

## Background

- ▶ General-purpose correct-by-construction synthesis methods rely on set reasoning
- ▶ Set manipulation intractable for large-scale systems. Must exploit symmetry when possible
- ▶ Best-case solution: propose analytical controller and prove that closed-loop system satisfies specification

## Motivating example: **TCL** scheduling

- ▶ Thermostatically controlled loads (TCLs) are air condition units, refrigerators etc, that operate within a temperature range
- ▶ Local temperature constraints vs aggregate power consumption constraints

## Counting problem: Definition

Large number  $i = 1, \dots, N$  of subsystems  $x_i \in \mathbb{R}^d$  with dynamics:

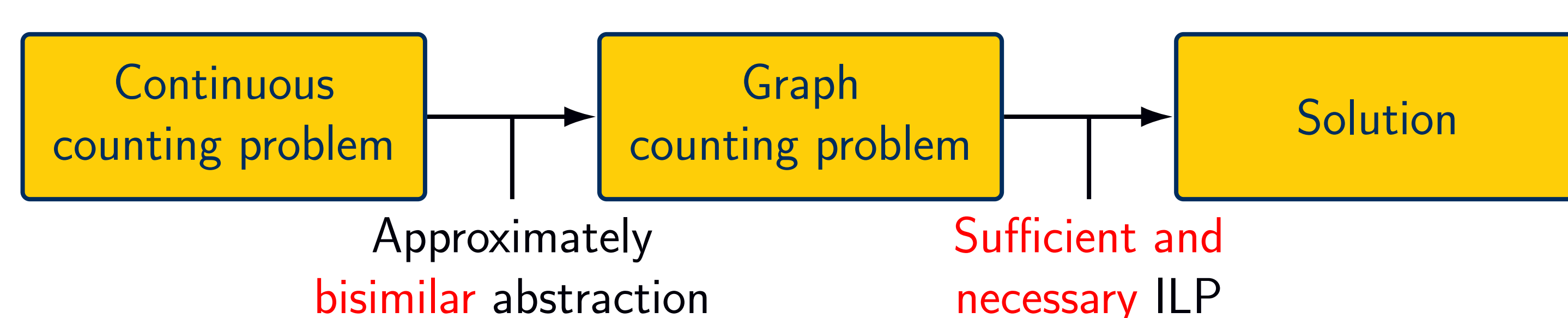
$$\dot{x}^i(t) = f_{\sigma^i(t)}(x^i(t), \delta^i(t)), \quad \sigma^i : \mathbb{R} \rightarrow \{1, \dots, M\}$$

Counting constraints  $(X^l, R^l)$ :

