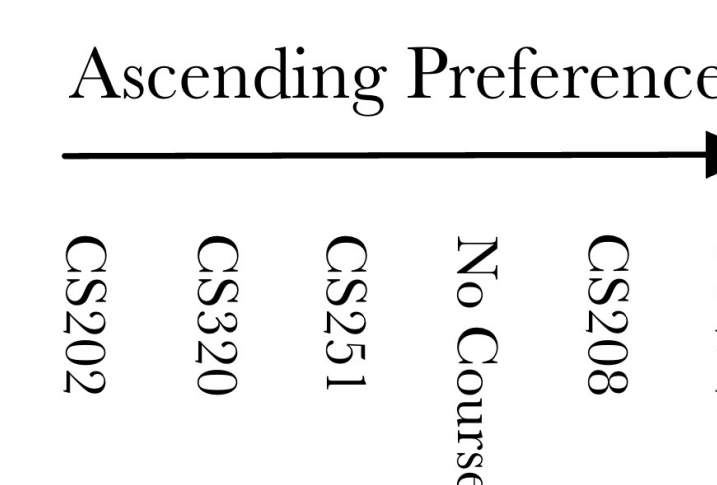
CS202	CS208	CS251	CS252	CS320
0.34	0.10	0.34	0.27	0.34

4.15 / 9 = 0.45 cutoff

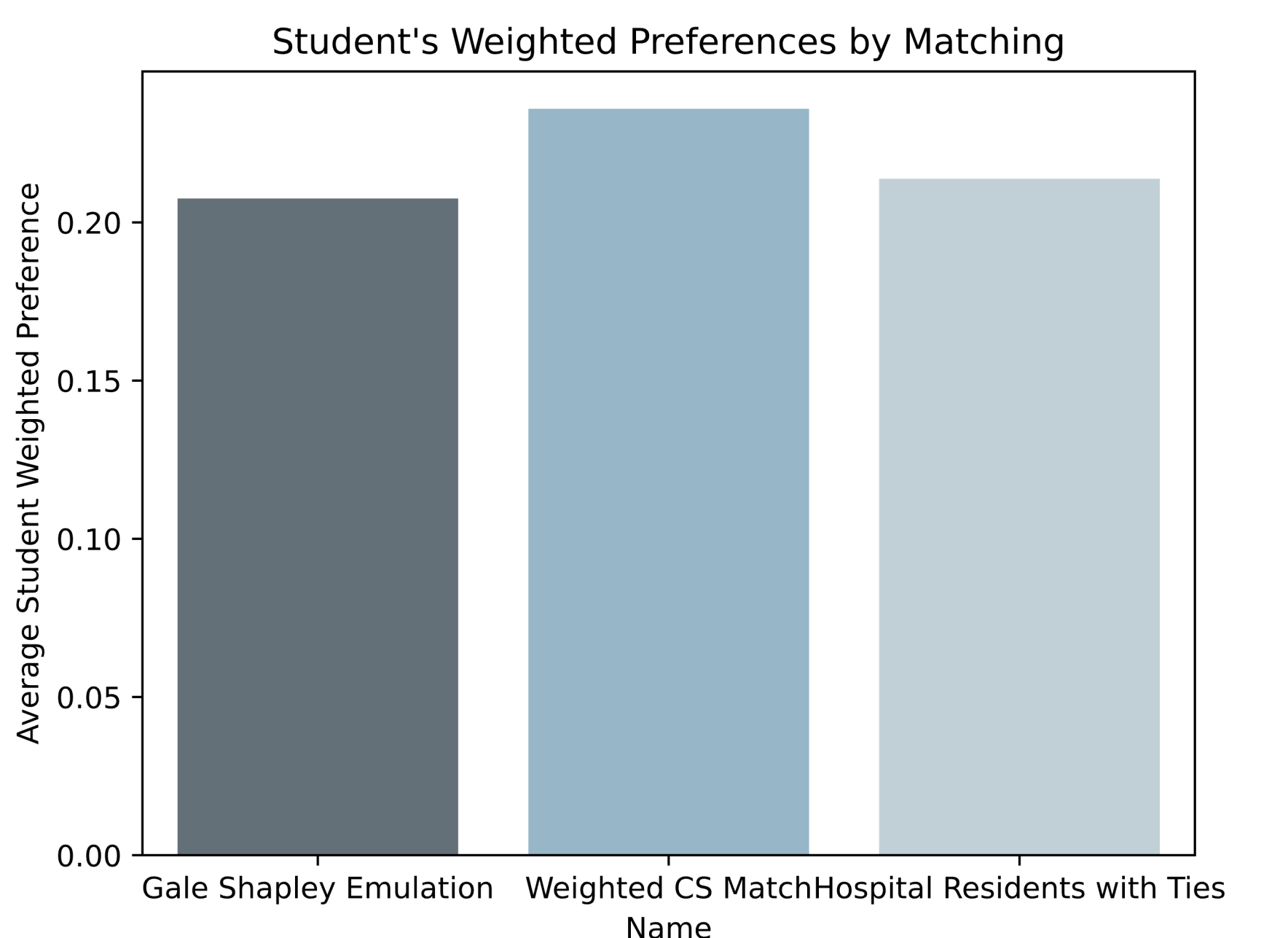
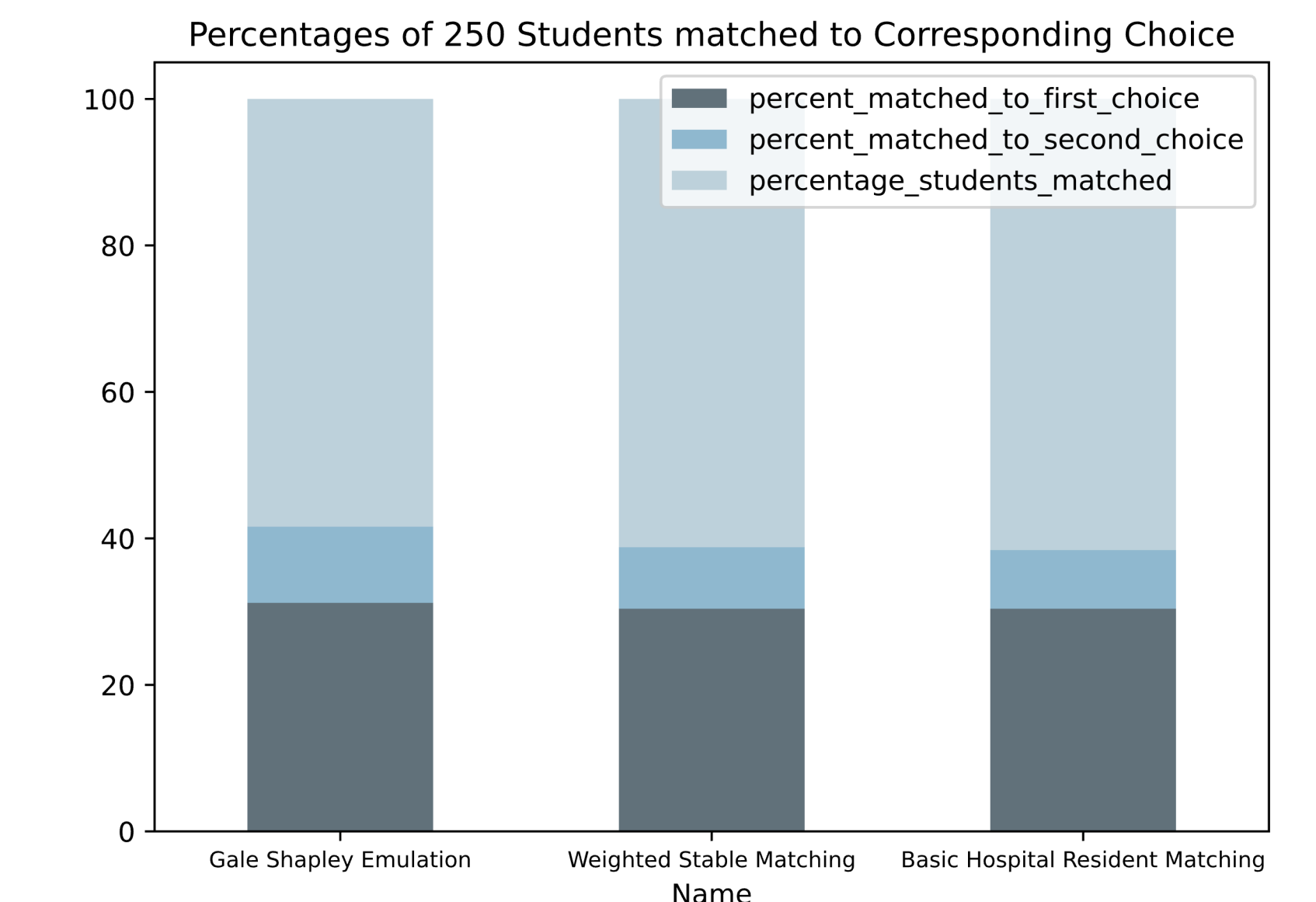
3.1 Generate normalized weights

CS202	CS208	CS251	CS252	CS320
0.00	0.27	0.0	0.73	0.00

3.2 Generate preference list



6. Preliminary Results



7. Conclusion

Integer Linear Programming is a flexible technique that can model not only trivially linear problems but also various optimization problems with non-trivial structures.

Adjusting the objective function provides us with even greater flexibility in defining what is the most optimal, consequently influencing the resulting solutions of the ILP algorithm.

Initially, we had intended to run our ILP models against Carleton's CS Match algorithm using the same generated dataset. We had to forego this plan due to time constraints. It is important for the study to running such comparisons in the future.

8. Acknowledgements

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9. References

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