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1 Basic

```
.1
     .vimrc
map jj <Esc>
v on
e sw=4 ts=4 sts=4 et nu sc hls cc=69
filet plugin indent on
m <F5> :!"./%<"<CR>
m <F6> :!"./%<" < input.txt<CR>
u FileType cpp no <F9> :!g++ % -o
\ -std=c++14 -O3 -Wall -Wextra
\ -Wshadow -Wno-unused-result<CR>
\langle \text{expr} \rangle \langle \text{silent} \rangle \langle \text{Home} \rangle \text{col}('.') = \\ \text{match}(\text{getline}('.'), '\S') + 1
m < silent > < Home > < C-O > < Home >
     Increase Stack Size
/stack resize
```

```
//stack resize
asm( "mov %0,%%esp\n" ::"g"(mem+10000000) );
//change esp to rsp if 64-bit system

//stack resize (linux)
#include <sys/resource.h>
void increase_stack_size() {
   const rlim_t ks = 64*1024*1024;
   struct rlimit rl;
   int res=getrlimit(RLIMIT_STACK, &rl);
   if(res==0){
      if(rl.rlim_cur=ks);
      res=setrlimit(RLIMIT_STACK, &rl);
   }
   res=setrlimit(RLIMIT_STACK, &rl);
}
```

1.3 digitDP

```
顷名思义,所谓的数位DP就是按照数字的个,十,百,千……
   位数进行的DP。
数位DP的题目有着非常明显的性质:
   询问[1,r]的区间内,有多少的数字满足某个性质
收法根据前缀和的思想,求出[0,1-1]和[0,r]中满足性质的数
   的个数,然后相减即可。
算法核心
 L dfs(int x, int pre, int bo, int limit);
一般需要以上参数(当然具体情况具体分析)
   x表示当前的数位 (一般都是从高位到低位)
   pre表示前一位的数字
   bo可以表示一些附加条件:是否有前项0,是否当前已经符
      合条件……
   limit 这个很重要! 它表示当前数位是否受到上一位的限
     制, 比较抽象, 举例说明
   如果上限是135,前两位已经是1和3了,现在到了个位,个
     位只能是5以下的数字
注: 如果当前受限,不能够记忆化,也不能返回记忆化的结果
为了避免多次调用时 每次上限不同 而导致的错
//http://acm.csie.org/ntujudge/view_code.php?id=106844
// Multiples
LL x;
int digit [100];
LL ten_pow[ 15 ];
bool ava[15];
LL \ dp \, [15][2][1000000];
LL dfs(int len, LL mod, bool bo, bool limit) {
```

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```
if ( len < 0 ) return mod == 0;
    if (!limit && dp[len][bo][mod]!= -1 ) return dp[
         len ] [bo] [mod];
    int up = limit? digit[len] : 9;
    LL ret = 0;
    for (int i = 0; i \le up; i++) if ( ava[i] || (i=0&&
        bo) ) {
         ret += dfs(len-1, (mod+ten_pow[len]*i)%x, bo
             &&(!i), limit&&(i=up));
    if( !limit ) dp[len][bo][mod] = ret;
    return ret;
LL solve (LL num) {
    int len = 0; digit [0] = 0;
    while( num ) {
        digit[len++] = num\%10;
        num = 10;
    return dfs (len -1, 0,1, 1);
bool check (LL num) {
    while ( num )
        if (!ava[ num%10 ] ) return false;
        num /= 10:
    return true;
int main() {
    LL A, B;
    cin>>x>>A>>B;
    ten_pow[0] = 1;
    mem(dp, -1);
    for (int i = 1; i < 15; i++)
        ten_pow[i] = (ten_pow[i-1]*10)\%x;
    string \ dig; \ cin>>dig;
    mem(ava, false);
    for (char c : dig) ava [c-'0'] = 1;
    if ( x <= 1000000 ) {
        cout << solve(B) - solve(A-1) << endl;
    }else {
        LL ans = 0;
        LL cur = 0;
         while ( cur < A ) cur += x;
         while ( cur <= B ) {
             i\dot{f} ( check(cur) ) ans++;
             cur += x;
        cout << ans << endl;
    }
}
```

1.4 DP(convex hull optimization)

```
// \texttt{http://codeforces.com/contest/311/problem/B}
struct line {
     LL slope,
                  inter;
     LL value(LL x) { return x*slope + inter; }
bool check(line x, line y, line z) {
     return (z.slope - y.slope) * (z.inter - x.inter )
              (z.slope - x.slope) * (z.inter - y.inter);
#define maxn 100005
int n, m, p;
 LL \ a \, [\, maxn ] \; , \ d \, [\, maxn ] \; , \ dp \, [\, 1\,0\,1\,] \, [\, maxn \,] \; , \ s \, [\, maxn \,] \; ; 
int main() {
     cin>\!\!> n>\!\!> m>\!\!> p;
     for(int i = 2; i <= n; ++i) {
          d[i] = getint();
          d\,[\,i\,] \; +\!\!= d\,[\,i\,-1\,]\,;
     for(int i = 1; i<=m; ++i) {
    int h; scanf("%d %lld", &h, a+i);
          a[i] -= d[h];
     sort(a+1,a+1+m);
```

```
\label{eq:formalized} \begin{array}{ll} \text{for}\,(\,int\ i \!=\! 1; i \!<\!\!=\! m;\, i \!+\! +\! )\ s\,[\,i\,] \ = \ a\,[\,i\,] \!+\! s\,[\,i\,\!-\! 1\,]\,; \end{array}
//start dp
for (int i=1; i < p; i++) {
      if(i = 1) {
            for (int j=1; j \le m; j++) dp[i][j] = j*a[j] - s
                  [j];
      }else {
           deque<line> dq;
           dq.pb(\{0, 0\});
            for (int j=1; j \le m; j++) {
                  while (dq. size() >= 2 \&\& dq[0]. value(-a[j]) > dq[1]. value(-a[j])) dq.
                        pop_front();
                  dp[i][j] = dq[0].value(-a[j]);
                  line newline{ j, dp[i-1][j]+s[j] };
                  \label{eq:while} \mbox{while} (\ \mbox{dq.size}() >= 2 \ \&\& \ \mbox{check}(\mbox{dq}[\mbox{dq}.
                        size()-2], dq.back(), newline)) dq
                        .pop_back();
                  dq.pb( newline );
                  if ( i==1 )
                       dp\,[\,i\,\,]\,[\,j\,] \;=\; j\,{}^*a\,[\,j\,] \;\; -\;\; s\,[\,j\,\,]\,;
                  }else {
                       LL mn = 0;
                        for (int \ k = 1; \ k < j; \ k+\!\!\!\!+\!\!\!\!+) \ \{
                             mn = min(mn, dp[i-1][k] + s[k]
                                    - a[j]*k);
                        dp[i][j] = mn + a[j]*j-s[j];
                        // apply convex hull optimization
                  dp[i][j] += a[j]*j - s[j];
      }
cout \ll dp[p][m] \ll endl;
```

1.5 simulated annealing

```
//http://mikucode.blogspot.tw/2015/03/algorithm.html
//尋找和所有點距離和最小的點
#include <cstdio>
#include <cstdlib>
#include <cmath>
#define F(n) Fi(i,n)
#define Fi(i,n) for (int i=0;i< n;i++)
#define N 1010
using namespace std;
int X[N], Y[N], n;
inline double pow2(double x){
    return x*x;
double check(double x, double y){
    double ans=0;
    F(n) ans += sqrt(pow2(x-X[i]) + pow2(y-Y[i]));
    return ans;
int main(){
     while (~scanf("%d",&n)) {
        F(n) scanf("%d%d",X+i,Y+i);
        double x=0,y=0,tx,ty,tans,l=10000,ans;
        ans=check(x,y);
        while(1>1e-4) {
            int tmp=rand();
             tx=x+l*cos(tmp); ty=y+l*sin(tmp);
            tans=check(tx,ty);
             if(tans < ans) ans=tans, x=tx, y=ty;
            else 1*=0.9;
        printf("%.9f\n",2*ans);
    }
}
//尋找兩個點使他們跟給定的四個點最小生成樹最小
#include <cstdio>
#include <cstdlib>
```

```
#include <cmath>
#include <algorithm>
#define F(n) Fi(i,n)
#define Fi(i,n) Fl(i,0,n)
#define Fl(i,l,n) for (int i=1; i<n; i++)
#define N 10
using namespace std;
int X[N], Y[N], n, F[N], e;
struct E{
    int a,b;
    double c;
G[N*2];
struct V{
    double x, y;
    V operator+(double 1){
         int tmp=rand();
         return (V) \{x+l*cos(tmp), y+l*sin(tmp)\};
}v[N];
int find(int x){
    return x = F[x]?x:F[x] = find(F[x]);
inline double pow2(double x){
    return x*x;
double check (V s1, V s2) {
    double ans=0;
    e=0;v[4]=s1,v[5]=s2;
    F(5)Fl(j,i+1,6)
         G[e++]=(E)\{i,j,sqrt(pow2(v[i].x-v[j].x)+pow2(v[i].x)\}
             i].y-v[j].y))};
    F(6)F[i]=i;
    sort(G,G+e,[](E a,E b)\{return a.c < b.c;\});
    F(e) {
         if (find (G[i].a)!=find (G[i].b)) {
             ans+=G[ i ] . c;
             F[find(G[i].a)]=find(G[i].b);
         }
    return ans;
int main() {
    scanf("%d",&n);
    while (n--) {
         F(4) scanf("%lf%lf",&v[i].x,&v[i].y);
         double ttans, tans, ans, 11=10000, 12;
         V s1=(V) \{0,0\}, s2=(V) \{0,0\}, ts1, ts2, tmp;
         ans=check(s1, s2);
         while (l1>1e-3) {
             12 = 10000;
             ts1=s1+l1;
             tans=check(ts1,s2);
             tmp=s2;
             while (12>1e-3) {
                  ts2=s2+l2;
                  ttans=check(ts1,ts2);
                  if (ttans<tans) tans=ttans, s2=ts2;
                  else 12*=0.9;
             if (tans<ans) ans=tans, s1=ts1;
             else 11 = 0.9, s2 = tmp;
         printf("%f \setminus n", 2*ans);
    }
}
```

2 Graph

2.1 HLD

```
//we can reference the problem Greatest graph
///http://acm.csie.org/ntujudge/problemdata/2582.pdf
//this template operate on edges
#define maxn 100005
struct segment_tree{
    #define right(x) x << 1 | 1
    #define left(x) x << 1
    int* arr;
    int m[4*maxn];</pre>
```

```
int tag[4*maxn];
     const int inf = 1e9;
    void init() {
           /\text{memset}(\text{tag}, -1, \text{sizeof}(\text{tag}));
          fill(tag, tag+4*maxn, inf);
     void pull(int ind) {
              m[ind] = min(m[right(ind)], m[left(ind)]);
     void push(int ind) {
         if(tag[ind] != inf) {
              tag[left(ind)] = min(tag[left(ind)], tag[
                   ind]);
              tag[right(ind)] = min(tag[right(ind)], tag[
                   ind]);
              m[left(ind)] = min( m[left(ind)], tag[left(
                   ind)])
              m[right(ind)] = min(m[right(ind)], tag[
                   right (ind)]);
              tag[ind] = inf;
         }
     /// \text{ root} \Rightarrow 1
     void build(int ind, int l, int r) {
          if( r - l == 1) {
              m[ind] = arr[l];
              return;
         \begin{array}{ll} {\bf int} \ \ {\rm mid} \ = \ (\ l{+}r \ ){>}{>}1; \end{array}
         build( left(ind), l, mid );
         build( right(ind), mid, r );
         pull(ind);
     int query_min(int ind, int L, int R, int ql, int qr
         if (L >= qr \mid \mid R <= ql) return 1e9;
         if (R \ll qr \&\& L >= ql) \{
              return m[ind];
         }
         push(ind);
         \quad \text{int} \ \operatorname{mid} = (L\!+\!R) >> 1;
         return min( query_min(left(ind), L, mid, ql, qr
              ), query_min(right(ind), mid, R, ql, qr));
     void modify(int ind, int L, int R, int ql, int qr,
         int x) {
         if (L >= qr \mid \mid R <= ql) return;
         if(R \le qr \&\& L = ql)
              m[ind] = min(m[ind], x);
              tag[ind] = min(tag[ind], x);
              return:
         push(ind);
         int mid = (L+R) >> 1;
         modify(left(ind), L, mid, ql, qr, x);
         modify(\,right\,(ind\,)\,,\,\,mid\,,\,\,R,\,\,ql\,,\,\,qr\,,\,\,x)\,;
         pull(ind);
    }
};
int seg_arr[maxn];
struct Tree{
    segment_tree seg;
    int n;
    struct Edge { int u, v, c; };
     vector<Edge> e;
     \begin{array}{c} void \ addEdge(int \ x, \ int \ y, \ int \ c) \ \{ \\ G[x].pb(\ SZ(e)\ ); \end{array} 
         G[y].pb(SZ(e));
         e.pb(Edge\{x, y, c\});
    int siz [maxn], max_son[maxn], pa[maxn], dep[maxn];
     /*size of subtree index of max_son, parent index >
          depth*/
    \verb|int link_top[maxn]|, \verb|link[maxn]|, \verb|timer|;
     /*chain top index in segtree ime stamp*/
    std::vector<int >G[maxn];
    void init(int N) {
```

