Contents

1	Basic	1	1.1 .vimrc
	1.1 .vimrc	1	
	1.2 Increase Stack Size	1	imap jj <esc></esc>
	1.3 digitDP	1	sy on
	1.4 DP(convex hull optimization)	2	se sw=4 ts=4 sts=4 et nu sc hls cc=69
	1.5 simulated annealing	2	filet plugin indent on
2	Graph	3	mm <f5> :!"./%<"<cr> mm <f6> :!"./%<" < input.txt<cr></cr></f6></cr></f5>
_	2.1 HLD	3	au FileType cpp no <f9>:!g++ % -o</f9>
	2.2 Hungarian	4	$\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ $
	2.3 KM	5	\ -Wshadow -Wno-unused-result <cr> no <expr> <silent> <home> col('.') ==</home></silent></expr></cr>
	2.4 Bi-vertex-connected Subgraph	5	$\begin{array}{c} \text{motch}(\text{getline}('.'), '\backslash S') + 1 \end{array}$
	2.5 Bi-edge-connected Subgraph	6	\? ',0', : ',^,
	2.6 SCC	6	im < silent > < Home > < C-O > < Home >
	2.7 Edmond's Matching Algorithm	6	
	2.8 Tree Decomposition	7	10 I C 1 C'
	2.9 Tree Longest Path	8	1.2 Increase Stack Size
	2.0 Tree Bongest Fault	O	1//
3	Flow	9	//stack resize asm("mov %0,%%esp\n" :: "g"(mem+10000000));
	3.1 Dinic Maxflow	9	//change esp to rsp if 64-bit system
	3.2 Sw Mincut	10	
			//stack resize (linux) #include <sys resource.h=""></sys>
4	Data Structure	10	void increase_stack_size() {
	4.1 Disjoint Set	10	const rlim_t ks = $64*1024*1024$;
	4.2 Djs + Seg \dots	10	struct rlimit rl;
	4.3 Sparse Table	11	<pre>int res=getrlimit(RLIMIT_STACK, &rl); if(res==0){</pre>
	4.4 Link Cut Tree	12	if (rl.rlim_cur <ks){< td=""></ks){<>
	4.5 Treap	12	rl.rlim_cur=ks;
			res=setrlimit(RLIMIT_STACK, &rl);
5	Math	12	}
	5.1 Prime Table	12	}
	5.2 Miller Rabin Prime Test	13	
	5.3 Extended Euclidean Algorithm	13	
	5.4 Gauss Elimination	13	1.3 digitDP
	5.5 FFT	14	
	5.6 NNT	15	/** / **
	5.7 Big Number	16	简介 顾名思义,所谓的数位DP就是按照数字的个,十,百,千
	5.8 Simplex	18	位数进行的DP。
6	string	19	数位DP的题目有着非常明显的性质:
U	6.1 Palindromic Tree	19	询问[1,r]的区间内,有多少的数字满足某个性质
	6.2 Suffix Array	19	做法根据前缀和的思想,求出[0,l-1]和[0,r]中满足性质的数
	6.3 Longest Palindromic Substring	19	的个数,然后相减即可。
	0.5 Longest I annuronne Substituig	13	 算法核心
7	geometry	20	并 仏 似 心
	7.1 Point Class	20	LL dfs(int x, int pre, int bo, int limit);
	7.2 Intersection of Circles/Lines/Segments	20	一般需要以上参数(当然具体情况具体分析)。
	7.3 Convex Hull	20	
	7.4 Half Plane Intersection	21	x表示当前的数位(一般都是从高位到低位) pre表示前一位的数字
			pre表示前一位的数子
			合条件
			limit这个很重要! 它表示当前数位是否受到上一位的限

Basic 1

```
map jj <Esc>
sy on
e sw=4 ts=4 sts=4 et nu sc hls cc=69
filet plugin indent on
nm <F5> :!"./%<"<CR>
m <F6> :!"./%<" < input.txt<CR>
u FileType cpp no <F9> :!g++ % -o
\ %< -std=c++14 -O3 -Wall -Wextra
\ -Wshadow -Wno-unused-result<CR>
no <expr> <silent> <Home> col('.') ==
\ match(getline('.'), '\S') + 1
? '0' : '^
m < silent > < Home > < C-O > < Home >
```

Increase Stack Size

```
/stack resize
sm( "mov %0,%%esp\n" :: "g"(mem+10000000) );
/change esp to rsp if 64-bit system
/stack resize (linux)
include <sys/resource.h>
roid increase\_stack\_size() + roid increase\_stack\_size()
 const rlim_t ks = 64*1024*1024;
 struct rlimit rl;
 int res=getrlimit(RLIMIT_STACK, &rl);
 if(res==0){
   if(rl.rlim_cur<ks){</pre>
     rl.rlim_cur=ks;
     res=setrlimit(RLIMIT_STACK, &rl);
 }
```

$\operatorname{digitDP}$ 1.3

```
位数进行的DP。
数 位DP的 题 目 有 着 非 常 明 显 的 性 质:
   询问[1,r]的区间内, 有多少的数字满足某个性质
做法根据前缀和的思想,求出[0,1-1]和[0,r]中满足性质的数
   的个数,然后相减即可。
算法核心
L dfs(int x, int pre, int bo, int limit);
一般需要以上参数(当然具体情况具体分析)
   x表示当前的数位 (一般都是从高位到低位)
   pre表示前一位的数字
   bo可以表示一些附加条件:是否有前项0,是否当前已经符
      合条件……
   limit 这个很重要! 它表示当前数位是否受到上一位的限
      制, 比较抽象, 举例说明
   如果上限是135,前两位已经是1和3了,现在到了个位,个
      位只能是5以下的数字
注: 如果当前受限,不能够记忆化,也不能返回记忆化的结果
为了避免多次调用时 每次上限不同 而导致的错
//http://acm.csie.org/ntujudge/view_code.php?id=106844
// Multiples
LL x;
int digit [100];
LL ten_pow[ 15 ];
bool ava[15];
LL \ dp \, [15][2][1000000];
LL dfs(int len, LL mod, bool bo, bool limit) {
```

```
if ( len < 0 ) return mod == 0;
    if (!limit && dp[len][bo][mod]!= -1 ) return dp[
        len ] [bo] [mod];
    int up = limit? digit[len] : 9;
    LL ret = 0;
    for (int i = 0; i \le up; i++) if (ava[i] || (i=0&&
        bo) ) {
        ret += dfs(len-1, (mod+ten_pow[len]*i)%x, bo
            &&(!i), limit&&(i=up));
    if( !limit ) dp[len][bo][mod] = ret;
    return ret;
LL solve (LL num) {
    int len = 0; digit [0] = 0;
    while( num ) {
        digit[len++] = num\%10;
        num = 10;
    return dfs (len -1, 0,1, 1);
bool check (LL num) {
    while ( num )
        if (!ava[ num%10 ] ) return false;
        num /= 10:
    return true;
int main() {
    LL A, B;
    cin>>x>>A>>B;
    ten_pow[0] = 1;
    mem(dp, -1);
    for (int i = 1; i < 15; i++)
        ten_pow[i] = (ten_pow[i-1]*10)\%x;
    string \ dig; \ cin>>dig;
    mem(ava, false);
    for (char c : dig) ava [c-'0'] = 1;
    if ( x <= 1000000 ) {
        cout << solve(B) - solve(A-1) << endl;
    }else {
        LL ans = 0;
        LL cur = 0;
        while ( cur < A ) cur += x;
        while ( cur <= B ) {
            i\dot{f} ( check(cur) ) ans++;
            cur += x;
        cout << ans << endl;
    }
}
```

