

# EXTENSIONS OF GRUNDY'S GAME

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## 1. GRUNDY'S GAME

Grundy's Game is well known in game theory. Two players begin with a pile of coins. The first player splits the pile into two uneven piles. The second player picks one of the piles and splits it into two uneven piles. The first player then picks one of the three piles and splits it into uneven piles. This goes on until no more uneven splits can be made. The last player to move wins.

Grundy's Game has been studied extensively and although some patterns have been found in its SG-values<sup>1</sup>, no one has been able to come up with an overall description of them that does not involve examining the whole game tree.

## 2. EXTENDING GRUNDY'S GAME

We present two ways extend Grundy's game by allowing more splits and solve some cases.

## 3. FIXED NUMBER OF SPLITS

We can extend the game such that for a given natural number  $M \geq 2$  at each turn the pile must be split into *exactly*  $M$  piles, no two of which have the same size.

**3.1. Upper bound.** It is obvious that if we have a pile of size  $n < \sum_{i=0}^M i = \frac{M(M+1)}{2}$  that it is not possible to split it into  $M$  uneven piles. Therefore, for such  $n$ ,  $g(n) = 0$ . However with  $n = \frac{M(M-1)}{2}$ , the split  $1, 2, \dots, M$  is valid and will have SG-value 1, since the SG-values of all the smaller piles are 0.

We now want to prove that the SG-values cannot increase faster than once for every  $\sum_{i=1}^{M-1} i = \frac{(M-1)M}{2}$  increases in  $n$ , the pile size. We know that the first value of 1 appears at pile size  $\frac{M(M+1)}{2}$ . To get the first value of 2 we must have at least one of the piles with a value of 1. Therefore at least one pile must be of size  $\frac{M(M-1)}{2}$ . The minimum sizes of the remaining piles that give a valid split are  $1, 2, \dots, M-1$ , so the first value 2 can appear at first with pile size  $\frac{M(M+1)}{2} + \frac{(M-1)M}{2}$ , and it does, since all the small piles have SG-value 0.

Let us now assume that for all piles up to a given  $n$ , the SG-values increase at most at the rate of one every  $\frac{(M-1)M}{2}$  increases in pile size. Further assume that  $n$  is the smallest pile that has a given SG-value  $k$ . Then we want to show that the smallest possible pile size  $m$  that has SG-value  $k+1$  is  $m = n + \frac{(M-1)M}{2}$ .

**3.2. Upper bound can be reached for  $M > 3$ .** Now we want to show that we can reach the upper bound for  $M > 3$ . Let  $n$  be some pile and  $k$  be it's upper bound according to the previous section. Let  $0 \leq l < k$ . We will show that we can split  $n$  such that the new piles together have SG-value  $l$ .

First find the smallest pile  $n'$  with SG-value  $l$ . That is our first pile. We will then split the remaining coins into  $M-1$  piles such that their total values are 0. If  $M-1$  is odd, then take one coin and put it in its own pile. It has SG-value 0. We now have an even number of piles remaining to fill. Split the remaining coins as evenly as possible into these piles. That means that the pile sizes are a sequence  $n_1, n_2, \dots, n_{M-2}$  such that  $n_i - n_{i-1} = 1$  for all  $i$  except possibly one, where the difference is 2.

These  $M-2$  piles can have at most 2 different SG-values according to the induction hypothesis, since  $M-1 < \frac{(M-1)M}{2}$  for all  $M > 3$  ( $M-1$  is the largest possible difference between the sizes). If all the piles have the same size, then their SG-values cancel out and we are done. Further, if they have two different SG-values but there's an even number of piles for each value, they also cancel out and we are done.

We will fix this by modifying the two largest piles with the lower SG-value. We remove  $\frac{(M-1)M}{2} - 3$  coins from the smaller of these two piles. That will change it from being the second biggest possible pile

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<sup>1</sup>See for example *Winning Ways for Your Mathematical Plays*, page 111.

with this SG-value to being the second smallest. The other, larger, pile will increase in SG-value. Thus we now have an even number of piles with each SG-value. Note that this operation will not change these piles into piles of a size already taken, since we move them far enough.

There are two cases to consider where we still have piles of even size. The first one is when one of the remainder piles is of size 1 and therefore collides with the special pile of size one. If this happens, then in fact the remainder piles must be of sizes  $1, \dots, M-2$ , which means we're trying to create an SG-value that is in fact not possible to create, ie  $l \geq k$ . That's nonsense so we don't need to worry about that.