```
n = N:
    e.clear();
    for(int i = 1; i <= n; i++) G[i].clear();
    timer=0;
    pa[1] = 1;
    dep[1] = 0;
void find_max_son(int x){
    siz[x]=1;
    \max_{son}[x] = -1;
    for(int e\_ind : G[x]) {
        int v = e[e\_ind].u == x ? e[e\_ind].v : e[
            e_ind].u
        find_max_son(v);
        if(max\_son[x] = -1 \mid \mid siz[v] > siz[max\_son]
             [x]])
            \max_{son}[x] = v;
        siz[x] += siz[v];
    }
void build_link(int x, int top){
    link[x] = timer++;/*記錄x點的時間戳*/
    link\_top[x] = top;
    if(\max_{x \in \mathbb{R}} |x| != -1)
        build_link( max_son[x], top);/*優先走訪最大
             孩子*/
    for(int e\_ind : G[x]) {
        int v = e[e\_ind].u == x ? e[e\_ind].v : e[
             e\_ind].u;
        if(v = pa[x])
            seg\_arr[link[x]] = e[e\_ind].c;
        if (v = \max_{son}[x] \mid | v = pa[x]) continue
        // edge from x \Rightarrow v
        build_link(v, v);
    }
inline int lca(int a, int b){
    /*求LCA, 可以在過程中對區間進行處理*/
    int ta=link_top[a],tb=link_top[b];
    while(ta != tb){
        if (dep[ta]<dep[tb]) {
            std::swap(ta,tb);
            std::swap(a,b);
        //interval [ link[ta], link[a] ]
        a = pa[ta];
        ta = link\_top[a];
    return dep[a] < dep[b] ? a:b;
}
int modify(int a, int b, int c){
    int ta=link_top[a], tb=link_top[b];
    while(ta != tb){
        if (dep[ta]<dep[tb]) {
            std::swap(ta,tb);
            std::swap(a,b);
        //interval [ link[ta], link[a] ]
        //same interval if operate on edges
        seg.modify(1, 1, n, link[ta], link[a]+1, c)
        a = pa[ta];
        ta = link\_top[a];
    //a, b are on the same chain
    if( a = b ); // interval [ link[a], link[a]], if operate on edges \Rightarrow no edge
        if (dep[a]>dep[b])
            swap(a,b)
        //interval [ link[a], link[b] ]
        // if operate on edges \Longrightarrow [ link[ max_son[
            a] ], link[b] ]
```

```
seg.modify(1, 1, n, link[max\_son[a]],
                   link[b]+1, c);
         }
    }
/*
     void modify(int a, int b, int c) {
         if ( a==b ) return;
         if( link\_top[a] = link\_top[b]) {
              if \, (\ dep \, [\, a\, ] \, > \, dep \, [\, b\, ] \ ) \ swap (a \, , \ b) \, ;
              seg.modify(1,\ 1,\ n,\ link\,[a]{+}1,\ link\,[b]{+}1,\ c
              assert(link[a]+1 = link[max_son[a]]);
              return;
         if(dep[link\_top[a]] < dep[link\_top[b]])
              swap(a, b);
         // a is the node with deeper link_top
         seg.modify(1, 1, n, link[link_top[a]], link[a]
               + 1. c):
         modify( pa[link_top[a]], b, c);
    /// Heavy Light Decomposition
    void HLD() {
          // root is indexed 1 here !
         find_max_son(1);
         build_link(1, 1);
}tree;
int main() {
    int T; cin>>T;
     while (T--) {
         int n,m;
         scanf("%d %d",&n, &m);
         int ans = 0;
         tree.init(n);
         for (int i=0; i< n-1; i++) {
              int a, b, c;
              {\rm scanf}\,(\,\text{``'}\!\text{d}\!\text{''}\!\text{d}\!\text{''}\!\text{d}\!\text{''},\!\&a,\!\&b,\!\&c\,)\;;
              //a--, b--; be careful here
              tree.addEdge(a, b, c);
              ans += c;
         tree.HLD();
         tree.seg.arr = seg_arr;
         tree.seg.build(1, 1, n);
    return 0;
2.2 Hungarian
```

```
// edge and node index starting from 0
  dfs version below
//complexity O ( V*E )
/* to do
#define ___maxNodes
num\_left = ?
struct Edge {
    int from:
    int to;
    int weight;
    Edge(int f, int t, int w):from(f), to(t), weight(w)
         {}
};
vector<int> G[__maxNodes]; /* G[i] 存储顶点 i 出发的边
    的编号 */
vector<Edge> edges;
int num_nodes;
int num_left;
int num_right;
int num_edges;
int matching[__maxNodes]; /* matching result */
int check [___maxNodes];
```

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```
bool dfs(int u) {
   for (auto i = G[u].begin(); i != G[u].end(); ++i) {
        // 对 u 的每个邻接点
       int v = edges[*i].to;
       if (!check[v]) { // 要求不在交替路中
           check[v] = true; // 放入交替路
           if (matching[v] == -1 || dfs(matching[v]))
              // 如果是未盖点,说明交替路为增广路,则
                  交换路径,并返回成功
              matching[v] = u;
              matching[u] = v;
              return true;
       }
   }
   return false; // 不存在增广路, 返回失败
int hungarian() {
   int ans = 0;
   memset(matching, -1, sizeof(matching));
   for (int u=0; u < num_left; ++u) {</pre>
       if (matching[u] = -1) {
           memset(check, 0, sizeof(check));
           if (dfs(u)) + ans;
       }
   return ans;
```

2.3 KM

```
// 最小帶權匹配~ km算法
//http://acm.csie.org/ntujudge/contest_view.php?id=836&
    contest\_id=449
#include <bits/stdc++.h>
using namespace std;
struct bipartite {
    #define maxn 602
    #define INF 0xfffffff
     int \ sx\left[maxn\right], \ sy\left[maxn\right], \ mat\left[maxn\right]\left[maxn\right]; 
    int x[maxn], y[maxn], link[maxn];
    int N, M, slack;
    int DFS(int t) {
        int tmp;
        sx[t] = 1;
        for (int i = 0; i < M; i++) {
             if (!sy[i]) {
                 tmp = x[t] + y[i] - mat[t][i];
                  if (tmp == 0) {
                      sy[i] = 1;
                      if (link[i] = -1 || DFS(link[i]))
                          \,l\,i\,n\,k\,\,[\,\,i\,\,]\,\,=\,\,t\,\,;
                          return 1;
                  else if (tmp < slack) slack = tmp;
             }
        }
        return 0;
    int KM() {
        for (int i = 0; i < N; i++) {
             x[i] = 0;
             for (int j = 0; j < M; j++) {
                  if (mat[i][j] > x[i]) x[i] = mat[i][j];
        for (int j = 0; j < M; j++) { y[j] = 0; }
        memset(link, -1, sizeof(link));
        for (int i = 0; i < N; i++) {
             while (1) {
                 memset(sx, 0, sizeof(sx));
                 memset(sy, 0, sizeof(sy));
                 slack = INF;
```

```
if (DFS(i)) break;
                 \quad \text{for (int } j \, = \, 0; \ j \, < \, N; \ j++) \, \, \{
                      if (sx[j]) x[j] = slack;
                 for (int j = 0; j < M; j++) {
                      if (sy[j]) y[j] += slack;
                 }
             }
        }
        int ans = 0;
        int cnt = 0;
        int t;
        for (int i = 0; i < M; i++)
        {
             t = link[i];
             if (t >= 0 \&\& mat[t][i] != -INF)
                 cnt ++:
                 ans += mat[t][i];
        // 最大權 : 沒有負號
        return -ans;
    void init(int n, int m) {
        N\,=\,n\,,\ M=m;
        for (int i = 0; i < N; i++)
            for (int j = 0; j < M; j++)
                 mat[i][j] = -INF;
    void input() {
        for(int i = 0; i < N; i++)
             for (int j = 0; j < M; j++) {
                 // fill in mat[i][j]
                 // stands for the weighting , but
                      negative sign !
                 // if 最大權 : 沒有負號
             }
}km;
int main(){
    int n,E;
    while (scanf("%d", &n) != EOF)
    {
        km.init(n, n);
        km.input();
        cout<< km.KM() <<endl;
    return 0;
```

2.4 Bi-vertex-connected Subgraph

```
#include <bits/stdc++.h>
using namespace std;
#ifdef DEBUG
   #define debug(...) printf(__VA_ARGS__)
   #define debug(...) (void)0
#endif
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair<long long, long long>
#define fi first
#define se second
#define all(x) (x).begin(),(x).end()
#define SZ(x) ((int)(x).size())
const int inf = 0x7ffffffff; //beware overflow
\#define mem(x, y) memset(x, (y), sizeof(x));
#define IOS ios_base::sync_with_stdio(0); cin.tie(0)
template<typename A, typename B>
ostream& operator <<(ostream &s, const pair<A,B> &p) {
    return s<<"("<<p.first<<","<<p.second<<")";
```

```
template<tvpename T>
ostream& operator <<(ostream &s, const vector<T> &c) {
    for (auto it : c) s << it << " ";
    s << "]";
                                                          }
   return s;
template<typename T>
ostream& operator << (ostream &o, const set<T> &st) {
    o << "{";
    for (auto it=st.begin(); it!=st.end(); it++) o << (
       it=st.begin()? "": ", ") << *it;
    return o << "}";
template < typename T1, typename T2>
ostream& operator << (ostream &o, const map<T1, T2> &mp
   ) {
   o << "{";
    for (auto it=mp.begin(); it!=mp.end(); it++) {
       o << (it=mp. begin()?"":", ") << it->fi << ":"
            << it->se;
    o << "}";
    return o;
     regard every vbcc as a set of edges
/** needed for tarjan **/
#define maxn 100005
\#define maxm 100005
int n, m;
struct Edge{int s, t;};
vector<Edge> edge;
int dfn[maxn], low[maxn];
stack<int> st;
bool vis[maxn];
int Time;
bool vis_e [maxm];
int bcnt, vbb[maxm];
vector<int> vb[maxm];
vector<int> G[maxn];
void tarjan(int s){
    dfn[s] = low[s] = ++Time;
    vis[s] = true;
    for (int e_ind : G[s]) {
        if (!vis_e[e_ind]) {
            vis_e[e_ind] = true; st.push(e_ind);
            int to = edge[e\_ind].s + edge[e\_ind].t - s;
            if (! vis[to]) {
                tarjan(to);
                low[s] = min(low[s], low[to]);
                if(low[to] >= dfn[s]){
                    vb[bcnt].clear();
                    while (1) {
                        int t = st.top(); st.pop();
                        vbb[t] = bcnt;
                        vb[bcnt].push_back(t);
                        if(t = e_ind) break;
                    bcnt++;
                }
            }else
                low[s] = min(low[s], dfn[to]);
        }
   }
void init_tarjan() {
   mem(vis, false); mem(vis_e, false);
    Time = bcnt = 0; edge.clear();
    for (int i = 1; i \le n; i++) G[i]. clear ();
int main() {
    cin >> n >> m;
    init_tarjan();
    edge.push\_back(Edge\{a,b\});
```

```
G[a].push_back((int)edge.size()-1);
G[b].push_back((int)edge.size()-1);
}
tarjan(1);
```