1.4 DP(convex hull optimization)

```
// \texttt{http://codeforces.com/contest/311/problem/B}
struct line {
     LL slope,
                  inter;
     LL value(LL x) { return x*slope + inter; }
bool check(line x, line y, line z) {
     return (z.slope - y.slope) * (z.inter - x.inter )
              (z.slope - x.slope) * (z.inter - y.inter);
#define maxn 100005
int n, m, p;
 LL \ a \, [\, maxn ] \; , \ d \, [\, maxn ] \; , \ dp \, [\, 1\,0\,1\,] \, [\, maxn \,] \; , \ s \, [\, maxn \,] \; ; 
int main() {
     cin>\!\!> n>\!\!> m>\!\!> p;
     for(int i = 2; i <= n; ++i) {
          d[i] = getint();
          d\,[\,i\,] \; +\!\!= d\,[\,i\,-1\,]\,;
     for(int i = 1; i<=m; ++i) {
    int h; scanf("%d %lld", &h, a+i);
          a[i] -= d[h];
     sort(a+1,a+1+m);
```

```
\label{eq:formalized} \begin{array}{ll} \text{for}\,(\,int\ i\!=\!1; i\!<\!\!=\!\!m;\, i\!+\!\!+\!\!)\ s\,[\,i\,] \ = \ a\,[\,i\,]\!+\!s\,[\,i\,-\,1\,]\,; \end{array}
//start dp
for (int i=1; i < p; i++) {
      if(i = 1) {
            for (int j=1; j \le m; j++) dp[i][j] = j*a[j] - s
                  [j];
      }else {
           deque<line> dq;
           dq.pb(\{0, 0\});
            for (int j=1; j \le m; j++) {
                  while (dq. size() >= 2 \&\& dq[0]. value(-a[j]) > dq[1]. value(-a[j])) dq.
                        pop_front();
                  dp[i][j] = dq[0].value(-a[j]);
                  line newline{ j, dp[i-1][j]+s[j] };
                  \label{eq:while} \mbox{while} (\ \mbox{dq.size}() >= 2 \ \&\& \ \mbox{check}(\mbox{dq}[\mbox{dq}.
                        size()-2], dq.back(), newline)) dq
                        .pop_back();
                  dq.pb( newline );
                  if ( i==1 )
                       dp\,[\,i\,\,]\,[\,j\,\,] \;=\; j\,{}^*a\,[\,j\,\,] \;\; -\;\; s\,[\,j\,\,]\,;
                  }else {
                       LL mn = 0;
                        for (int \ k = 1; \ k < j; \ k+\!\!\!\!+\!\!\!\!+) \ \{
                             mn = min(mn, dp[i-1][k] + s[k]
                                    - a[j]*k);
                        dp[i][j] = mn + a[j]*j-s[j];
                        // apply convex hull optimization
                  dp[i][j] += a[j]*j - s[j];
      }
cout \ll dp[p][m] \ll endl;
```

1.5 simulated annealing

```
//http://mikucode.blogspot.tw/2015/03/algorithm.html
//尋找和所有點距離和最小的點
#include <cstdio>
#include <cstdlib>
#include <cmath>
#define F(n) Fi(i,n)
#define Fi(i,n) for (int i=0;i< n;i++)
#define N 1010
using namespace std;
int X[N], Y[N], n;
inline double pow2(double x){
    return x*x;
double check(double x, double y){
    double ans=0;
    F(n) ans += sqrt(pow2(x-X[i]) + pow2(y-Y[i]));
    return ans;
int main(){
     while (~scanf("%d",&n)) {
        F(n) scanf("%d%d",X+i,Y+i);
        double x=0,y=0,tx,ty,tans,l=10000,ans;
        ans=check(x,y);
        while(1>1e-4) {
            int tmp=rand();
             tx=x+l*cos(tmp); ty=y+l*sin(tmp);
            tans=check(tx,ty);
             if(tans < ans) ans=tans, x=tx, y=ty;
            else 1*=0.9;
        printf("%.9f\n",2*ans);
    }
}
//尋找兩個點使他們跟給定的四個點最小生成樹最小
#include <cstdio>
#include <cstdlib>
```

```
#include <cmath>
#include <algorithm>
#define F(n) Fi(i,n)
#define Fi(i,n) Fl(i,0,n)
#define Fl(i,l,n) for (int i=1; i<n; i++)
#define N 10
using namespace std;
int X[N], Y[N], n, F[N], e;
struct E{
    int a,b;
    double c;
G[N*2];
struct V{
    double x, y;
    V operator+(double 1){
         int tmp=rand();
         return (V) \{x+l*cos(tmp), y+l*sin(tmp)\};
}v[N];
int find(int x){
    return x = F[x]?x:F[x] = find(F[x]);
inline double pow2(double x){
    return x*x;
double check (V s1, V s2) {
    double ans=0;
    e=0;v[4]=s1,v[5]=s2;
    F(5)Fl(j,i+1,6)
         G[e++]=(E)\{i,j,sqrt(pow2(v[i].x-v[j].x)+pow2(v[i].x)\}
             i].y-v[j].y))};
    F(6)F[i]=i;
    sort(G,G+e,[](E a,E b)\{return a.c < b.c;\});
    F(e) {
         if (find (G[i].a)!=find (G[i].b)) {
             ans+=G[ i ] . c;
             F[find(G[i].a)]=find(G[i].b);
         }
    return ans;
int main() {
    scanf("%d",&n);
    while (n--) {
         F(4) scanf("%lf%lf",&v[i].x,&v[i].y);
         double ttans, tans, ans, 11=10000, 12;
         V s1=(V) \{0,0\}, s2=(V) \{0,0\}, ts1, ts2, tmp;
         ans=check(s1, s2);
         while (l1>1e-3) {
             12 = 10000;
             ts1=s1+l1;
             tans=check(ts1,s2);
             tmp=s2;
             while (12>1e-3) {
                  ts2=s2+l2;
                  ttans=check(ts1,ts2);
                  if (ttans<tans) tans=ttans, s2=ts2;
                  else 12*=0.9;
             if (tans<ans) ans=tans, s1=ts1;
             else 11 = 0.9, s2 = tmp;
         printf("%f \setminus n", 2*ans);
    }
}
```

