

Preliminary Results on Correct-by-Construction Control Software Synthesis for Adaptive Cruise Control

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Outline

Introduction

Problem setup

Solution: Continuous reachability computation

Solution: Discrete abstraction [PESSOA]

Simulation results

Comparison and Conclusion

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Adaptive Cruise Control

- ▶ Maintains set speed when no lead car present.
- ▶ Follows lead car at a safe distance.

Formal methods in automotive safety

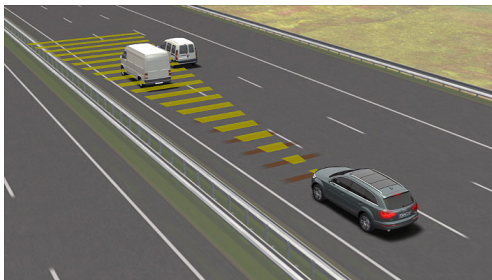


Image credit: Audi

Why formal methods?

- ▶ Safety and dynamics intimately connected.
- ▶ Explicit characterization of performance limits.
- ▶ Enables formal composition of components.

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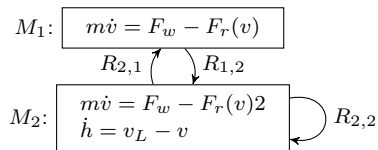
Comparison and Conclusion

System model

Car model:

$$m\dot{v} = F_w - F_r(v), \quad F_r \text{ polynomial.} \quad (1)$$

Two-car system modeled as hybrid system with two modes M_1 (no lead car) and M_2 (lead car):



Reset maps:

- ▶ Lead car leaves: use reset map $R_{2,1}$.
- ▶ Lead car cuts in: use reset map $R_{1,2}$.
- ▶ New lead car: use reset map $R_{2,2}$.

Specifications

English specifications:

Set/LTL specifications:

¹The time headway, or time to collision, is $\tau = h/v$.

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1. $\square(F_w \in \underbrace{[-0.3mg, 0.2mg]}_{S_2})$

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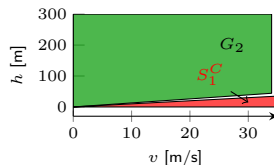
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3. $\square(\square M_2 \implies \diamond \square(v, h) \in G_2) \wedge \square(M_2 \implies (v, h) \notin S_1^C)$



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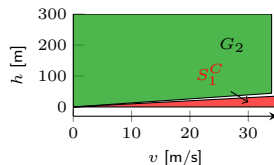
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Complete specification:

$$\square((M_1 \vee S_1) \wedge S_2) \wedge \square\left(\bigwedge_{i=1}^2 \square M_i \implies \diamond \square G_i\right) \quad (2)$$

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Assumptions and Problem Statement

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Problem statement

Under the preceding assumptions, synthesize a domain \mathcal{D} and a controller for the hybrid dynamics, which enforces the specifications given that the reset maps are restricted to \mathcal{D} .

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Continuous computation: Idea

Focus on the lead car mode M_2 . To satisfy the given “reach-stay-while avoiding” specification:

- ▶ Find a **controlled-invariant** set $C_1 \subset G_2$.
- ▶ Compute a set of chains that reaches C_1 while staying in S_1 .

Definition

Let Σ be the system $x(t+1) = f(x(t), u(t))$ with the constraint $u(t) \in U$. A set C is *controlled-invariant* for Σ if for all $x_0 \in C$ there is a $u_0 \in U$ such that $f(x_0, u_0) \in C$.