The other case is where the original pile is equal in size with one of the remainder piles. Then these piles are the smallest piles with this SG-value. We cancel the whole process and start over, but this time we pick the first pile to be the largest one with its SG-value. That guarantees that the remainder piles cannot collide with it. The rest of the process is the same.

**3.3.  $M = 3$ .** Clearly from the data, the previous does not hold for  $M = 3$ . The proof fails when the two remainder piles are of an uneven size and do not have the same SG-value. We cannot as easily move coins around in that case.

However, there is an apparent pattern by looking at the plot of the first 2000 values. Most of the values below the main line appear to be on a line of their own. I haven't been able to figure out what the pattern is.

#### 4. BOUNDED NUMBER OF SPLITS

We can also extend the game such that we allow the players to choose how many piles to split into, with an upper limit  $M$  on the number of piles.

**4.1.  $M > 4$ .** The SG-value for a pile of size  $n \geq 9$  appears to be  $n - 4$ . For smaller piles the SG-values can be read from the graphs.

Note that it is enough to prove this for the case  $M = 5$ . The larger cases follow, since they can simply choose the splits in the  $M = 5$  case.

**4.2.  $M = 4$ .** The SG-value for a pile of size  $n \geq 17$  appears to be  $n - 5$ . The smaller values can be read from the graph. This case is different from the larger ones due to  $n = 17$  where a 5-way split is needed to achieve the SG-value 0.

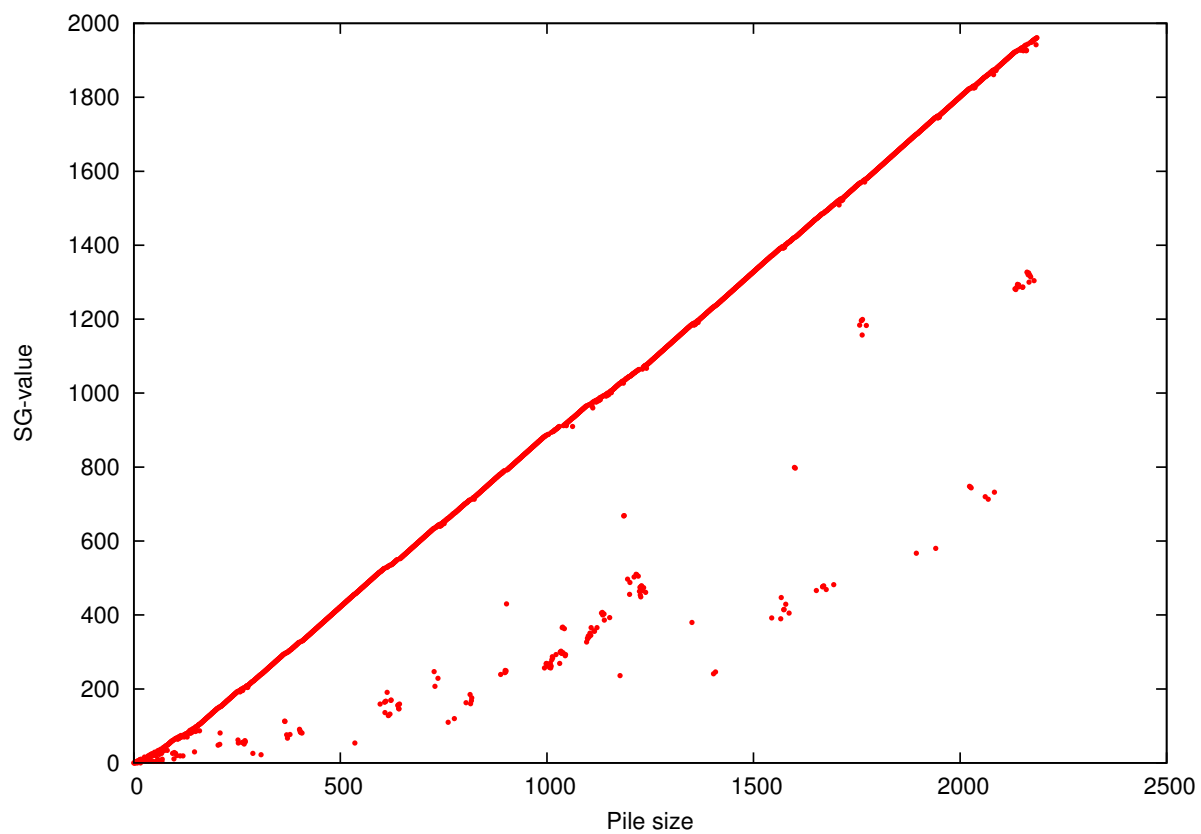


FIGURE 1. Any 3