2 Graph

2.1 HLD

```
//we can reference the problem Greatest graph
///http://acm.csie.org/ntujudge/problemdata/2582.pdf
//this template operate on edges
#define maxn 100005
struct segment_tree{
    #define right(x) x << 1 | 1
    #define left(x) x << 1
    int* arr;
    int m[4*maxn];</pre>
```

```
int tag[4*maxn];
     const int inf = 1e9;
    void init() {
           /\text{memset}(\text{tag}, -1, \text{sizeof}(\text{tag}));
          fill(tag, tag+4*maxn, inf);
     void pull(int ind) {
              m[ind] = min(m[right(ind)], m[left(ind)]);
     void push(int ind) {
         if(tag[ind] != inf) {
              tag[left(ind)] = min(tag[left(ind)], tag[
                   ind]);
              tag[right(ind)] = min(tag[right(ind)], tag[
                   ind]);
              m[left(ind)] = min( m[left(ind)], tag[left(
                   ind)])
              m[right(ind)] = min(m[right(ind)], tag[
                   right (ind)]);
              tag[ind] = inf;
         }
     /// \text{ root} \Rightarrow 1
     void build(int ind, int l, int r) {
          if( r - l == 1) {
              m[ind] = arr[l];
              return;
         \begin{array}{ll} {\bf int} \ \ {\rm mid} \ = \ (\ l{+}r \ ){>}{>}1; \end{array}
         build( left(ind), l, mid );
         build( right(ind), mid, r );
         pull(ind);
     int query_min(int ind, int L, int R, int ql, int qr
         if (L >= qr \mid \mid R <= ql) return 1e9;
         if (R \ll qr \&\& L >= ql) \{
              return m[ind];
         }
         push(ind);
         \quad \text{int} \ \operatorname{mid} = (L\!+\!R) >> 1;
         return min( query_min(left(ind), L, mid, ql, qr
              ), query_min(right(ind), mid, R, ql, qr));
     void modify(int ind, int L, int R, int ql, int qr,
         int x) {
         if (L >= qr \mid \mid R <= ql) return;
         if(R \le qr \&\& L = ql)
              m[ind] = min(m[ind], x);
              tag[ind] = min(tag[ind], x);
              return:
         push(ind);
         int mid = (L+R) >> 1;
         modify(left(ind), L, mid, ql, qr, x);
         modify(\,right\,(ind\,)\,,\,\,mid\,,\,\,R,\,\,ql\,,\,\,qr\,,\,\,x)\,;
         pull(ind);
    }
};
int seg_arr[maxn];
struct Tree{
    segment_tree seg;
    int n;
    struct Edge { int u, v, c; };
     vector<Edge> e;
     \begin{array}{c} void \ addEdge(int \ x, \ int \ y, \ int \ c) \ \{ \\ G[x].pb(\ SZ(e)\ ); \end{array} 
         G[y].pb(SZ(e));
         e.pb(Edge\{x, y, c\});
    int siz [maxn], max_son[maxn], pa[maxn], dep[maxn];
     /*size of subtree index of max_son, parent index >
          depth*/
    \verb|int link_top[maxn]|, \verb|link[maxn]|, \verb|timer|;
     /*chain top index in segtree ime stamp*/
    std::vector<int >G[maxn];
    void init(int N) {
```

```
n = N:
    e.clear();
    for(int i = 1; i <= n; i++) G[i].clear();
    timer=0;
    pa[1] = 1;
    dep[1] = 0;
void find_max_son(int x){
    siz[x]=1;
    \max_{son}[x] = -1;
    for(int e\_ind : G[x]) {
        int v = e[e\_ind].u == x ? e[e\_ind].v : e[
            e_ind].u
        find_max_son(v);
        if(max\_son[x] = -1 \mid \mid siz[v] > siz[max\_son]
             [x]])
            \max_{son}[x] = v;
        siz[x] += siz[v];
    }
void build_link(int x, int top){
    link[x] = timer++;/*記錄x點的時間戳*/
    link\_top[x] = top;
    if(\max_{x \in \mathbb{R}} |x| != -1)
        build_link( max_son[x], top);/*優先走訪最大
             孩子*/
    for(int e\_ind : G[x]) {
        int v = e[e\_ind].u == x ? e[e\_ind].v : e[
             e\_ind].u;
        if(v = pa[x])
            seg\_arr[link[x]] = e[e\_ind].c;
        if (v = \max_{son}[x] \mid | v = pa[x]) continue
        // edge from x \Rightarrow v
        build_link(v, v);
    }
inline int lca(int a, int b){
    /*求LCA, 可以在過程中對區間進行處理*/
    int ta=link\_top[a],tb=link\_top[b];
    while(ta != tb){
        if (dep[ta]<dep[tb]) {
            std::swap(ta,tb);
            std::swap(a,b);
        //interval [ link[ta], link[a] ]
        a = pa[ta];
        ta = link\_top[a];
    return dep[a] < dep[b] ? a:b;
}
int modify(int a, int b, int c){
    int ta=link_top[a], tb=link_top[b];
    while(ta != tb){
        if (dep[ta]<dep[tb]) {
            std::swap(ta,tb);
            std::swap(a,b);
        //interval [ link[ta], link[a] ]
        //same interval if operate on edges
        seg.modify(1, 1, n, link[ta], link[a]+1, c)
        a = pa[ta];
        ta = link\_top[a];
    //a, b are on the same chain
    if( a = b ); // interval [ link[a], link[a]], if operate on edges \Rightarrow no edge
        if (dep[a]>dep[b])
            swap(a,b)
        //interval [ link[a], link[b] ]
        // if operate on edges \Longrightarrow [ link[ max_son[
            a] ], link[b] ]
```

```
seg.modify(1, 1, n, link[max\_son[a]],
                   link[b]+1, c);
         }
    }
/*
     void modify(int a, int b, int c) {
         if ( a==b ) return;
         if( link\_top[a] = link\_top[b]) {
              if \, (\ dep \, [\, a\, ] \, > \, dep \, [\, b\, ] \ ) \ swap (a \, , \ b) \, ;
              seg.modify(1,\ 1,\ n,\ link\,[a]{+}1,\ link\,[b]{+}1,\ c
              assert( link[a]+1 = link[max_son[a]]);
              return;
         if(dep[link\_top[a]] < dep[link\_top[b]])
              swap(a, b);
         // a is the node with deeper link_top
         seg.modify(1, 1, n, link[link_top[a]], link[a]
               + 1. c):
         modify( pa[link_top[a]], b, c);
    /// Heavy Light Decomposition
    void HLD() {
          // root is indexed 1 here !
         find_max_son(1);
         build_link(1, 1);
}tree;
int main() {
    int T; cin>>T;
     while (T--) {
         int n,m;
         scanf("%d %d",&n, &m);
         int ans = 0;
         tree.init(n);
         for (int i=0; i< n-1; i++) {
              int a, b, c;
              {\rm scanf}\,(\,\text{``'}\!\text{d}\!\text{''}\!\text{d}\!\text{''}\!\text{d}\!\text{''},&a,&b,&c\,)\;;
              //a--, b--; be careful here
              tree.addEdge(a, b, c);
              ans += c;
         tree.HLD();
         tree.seg.arr = seg_arr;
         tree.seg.build(1, 1, n);
    return 0;
2.2 Hungarian
```

```
// edge and node index starting from 0
  dfs version below
//complexity O ( V*E )
/* to do
#define ___maxNodes
num\_left = ?
struct Edge {
    int from:
    int to;
    int weight;
    Edge(int f, int t, int w):from(f), to(t), weight(w)
         {}
};
vector<int> G[__maxNodes]; /* G[i] 存储顶点 i 出发的边
    的编号 */
vector<Edge> edges;
int num_nodes;
int num_left;
int num_right;
int num_edges;
int matching[__maxNodes]; /* matching result */
int check [___maxNodes];
```

```
bool dfs(int u) {
   for (auto i = G[u]. begin(); i != G[u]. end(); ++i) {
        // 对 u 的每个邻接点
       int v = edges[*i].to;
       if (!check[v]) { // 要求不在交替路中
          check[v] = true; // 放入交替路
           if (matching[v] == -1 || dfs(matching[v]))
              // 如果是未盖点,说明交替路为增广路,则
                  交换路径,并返回成功
              matching[v] = u;
              matching[u] = v;
              return true;
       }
   }
   return false; // 不存在增广路,返回失败
int hungarian() {
   int ans = 0;
   memset(matching, -1, sizeof(matching));
   for (int u=0; u < num_left; ++u) {
       if (matching[u] = -1) {
          memset(check, 0, sizeof(check));
          if (dfs(u)) + ans;
       }
   return ans;
```