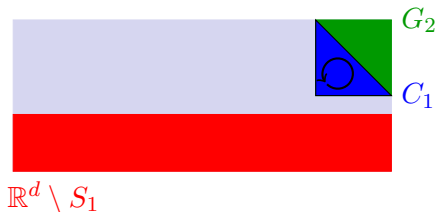
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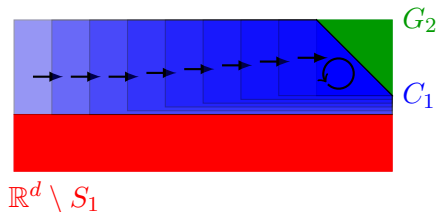
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Continuous computation: Solution overview

1. Linearize and time-discretize system dynamics in a way that correctness is inherited by the original system.
 - ▶ Reachability in linear system must imply reachability in original system.
2. Use linear reachability computations to synthesize a correct-by-construction controller for the linearized system.
3. Implement controller on original system.

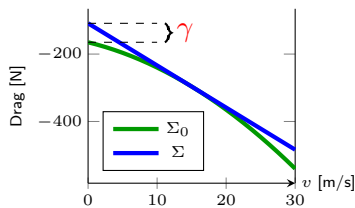
Linearization and time-discretization of system dynamics

We linearize and integrate the dynamics

$$\Sigma_0 : \begin{aligned} m\dot{v} &= F_w - f_0 - f_1 v - f_2 v^2, \\ \dot{h} &= v_L - v \end{aligned}$$

for a time Δ to obtain

$$\Sigma : \begin{bmatrix} v(t + \Delta) \\ h(t + \Delta) \\ v_L(t + \Delta) \end{bmatrix} = A \begin{bmatrix} v(t) \\ h(t) \\ v_L(t) \end{bmatrix} + B\bar{F}_w(t) + K.$$



Proposition

For a fixed $d : [0, \tau] \rightarrow D$, if there exists an input $\bar{F}_w : [0, \tau] \rightarrow [-0.3mg, 0.2mg - \gamma]$ that steers the state of Σ from (v^0, h^0, v_L^0) to (v^1, h^1, v_L^1) , then there exists an input $F_w : [0, \tau] \rightarrow [-0.3mg, 0.2mg]$ that steers the state of Σ_0 from (v^0, h^0, v_L^0) to (v^1, h^1, v_L^1) .

Reachability computations for affine systems I

Problem

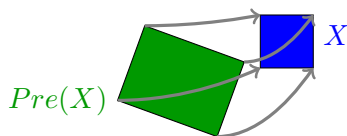
Given a discrete-time affine system

$$x(t+1) = Ax(t) + Bu(t) + K,$$

with (possibly state-dependent) input constraints

$$H_x^u x(t) + H_u^u u(t) \leq h^u$$

and a final set $X = \{x : Hx \leq h\}$, we want to find the set of initial states $Pre(X)$ from where the system can be steered to X .



Reachability computations for affine systems II

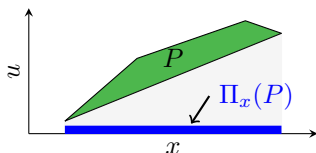
Solution

Write all constraints as a polyhedron in $x - u$ -space:

$$P = \left\{ (x, u) : \begin{bmatrix} HA & HB \\ H_x^u & H_u^u \end{bmatrix} \begin{bmatrix} x \\ u \end{bmatrix} \leq \begin{bmatrix} h - HK \\ h^u \end{bmatrix} \right\}.$$

Then $Pre(X)$ can be characterized as the x -projection of this polyhedron:

$$Pre(X) = \Pi_x(P).$$



Can be extended to multiple time steps and/or systems with disturbance. Projection can be computed with MPT [Herceg et al. (2013)].

Algorithms

With the reachability operator Pre , the required sets can be found for a 'reach-and-stay while avoiding' type specification as follows.

- ▶ A controlled-invariant set C_1 inside a goal set G is the fixed point of the operator

$$X \mapsto G \cap Pre(X)$$

initialized with G itself.