2.5 Bi-edge-connected Subgraph

```
/** needed for tarjan **/
#define maxn 100005
#define maxm 100005
int n, m;
int dfn[maxn], low[maxn];
stack<int> st;
int Time;
int bcnt;
vector<int> G[maxn];
bool in_cyc[maxn];
void tarjan(int s, int p){
    dfn[s] = low[s] = ++Time;
    st.push(s);
    for (int to : G[s]) if ( to != p ){
        if (!dfn[to]) {
            tarjan(to, s);
            low[s] = min(low[s], low[to]);
            if(low[to] > dfn[s]) {
                // is cut_edge
                 // pop stack 的過程也可以寫在這
                 // 但最後(after tarjan)還要多判stack
                     not empty的情况
                 if ( low [to] > dfn [s]) {
                 in\_cyc[bcnt] = st.top()!=to;
                 while (1) {
                     int t = st.top(); st.pop();
                     id[t] = bcnt;
                     if (t == to) break;
                 bcnt++;
            }
            }
        }else
            low[s] = min(low[s], dfn[to]);
    if(low[s] = dfn[s]){
        in\_cyc[bcnt] = st.top()!=s;
        while (1) {
            int t = st.top(); st.pop();
            id[t] = bcnt;
            if(t = s) break;
        bcnt++;
    }
void init_tarjan() {
    Time = bcnt = 0;
int main() {
  cin >> n >> m;
  init_tarjan();
  for (int i = 0; i < m; i++) {
        int a, b; scanf("%d %d", &a, &b);
        G[a].pb(b), G[b].pb(a);
  mem( in_cyc , false);
  tarjan(1, 1);
}
```

2.6 SCC

```
#include <bits/stdc++.h>
using namespace std;
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair<long long, long long>
#define fi first
#define se second
const int inf = 1e9;
const LL INF = 1e18;
const int mod = 1e9 + 7;
#define maxn 100050
int n, m;
vector<int> g[maxn];
stack<int> Stack;
int scnt, Time;
int belong[maxn], dfn[maxn], low[maxn], indegree[maxn];
bool instack [maxn];
void input(){
  cin >> n >> m;
  for (int i = 0; i < m; i++){}
    int a, b; scanf("%d%d", &a, &b);
    g[a].pb(b);
void init() {
  scnt = Time = 0;
  for (int i = 1; i <= n; i++)
    g[i].clear();
  while (!Stack.empty()) Stack.pop();
  memset(indegree, 0, sizeof(indegree));
  memset(dfn, 0, sizeof(dfn));
  memset(instack, false, sizeof(instack));
void dfs(int u) {
  dfn\left[u\right] \ = \ low\left[u\right] \ = \ +\!\!\!+\!\! Time;
  Stack.push(u); instack[u] = true;
  for(int v : g[u]) {
    if (!dfn[v]) {
      dfs(v);
      low[u] = min(low[u], low[v]);
    else if (instack [v])
      low[u] = min(low[u], dfn[v]);
  if(low[u] = dfn[u]) {
    scnt++;
    int tp;
      tp = Stack.top(); Stack.pop();
      instack[tp] = false;
      belong[tp] = scnt;
    }  while (tp != u);
void tarjan() {
  for (int i = 1; i \le n; i++)
    if (!dfn[i])
      dfs(i);
int main(){
  int T; cin >> T;
  while (T--) {
    init();
    input();
    tarjan();
    for (int i = 1; i \le n; i++) {
      for(int v : g[i]) {
  if(belong[v] != belong[i])
          indegree [belong [v]]++;
      }
    LL ans = 0;
    for (int i = 1; i \le scnt; i++)
```

if (!indegree[i]) ans++;

```
cout << ans << endl;
}
return 0;
}</pre>
```

2.7 Steiner Tree(PECaveros)

```
// Minimum Steiner Tree
^{'}// O(V 3^{^{'}}T + V^{^{'}}2 2^{^{'}}T)
struct SteinerTree{
#define V 33
#define T 8
#define INF 1023456789
  int n , dst [ V ][ V ] , dp[ 1 << T ][ V ] , tdst [ V
   \begin{tabular}{ll} \bf void & init ( & int & \_n \ ) \{ \end{tabular} 
    n\,=\,\underline{}\,n;
     for (int i = 0 ; i < n ; i ++){
      for ( int j = 0 ; j < n ; j ++ )

dst[i][j] = INF;

dst[i][i] = 0;
    }
  void add_edge( int ui , int vi , int wi ){
    void shortest_path(){
     for (int k = 0 ; k < n ; k ++ )
       for (int i = 0 ; i < n ; i ++)
         for(int j = 0 ; j < n ; j +++)
           ][ j ]);
  int solve( const vector<int>& ter ){
    int t = (int) ter.size();
    for( int i = 0 ; i < (1 << t) ; i ++)
for( int j = 0 ; j < n ; j ++)
         dp[i][j] = INF;
     for (int i = 0 ; i < n ; i ++)
      dp[0][i] = 0;
     for ( int msk = 1 ; msk < (1 << t) ; msk ++ ){}
       i\dot{f} ( msk == (msk \& (-msk)) ){
         int who = \underline{\hspace{1cm}} \lg(msk);
         for (int i = 0); i < n; i ++)
           dp[msk][i] = dst[ter[who]][i];
       for (int i = 0 ; i < n ; i ++)
         for(int submsk = (msk - 1) \& msk ; submsk ;
                  submsk = (submsk - 1) \& msk)
             dp[msk][i] = min(dp[msk][i],
                                    dp[ submsk ][ i ] +
   dp[ msk ^ submsk
                                         dp[ msk
                                         ][ i ]);
       for (int i = 0 ; i < n ; i ++) {
         tdst[i] = INF;
         for (int j = 0 ; j < n ; j ++)
           tdst[i] = min(tdst[i])
                             dp[msk][j] + dst[j][i]
                                  ] );
       for (int i = 0 ; i < n ; i ++)
         dp[msk][i] = tdst[i];
     int ans = INF;
     for (int i = 0 ; i < n ; i ++)
       ans = \min( ans , dp[ ( 1 << t ) - 1 ][ i ] );
     return ans;
} solver;
```

2.8 Edmond's Matching Algorithm

2.9 Tree Decomposition

```
//codeforces Digit Tree
//http://codeforces.com/problemset/problem/715/C
#include <bits/stdc++.h>
using namespace std;
#ifdef DEBUG
   #define debug(...) printf(__VA_ARGS__)
#else
    #define debug(...) (void)0
#endif
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair<long long, long long>
#define fi first
#define se second
#define all(x) (x).begin(),(x).end()
#define SZ(x) ((int)(x).size())
const int inf = 0x7ffffffff; //beware overflow
\#define mem(x, y) memset(x, (y), sizeof(x));
#define IOS ios_base::sync_with_stdio(0); cin.tie(0)
template {<} typename \ A, \ typename \ B {>}
ostream& operator <<(ostream &s, const pair<A,B> &p) {
    return s<<"("<<p.first<<","<<p.second<<")";
template<typename T>
ostream& operator <<(ostream &s, const vector<T> &c) {
    s \ll "[";
    for (auto it : c) s << it << " ";
    s << "]";
    return s;
template<typename T>
ostream& operator << (ostream &o, const set<T> &st) {
    o << "{";
    for (auto it=st.begin(); it!=st.end(); it++) o << (
    it==st.begin() ? "" : ", ") << *it;</pre>
    return o << "}";
template<typename T1, typename T2>
ostream& operator << (ostream &o, const map<T1, T2> &mp
    ) {
    o << "{";
     for (auto it=mp.begin(); it!=mp.end(); it++) \{
        o << (it=mp.begin()?"":", ") << it->fi << ":"
             << it->se:
    o << "}";
    return o;
typedef long long 11;
bool isprime [100005];
vector<LL> primes;
LL M, PHI;
#define MOD M
ll modpow(ll a, ll b) {
 11 r = 1;
  while(b) {
    if(b\&1) r=(r*a)\%MOD;
    a=(a*a)MOD;
    b >>= 1;
  }
  return r;
void Sieve(int n) {
  memset(isprime, 1, sizeof(isprime));
  isprime[1] = false;
  for(int i = 2; i <= n; i++) {
    if(isprime[i]) {
      primes.pb(i);
      for (int j = 2*i; j \le n; j += i)
        isprime[j] = false;
  }
LL phi(LL n) {
  11 \text{ num} = 1; 11 \text{ num} 2 = n;
  \quad \quad for(\,l\,l\ i\,=\,0\,;\ primes\,[\,i\,]\,^*\,primes\,[\,i\,]\,<=\,n\,;\ i++)\ \{
```

```
if(n\%primes[i]==0) {
      num2/=primes[i];
      num^* = (primes [i] - 1);
    while (n\%primes[i]==0) {
      n/=primes[i];
  if(n>1) {
    num2/=n; num*=(n-1);
  n = 1;
  num *= num2;
  return num;
ll inv(ll a) {
  return modpow(a, PHI-1);
#define maxn 100005
struct edge{
    int u, v, dig;
    int no(int x) {
         }:
vector<edge> e;
vector<int> G[maxn];
LL n, ans;
bool vis[maxn];
int sz [maxn], dep [maxn];
LL tenPow[maxn];
int dfs(int u, int p, int d) {
    sz[u] = 1;
    dep[u] = d;
    for (int eind : G[u] ) {
         int v = e[eind] \cdot no(u);
if (v == p \mid \mid vis[v]) continue;
         sz[u] += dfs(v, u, d+1);
    return sz[u];
int findCenter(int u, int p, int treesize) {
    for (int eind : G[u] ) {
         int v = e[eind].no(u);
         if (v = p \mid | vis[v]) continue;
         if (sz[v]*2 > treesize)
             return findCenter( v, u, treesize);
    return u;
}
LL up[maxn], down[maxn];
int belong[maxn];
map<LL, LL> tot;
vector < map < LL, LL > vec;
vector<int> pt;
void calc(int u, int p, int b, int d) {
    pt.pb( u );
    belong[u] = b;
    dep[u] = d;
    int \ id = \ find\_if(\ all(G[u]) \ , [u,p](int \ x) \ \{ \ return
          e[x].no(u) == p; ) - G[u].begin();
    down[u] = (down[p]*10 + e[G[u][id]].dig)M;
    up[u] = (tenPow[d-1]*e[G[u][id]].dig + up[p])
        %M:
    for(int eind : G[u]) {
         int v = e[eind].no(u);
         if(\ vis[v]\ ||\ v = p\ )\ continue;
         calc(v, u, b, d+1);
    vec[b][up[u]]++;
    tot[ up[u] ]++;
LL solve(int cent) {
    //cent is the root now
    vector<int> L;
    \quad \quad \text{for} \, (\, \text{int eind } : \, G[\, \text{cent} \,] \,) \  \, \{ \,
```

```
int v = e[eind].no(cent);
          if( !vis[v]) {
               L.pb( v );
     vec.clear();
     vec.resize(SZ(L), {});
     tot.clear();
     up[cent] = down[cent] = 0;
     dep[cent] = 0;
     pt.clear();
     for (int i = 0; i < SZ(L); i++)
          calc( L[i], cent, i, 1);
     LL ret = 0;
     for(int u : pt) {
           LL \ tmp \ = \ (\ -down \ [\ u] +\! M) \%\! M; 
          tmp = (tmp*inv(tenPow[dep[u]]))%M;
          ret += tot [ tmp ] - vec [ belong [u] ] [ tmp ];
     assert( (LL)count_if(all(pt), [] (int x) { return
     \begin{array}{c} \operatorname{up}[x] = 0; \; \} \; ) = \operatorname{tot}[0]); \\ \operatorname{LL} \; \operatorname{tmp} = \operatorname{tot}[0] \; + \; (\operatorname{LL})\operatorname{count\_if}(\operatorname{all}(\operatorname{pt}), \; [] \; (\operatorname{int} \; x) \end{array}
           \{ return down[x] = 0; \} );
     debug("%lld\n", tmp);
     return ret+tmp;
void solveAll(int node) {
     dfs(node, -1, 0);
     int cent = findCenter(node, -1, sz[node]);
     ans += solve ( cent );
     \label{eq:debug_norm} \mbox{debug("%d \%lld\n", cent, ans);}
     vis[cent] = true;
     for (int eind : G[cent] ) {
          int v = e[eind].no(cent);
          if( vis[v] ) continue;
          solveAll(v);
     }
int main() {
     cin>>n>>M;
  Sieve ( 100000 );
     PHI = phi(M);
     for(int i = 0; i < n-1; i++) {
           int a, b, c; scanf("%d %d %d", &a, &b, &c);
          G[a].pb(SZ(e)); G[b].pb(SZ(e));
          e.pb( edge{a, b, c} );
     //init
     tenPow[0] = 1;
     for(int i = 1; i < maxn; i++) tenPow[i] = (tenPow[i])
          -1]*10)%M;
     ans = 0;
     mem( vis, false);
     solveAll(0);
     cout << ans << endl;
}
```