2.3 KM

```
// 最小帶權匹配~ km算法
//http://acm.csie.org/ntujudge/contest_view.php?id=836&
     contest\_id{=}449
#include <bits/stdc++.h>
using namespace std;
struct bipartite {
    #define maxn 602
    #define INF 0xfffffff
     \quad \text{int } \operatorname{sx}\left[\operatorname{maxn}\right], \ \operatorname{sy}\left[\operatorname{maxn}\right], \ \operatorname{mat}\left[\operatorname{maxn}\right]\left[\operatorname{maxn}\right];
     int x[maxn], y[maxn], link[maxn];
     int N, M, slack;
     int DFS(int t) {
          int tmp;
          sx[t] = 1;
          for (int i = 0; i < M; i++) {
               if (!sy[i]) {
                    tmp = x[t] + y[i] - mat[t][i];
                    if (tmp == 0) {
                         sy[i] = 1;
                          if (link[i] = -1 || DFS(link[i]))
                              \,l\,i\,n\,k\,\,[\,\,i\,\,]\,\,=\,\,t\,\,;
                              return 1;
                    else if (tmp < slack) slack = tmp;
               }
          }
          return 0;
     int KM() {
          for (int i = 0; i < N; i++) {
               x[i] = 0;
               for (int j = 0; j < M; j++) {
                    \inf [\max[i][j] > x[i]) x[i] = \max[i][j];
          for (int j = 0; j < M; j++) { y[j] = 0; }
          memset(link, -1, sizeof(link));
          for (int i = 0; i < N; i++) {
               while (1) {
                    memset(sx, 0, sizeof(sx));
                    memset(sy\,,\ 0\,,\ {\tt sizeof}(sy))\,;
                    slack = INF;
```

```
if (DFS(i)) break;
                 \quad \text{for (int } j \, = \, 0; \ j \, < \, N; \ j++) \, \, \{
                      if (sx[j]) x[j] = slack;
                 for (int j = 0; j < M; j++) {
                      if (sy[j]) y[j] += slack;
                 }
             }
        }
        int ans = 0;
        int cnt = 0;
        int t;
        for (int i = 0; i < M; i++)
        {
             t = link[i];
             if (t >= 0 \&\& mat[t][i] != -INF)
                 cnt ++:
                 ans += mat[t][i];
        // 最大權 : 沒有負號
        return -ans;
    void init(int n, int m) {
        N\,=\,n\,,\ M=m;
        for (int i = 0; i < N; i++)
            for (int j = 0; j < M; j++)
                 mat[i][j] = -INF;
    void input() {
        for(int i = 0; i < N; i++)
             for (int j = 0; j < M; j++) {
                 // fill in mat[i][j]
                 // stands for the weighting , but
                      negative sign !
                 // if 最大權 : 沒有負號
             }
}km;
int main(){
    int n,E;
    while (scanf("%d", &n) != EOF)
    {
        km.init(n, n);
        km.input();
        cout<< km.KM() <<endl;
    return 0;
```

2.4 Bi-vertex-connected Subgraph

```
regard every vbcc as a set of edges
      so vb[i] is a vector that contains a set of edge
    indexes
/** needed for tarjan **/
#define maxn 100005
#define maxm 100005
int n, m;
struct Edge{int s, t;};
vector<Edge> edge;
\quad \quad \text{int} \ dfn\left[maxn\right], \ low\left[maxn\right];
stack<int> st
bool vis [maxn];
int Time;
bool vis_e [maxm];
int bcnt, vbb[maxm];
vector<int> vb[maxm];
vector<int> G[maxn];
/** **/
void tarjan(int s){
    dfn[s] = low[s] = ++Time;
     vis[s] = true;
    for(int e_ind : G[s]){
         if (!vis_e[e_ind]) {
```

```
vis_e [e_ind] = true; st.push(e_ind)
                  \label{eq:int_to_edge} \begin{array}{l} \texttt{int} \hspace{0.1cm} \texttt{to} \hspace{0.1cm} = \hspace{0.1cm} \texttt{edge} \hspace{0.1cm} [\hspace{0.1cm} \texttt{e} \hspace{0.1cm} \texttt{ind} \hspace{0.1cm}] \hspace{0.1cm} \texttt{.} \hspace{0.1cm} \texttt{s} \hspace{0.1cm} \texttt{;} \end{array}
                  if (! vis [to]) {
                       tarjan(to);
                       low[s] = min(low[s], low[to]);
                        if(low[to] >= dfn[s]) {
                             vb[bcnt].clear();
                              \mathbf{while}(1){
                                   int t = st.top(); st.pop();
                                   vbb[t] = bcnt;
                                   vb[bcnt].push_back(t);
                                   if(t == e_ind) break;
                             bcnt++;
                       }
                 }else
                       low[s] = min(low[s], dfn[to]);
           }
      }
void init_tarjan() {
     mem(vis, false); mem(vis_e, false);
     Time = bcnt = 0; edge.clear();
      for (int i = 1; i \le n; i++) G[i]. clear ();
int main() {
      cin >> n >> m;
      init_tarjan();
      for (int i = 0; i < m; i++) {
           int a, b; scanf("%d %d"
                                               &a, &b);
           edge.push\_back(Edge\{a,b\});
           G[a].push\_back((int)edge.size()-1);
           G[b].push_back((int)edge.size()-1);
      tarjan(1);
}
```

2.5 Bi-edge-connected Subgraph

```
/** needed for tarjan **/
#define maxn 100005
#define maxm 100005
int n, m;
int dfn[maxn], low[maxn];
stack<int> st;
int Time;
int bcnt;
vector<int> G[maxn];
bool in_cyc[maxn];
void tarjan(int s, int p){
    dfn[s] = low[s] = ++Time;
    st.push(s);
    for(int to : G[s]) if( to != p ){
    if(!dfn[to]) {
             tarjan(to, s);
             low\,[\,s\,] \;=\; min\,(\,low\,[\,s\,]\,\,,\;\; low\,[\,to\,]\,)\;;
             if(low[to] > dfn[s]) {
                 // is cut_edge
                  // pop stack 的過程也可以寫在這
                  // 但最後(after tarjan)還要多判stack
                      not empty的情况
                  if ( low [to] > dfn [s]) {
                  in\_cyc[bcnt] = st.top()!=to;
                  while (1) {
                      int t = st.top(); st.pop();
                      id[t] = bcnt;
                      if (t == to) break;
                  bcnt++;
```

```
}
             }
         }else
             low[s] = min(low[s], dfn[to]);
    if(low[s] = dfn[s]){
         in\_cyc\,[\,bcnt\,] \;=\; st\,.\,top\,(\,)\,!{=}\,s\,;
         while (1) {
             int t = st.top(); st.pop();
             id[t] = bcnt;
             if(t == s) break;
         bcnt++;
    }
void init_tarjan() {
    Time = bcnt = 0;
int main() {
  cin >> n >> m;
  init_tarjan();
  for (int i = 0; i < m; i++) {
         int a, b; scanf("%d %d", &a, &b);
        G[a].pb(b), G[b].pb(a);
  mem( in_cyc , false);
  tarjan(1, 1);
}
```