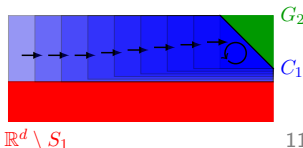
- ▶ To reach C_1 while staying in a safe set S , iterating

$$C_{i+1} = S \cap Pre(C_i)$$

will give a list of sets $C_1 \leftarrow C_2 \leftarrow C_3 \leftarrow \dots$ such that C_i can be reached from C_{i+1} while staying in S .

Correct control strategy:

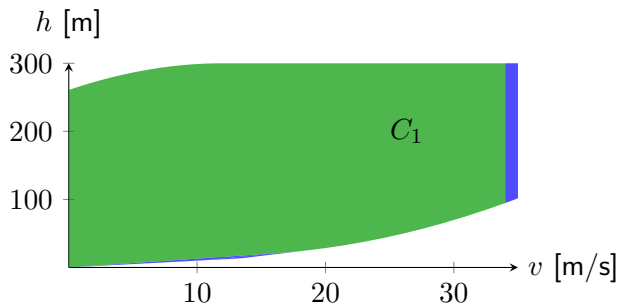
When in C_i , steer to $C_{\max(1,i-1)}$ in finite time.



Result for ACC

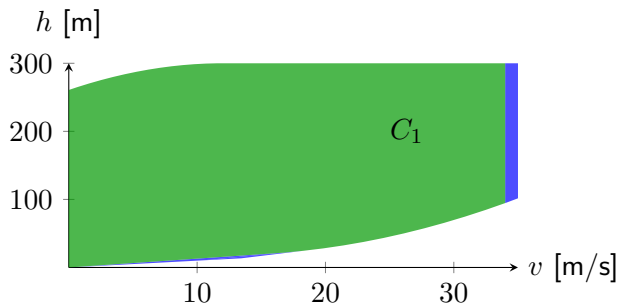
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Controlled-invariant set C_1 in $v - h$ space:



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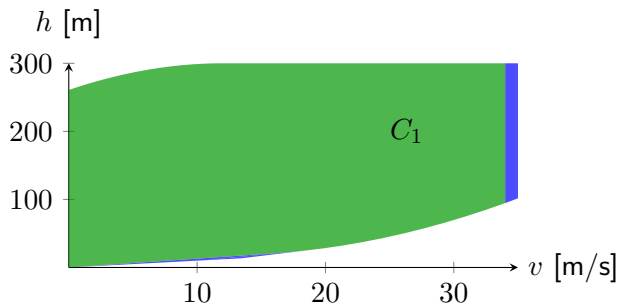
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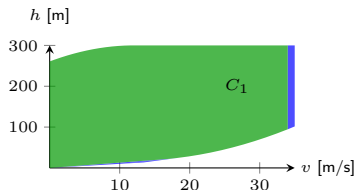
- How to move between the sets returned by the reachability algorithms?

Control strategy implementation: MPC

Use N -step horizon Model-Predictive Control (MPC) to implement low-level controller:

$$u^* = \begin{cases} \min & L(x, u), \\ \text{s.t.} & H_x^u x + H_u^u u \leq h^u, \\ & x \in C_i \end{cases}$$

Feasibility of the QP is guaranteed by construction.

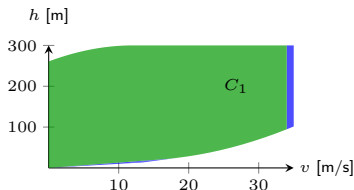


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Composition with CCC

- ▶ Using any conventional cruise controller (CCC) for mode M_1 together with the presented controller for M_2 is correct if reset maps are restricted.
- ▶ For a general hybrid system, control domains need to be propagated back and forth, but in this case there is convergence in the first time step (thanks to the simplicity of CCC).