2.10 Tree Longest Path

```
/** codeforces 592D - Super M **/
#include <bits/stdc++.h>
using namespace std;
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair < long long, long long>
#define fi first
#define se second
const int inf = 1e9;
const LL INF = 1e18;
const int mod = 1e9 + 7;
#define maxn 123460
int n, m;
vector<int> g[maxn];
```

```
bool is [maxn];
int dep[maxn], R, max_depth, A;
int cnt[maxn], parent[maxn];
bool dfs(int u, int par = 0){
  parent[u] = par;
  \operatorname{dep}\left[u\right] = \operatorname{dep}\left[\operatorname{par}\right] \,+\, 1;
   if(dep[u] > max_depth \&\& is[u])
    \max_{\mathbf{depth}} = \operatorname{dep}[\mathbf{u}], \ \mathbf{R} = \mathbf{u};
  bool ret = is[u];
   for(int v : g[u])
     i\dot{f}(v != par)
       ret = dfs(v, u);
  if (ret) A++;
  return ret;
int find_center(int start) {
  R = start; dep[0] = -1; max_depth = 0;
  dfs(start);
  \max_{\text{depth}} = 0; \text{ dep}[R] = -1;
   dfs(R, R);
   int ret = R, d = max_depth/2;
   while(d>0)
     d--;
     ret = parent[ret];
  return ret;
int S, dis, max_length;
bool \ dfs1(int \ u, \ int \ par = 0) \ \{
  dep[u] = dep[par] + 1;
  if (is [u])
     if (dep[u] > max_length)
       \max_{\text{length}} = \text{dep}[u], S = u;
     else if (dep[u] = max_length && u < S)
       S = u;
  bool c = false;
   \quad \quad for (int \ v \ : \ g[u])
     if ( v != par )
       dfs1(v, u);
int main(){
   cin >> n >> m;
   for (int i = 0; i < n-1; i++){
     int a, b; scanf("%d%d",&a, &b);
     g[a].pb(b), g[b].pb(a);
  memset(is, false, sizeof(is));
  int tmp;
   for (int i = 0; i < m; i++){
     cin>>tmp; is [tmp] = true;
  int C = find_center(tmp);
  dep[0] = -1; S = inf; dis = (max_depth+1)/2;
  // distance(center, any other node) <= (longestpath +
        1) / 2
   dfs1(C);
  if(max\_depth \& 1)
     dfs1(parent[C]);
  cout << S << endl << A-2-max_depth << endl;
  return 0;
}
```

3 Flow

3.1 Dinic Maxflow

```
//http://acm.csie.org/ntujudge/problem.php?id=2581
//French Fries Festival
//dinic runs in O( V^2*E )
#include <bits/stdc++.h>
using namespace std;
#ifdef DEBUG
    #define debug(...) printf(_VA_ARCS_)
#else
    #define debug(...) (void)0
#endif
```

```
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair<long long, long long>
#define fi first
#define se second
\#define all(x) (x).begin(),(x).end()
#define SZ(x) ((int)(x).size())
const int inf = 0x7ffffffff; //beware overflow
\#define mem(x, y) memset(x, (y), sizeof(x));
\#define IOS ios_base::sync_with_stdio(0); cin.tie(0)
template<typename A, typename B>
template<typename T>
ostream& operator <<(ostream &s, const vector<T> &c) {
    s << "[ ";
    for (auto it : c) s \ll it \ll ";
    s\,<<\,"\,|\,"\,;
    return s;
template<typename T>
ostream& operator << (ostream &o, const set<T> &st) {
    o << "{";
     for (auto it=st.begin(); it!=st.end(); it++) o << (
        it=st.begin() ? "" : ", ") << *it;
    return o << "}";
template<typename T1, typename T2>
ostream& operator << (ostream &o, const map<T1, T2> &mp
   ) {
    o << "{";
    \quad \text{for (auto it=mp.begin(); it!=mp.end(); it++) } \{
        o << (it=mp.begin()?"":", ") << it->fi << ":"
            << it->se;
   o << "}";
   return o:
#define maxn 500
struct Edge{ int to, cap, rev; };
struct Dinic{
    vector < Edge > G[maxn];
    int dis[maxn], iter[maxn];
    void init(int n) {
        //zero based
      for (int i = 0; i < n; i++) G[i]. clear ();
    void addEdge(int from, int to, int cap) {
        vector < Edge > :: iterator it;
        if( ( it=find_if( all(G[from]), [to](Edge& e) {
             return e.to == to; } )) != G[from].end() )
            it -> cap += cap;
            return;
      G[from].pb(Edge\{to, cap, (int)G[to].size()\});
      G[to].pb(Edge\{from, 0, (int)G[from].size()-1\});
        //if undirected 0 will be cap
    bool bfs(int s, int t) {
      memset(dis, -1, sizeof(dis));
      queue<int> que;
      que.push(s); dis[s] = 0;
      while (!que.empty()) {
        int tp = que.front(); que.pop();
        for (Edge &e : G[tp])
          if(e.cap > 0 \&\& dis[e.to] = -1)
            dis[e.to] = dis[tp] + 1, que.push(e.to);
        }
      return dis[t] != -1;
    int dfs(int v, int t, int f) {
      if(v = t) return f;
      for(int \& i = iter[v]; i < G[v].size(); i++) {
        Edge &e = G[v][i];
        if(e.cap > 0 \&\& dis[v] < dis[e.to]) {
          int d = dfs(e.to, t, min(f, e.cap));
```

```
if(d > 0) {
              e.cap -= d;
              G[e.to][e.rev].cap += d;
              f += d:
              return d;
         }
       }
       return 0;
     int maxFlow(int s, int t) {
       int ret = 0;
       while (bfs(s, t))
         memset(iter, 0, sizeof(iter));
         while ((f = dfs(s, t, inf)) > 0)
           ret += f:
       return ret;
}dinic, dinic2;
void solve() {
    // flow problem with lower bounds;
    int s = 0, t = n+2, ss = n+3, tt = n+4;
     dinic.init (n+5);
     dinic.addEdge(s, 1, k);
    {\tt dinic.addEdge(n+1,\ t\,,\ k)\,;}
    int slb = 0;
     while (m--) {
         int l, r, a, b; scanf("%d %d %d %d", &l, &r, &a
              , &b);
         slb += a;
         r++;
         dinic.addEdge(l, r, b-a);
         \label{eq:dinic.addEdge(ss, r, a);} \\ \text{dinic.addEdge(ss, r, a);}
         \label{eq:dinic.addEdge(l, tt, a);} \\ \text{dinic.addEdge(l, tt, a);}
    dinic2 = dinic;
     dinic.addEdge(t, s, k);
    int f1 = dinic.maxFlow(ss, tt);
     if (!all_of (all (dinic.G[ss]), [] (Edge x) { return
         x. cap = 0;  \} ) ) { puts("-1"); return; }
     dinic2.addEdge(ss, s, 1e9);
     dinic2.addEdge(t, tt, 1e9);
    int f2 = dinic2.maxFlow(ss, tt);
    // maxflow in current graph is f2 - slb
    printf("%d\n", (f2 - slb)*n);
int main() {
    \begin{array}{ll} \textbf{int} & t \; ; cin >\!\!> t \; ; \end{array}
     while(t--)
         solve();
```

4 Data Structure

4.1 Disjoint Set

```
struct Disjoint_set {
    #define MAX_N 500005
    // define MAX_N
    int pa[MAX_N], Rank[MAX_N];
    int sz[MAX_N];
    void init_union_find(int V) {
        for(int i=0; i<V; i++) {
            pa[i] = i;
            Rank[i] = 0;
            sz[i] = 1;
        }
    }
    int find(int x) {</pre>
```

```
return x == pa[x] ? x : pa[x] = find(pa[x]);
    int unite(int x, int y) {
        x = find(x), y = find(y);
        int S = sz[x] + sz[y];
         if(x != y){
             if(Rank[x] < Rank[y]) {
                 pa[x] = y;
                 sz[y]=S;
                 return y;
             }
             else{
                 pa[y] = x;
                 \operatorname{sz}[x] = S;
                 if(Rank[x] = Rank[y]) Rank[x] ++;
                 return x;
        }
    bool same(int x, int y) {
        return find(x) = find(y);
}
```

$4.2 \quad \text{Djs} + \text{Seg}$

```
// demo \Longrightarrow undo djs + segtree with offline // this program doesn't consider the problem of
     overflowing varaible ans
// http://acm.csie.org/ntujudge/view_code.php?id
     =108190&contest_id=472
#define maxn 100005
#define maxm 500005
//can be used to solve dynamic connectivity problem
//can be used with segment tree \Longrightarrow offline
struct DisjointSet {
 // save() is like recursive
  // undo() is like return
  int n, fa[maxn], sz[maxn];
vector<pair<int*,int>>> h;
  vector<int> sp;
  int ans;
  void init(int tn) {
    ans = 0;
     n=tn:
     for (int i=0; i<n; i++) {
       fa[i]=i;
       sz[i]=1;
     sp.clear(); h.clear();
  void assign(int *k, int v) {
     h.\, pb(\{k\,,\ ^{*}k\})\,;
     *k=v;
  void save() { sp.pb(SZ(h)); }
  void undo() {
     assert(!sp.empty());
     int last=sp.back(); sp.pop_back();
     while (SZ(h)!=last) {
       auto x=h.back(); h.pop_back();
        *x.fi=x.se;
     }
  int f(int x) {
     while (fa[x]!=x) x=fa[x];
     return x;
  void uni(int x, int y) {
     x=f(x); y=f(y);
     if (x=y) return ;
     \begin{array}{ll} \textbf{if} & (\,sz\,[\,x]\!<\!sz\,[\,y\,]\,) & swap\,(\,x\,,\,\,y\,)\,; \end{array}
     //nans stands for new answer
     int t = sz[x] + sz[y];
     int nans = ans - (sz[x]*sz[x]-sz[x]) - (sz[y]*sz[y]
          ]-sz[y]) + t*t-t;
     \operatorname{assign}(\&\operatorname{sz}\left[\,x\,\right]\,,\ \operatorname{sz}\left[\,x\right]\!+\!\operatorname{sz}\left[\,y\,\right]\,)\;;
     assign(&fa[y], x);
```

```
assign(&ans, nans);
}djs;
int n, m;
map<int , int> ma[maxn];
vector<pii> seg[4*maxm];
LL ans [maxm];
void add(int ql, int qr, int a, int b, int id=1, int l
     =0, int r=m) {    if( qr <= l || ql >= r ) return;    if( l >= ql && r <= qr ) {
           seg[id].pb(mp(a, b));
          return ;
     int mid = (l+r) >> 1;
     add( ql, qr, a, b, id*2, l, mid);
add( ql, qr, a, b, id*2+1, mid, r);
void dfs(int u=1, int l=0, int r=m) {
     djs.save();
     \quad \quad \text{for} \, (\, \text{pii} \ v \ : \ \text{seg} \, [\, u\, ] \ ) \quad djs.\, uni \, (\ v.\, \text{fi} \ , \ v.\, \text{se} \ ) \, ;
     if(r-l > 1)
           int mid = (l+r) >> 1;
           dfs(u*2, 1, mid);
           dfs(u*2+1, mid, r);
     }else {
           // do sth here
           ans[l] = djs.ans;
     djs.undo();
int main() {
    scanf("%d %d", &n, &m);
     for (int i = 0; i < m; i++) {
           int a, b; scanf("%d %d",&a, &b);
          a--, b--; if (b < a) swap(a, b);
           if (ma[a].count(b)) {
                add(ma[a][b], i, a, b);
                ma[a].erase(b);
           else ma[a][b] = i;
     for(int i = 0; i < n; i++) if(!ma[i].empty()) 
           for(auto p : ma[i])
                add( p.se, m, i, p.fi);
     djs.init(n);
     dfs();
     \label{eq:for_int} \begin{array}{lll} \mbox{for(int $i = 0$; $i < m$; $i++)$ $printf("\%lld\n"$, $ans[i])$;} \end{array}
```