$$\sum_{i=1}^N \mathbb{1}_{X^l}(x^i(t), \sigma^i(t)) \leq R^l.$$

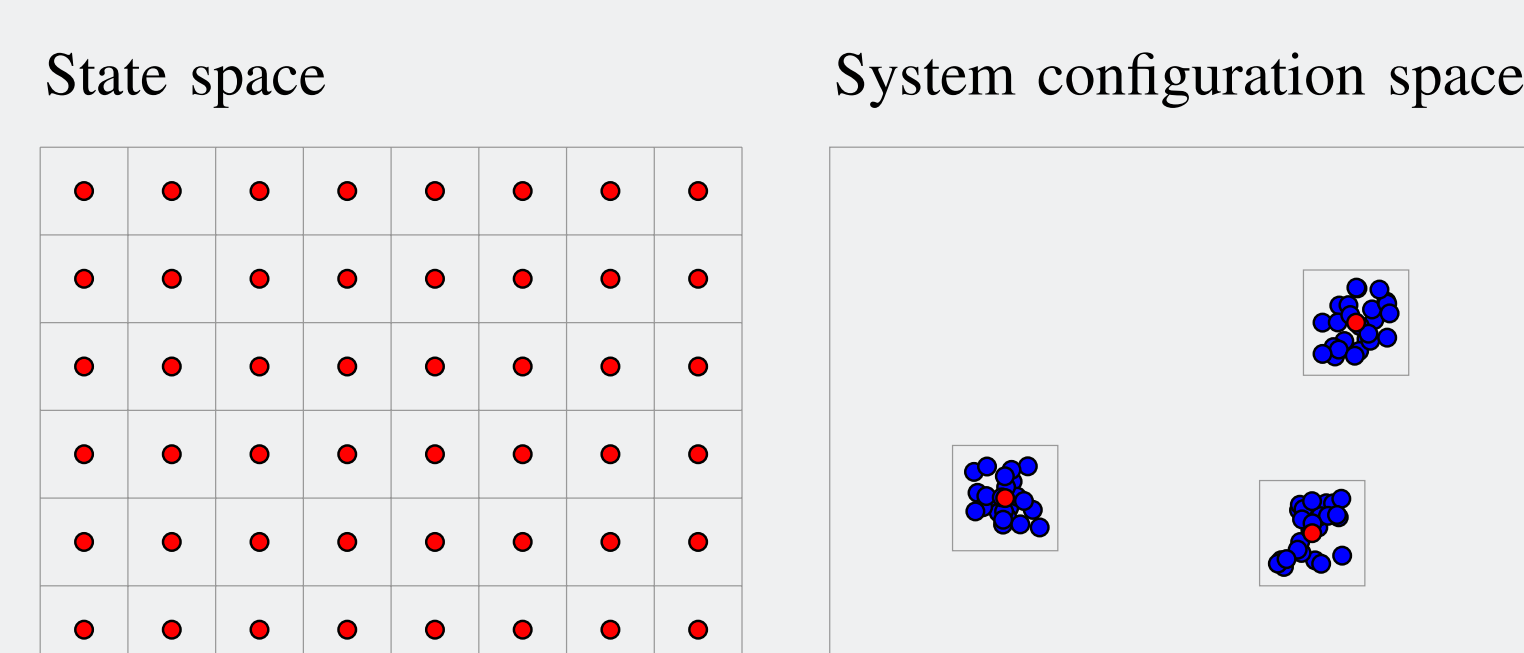
- ▶ Counting constraints for safety, aggregate consumption, balancing, etc
- ▶ Counting constraints are **permutation invariant**
- ▶ **Problem:** given initial state, construct a switching protocol  $\{\sigma^i(t)\}_{n=1}^N$  such that the constraints are satisfied

## Synthesis for symmetric problem via abstraction and integer programming



## Re-use abstraction

- For **mild heterogeneity**, use same abstraction for similar subsystems. Can “abstract away” mild heterogeneity



- ▶ Standard abstraction:  $\mathcal{O}\left(\frac{1}{\eta}\right)^{Nd}$
- ▶ Counting abstraction:  $\mathcal{O}\left(\frac{1}{\eta}\right)^d$

- Necessary and sufficient conditions for feasibility of continuous problem via approximate simulation relations