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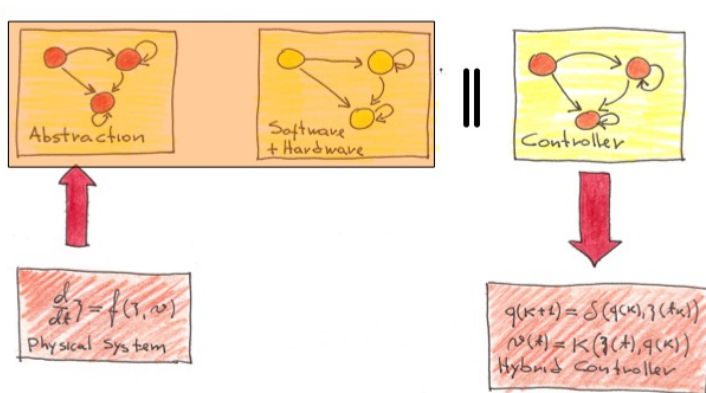
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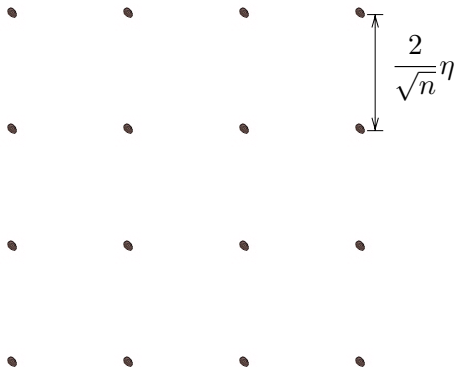
Comparison and Conclusion

1. Finite-state abstraction;
2. Finite-state controller synthesis;
3. Controller refinement.



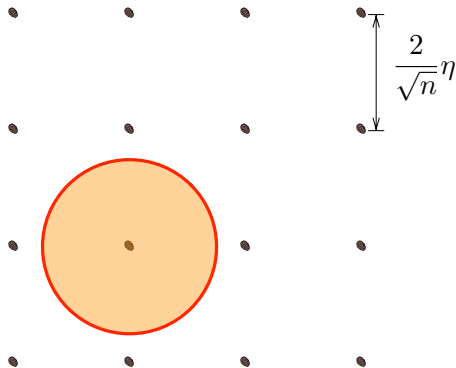
Abstraction Computation

Finite-state abstractions can be extracted from control systems using a very simple construction.