2.6 SCC

```
struct SCC{
    \#define maxn 1005
     vector<int> G[maxn];
     stack<int> Stack;
     int scnt, Time;
     int belong [maxn], dfn[maxn], low[maxn];
     bool instack [maxn];
     void init(int n) {
          scnt = Time = 0;
          \label{eq:formula} \mbox{for(int $i = 0$; $i < n$; $i++$) $G[i].clear()$;}
          while(!Stack.empty()) Stack.pop();
         memset(dfn, 0, sizeof(dfn));
         memset(instack, false, sizeof(instack));
     void dfs(int u) -
          dfn[u] = low[u] = ++Time;
          Stack.push(u)\,;\ instack\,[u]\,=\, {\tt true}\,;
          for(int v : G[u]) {
              if ( !dfn[v] ) {
                   dfs(v);
                   low[u] = min(low[u], low[v]);
              else if (instack[v])
                   low\left[u\right] \,=\, min(low\left[u\right],\ dfn\left[v\right])\,;
          if(low[u] = dfn[u]) {
              scnt++;
              int \ {\rm tp}\,;
              do{
                   tp = Stack.top(); Stack.pop();
                   instack[tp] = false;
                   belong[tp] = scnt;
              \} while (tp != u);
          }
     void tarjan(int n) {
          //zero based here
          for (int i = 0; i < n; i++)
              if (!dfn[i])
                   dfs(i);
};
```

2.7 Edmond's Matching Algorithm

```
//http://acm.csie.org/ntujudge/contest_view.php?id=370&
    contest_id=466
#include <bits/stdc++.h>
using namespace std;
//带花树, Edmonds's matching algorithm, 一般图最大匹配
// have to be a undirected graph
#define MAXN 505
vector<int>G[MAXN];//用vector存圖
int pa [MAXN], match [MAXN], st [MAXN], S [MAXN], vis [MAXN];
int t,n;
inline int lca(int u, int v){//找花的花托
    for(++t; swap(u,v))
        if (u==0)continue;
        if(vis[u]==t)return u;
        vis [u]=t;//這種方法可以不用清空vis 陣列
        u=st[pa[match[u]]];
#define qpush(u) q.push(u),S[u]=0
inline void flower (int u, int v, int l, queue<int> &q) {
    while (st[u]!=1){
        pa[u]=v; //所有未匹配邊的pa都是雙向的
        if (S[v=match[u]]==1)qpush(v); //所有奇點變偶點
        st[u]=st[v]=l,u=pa[v];
inline bool bfs(int u){
    for (int i=1;i<=n;++i)st[i]=i;//st[i]表示第i個點的集
    memset(S+1,-1,sizeof(int)*n);//-1:沒走過 0:偶點 1:
        奇點
    queue<int>q;qpush(u);
    while (q. size()) {
        u=q. front(), q. pop();
        for (size_t i=0; i < G[u]. size();++i){
            int v=G[u][i];
            if(S[v]==-1){
                pa[v]=u,S[v]=1;
                if (!match[v]) {//有增廣路直接擴充
                    for(int lst;u;v=lst,u=pa[v])
                        lst=match[u], match[u]=v, match[v]
                            ]=u;
                    return 1;
                qpush (match [v]);
            else\ if\ (!S[v]\&\&st[v]!=st[u])
                int l=lca(st[v],st[u]);//遇到花, 做花的
                flower(v,u,l,q), flower(u,v,l,q);
        }
    return 0;
inline int blossom(){
   memset(pa+1,0,sizeof(int)*n);
    memset(match+1,0,sizeof(int)*n);
    int ans=0;
    for (int i=1; i \le n; ++i)
        if (!match[i]&&bfs(i))++ans;
    return ans;
void solve() {
    cin>>n;
    int m; cin>>m;
    while (m--) {
        int a,b;
        scanf("%d %d", &a, &b);
        a++, b++;
        // since node indexed [ 1 .. n ] in this
            template
        #define pb push_back
        //Multiedge and self-cycles are not forbidden
        G[a].pb(b);
        G[b].pb(a);
    cout<< blossom() <<endl;</pre>
    for (int i = 1; i \le n; i++) G[i]. clear();
int main() {
```

```
int t; cin>>t;
while(t--) solve();
```