4.3 Sparse Table

```
//codeforces 689D
#define maxn 200005
template< typename T, typename Cmp = less<T>>
struct RMQ {
    T d[maxn][20];
     Cmp cmp;
     int w[maxn], sz;
     void init (const T *a, int n) {
          int i, j;
          for (w[0] = -1, i = 1; i \le n; ++i) w[i] = (i & 
                 (i-1)) ? w[i-1] : w[i-1]+1;
           \mbox{for } (sz = n, \ i = 0; \ i < n; \ +\!\!+\!\! i) \ d[\,i\,][\,0\,] = a[\,i\,]; 
          for (j = 1; (1 \ll j) \ll n; ++j) {
                for (i = 0; i + (1 \ll j) \ll n; ++i) {
                    \begin{array}{l} d[i][j] = cmp(d[i][j-1], \ d[i+(1<<\\(j-1))][j-1]) \ ? \ d[i][j-1] \ : \ d\\[i+(1<<(j-1))][j-1]; \end{array}
          }
     }
```

```
// index of a [l \dots r] const T &query(int l, int r) const {
        int x = w[r - l + 1];
        [1][x] : d[r - (1 \ll x) + 1][x];
};
int a [maxn], b [maxn];
int n:
RMQ < int > s;
RMQ<int, greater<int>> t;
int main() {
    cin>>n;
    for(int i = 0; i < n; i++) scanf("%d", &a[i]);
     for(int \ i = 0; \ i < n; \ i++) \ scanf("%d", \&b[i]); 
    s.init(b, n);
    t.init(a, n);
    int \ c\,,\ d\,;
    LL ans = 0;
    for (int i=0; i< n; i++) {
        if (a[i] > b[i]) continue;
        int ub = n+1, lb = i;
        while (ub-lb>1) {
            int mid = (ub+lb) >> 1;
             if(t.query(i, mid-1) - s.query(i, mid-1) >
                  0) ub = mid;
            else lb = mid;
        int up = ub:
        ub = n+1, lb = i;
         while (ub-lb>1) {
            int mid = (ub+lb) >> 1;
            if( t.query(i, mid-1) - s.query(i, mid-1)
                >= 0) ub = mid;
            else lb = mid;
        int down = ub;
        ans += up-down;
    cout << ans << endl;
    return 0;
}
```

4.4 Link Cut Tree

```
//\,https://\,github.com/\,yzgysjr/\!ACM-ICPC-\,Templates/\,blob/
    master/Data\%20Structure/Link\%20Cut\%20Tree.cpp
struct node { int rev; node *pre, *ch[2]; } base[MAXN],
     nil, *null;
typedef node *tree;
#define isRoot(x) (x->pre->ch[0] != x && x->pre->ch[1]
    != x)
#define isRight(x) (x->pre->ch[1] == x)
inline void MakeRev(tree t) { if (t != null) { t->rev
    \hat{} = 1; \text{ swap}(t->\text{ch}[0], t->\text{ch}[1]); \}
inline void PushDown(tree t) { if (t->rev) { MakeRev(t
    ->ch[0]); MakeRev(t->ch[1]); t->rev = 0; }
inline void Rotate(tree x) {
  tree y = x->pre; PushDown(y); PushDown(x);
  int d = isRight(x);
  if (!isRoot(y)) y->pre->ch[isRight(y)] = x; x->pre =
  if ((y->ch[d] = x->ch[!d]) != null) y->ch[d]->pre = y
  x->ch[!d] = y; y->pre = x; Update(y);
inline void Splay(tree x) {
  PushDown(x)\,;\;\; \textbf{for}\;\; (\,t\, ree\;\; y\,;\;\; !\, isRoot(x)\,;\;\; Rotate(x)\,) \;\; \{
    y = x-pre; if (!isRoot(y)) Rotate(isRight(x))!= isRight(y) ? x : y);
  } Update(x);
inline void Splay(tree x, tree to) {
```

```
Update(x);
inline tree Access(tree t) {
  tree \ last = null; \ for \ (; \ t \ != \ null; \ last = t \, , \ t = t \text{-}>
      pre) Splay(t), t->ch[1] = last, Update(t);
  return last:
inline void MakeRoot(tree t) { Access(t); Splay(t);
    MakeRev(t); }
inline tree FindRoot(tree t) { Access(t); Splay(t);
     tree last = null;
  for ( ; t \stackrel{!}{=} null; last = t, t = t->ch[0]) PushDown(t)
       ); Splay(last); return last;
inline void Join(tree x, tree y) { MakeRoot(y); y->pre
    = x; 
inline void Cut(tree t) {Access(t); Splay(t); t->ch
    [0]->pre = null; t->ch[0] = null; Update(t);
inline void Cut(tree x, tree y) {
  tree upper = (Access(x), Access(y));
  if (upper == x) { Splay(x); y->pre = null; x->ch[1] =
        null; Update(x); }
  else if (upper == y) \{ Access(x); Splay(y); x->pre =
  null; y->ch[1] = null; Update(y); }
else assert(0); // 'impossible to happen'
inline int Query(tree a, tree b) { // 'query the cost
    in path a <-> b, lca inclusive
  Access(a); tree c = Access(b); // c is lca
  int v1 = c - ch[1] - maxCost; Access(a);
  int v2 = c - sch[1] - smaxCost;
  return \max(\max(v1, v2), c->cost);
void Init() {
  null = &nil; null -> ch[0] = null -> ch[1] = null -> pre =
       null; null->rev = 0;
  Rep(i, 1, N)  { node &n = base[i]; n.rev = 0; n.pre = 0
       n.ch[0] = n.ch[1] = null;
//compressed version
// \texttt{http://trinklee.blog.163.com/blog/static}
     /238158060201521101957375/
const int N=30010;
int n, fa [N], son [N][2], val[N], siz[N], stmp, rev[N];
\#define swap(a,b) (stmp=a,a=b,b=stmp)
void pu(int t) \{ siz[t] = siz[son[t][0]] + siz[son[t][1]] + 1; \}
void pd(int t) \{rev[t]? rev[t]=0, rev[son[t][0]]^=1, rev[
    son[t][1]]^=1, swap(son[t][0], son[t][1]), 1:1;
bool \operatorname{nr}(\operatorname{int} t) \{ \operatorname{return} \operatorname{son} [\operatorname{fa}[t]][0] == t \mid | \operatorname{son} [\operatorname{fa}[t]][1] == t \}
    t;
void rtt(int t,int f=0,bool p=0){
    p=son[f=fa[t]][1]==t,
    fa[t] = fa[f], nr(f)? son[fa[f]][son[fa[f]][1] == f] = t:1,
    (son[f][p]=son[t][!p])?fa[son[f][p]]=f:1,
    pu(son[fa[f]=t][!p]=f);
void pv(int t){if(nr(t))pv(fa[t]);pd(t);}
void splay(int t, int f=0){
     for (pv(t); nr(t); rtt(t)) nr(f=fa[t])?
    rtt(son[f][1] == t^son[fa[f]][1] == f?t:f), 1:1; pu(t);
void access (int t, int la=0) {for (;t; splay(t), son [t][1]=
    la , la=t , t=fa [ t ] ) ;}
void makeroot(int t){access(t), splay(t), rev[t]^=1;}
void link(int u, int v){makeroot(u), fa[u]=v;}
void cut(int u,int v){makeroot(u),access(v),splay(v),
     son[v][0] = fa[u] = 0;
4.5 Treap
#include <bits/stdc++.h>
using namespace std;
struct Treap{
    Treap *1, *r;
    int pri, key, val;
    Treap(int _val, int _key):
```

PushDown(x); for (tree y; (y = x->pre) != to; Rotate(

Rotate(isRight(x) != isRight(y) ? x : y);

x)) if (y-pre != to)

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```
val(val), key(key), l(NULL), r(NULL), pri(key)
             rand()){}
};
/// We assure that key value in A treap is greater than
     that in treap B
Treap *merge( Treap *a, Treap *b){
    if(a=NULL || b=NULL) return (!a) ? b : a;
    if(a->pri > b->pri){
        a \rightarrow r = merge(a \rightarrow r, b);
        return a;
    }else{
        b->l = merge(a, b->l);
        return b;
    }
void split(Treap *t, int k, Treap *&a, Treap *&b){
    if(!t) a = b = NULL;
    else if (t->key <= k) {
       a = t;
        split(t->r, k, a->r, b);
    } else {
        b = t:
        split(t->l, k, a, b->l);
Treap* insert ( Treap *t, int k, int _val) {
   Treap *tl, *tr;
    split(t, k, tl, tr);
    return merge(tl, merge(new Treap(_val, k), tr));
Treap* remove( Treap* t, int k){
    Treap *tl, *tr;
    split(t, k-1, tl, t);
    split(t, k, t, tr);
    return merge(tl, tr);
int main(){
    return 0;
```

5 Math

5.1 Prime Table

```
#include <bits/stdc++.h>
using namespace std;
{\tt struct \ Prime\_table \ \{}
    int prime [1000000] = \{2,3,5,7\};
    int sz=4;
    // biggest prime < ub
    int ub=(1<<20);
    int check(int num){
         int k = 0;
         for(k = 0; k < sz \&\& prime[k]*prime[k] <= num;
             k++){}
              if ( num % prime [k]==0) return 0;
         return 1;
    void buildprime(){
         int currentPrime=7;
         int i=4;
         for(sz=4,j=4; currentPrime < ub; sz++, j=6-j)
               currentPrime \!\!=\!\! currentPrime \!\!+\!\! j \; ;
               if (check(currentPrime)) {
                  prime[sz] = currentPrime;
               else{
                  SZ - -;
         }
}ptable;
```

5.2 Miller Rabin Prime Test

```
#include <bits/stdc++.h>
using namespace std;
typedef long long LL;
LL mul(LL a, LL b, const LL mod) {
    LL x = 0, y = a \% \text{ mod};
    while (b > 0) {
         if (b&1)
        x = (x + y) \% \text{ mod};

y = (y * 2) \% \text{ mod};
         b >>= 1;
    return x % mod;
}
LL mul(LL lhs, LL rhs, const LL mod) {
    return ( lhs * rhs ) % mod;
LL mypow(LL b, LL e, const LL mod) {
    LL x = 1;
    LL y = b;
    while (e) {
         if ( e\&1 ) x = mul(x, y, mod);
        y = \operatorname{mul}(y\,,\ y\,,\ \operatorname{mod})\,;
         e >>= 1;
    return x;
const int testbase [] = \{2, 3, 5, 7, 11, 13, 17, 19, 23,
     29, 31, 37};
bool isprime (const LL p) {
    if (p < 2) return false;
    if (p != 2 && !(p&1) ) return false;
    LL d = p - 1;
    while (!(d\&1)) d >>= 1;
    for( int a : testbase ) {
        LL td = d:
         if ( a >= p-1 ) return true;
        LL st = mypow(a, td, p);
         while (td != p - 1 && st != 1 & st != p - 1)
             st = mul(st, st, p);
             td \ll 1;
         if ( st != p - 1 && !(td&1) ) return false;
    return true;
int main() {
    int T;
    scanf("%d",&T);
    while (T--) {
         scanf("%lld",&q);
         puts(isprime(q)?"YES":"NO");
    return 0;
```

5.3 Extended Euclidean Algorithm

```
/** normal gcd function using recursion **/
int gcd(int a, int b){
    if(b == 0) return a;
    return gcd(b, a%b);
}

// Find solution of ax + by = gcd(a, b)
// ps : x, y may be negative
int extgcd(int a, int b, int& x, int& y){
    int d = a;
    if(b!= 0) {
        d = extgcd(b, a%b, y, x);
        y -= (a/b) * x;
    }else {
        x = 1, y = 0;
    }
}
```