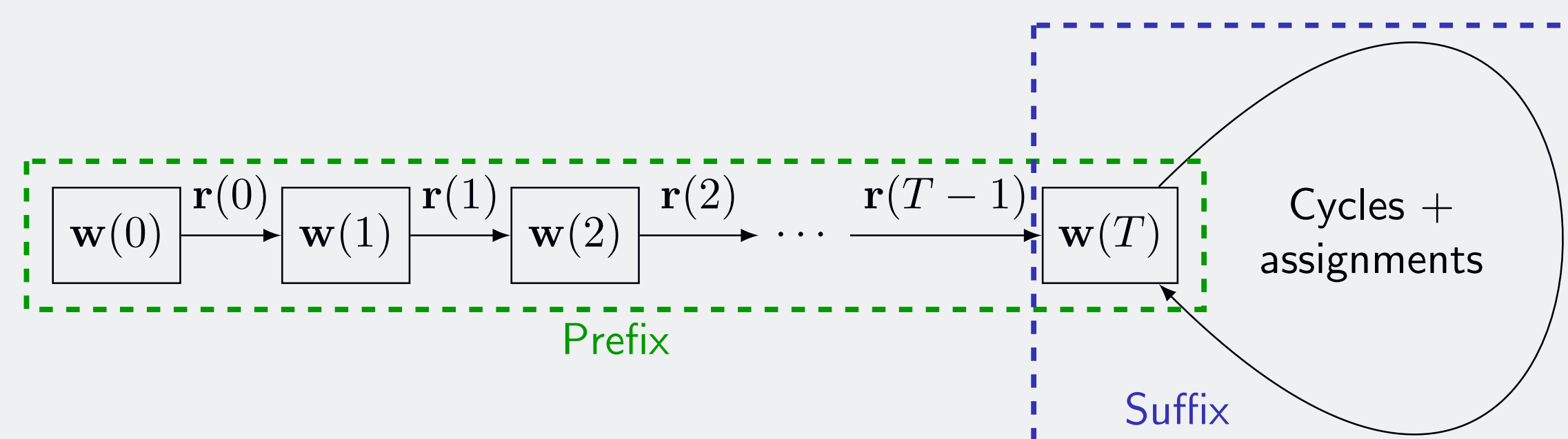
$$\beta(\epsilon, \tau) + \frac{\delta}{K} (e^{K\tau} - 1) + \frac{\eta}{2} \leq \epsilon,$$

## Aggregate counting dynamics and ILP solution

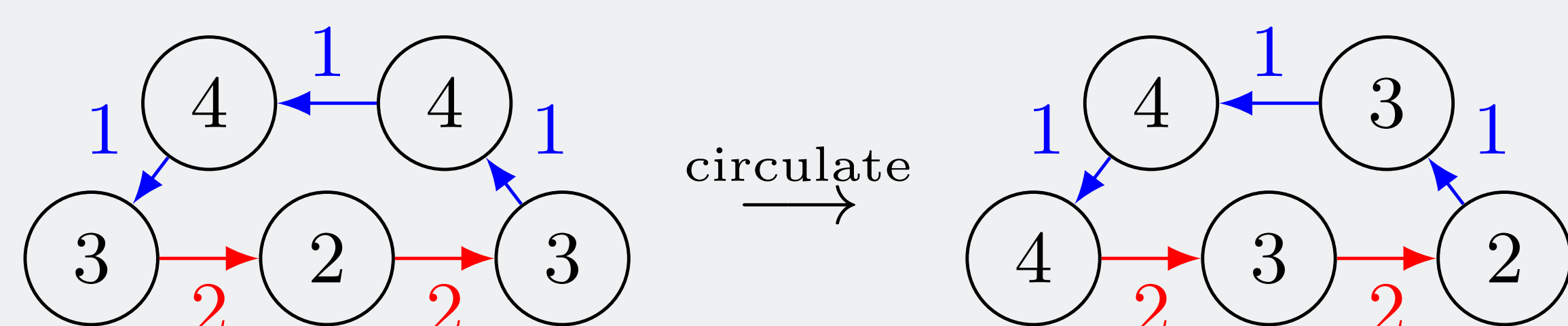
- Abstraction graph has nodes  $\{\nu_k\}_{k=1}^K$ , let  $w_k(s) \in \mathbb{N}$ : number of subsystems at node  $\nu_k$ . Aggregate dynamics:

$$w_k(s+1) = \sum_{m \in [M]} \sum_{k' \in \mathcal{N}_k^m} r_{k'}^m(s), \quad k \in [K],$$

- ▶ Search for prefix-suffix trajectories through ILP formulation



- ▶ Prefix is finite-horizon trajectory, suffix consists of **cycle-assignment pairs**



- ▶ Sufficient if prefix-horizon and cycle set is large enough
- ▶ If ILP too large: solve as **relaxed LP** and round solution