}

2.8 Tree Decomposition

```
//codeforces Digit Tree
//http://codeforces.com/problemset/problem/715/C
typedef long long 11;
bool isprime [100005];
vector<LL> primes;
LL M, PHI;
#define MOD M
ll modpow(ll a, ll b) {
  11 r = 1;
   while(b) {
    if(b\&1) r=(r*a)\%MOD;
     a=(a*a)MOD;
    b >>= 1:
  return r;
}
void Sieve(int n) {
  memset(isprime, 1, sizeof(isprime));
  isprime[1] = false;
  for (int i = 2; i \le n; i++) {
     if(isprime[i]) {
       primes.pb(i);
       for (int j = 2*i; j \le n; j += i)
         isprime[j] = false;
  }
}
LL phi(LL n) {
    ll num = 1;    ll num2 = n;
   for(11 i = 0; primes[i]*primes[i] \le n; i++) {
     if (n%primes [i]==0) {
       num2/=primes[i];
       num*=(primes[i]-1);
     while (n\%primes[i]==0) {
       n/=primes[i];
   if(n>1) {
    num2/=n; num*=(n-1);
  n = 1;
  num \ *= \ num2;
  return num;
ll inv(ll a) {
  return modpow(a, PHI-1);
#define maxn 100005
struct edge{
     int u, v, dig;
     int no(int x) {
         };
vector<edge> e;
vector < int > G[maxn];
LL n, ans;
bool vis [maxn];
int sz[maxn], dep[maxn];
LL \ tenPow[maxn];
int dfs(int u, int p, int d) {
     sz[u] = 1;
     dep[u] = d;
     for (int eind : G[u] ) {
         \begin{array}{ll} \textbf{int} & v \, = \, e \, [\, eind \, ] \, . \, no(u) \, ; \end{array} \label{eq:varphi}
              v = p \mid | vis[v] ) continue;
         sz[u] += dfs(v, u, d+1);
     return sz[u];
```

```
cin >> n >> M:
int findCenter(int u, int p, int treesize) {
                                                                     Sieve( 100000 );
     for (int eind : G[u] ) {
                                                                       PHI = phi(M);
                                                                       for (int i = 0; i < n-1; i++) {
    int a, b, c; scanf(``%d %d %d ``, &a, &b, &c);
    G[a].pb(SZ(e)); G[b].pb(SZ(e));
         int v = e[eind].no(u);
         if (v = p || vis[v]) continue;
if (sz[v]*2 > treesize)
              return findCenter( v, u, treesize);
                                                                           e.pb(edge{a, b, c});
    return u;
                                                                       //init
                                                                       tenPow[0] = 1;
                                                                       for (int i = 1; i < maxn; i++) tenPow[i] = (tenPow[i]
                                                                           -1]*10)%M;
LL up [maxn], down [maxn];
int belong[maxn];
                                                                       ans = 0;
map<LL, LL> tot;
                                                                       mem( vis, false);
                                                                       solveAll(0);
vector<int> pt;
                                                                       cout << ans << endl;
void calc(int u, int p, int b, int d) {
    pt.pb(u);
    belong[u] = b;
                                                                  2.9
                                                                        Tree Longest Path
    \mathrm{dep}\,[\,u\,]\ =\ d\,;
    \begin{array}{lll} int & id = & find\_if( & all(G[u]) &, [u,p](int & x) & \{ & return \\ & e \, [\, x\,] \,.\, no(u) == p; & \}) & - & G[u] \,.\, begin(); \end{array}
                                                                  /** codeforces 592D - Super M **/
                                                                  #include <bits/stdc++.h>
    down[u] = (down[p]*10 + e[G[u][id]].dig)M;
    up[u] = (tenPow[d-1]*e[G[u][id]].dig + up[p])
                                                                  using namespace std;
                                                                  #define mp make_pair
    for(int eind : G[u]) {
                                                                  #define pb push_back
         \begin{array}{ll} int \ v = e \, [\, eind \, ]. \, no(u) \, ; \\ if ( \ vis \, [v] \ || \ v == p \, ) \ continue \, ; \end{array}
                                                                  #define LL long long
                                                                  #define pii pair<int,int>
         calc(v, u, b, d+1);
                                                                  #define PII pair < long long, long long>
    }
                                                                  #define fi first
                                                                  #define se second
     vec[b][ up[u] ]++;
    tot[ up[u] ]++;
                                                                  const int inf = 1e9;
                                                                  const LL INF = 1e18;
                                                                  const int mod = 1e9 + 7;
LL solve(int cent) {
                                                                  #define maxn 123460
    //cent is the root now
     vector<int> L;
                                                                  int n, m;
     for(int eind : G[cent]) {
                                                                  vector<int> g[maxn];
         int v = e[eind].no(cent);
                                                                   bool is [maxn];
         if (!vis[v]) {
                                                                  int dep[maxn], R, max_depth, A;
             L.pb(v);
                                                                  int cnt[maxn], parent[maxn];
         }
                                                                  bool dfs(int u, int par = 0){
     vec.clear();
                                                                    parent[u] = par;
    vec.resize(SZ(L), {});
                                                                     dep[u] = dep[par] + 1;
                                                                     if(dep[u] > max\_depth \&\& is[u])
     tot.clear();
    up[cent] = down[cent] = 0;
                                                                       \max_{\underline{\phantom{a}}} depth = dep[u], R = u;
    dep[cent] = 0;
                                                                     bool ret = is[u];
     pt.clear();
                                                                     for(int v : g[u])
     for (int i = 0; i < SZ(L); i++)
                                                                       if (v != par)
         calc( L[i], cent, i, 1);
                                                                         ret = dfs(v, u);
                                                                     if (ret) A++;
    LL ret = 0;
                                                                     return ret;
                                                                  }
    for(int u : pt) {
         LL tmp = (-down[u]+M)\%M;
         tmp = (tmp*inv(tenPow[dep[u]]))%M;
                                                                  int find_center(int start) {
         ret += tot[ tmp ] - vec[ belong[u] ][ tmp ];
                                                                    R = start; dep[0] = -1; max_depth = 0;
                                                                     dfs(start);
                                                                     max\_depth = 0; \ dep[R] = -1;
    assert( (LL)count_if(all(pt), [] (int x) { return
                                                                     dfs(R, R);
         up[x] = 0; \} ) = tot[0]);
    LL tmp = tot[0] + (LL)count_if(all(pt), [] (int x)
                                                                     int ret = R, d = max_depth/2;
         \{ return down[x] = 0; \} );
                                                                     while(d>0)
    debug("\%lld \n", tmp);
                                                                       d--;
    return ret+tmp;
                                                                       ret = parent[ret];
void solveAll(int node) {
                                                                     return ret;
    dfs(node, -1, 0);
     int cent = findCenter(node, -1, sz[node]);
                                                                  int S, dis, max_length;
    ans += solve( cent );
                                                                  bool dfs1(int u, int par = 0) {
    debug("\%d \%lld \n", cent, ans);
                                                                     dep[u] = dep[par] + 1;
     vis[cent] = true;
                                                                     if ( is [u] )
    for(int eind : G[cent] ) {
                                                                       if(dep[u] > max\_length)
         int v = e[eind].no(cent);
                                                                         max\_length = dep[u], S = u;
         if (vis[v]) continue;
                                                                       else if (dep[u] = max_length & u < S)
         solveAll(v);
                                                                         S = u:
    }
                                                                     bool c = false;
int main() {
                                                                     for(int v : g[u])
```

```
if(v!=par)
      dfs1(v, u);
int main(){
  cin >> n >> m;
  for (int i = 0; i < n-1; i++){
   int a, b; scanf("%d%d",&a, &b);
   g[a].pb(b), g[b].pb(a);
 memset(is, false, sizeof(is));
  for (int i = 0; i < m; i++){
   cin>>tmp; is [tmp] = true;
  int C = find_center(tmp);
  dep[0] = -1;S = inf; dis = (max_depth+1)/2;
  // distance(center, any other node) <= (longestpath +
       1) / 2
  dfs1(C);
 if ( max_depth & 1)
    dfs1(parent[C]);
  cout << S << endl << A-2-max_depth << endl;
  return 0;
```

3 Flow

```
//Circulation problems
http://www.win.tue.nl/~nikhil/courses/2013/2WO08/
    scribenotes 26 febv 02.pdf
Flow problems with boundary
1. feasible flow in a network with both upper and lower
    capacity constraints, no source or sink:
  capacities are changed to upper bound - lower bound.
  Add a new source and a sink.
  let M[v] = (sum of lower bounds of ingoing edges to v
      ) - (sum of lower bounds of outgoing edges from
      v).
  For all v,
    if M[v]>0 then add edge (S,v) with capacity M,
    otherwise add (v,T) with capacity -M.
   Actually, this equals to doing the following steps
       for every edges
   edges u --> v with lb
   addEdge(s, v, lb) and addEdge(u, t, lb)
   If all outgoing edges from S are full, then a
       feasible flow exists, it is the flow plus the
       original lower bounds.
2.maximum flow in a network with both upper and lower
    capacity constraints, with source s and sink t:
  //referenced from the book 挑戰城市競賽acm-icpc and
  //http://web.engr.illinois.edu/~jeffe/teaching/
      algorithms/2009/notes/18-maxflowext.pdf
  a. add edge (t,s) with capacity infinity.
  // Binary search for the lower bound, check whether a
       feasible exists for a network WITHOUT source or
      sink ??
  b.\ add\ new\ source\ and\ sink\,,\ ss\ and\ tt
      for all edges u --> v with lb
      addEdge(\ ss\ ,\ v\ ,\ lb\,)\ \ \text{and}\ \ addEdge(\ u\ ,\ tt\ ,\ lb\,)
  c. f1 = maxFlow of current graph
  d. if ss -> other verticles aren't all used \Longrightarrow no
      feasible solution
  e. \ addEdge(ss\,,\ s\,,\ inf)\,,\ addEdge(t\,,\ tt\,,\ inf)\,,
      removeEdge(t, s) (not necessary)
  f. f2 = maxFlow of current graph
  c. final answer will be f2 - the sum of edge demands
```