```
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     return d;
                                                                       // Read input data
|}
                                                                       for (int i=0; i<n; i++) {
                                                                                cin >> A[i][j];
       Gauss Elimination
                                                                           }
                                                                       }
// solving linear equations with gauss elimination
#include <iostream>
                                                                       for (int i=0; i< n; i++) {
#include <cmath>
                                                                           cin \gg A[i][n];
#include <vector>
using namespace std;
                                                                       // Print input
                                                                       print(A);
void print(vector< vector<double> > A) {
                                                                       // Calculate solution
     int n = A. size();
     for (int i=0; i<n; i++) {
                                                                       vector < double > x(n);
                                                                       x = gauss(A);
          for (int j=0; j< n+1; j++) {
              cout << A[i][j] << "\t";
              // Print result
                  cout << "|
                                                                       cout << "Result:\t";</pre>
                                                                        \  \  \, \text{for} \  \, (\, i\text{nt} \  \, i\!=\!0; \  \, i\!<\!\!n\,; \  \, i\!+\!+\!) \,\, \{ \,
                                                                           cout << x[i] << "
         cout \ll "\n";
                                                                       cout << endl;
     cout << endl;
vector<double> gauss(vector< vector<double> > A) {
                                                                  5.5 FFT
     int n = A. size();
                                                                  typedef long double ld;
     for (int i=0; i< n; i++) {
          // Search for maximum in this column
                                                                   * \ \mathrm{FFT}(\ \mathrm{a}\ )\ ;
          double maxEl = abs(A[i][i]);
                                                                   * \ \mathrm{FFT}(\dot{b} \ ) \ ;
         int \max Row = i;
          for (int k=i+1; k< n; k++)
                                                                   * FFT(c);
              if (abs(A[k][i]) > maxEl) {
                  maxEl = abs(A[k][i]);
                  \max Row = k;
         }
                                                                       int N = v.size();
         // Swap maximum row with current row (column by
               column)
                                                                           if(i>j) swap(v[i],v[j]);
          for (int k=i; k<n+1;k++) {
              double tmp = A[maxRow][k];
              A[\max Row][k] = A[i][k];
                                                                           ld w = -2.0*pi/k;
             A[i][k] = tmp;
         // Make all rows below this one 0 in current
          for (int k=i+1; k< n; k++) {
              double c = -A[k][i]/A[i][i];
                                                                                    v[i] = a+b;
              \quad \  \   \text{for (int } j{=}i\;;\;\;j{<}n{+}1;\;\;j{+}{+})\;\;\{
                  if (i==j) {
                                                                                    theta *= deg;
                       A[k][j] = 0;
                                                                                }
                  } else {
                                                                           }
                       A[k][j] += c * A[i][j];
                                                                       }
                                                                  }
              }
         }
     }
                                                                  struct Complex {
     // Solve equation Ax=b for an upper triangular
                                                                       double x , y;
         matrix A
     vector < double > x(n);
                                                                           x = \_x , y = \_y;
     for (int i=n-1; i>=0; i--) {
         x[i] = A[i][n]/A[i][i];
         for (int k=i-1; k>=0; k--)
             A[k][n] -= A[k][i] * x[i];
```

return x;

vector < double > line(n+1,0);

int main() {

int n: cin >> n;

```
vector < vector < double > > A(n, line);
          for (int j=0; j<n; j++) {
/* N must be 2^k and greater than array.size()
 * for (int i = 0; i < N; ++i) c[i] = conj(a[i] * b[i]);
 * for (int i = 0; i < N; ++i) c[i] = conj(c[i]);
 * for (int i = 0; i < N; ++i) c[i] /= N;
void FFT(vector< complex<ld>>& v) {
     \begin{array}{lll} & \text{for}\,(\,\text{int}\ i\,=\,1\,,\ j\,=\,0;\ i\!<\!\!N;\ +\!\!\!+\!\!\!i\,)\ \{\\ & \text{for}\,(\,\text{int}\ k\,=\,N\!\!>\!\!>\!\!1;\ !((\,j\!\!\stackrel{-}{=}\!\!k)\&k)\,;\ k\!\!>\!\!=\!\!1); \end{array}
     for (int k = 2; k \le N; k < = 1) {
          complex{<}ld{>}\ deg\left(\,cos\left(w\right)\,,sin\left(w\right)\,\right);
          for(int j = 0; j < N; j + = k) {
                complex < ld > theta(1,0);
                for(int i = j; i < j+k/2; ++i) {
                     complex < ld > a = v[i];
                     complex< ld > b = v[i+k/2]*theta;
                     v[i+k/2] = (a-b);
//http://sd-invol.github.io/2016/02/13/FFT-mod-prime/
     Complex (double _x = 0, double _y = 0) {
     Complex operator + (const Complex &r) const {
          return Complex(x + r.x , y + r.y);
     Complex operator - (const Complex &r) const {
          return Complex(x - r.x , y - r.y);
     Complex operator * (const Complex &r) const {
          return Complex(x * r.x - y * r.y , x * r.y + y
               * r.x);
     Complex conj () const {
          return Complex(x , -y);
```

```
double operator = (const double a) {
        *this = Complex(a, 0);
        return a;
    }
                                                                                + z) \% Mod;
const double pi = acos(-1.0);
                                                                      }
//\,\mathrm{fft} with modulo, code referenced from the internet
                                                                  }
                                                             };
    fftMod::fftPrepare(len);
    fftMod::convolution(res, le, ri, len, r-l);
                                                                  NNT
                                                             5.6
namespace fftMod{
    const int N = 1 \ll 18;
    const int Mod = 1e9 + 7;
    // to do, M should be about sqrt (Mod)
                                                             NTT( a );
    const int M = 32768;
                                                             NTT(b);
    for (int i = 0; i < N: ++i)
                                                                  c\,[\,i\,] \,=\, (\,long\ long\,)\ a\,[\,i\,]\ *\ b\,[\,i\,]\ \%\ mod\,;
                                                             NTT(c, true);
    Complex w[N];
                                                              for (int i = 0; i < N; ++i)
                                                                  c\,[\,i\,]\ =\ (786433LL\text{--}12)^{'}\ *\ c\,[\,i\,]\ \%\ mod\,;
    int rev[N];
    void fftPrepare(int n) {
        int LN = __builtin_ctz(n);
                                                              constexpr int mod = 786433;
        for (int i = 0; i < n; ++ i) {
    double ang = 2 * pi * i / n;
                                                              constexpr int N = 65536;
            void NTT(vector< int >& v, bool flag = false)
                LN - 1));
                                                                  for(int i = 1, j = 0; i < N; ++i)
                                                                      for (int k = N > 1; !((j^=k)&k); k > = 1);
    void FFT(Complex P[], int n, int oper) {
                                                                      if(i>j) swap(v[i],v[j]);
        for (int i = 0; i < n; i ++) {
    if (i < rev[i]) {
                                                                  for (int k = 2; k \le N; k < = 1)
                swap(P[i], P[rev[i]]);
                                                                      int deg = mypow(flag ? 524289 : 3, N / k);
                                                                      for (int j = 0; j < N; j + = k)
        for (int d = 0; (1 \ll d) < n; d++) {
                                                                      {
            int m = 1 \ll d, m2 = m * 2, rm = n / m2;
                                                                          int theta = 1:
            for (int i = 0; i < n; i += m2) {
                                                                           for (int i = j; i < j+k/2; ++i)
                 for (int j = 0; j < m; j++) {
                     Complex &P1 = P[i + j + m], &P2 = P
                                                                               int a = v[i];
                          [i + j];
                                                                               int b = (long long) v[i+k/2]*theta%mod;
                     Complex t = w[rm * j] * P1;
                                                                               v[i] = (a+b) \% mod;
                     P1 = P2 - t;
                                                                               v[i+k/2] = (a-b+mod)\%mod;
                     P2 = P2 + t;
                                                                               theta = (long long) theta * deg \% mod;
                 }
                                                                          }
            }
                                                                      }
        }
                                                                  }
    }
                                                             }
    Complex \ A[N] \ , \ B[N] \ , \ C1[N] \ , \ C2[N] \, ;
                                                             constexpr int mod = 1e9+7;
    void convolution (vector<int> &res, vector<int> &a,
                                                              typedef vector<int> VEC;
        vector<int> &b, int len, int K) {
                                                              // ntt + Crt, code referenced from the internet
        // a[ 0 .. len ) and b[ 0 .. len ) 's convolution % Mod
                                                             namespace nttCrt {
                                                                  constexpr int magic [3] = \{1004535809, 998244353,
        // stored in res[ 0 .. K+1 )
        for (int i = 0 ; i < len ; ++ i) {
    A[i] = Complex(a[i] / M , a[i] % M);
    B[i] = Complex(b[i] / M , b[i] % M);
                                                                      104857601};
                                                                  {\rm constexpr} \ \ {\rm int} \ \widetilde{M\!O\!D} = 1000000007;
                                                                  constexpr int G = 3;
                                                                  int P:
                                                                  inline int quick_mod(int x, int k, int MOD) {
        FFT(A, len, 1); FFT(B, len, 1);
                                                                      int ans = 1;
                                                                      while (k) {
        for (int i = 0; i < len; ++ i) {
                                                                          if (k\&1) ans = 1LL * ans * x % MOD;
            int j = i ? len - i : i;
                                                                          x = 1LL * x * x % MOD;
            Complex a1 = (A[i] + A[j].conj()) * Complex
                                                                          k >>= 1;
                 (0.5, 0);
            Complex a2 = (A[i] - A[j].conj()) * Complex
                                                                      return ans;
                 (0, -0.5)
            Complex b1 = (B[i] + B[j].conj()) * Complex
                                                                  inline void change(int *y, int len) {
                 (0.5, 0);
                                                                      for(int i = 1, j = len / 2; i < len - 1; i++) {
            Complex b2 = (B[i] - B[j].conj()) * Complex
                                                                           if(i < j) swap(y[i], y[j]);
                 (0, -0.5);
            Complex c11 = a1 * b1 , c12 = a1 * b2;
                                                                          //交换互为小标反转的元素, i<j保证交换一次
            Complex c21 = a2 * b1 , c22 = a2 * b2;
                                                                          //i做正常的+1, j左反转类型的+1,始终保持i和j
            C1[j] = c11 + c12 * Complex(0, 1);
                                                                               是反转的
            C2[j] = c21 + c22 * Complex(0, 1);
                                                                           int k = len / 2;
                                                                          while(j >= k) {
        FFT(C1 , len , -1); FFT(C2 , len , -1);
                                                                              j -= k;
                                                                               k /= 2;
        for (int i = 0; i \le K; ++ i) {
            int x = (LL)(C1[i].x / len + 0.5) \% Mod;
                                                                           if(j < k) j += k;
```

```
inline void ntt(int *y, int len, int on) {
         change(y, len);
          for (int h = 2; h \le len; h \le 1) {
              int wn = quick_{\text{mod}}(G, (P - 1) / h, P);
              for (int j = 0; j < len; j += h) {
                   int w = 1;
                   for(int k = j; k < j + h / 2; k++) {
                       \begin{array}{ll} \text{int } u = y[k] \ \% \ P; \\ \text{int } t = 1LL \ ^* w \ ^* \ y[k+h \ / \ 2] \ \% \ P; \end{array}
                       y[k] = (u + t) \% P;
                       y[k + h / 2] = ((u - t) \% P + P) \%
                       w = 1LL * w * wn \% P;
                  }
             }
          if (on == -1) {
              for (int i = 1; i < len / 2; i++)
                  swap(y[i], y[len - i]);
              int inv = quick\_mod(len, P - 2, P);
              for (int i = 0; i < len; i++)

y[i] = 1LL * y[i] * inv % P;
         }
     }
     int n;
     int r[3][3];
     inline int CRT(int *a) {
         int sb[3] = \{a[0], a[1], a[2]\};
         for(int i = 0; i < 3; ++i) {
              for(int j = 0; j < i; ++j) {
                  int t = (sb[i] - sb[j]) \% magic[i];
                  int mul = 1, ans = sb[0] \% MOD;
         for (int i = 1; i < 3; ++i) {
    mul = 1LL * mul * magic[i - 1] % MOD;
              ans = (ans + 1LL * sb[i] * mul) % MOD;
         }
         return ans;
     int tmp[maxn][3];
     int x1 [maxn*2], x2 [maxn*2];
     inline void gao(vector<int>& res, vector<int> &a,
          vector < int > \&b, int len, int kk) {
          for (int ti = 0; ti < 3; ti++) {
              P = magic[ti];
              int k;
              for (k = 0; k < SZ(a) \&\& k < len; k++) x1
                  k = a[k];
              \quad \  \  \text{for}\ (\ ;k<\,len\,;\ k+\!+\!)\ x1[\,k\,]\ =\ 0\,;
              for ( k = 0; k < SZ(b) && k < len; k++) x2[
                  k = b[k];
              for (; k < len; k++) x2[k] = 0;
              ntt(x1, len, 1); ntt(x2, len, 1);
              for (int i = 0; i < len; i++) x1[i] = 1LL *
              for (int i = 0; i \le kk; i++) tmp[i][ti] =
                   x1[i];
          for (int i = 0; i \le kk; i++) res[i] = CRT(tmp
              [ i ] )
     inline void init() {
         for (int i = 0; i < 3; i++) {
              for (int j = 0; j < 3; j++) {
                  r[i][j] = quick\_mod(magic[i], magic[j])
                         2, magic[j]);
         }
     }
};
```

5.7 Big Number