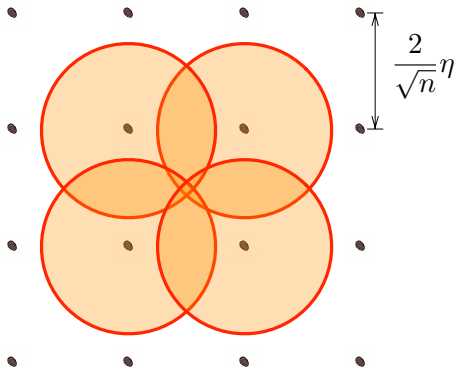
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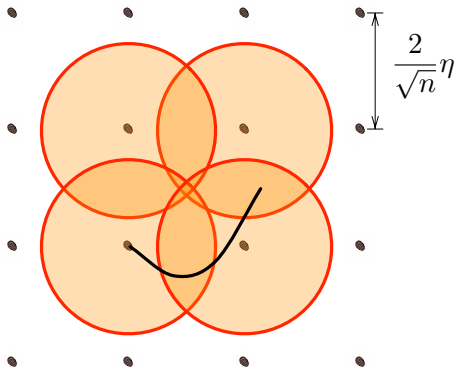
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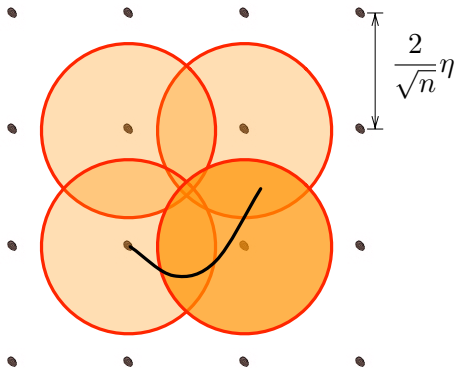
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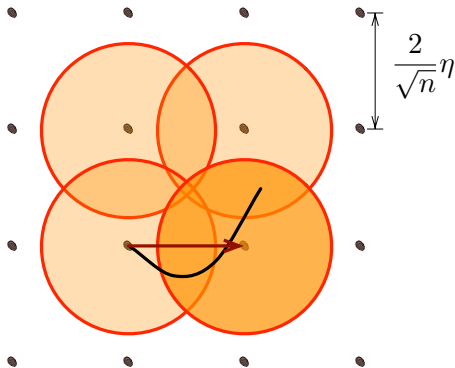
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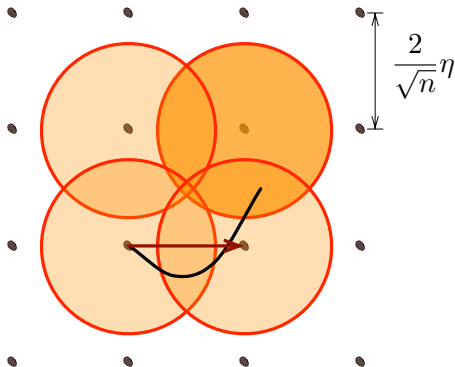
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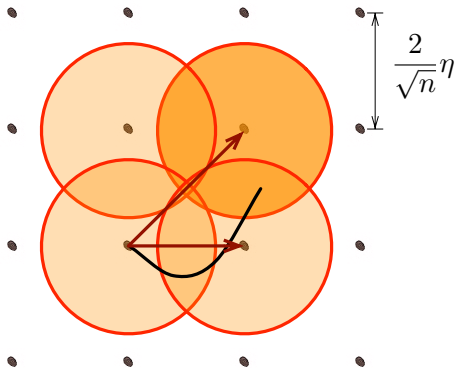
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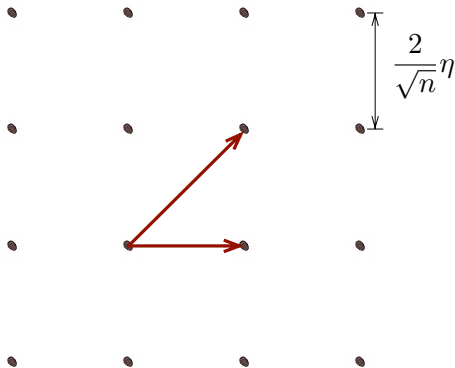
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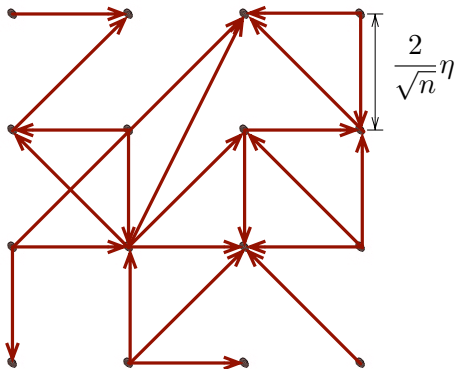
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Controller Synthesis

- ▶ Controller synthesis problem can be modeled as a game.
 - ▶ A game consists of
 1. Transition system **abstraction**
 2. Wining condition **specification**
 3. Players **controller and environment**

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- ▶ Fixed-point algorithms are used to compute the wining set (controller domain).
- ▶ Termination is guaranteed by finiteness of the set of states.