3.1 Dinic Maxflow

```
//http://acm.csie.org/ntujudge/problem.php?id=2581
//French Fries Festival
//dinic runs in O( V^2*E )
#define maxn 500
struct Edge{ int to, cap, rev; };
struct Dinic{
    vector < Edge > G[maxn];
    int dis[maxn], iter[maxn];
    void init(int n) {
        //zero based
      for(int i = 0; i < n; i++) G[i].clear();
    void addEdge(int from, int to, int cap) {
        vector < Edge > :: iterator it;
        if( ( it=find_if( all(G[from]), [to](Edge& e) {
              return e.to == to; } )) != G[from].end() )
             i\,t\,\text{-}\!>\!\mathrm{cap}\;+\!\!=\;\mathrm{cap}\,;
             return;
      G[from].pb(Edge\{to\,,\ cap\,,\ (int)G[to].size\,()\,\})\,;
      G[to].pb(Edge\{from, 0, (int)G[from].size()-1\});
        //if undirected 0 will be cap
    bool bfs(int s, int t) {
      memset(dis, -1, sizeof(dis));
      queue<int> que;
      que.push(s); dis[s] = 0;
      while (!que.empty())
        int tp = que.front(); que.pop();
        for (Edge &e : G[tp]) {
           if(e.cap > 0 \&\& dis[e.to] = -1)
             dis[e.to] = dis[tp] + 1, que.push(e.to);
        }
      }
      return dis[t] != -1;
    int dfs(int v, int t, int f) {
      if(v) = t) return f;
      for (int &i = iter[v]; i < G[v]. size(); i++) {
        Edge &e = G[v][i];
        if(e.cap > 0 \&\& dis[v] < dis[e.to]) {
           int d = dfs(e.to, t, min(f, e.cap));
           if(d > 0) {
             e.cap -= d;
             G[e.to][e.rev].cap += d;
            f += d;
             return d;
        }
      }
      return 0:
    int maxFlow(int s, int t) {
      int ret = 0;
      while(bfs(s, t))  {
        memset(iter, 0, sizeof(iter));
        int f;
        while ((f = dfs(s, t, inf)) > 0)
           ret += f:
      }
      return ret;
}dinic , dinic2;
void solve() {
    int n, m, k; cin>>n>>m>>k;
    // flow problem with lower bounds;
    int s = 0, t = n+2, ss = n+3, tt = n+4;
    {\tt dinic.init(n+5)};
    dinic.addEdge(s, 1, k);
    dinic.addEdge(n+1, t, k);
    int slb = 0;
    while (m--) {
        int l, r, a, b; scanf("%d %d %d %d", &l, &r, &a
             , &b);
        slb += a;
        r++;
```

```
\label{eq:dinic.addEdge(l, r, b-a);} \\ \text{dinic.addEdge(l, r, b-a);} \\
          dinic.addEdge(ss, r, a);
          dinic.addEdge(l, tt, a);
     dinic2 = dinic;
     dinic.addEdge(t, s, k);
     int f1 = dinic.maxFlow(ss, tt);
     if (!all_of (all (dinic.G[ss]), [] (Edge x) { return
          x.cap = 0;  } ) } {
          puts("-1"); return;
     {\tt dinic2.addEdge(ss\,,\,\,s\,,\,\,1e9)}\,;
     dinic2.addEdge(t, tt, 1e9);
     int f2 = dinic2.maxFlow(ss, tt);
     // maxflow in current graph is f2 - slb
     printf("%d\n", (f2 - slb)*n);
int main() {
     \begin{array}{ll} \textbf{int} & t \; ; cin >\!\!> t \; ; \end{array}
     while (t - -)
          solve():
```

3.2 Sw Mincut

```
//referenced from bcw's codebook
#include <cstdio>
#include <iostream>
#include <algorithm>
using namespace std;
{\color{red} \textbf{struct SW} \{ \ // \ O(V^3) \ 0\text{-base} }
     {\tt static \ const \ int \ MXN=514;}
     int n, vst [MXN], del [MXN];
     int edge [MXN] [MXN] , wei [MXN] ;
     void init(int _n){
         n\,=\,\underline{}\,n;
         for (int i=0; i< n; i++) {
              for (int j=0; j< n; j++)
                   {\rm edge}\,[\,i\,\,]\,[\,j\,]\,=\,0\,;
              del[i] = 0;
         }
     void add_edge(int u, int v, int w){
         edge[u][v] += w;
         void search (int &s, int &t) {
         for (int i=0; i< n; i++)
              vst[i] = wei[i] = 0;
         s = t = -1;
          while (true){
              int mx=-1, cur=0;
              for (int i=0; i<n; i++)
                   if (!del[i] && !vst[i] && mx<wei[i])
                        cur = i, mx = wei[i];
              if (mx == -1) break;
              vst[cur] = 1;
              s = t;
              t = cur;
              for (int i=0; i<n; i++)
                   if (!vst[i] && !del[i]) wei[i] += edge[
         }
     int solve(){
         int res = 2147483647;
          for (int i=0,x,y; i< n-1; i++){
              search(x,y);
              \mathtt{res} \, = \, \min(\,\mathtt{res}\,\,,\mathtt{wei}\,[\,\mathtt{y}\,]\,) \,\,;
              del[y] = 1;
              for (int j=0; j< n; j++)
                   edge[x][j] = (edge[j][x] += edge[y][j])
          return res;
```

```
}
}graph;
int main() {
    int n, m;
    while(cin>> n >> m ) {
        graph.init(n);
        while(m--) {
            int a, b, c; scanf("%d %d %d", &a, &b, &c);
            graph.add_edge(a, b, c);
        }
        cout << graph.solve() << endl;
    }
}</pre>
```

4 Data Structure

4.1 Disjoint Set

```
struct Disjoint_set {
    #define MAX_N 500005
    // define MAX_N
    int pa [MAX_N], Rank [MAX_N];
    int sz [MAX_N];
    void init_union_find(int V) {
        for(int i=0; i<V; i++) {
            pa[i] = i;
            Rank[i] = 0;
            sz[i] = 1;
    int find(int x) {
         return x == pa[x] ? x : pa[x] = find(pa[x]); 
    int unite(int x, int y) {
        x = find(x), y = find(y);
        int S = sz[x]+sz[y];
        if(x != y)
            if(Rank[x] < Rank[y]) {
                pa\left[\,x\,\right] \;=\; y\,;
                 sz[y]=S;
                 return y;
            }
             else{
                 pa[y] = x;
                 sz[x] = S:
                 if(Rank[x] = Rank[y]) Rank[x] ++;
                 return x;
            }
        }
    bool same(int x, int y) {
        return find(x) = find(y);
```