```
//http://blog.csdn.net/hackbuteer1/article/details
    /6595881
#include<iostream>
#include<string>
#include<iomanip>
#include < algorithm >
using namespace std;
#define MAXN 9999
#define MAXSIZE 10
#define DLEN 4
class BigNum
{
private:
 int a[500];
               //可以控制大数的位数
 int len;
               //大数长度
public:
 BigNum() { len = 1; memset(a,0,sizeof(a)); }
                       //将一个int类型的变量转化为
 BigNum(const int);
     大数
                       //将一个字符串类型的变量转化
 BigNum(const char*);
     为大数
 BigNum(const BigNum &); //拷贝构造函数
 BigNum & operator=(const BigNum &); //重载赋值运算
     符,大数之间进行赋值运算
  friend istream& operator>>(istream&, BigNum&);
                                              //
     重载输入运算符
  friend ostream& operator<<(ostream&, BigNum&);</pre>
     重载输出运算符
 BigNum operator+(const BigNum &) const;
                                       //重载加法
     运算符,两个大数之间的相加运算
 BigNum operator - (const BigNum &) const;
                                       //重载减法
     运算符,两个大数之间的相减运算
 BigNum operator*(const BigNum &) const;
                                       //重载乘法
     运算符, 两个大数之间的相乘运算
 BigNum operator/(const int &) const;
                                       //重载除法
     运算符,大数对一个整数进行相除运算
 BigNum operator (const int &) const;
                                      //大数的n次
     方运算
       operator%(const int &) const;
                                      //大数对一个
     int类型的变量进行取模运算
       operator > (const BigNum & T) const;
                                        //大数和另
     一个大数的大小比较
 bool operator>(const int & t)const;
                                        //大数和一
     个int类型的变量的大小比较
                   //输出大数
 void print();
BigNum::BigNum(const int b)
                           //将一个int类型的变量转
    化为大数
 int c, d = b;
 len = 0;
 memset(a, 0, sizeof(a));
  while (d > MAXN)
   c = d - (d / (MAXN + 1)) * (MAXN + 1);
   d = d / (MAXN + 1);
   a[len++] = c;
 a[len++] = d;
BigNum::BigNum(const char*s)
                             //将一个字符串类型的变
    量转化为大数
 int t,k,index,l,i;
 memset(a, 0, sizeof(a));
 l=strlen(s);
 len=l/DLEN;
 if (1%DLEN)
```

```
t.a[i] -=MAXN+1;
    len++:
  index=0:
  for (i=l-1; i>=0; i-=DLEN)
                                                                      if (t.a[big] != 0)
    t = 0;
                                                                       t.len = big + 1;
    k=i -DLEN+1;
    if(k<0)
                                                                       t.len = big;
      k=0;
                                                                     return t;
    for (int j=k; j <= i; j++)
      t=t*10+s[j]-'0';
                                                                   BigNum BigNum::operator-(const BigNum & T) const
    a [index++]=t;
                                                                        个大数之间的相减运算
  }
                                                                     int i,j,big;
BigNum::BigNum(const BigNum & T) : len(T.len) //拷贝构
                                                                     bool flag;
                                                                     BigNum t1, t2;
                                                                     if(*this>T)
  int i;
  memset(a,0,sizeof(a));
                                                                        t1=*this;
  for(i = 0 ; i < len ; i++)
                                                                        t2=T:
    a[i] = T.a[i];
                                                                        flag = 0;
BigNum & BigNum::operator=(const BigNum & n)
                                                      //重载赋
                                                                      else
                                                                     {
    值运算符, 大数之间进行赋值运算
                                                                       t1=T:
                                                                        t2 = *this;
  int i;
                                                                        flag = 1;
  len = n.len;
  memset(a, 0, sizeof(a));
                                                                     big=t1.len;
  for(i = 0 ; i < len ; i++)
                                                                     for (i = 0 ; i < big ; i++)
   a[i] = n.a[i];
  return *this;
                                                                        if(t1.a[i] < t2.a[i])</pre>
istream& operator>>(istream & in, BigNum & b)
                                                        //重载
                                                                          j = i + 1;
    输入运算符
                                                                          \textcolor{red}{\textbf{while}\,(\,t\,1\,.\,a\,[\,j\,]} \; = \!\!\!= \; 0)
                                                                            j++;
  char ch[MAXSIZE*4];
                                                                          t1.a[j--]--;
  int i = -1;
                                                                          while(j > i)
  in>>ch;
                                                                            t1.a[j--] += MAXN;
  int l=strlen(ch);
                                                                          t1.a[i] += MAXN + 1 - t2.a[i];
  int count=0,sum=0;
  for (i=l-1; i>=0;)
                                                                        else
                                                                         t1.a[i] -= t2.a[i];
    sum = 0;
    int t=1:
                                                                     t1.len = big;
    for (int j=0; j<4&&i>=0; j++,i--,t*=10)
                                                                     while (t1.a[len - 1] = 0 \&\& t1.len > 1)
      sum + = (ch[i] - '0') *t;
                                                                        t1.len - -;
                                                                       big - -;
    b.a[count]=sum;
    count++;
                                                                     if (flag)
                                                                       t1.a[big-1]=0-t1.a[big-1];
  b.len =count++;
                                                                     return t1;
  return in;
                                                                   BigNum BigNum::operator*(const BigNum & T) const
ostream& operator << (ostream& out, BigNum& b)
                                                                        个大数之间的相乘运算
    输出运算符
                                                                     BigNum ret;
  int i;
                                                                     int i,j,up;
  cout << b.a[b.len - 1];
                                                                     int temp, temp1;
  for(i = b.len - 2 ; i >= 0 ; i--)
                                                                     for (i = 0 ; i < len ; i++)
    cout.width(DLEN);
                                                                        up = 0;
    cout. fill(',0');
                                                                        for(j = 0 ; j < T.len ; j++)
    cout << b.a[i];
  }
                                                                          temp \, = \, a \, [\, i \, ] \ * \ T. \, a \, [\, j \, ] \ + \ ret. \, a \, [\, i \ + \ j \, ] \ + \ up \, ;
  return out;
                                                                          if(temp > MAXN)
                                                                          {
                                                                            temp1 \, = \, temp \ - \ temp \ / \ \left( \text{MAXN} \, + \, 1 \right) \ * \ \left( \text{MAXN} \, + \, 1 \right) \, ;
BigNum BigNum::operator+(const BigNum & T) const
                                                           //两
                                                                            up = temp / (MAXN + 1);

ret.a[i + j] = temp1;
     个大数之间的相加运算
                                                                          }
 BigNum t(*this);
                                                                          else
                    //位数
  int i, big;
  big = T.len > len ? T.len : len;
                                                                            up = 0;
  for(i = 0 ; i < big ; i++)
                                                                            \mathtt{ret.a}\,[\,\mathtt{i}\,+\,\mathtt{j}\,]\,=\,\mathtt{temp}\,;
    t.a[i] +=T.a[i];
    if(t.a[i] > MAXN)
                                                                        if(up != 0)
                                                                          ret.a[i + j] = up;
      t.a[i + 1]++;
```

```
ret.len = i + i:
  while (ret.a[ret.len - 1] = 0 \&\& ret.len > 1)
    ret.len--:
  return ret;
BigNum BigNum::operator/(const int & b) const
                                                  //大数
    对一个整数进行相除运算
 BigNum ret;
  int i, down = 0;
  for (i = len - 1 ; i >= 0 ; i--)
    ret.a[i] = (a[i] + down * (MAXN + 1)) / b;
    down = a[i] + down * (MAXN + 1) - ret.a[i] * b;
  ret.len = len;
  while (ret.a[ret.len - 1] = 0 \&\& ret.len > 1)
    ret.len - -;
  return ret;
int BigNum::operator %(const int & b) const
                                                 //大数对
    一个int类型的变量进行取模运算
  int i, d=0;
  for (i = len -1; i>=0; i--)
  {
   d = ((d * (MAXN+1))\% b + a[i])\% b;
  return d:
BigNum BigNum::operator^(const int & n) const
                                                   //大数
    的n次方运算
 BigNum t, ret(1);
  int i:
  if(n<0)
    exit(-1);
  if(n==0)
    return 1;
  if(n==1)
    return *this;
  int m=n;
  while (m>1)
    t=*this;
    \quad \  \  \textbf{for} \ (\quad i=\!1; i<\!\!<\!\!1\!\!<\!\!m; i<\!\!<\!\!=\!\!1)
    {
      t=t*t;
    }
   m=i;
    ret=ret*t;
    if (m==1)
      ret=ret*(*this);
  }
bool BigNum::operator>(const BigNum & T) const
                                                   //大数
    和另一个大数的大小比较
  int ln;
  if(len > T.len)
   return true;
  else if(len = T.len)
    ln = len - 1:
    while (a[ln] = T.a[ln] & ln >= 0)
    if(\ln >= 0 \&\& a[\ln] > T.a[\ln])
      return true;
    else
      return false;
  }
  else
    return false;
bool BigNum::operator >(const int & t) const
                                                   //大数
    和一个int类型的变量的大小比较
 BigNum b(t);
  return *this>b;
```

```
void BigNum::print()
                            //输出大数
  int i;
  cout << a[len - 1];
  for (i = len - 2 ; i >= 0 ; i--)
    cout.width(DLEN);
    cout. fill('0');
    cout <\!\!< a\,[\,i\,]\,;
  cout << endl;
int main(void)
  int i,n;
  BigNum x[101];
                         //定义大数的对象数组
  x[0] = 1;
  for (i=1; i<101; i++)
  x[i]=x[i-1]*(4*i-2)/(i+1);
while (scanf("%d",&n)==1 && n!=-1) {
    x[n].print();
}
```

6 string

6.1 Palindromic Tree

n=0:

```
回文自動機包含以下元素:
   狀態St, 所有節點的集合, 一開始兩個節點, 0表示偶數長
      度串的根和1表示奇數長度串的根
   last 新增加一個字符後所形成的最長回文串的節點編號
   s 當前的字符串(一開始設s[0]=-1(可以是任意一個在串S
      中不會出現的字符))
  n 表示添加的字符個數
每個節點代表一個不同的回文子字串,我們在每個節點會儲存
   一些數值:
   len 表示所代表的回文子字串長度
   next[c] 表示所代表的回文子字串在頭尾各增加一個字符c
      後的回文字串其節點編號
   sufflink 表示所代表的回文子字串不包括本身的最長後綴
      回文子串的節點編號
   cnt(非必要) 表示所代表的回文子字串在整體字串出現的
      次數(在建構完成後呼叫count()才能計算)
   //num(非必要) 表示所代表的回文子字串其後綴為回文字
      串的個數 <━ not included
struct palindromic_tree{
   struct node{
      int next[26], sufflink, len; /*這些是必要的元素*/
      int l, r; // this node is s[ l .. r ]
      int cnt, num;
                        /*這些是額外維護的元素*,
      node(int l=0): sufflink(0), len(l), cnt(0), num(0)
         for (int i=0; i<26;++i) next [i]=0;
   std::vector<node> St;
   std::string s; //current string [ 1 .. n ]
   int last, n;
   palindromic\_tree():St(2), last(1), n(0){
      St[0].sufflink=1;
      St[1].len=-1;
      s.push\_back(-1);
   inline void clear(){
     St.clear();
      s.clear();
      last=1;
```