## Analytic approach for heterogeneous problem with one-dimensional subsystems (e.g. TCLs)

## Heterogeneous problem

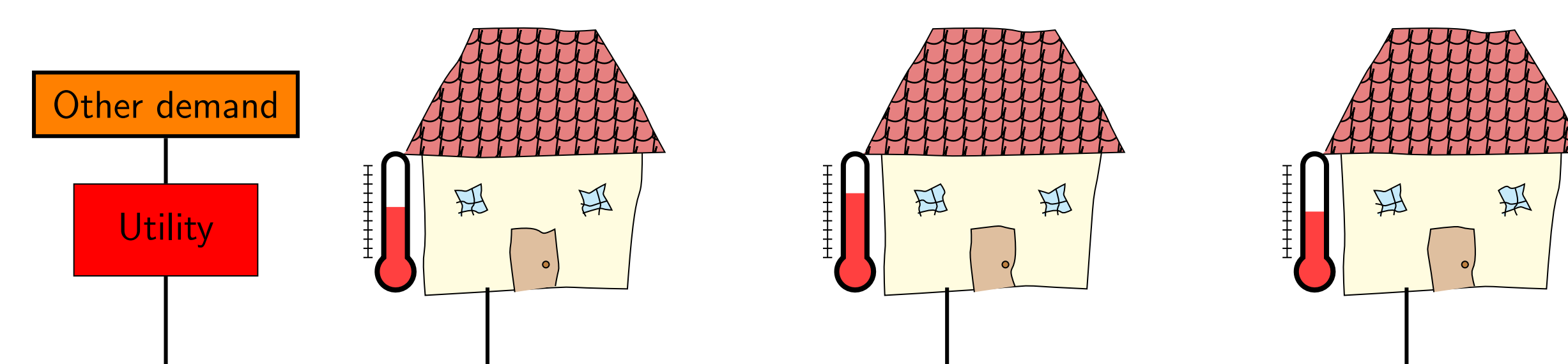
- Strongly heterogeneous family of one-dimensional switched systems

$$\frac{d}{dt}x^i(t) = \begin{cases} f_{\text{off}}^i(x^i(t)) & \text{if } \sigma^i(t) = \text{off}, \\ f_{\text{on}}^i(x^i(t)) & \text{if } \sigma^i(t) = \text{on}, \end{cases}$$

- ▶ Local constraints  $\underline{a}^i \leq x^i(t) \leq \bar{a}^i$
- ▶ Mode-on-counting constraint:  $\underline{K}_{\text{on}} \leq \sum_{i: \sigma^i(t)=\text{on}} 1 \leq \bar{K}_{\text{on}}$