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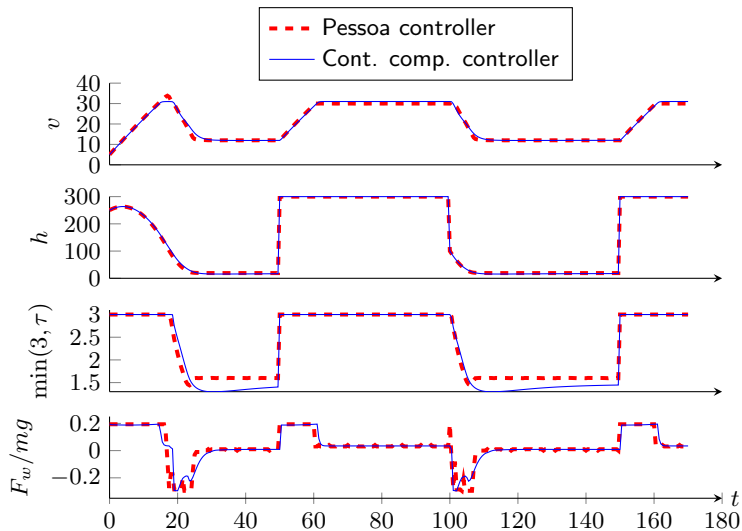
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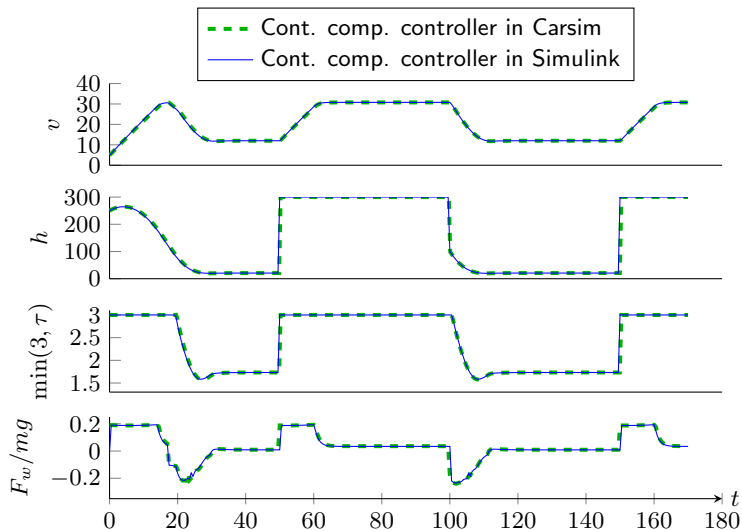
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Simulink simulation

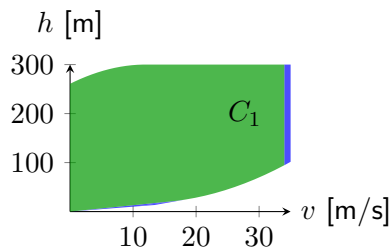
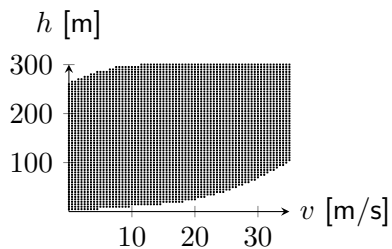


Carsim simulation



Carsim is a standard car simulator for Automotive industry.

Carsim simulation videos



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Comparison

- ▶ Both methods employ fixed point algorithms on a continuous and discrete state space, respectively.
- ▶ Discrete state space: no linearization required, algorithm guaranteed termination.
- ▶ Continuous state space: no abstraction required, tuning flexibility.

Summary

Conclusions:

- ▶ Synthesized two correct-by-construction adaptive cruise controllers for a low-order model and a high-order (~ 30 states), industry-standard model.
- ▶ Used industry-accepted model and realistic parameters.
- ▶ Compared abstraction-based method with continuous reachability computation method.

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Future work:

- ▶ Allow lead car to have variable speed (to appear).
- ▶ Implement on model cars.
- ▶ HIL simulation.
- ▶ Formal composition of interconnected systems equipped with correct-by-construction controllers.

Thank you for your attention.

References



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Multi-Parametric Toolbox 3.0.
In Proc. of the ECC, 2013.