```
St.push\_back(0);
        St.push\_back(-1);
        St[0].sufflink=1;
        s.push\_back(-1);
    inline int get_sufflink(int x){
        while (s[n-St[x].len-1] != s[n]) x=St[x].
           sufflink;
        return x:
    inline void add(int c){
       s.push_back(c-= 'a',);
        int cur=get_sufflink(last);
        if (!St[cur].next[c]) {
            int now=St.size();
            St.push\_back(St[cur].len+2);\\
            St [now]. sufflink=St [get_sufflink(St [cur].
                sufflink)].next[c];
            /*不用擔心會找到空節點,由證明的過程可知*/
            St[cur].next[c]=now;
            St[now].num=St[St[now].sufflink].num+1;
            St[now] \cdot l = n - St[now] \cdot len + 1, St[now] \cdot r
                = n;
        last=St [cur].next[c];
       ++St[last].cnt;
    inline void count(){/*cnt必須要在構造完後呼叫count
        ()去計算*/
        std::vector<node>::reverse_iterator i=St.rbegin
            ();
        for(;i!=St.rend();++i) {
            St[i->sufflink].cnt+=i->cnt;
    inline int size(){/*傳回其不同的回文子串個數*/
        return St. size()-2;
    }
}ptree;
```

6.2 Suffix Array

6.3 Longest Palindromic Substring

```
//ntu judge Earse
#define maxn 200001
char t[maxn];
char s[maxn * 2];
int z [maxn * 2];
int N;
int longest_palindromic_substring() {
    // t穿插特殊字元, 存放到s。
    int n = strlen(t);
    N = n * 2 + 1;
memset(s, '.', N);
    for (int i=0; i < n; ++i) s[i*2+1] = t[i];
    s[N] = ' \setminus 0';
    z[0] = 1;
               // if無須使用, then無須計算。
    int L = 0, R = 0;
    for (int i=1; i<N; ++i) // 從z[1]開始
    {
        z[i] = (R > i) ? min(z[2*L-i], R-i) : 1;
        s[i-z[i]] = s[i+z[i]]) z[i]++;
         \mbox{if} \ (\, i + z \, [\, i \, ] \, > R) \ L \, = \, i \, , \ R \, = \, i + z \, [\, i \, ] \, ; 
    }
    // 尋找最長迴文子字串的長度
    n = 0;
    int p = 0;
    for (int i=1; i<N; ++i) // 從z[1]開始
        if (z[i] > n)
            n = z[p = i];
    // longest 從中心到外端的長度 => (n-2)/2
```

```
//cout << "最長迴文子字串的長度是" << (2*n-1) / 2;
     // 印出最長迴文子字串, 記得別印特殊字元。
          for (int i=p-z[p]+1; i \le p+z[p]-1; ++i)
              if (i & 1) {
                   cout << s[i];
     return (2*n-1)/2;
int nxt[maxn * 2];
int main() {
     int T; cin>>T;
     while (T--)
         scanf("%s", t);
         #ifdef DEBUG
              cout << longest_palindromic_substring() <<</pre>
                   endl:
         #else
              longest_palindromic_substring();
         #endif
         \begin{array}{lll} memset(nxt, -1, \ sizeof(nxt)); \\ for(int \ i = 0; \ i < N; \ i++) \end{array} \{
              nxt \left[ \ i \cdot z \left[ \ i \right] + 1 \ \right] \ = \ i + 1;
          int leftmost = 0;
          for (int i = 0; i < N; i++) {
              leftmost = max(leftmost, nxt[i]);
              nxt[i] = max(leftmost, nxt[i]);
          int ans = 0:
          for (int cur = 0; cur < N-1;) {
              cur = nxt[cur];
              ans++;
          cout << ans << endl;
     return 0;
}
```

7 geometry

7.1 Point Class

```
const double eps = 1e-10;
#define N 100
struct P {
    {\color{red} \textbf{double}} \ x\,,\ y\,;
    P(double \_x=0, double \_y=0) : x(\_x), y(\_y) {};
     void read() {
         scanf("%lf%lf",&x,&y);
    void print()
         printf("%f %f\n", x, y);
} p[N];
bool operator <( Pa, Pb) { return tie(a.x,a.y)<tie(b
    .x,b.y);
P operator +( P a, P b ) { return P{a.x+b.x,a.y+b.y}; }
P operator -( P a, P b ) { return P{a.x-b.x,a.y-b.y}; }
P \ operator \ *( \ P \ b, \ double \ a \ ) \ \{ \ return \ P\{a*b.x,a*b.y\};
P operator /( P a, double b ) { return P{a.x/b,a.y/b};
P& operator /=( P &a, double b ) { return a=a/b; }
double operator *( Pa, Pb) { return a.x*b.x+a.y*b.y;
double operator ^( P a, P b ) { return a.x*b.y-a.y*b.x;
double x( Po, Pa, Pb) { return (a-o)^(b-o); }
double dot(Po, Pa, Pb) { return (a-o)*(b-o); }
```

7.2 Intersection of Circles/Lines/Segments

//PECaveros

```
vector<P> interCircle( P o1 , double r1 , P o2 , double
  double d2 = (01 - 02) * (01 - 02);
  double d = sqrt(d2);
  if (d > r1 + r2) return \{\};
  P \dot{u} = (o1+o2)*0.5 + (o1-o2)*((r2*r2-r1*r1)/(2*d2));
  double A = sqrt((r1+r2+d)*(r1-r2+d)*(r1+r2-d)*(-r1+r2
     +d));
  P\ v = P(\ o1.y \hbox{-} o2.y\ ,\ \hbox{-} o1.x \ +\ o2.x\ )\ *\ A\ /\ (2*d2)\,;
  return \{u+v, u-v\};
| P interPnt( P p1, P p2, P q1, P q2) {
  \frac{\text{double}}{\text{double}} \text{ f1} = (p2 - p1)
  double f = (f1 + f2);
  \label{eq:fabs}  \text{if ( fabs( f ) < eps ) } \\  \text{return Pt( } \\ \text{nan("") , } \\ \text{nan("") );} 
  return q1 * ( f2 / f ) + q2 * ( f1 / f );
// p1 == p2 || q1 == q2 need to be handled
bool banana (const PLL& p1 , const PLL& p2
            ((q1 - p2) * (q2 - p2)) <= 0;
  return (ori( p1, p2, q1 ) * ori( p1, p2, q2 )<=0) &&
         (ori(q1, q2, p1) * ori(q1, q2, p2)<=0);
```

7.3 Convex Hull

```
#define REP(i,n) for ( int i=0; i<int(n); i++)
void input() {
    scanf("%d",&n);
    REP(i,n) p[i].read();
P findCenter() {
    p[n]=p[0];
    P center=P\{0,0\};
    REP(i,n) {
        double v=p[i]*p[i+1];
        center.x += (p[i].x+p[i+1].x)*v;
        center.y += (p[i].y+p[i+1].y)*v;
    double area=0;
    REP(\,i\,\,,n)\ area+\!\!=\!\!p\,[\,i\,\,]\,^*\,p\,[\,\,i\,+1\,];
    area \neq 2;
    center /= 6*area;
    return center;
}
P q1[N], q2[N], q[N];
void convex() {
    sort(p,p+n);
    int m1=0,m2=0;
    REP(i,n) {
         while (m1>=2 \&\& X(q1[m1-2],q1[m1-1],p[i]) >= 0
         while (m2>=2 \&\& X(q2[m2-2],q2[m2-1],p[i]) <= 0
              ) m2- -:
        q1 [m1++]=q2 [m2++]=p[i];
    int m=0;
    REP(i, m1) q[m++]=q1[i];
    for ( int i=m2-2; i>=1; i-- ) q[m++]=q2[i];
    q[m]=q[0];
void solve() {
```

```
convex();
// continue ...
```

Point v[maxn];

7.4 Half Plane Intersection

```
// \texttt{http://acm.csie.org/ntujudge/problemdata/2575.pdf}
//http://www.csie.ntnu.edu.tw/~u91029/Half-
    planeIntersection.html
預先使用四個半平面, 設定一個極大的正方形邊界, 讓半平面
    交集擁有邊界。
二、逐一加入每個半平面, 求出當下的半平面交集(凸多邊
    形)。
online 演算法, 隨時維護一個半平面交集。每次更新需時 O(N
    ),總時間複雜度為 O(N<sup>2</sup>), N 是半平面數目。
#include <bits/stdc++.h>
using namespace std;
#define mp make_pair
typedef complex<double> Point;
typedef vector<Point> Polygon;
typedef pair < Point , Point > Line;
#define x real()
#define y imag()
// 兩向量叉積
double cross (Point& a, Point& b) {
    return a.x * b.y - a.y * b.x;
// 向量oa與向量ob進行叉積
double cross (Point& o, Point& a, Point& b) {
return (a.x-o.x) * (b.y-o.y) - (a.y-o.y) * (b.x-o.x)
        );
// 多邊形面積
double area (Polygon& p) {
    double a = 0;
    _{\hbox{\scriptsize int}}\ n=p.\,\hbox{\scriptsize size}\,(\,)\,;
    for (int i=0; i< n; ++i)
       a += cross(p[i], p[(i+1)\%n]);
    \begin{array}{lll} \textbf{return} & fabs\,(a) & / & 2; \end{array}
}
// 兩線交點
Point intersection (Point& a1, Point& a2, Point& b1,
    Point& b2) {
    Point a = a2 - a1, b = b2 - b1, s = b1 - a1;
    // 一個凸多邊形與一個半平面的交集
Polygon halfplane_intersection(Polygon& p, Line& line)
    Polygon q;
    Point \ p1 = \ line. \, first \; , \ p2 = \ line. \, second \, ;
    // 依序窮舉凸多邊形所有點,判斷是否在半平面上。
    // 如果凸多邊形與半平面分界線有相交, 就求交點。
    int n = p.size();
    for (int i=0; i< n; ++i)
        double c = cross(p1, p2, p[i]);
        double d = cross(p1, p2, p[(i+1)\%n]);
        if (c >= 0) q.push_back(p[i]);
        if (c * d < 0) q.push_back(intersection(p1, p2,
             p[i], p[(i+1)\%n]);
    return q;
#define maxn 550
//Line line[maxn];
```

```
double ans[maxn];
 int main() {
             int T; cin>>T;
             while (T--) {
                          int n;
                          double w, h;
                          scanf("%d %lf %lf", &n, &w, &h);
                           // 預先設定一個極大的正方形邊界
                          Polygon p, org;
                          /** initialize
                            p.push_back(Point(-1e9,-1e9));
                          p.push\_back(Point(-1e9,+1e9));
                          p.push_back(Point(+1e9,-1e9));
                         p.push_back(Point(+1e9,+1e9));
**/
                          p.push\_back(Point(0,0));
                          p.push\_back(Point(0,h));\\
                          p.push_back(Point(w,h));
                          p.push_back(Point(w,0));
                          org = p;
                          for (int i = 0; i < n; i ++) {
                                      double a, b;
                                       scanf("%lf %lf", &a, &b);
                                       v[i] = Point(a, b);
                           // 每一個半平面都與目前的半平面交集求交集
                          for (int i=0; i<n; ++i)
                                       p = org;
                                       for(int j = 0; j < n; j++) {
                                                    if(i=j) continue;
                                                   Line line;
                                                    // find perpendicular line to line i_j
                                                   Point a( (v[i].x+v[j].x)/2, (v[i].y+v[j
                                                                 ].y)/2);
                                                    Point \ b(a.x+(v[i].y-v[j].y) \,, \ a.y-(v[i].y) \,, \ a.y-(v[i].y
                                                                x - v[j].x));
                                                   line = cross(a, b, v[i]) >= 0 ? mp(a, b)
                                                                : mp(b,a);
                                                   p = halfplane_intersection(p, line);
                                                    if (area(p) == 0) break; // 退化或者
                                       ans[i] = area(p);
                          for (int i = 0; i < n; i ++) printf ("%.9f\n", ans [
                                        i]);
             }
}
/*
10
3 4 4
^1 1 2 2 3 3 */
```