- Feasibility condition

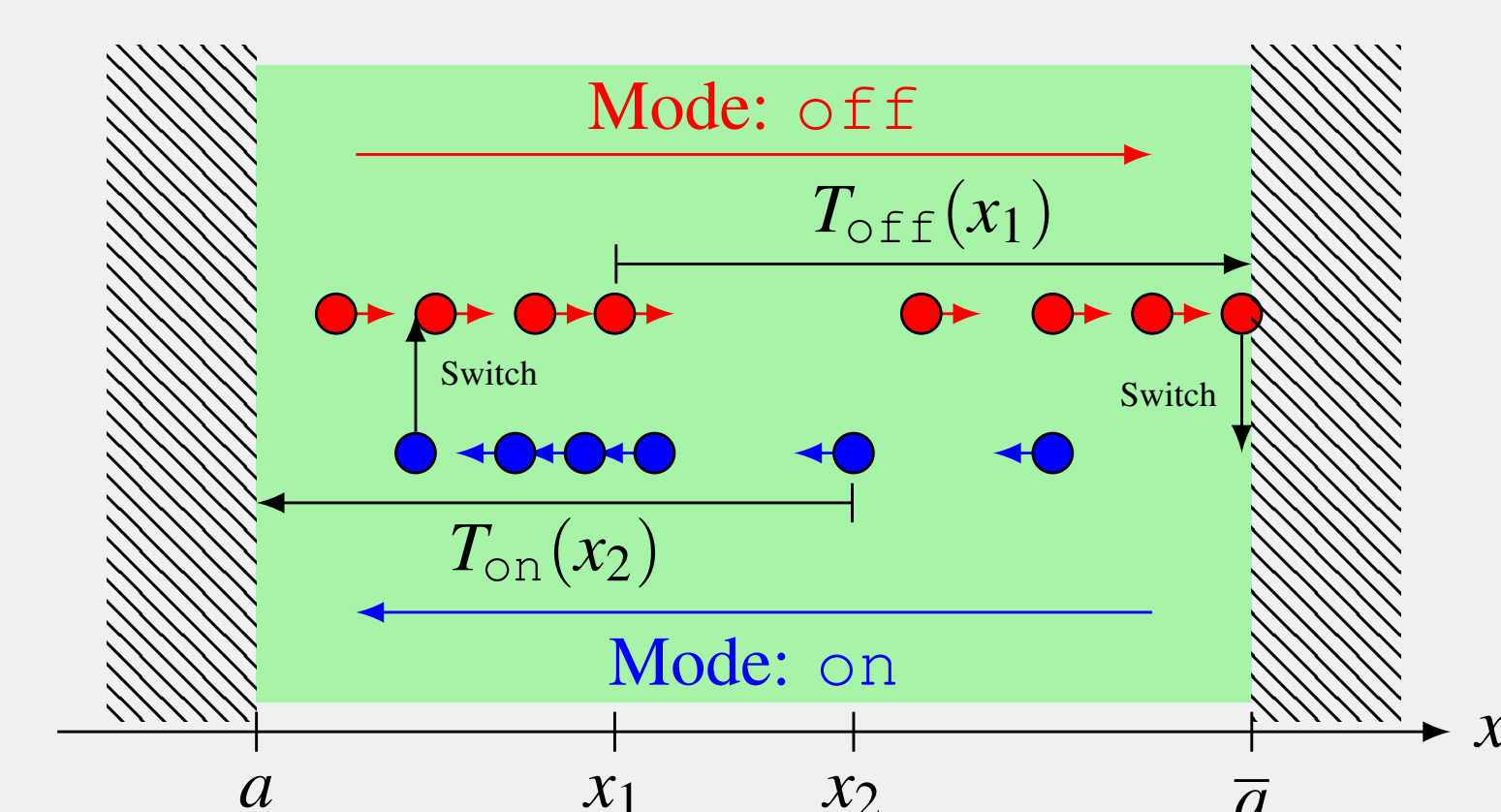
$$\sum_{i=1}^N \frac{\mathcal{L}_{\text{off}}^i T_{\text{on}}^i(\underline{a}^i)}{1 + \mathcal{L}_{\text{off}}^i T_{\text{on}}^i(\underline{a}^i)} > \underline{K}_{\text{on}}, \quad \sum_{i=1}^N \frac{\mathcal{L}_{\text{on}}^i T_{\text{off}}^i(\overline{a}^i)}{1 + \mathcal{L}_{\text{on}}^i T_{\text{off}}^i(\overline{a}^i)} > N - \overline{K}_{\text{on}} \quad (1)$$



## “Lazy” feedback control strategy

- ▶ Only switch subsystems when a constraint is about to be violated

- ▶ Select subsystems based on times to exit  $T_{\text{off}}^i$  and  $T_{\text{on}}^i$



## Theorem

Sufficient: If (1) holds, then the proposed strategy solves the problem for any non-degenerate initial condition.

Necessary: If  $f_{\text{on}}^i$  and  $f_{\text{off}}^i$  are monotonically decreasing in  $[\underline{a}^i, \bar{a}^i]$  and (1) is **strictly** violated, then the problem has no solution for any initial condition.

Simulation of 10,000 TCLs