$4.2 \quad \text{Djs} + \text{Seg}$

```
demo => undo djs + segtree with offline
// this program doesn't consider the problem of
     overflowing varaible ans
   http://acm.csie.org/ntujudge/view_code.php?id
     =\!108190\&\mathtt{contest\_id}\!=\!472
#define maxn 100005
#define maxm 500005
//can be used to solve dynamic connectivity problem
//can be used with segment tree \Longrightarrow offline
struct DisjointSet {
  // save() is like recursive
  // undo() is like return
  int n, fa[maxn], sz[maxn];
vector<pair<int*,int>>> h;
  vector < int > sp;
  int ans;
  void init(int tn) {
    ans = 0;
```

```
n=tn:
    for (int i=0; i<n; i++) {
       fa[i]=i;
       sz\ [\ i\ ]\!=\!1;
    sp.clear(); h.clear();
  void assign(int *k, int v) {
    h.\, pb(\{k\,,\ ^{*}k\})\,;
    *k=v;
  void save() { sp.pb(SZ(h)); }
  void undo() {
    assert(!sp.empty());
    int last=sp.back(); sp.pop_back();
    while (SZ(h)!=last) {
       auto x=h.back(); h.pop_back();
       *x.fi=x.se;
    }
  int f(int x) {
    while (fa[x]!=x) x=fa[x];
    return x;
  void uni(int x, int y) {
    x=f(x); y=f(y);
    if (x=y) return ;
    if (sz[x] < sz[y]) swap(x, y);
    //nans stands for new answer
    int t = sz[x]+sz[y];
    int nans = ans - (sz[x]*sz[x]-sz[x]) - (sz[y]*sz[y]
         ]-sz[y]) + t*t-t;
    assign(\&sz\left[\,x\,\right]\,,\ sz\left[\,x\right]\!+\!sz\left[\,y\,\right]\,)\;;
    assign(&fa[y], x);
    assign(&ans, nans);
\}djs;
int n, m;
map < int, int > ma[maxn]:
vector<pii> seg[4*maxm];
LL ans [maxm];
void add(int ql, int qr, int a, int b, int id=1, int l
    =0, int r=m) {
    if (qr \ll l | | ql \gg r) return;
    if ( l >= ql && r <= qr )
         seg\left[\,id\,\right].\,pb\left(\ mp(\,a\,,\ b\,)\ \right)\,;
         return ;
    int mid = (l+r) >> 1;
    {\rm add} \left( \ ql \, , \ qr \, , \ a \, , \ b \, , \ id \, ^{*}2 \, , \ l \, , \ mid \right);
    add(ql, qr, a, b, id*2+1, mid, r);
void dfs(int u=1, int l=0, int r=m) {
    djs.save();
    for(pii v : seg[u] ) djs.uni( v.fi , v.se );
     if(r-l > 1)
         int mid = (l+r) >> 1;
         dfs(u*2, l, mid);
         dfs(u*2+1, mid, r);
         // do sth here
         ans[l] = djs.ans;
    djs.undo();
int main() {
    scanf("%d %d", &n, &m);

for(int i = 0; i < m; i++) {
         int a, b; scanf("%d %d",&a, &b);
         a--, b--; if(b < a) swap(a, b);
         if (ma[a].count(b)) {
              add(ma[a][b], i, a, b);
              ma[a].erase(b);
         else ma[a][b] = i;
    for(int i = 0; i < n; i++) if( !ma[i].empty() ) {
         for(auto p : ma[i])
```

```
add( p.se, m, i, p.fi);
djs.init(n);
dfs();
for (int i = 0; i < m; i++) printf("%lld\n", ans[i]);
```

Sparse Table

```
//codeforces 689D
#define maxn 200005
template < typename T, typename Cmp = less < T > >
struct RMQ {
     T d[maxn][20];
     Cmp cmp;
     int w[maxn], sz;
     void init (const T *a, int n) {
           int i, j;
           \mbox{for } (sz = n, \ i = 0; \ i < n; \ +\!\!\!+\!\! i) \ d[\,i\,][\,0\,] = a[\,i\,];
           for (j = 1; (1 << j) <= n; ++j) {
                 for (i = 0; i + (1 << j) <= n; ++i) {
                     \begin{array}{l} d[\,i\,][\,j\,] = cmp(d[\,i\,][\,j\,-\,1]\,,\; d[\,i\,+\,(1 <\!< \\ (\,j\,-\,1))\,][\,j\,-\,1])\,\,?\,\,d[\,i\,][\,j\,-\,1]\,:\; d \end{array}
                            [i + (1 \ll (j - 1))][j - 1];
           }
     // index of a [l .. r]
     const T &query(int l, int r) const {
           int x = w[r - l + 1];
            \begin{array}{c} \textbf{return} \ \text{cmp}(d[1][x]\,, \ d[r - (1 <\!\!< x) + 1][x]) \ ? \ d \\ [1][x] \ : \ d[r - (1 <\!\!< x) + 1][x]; \end{array} 
     }
int a[maxn], b[maxn];
int n;
RMQ \leq int > s;
RMQ<int, greater<int>>t;
int main() {
     cin>>n;
     for (int i = 0; i < n; i++) scanf("%d", &a[i]);
     for (int i = 0; i < n; i++) scanf ("%d", &b[i]);
     s.init(b, n);
     t.init(a, n);
     int c, d;
     LL ans = 0;
     for (int i=0; i< n; i++) {
           if(a[i] > b[i]) continue;
           int ub = n+1, lb = i;
           while (ub-lb>1) {
                 int mid = (ub+lb)>>1;
                 if(t.query(i, mid-1) - s.query(i, mid-1) >
                        0) ub = mid;
                \begin{array}{ll} \textbf{else} & \textbf{lb} = \textbf{mid}; \end{array}
           int up = ub;
           ub = n+1, lb = i;
            while (ub-lb>1) {
                 \begin{array}{ll} \hbox{int} & \hbox{mid} = (ub + lb) >> 1; \\ \end{array} 
                 if( t.query(i, mid-1) - s.query(i, mid-1)
                     >= 0) ub = mid;
                \begin{array}{ll} {\bf else} & {\bf lb} \ = \ {\rm mid} \, ; \end{array}
           int down = ub;
           ans += up-down;
     cout << ans << endl;
     return 0;
```

4.4 Link Cut Tree

```
//\,https://\,github.com/\,yzgysjr/\!ACM-ICPC-\,Templates/\,blob/
     master/Data%20Structure/Link%20Cut%20Tree.cpp
struct node { int rev; node *pre, *ch[2]; } base [MAXN],
      nil, *null;
typedef node *tree;
#define isRoot(x) (x->pre->ch[0] != x && x->pre->ch[1]
     != x)
#define isRight(x) (x->pre->ch[1] == x)
inline void MakeRev(tree t) { if (t != null) { t->rev
     \hat{} = 1; \text{ swap}(t->\text{ch}[0], t->\text{ch}[1]); \}
inline void PushDown(tree t) { if (t->rev) { MakeRev(t
     ->ch[0]); MakeRev(t->ch[1]); t->rev = 0; }
inline void Rotate(tree x) {
  \label{eq:tree_y} \texttt{tree} \ \ y \, = \, x\text{-}\!\!>\!\! pre\,; \ \ PushDown(\,y\,)\,; \ \ PushDown(\,x\,)\,;
  int d = isRight(x);
  if (!isRoot(y)) y->pre->ch[isRight(y)] = x; x->pre =
       y->pre
  if ((y-ch[d] = x-ch[!d]) != null) y-ch[d]-pre = y
  x->ch[!d] = y; y->pre = x; Update(y);
inline void Splay(tree x) {
  PushDown(x)\,;\;\; \textbf{for}\;\; (\,t\, ree\;\; y\,;\;\; !\, isRoot(x)\,;\;\; Rotate(x)\,) \;\; \{
    y = x->pre; if (!isRoot(y)) Rotate(isRight(x) !=
          isRight(y) ? x : y);
  } Update(x);
inline void Splay(tree x, tree to) {
  PushDown(x); for (tree y; (y = x->pre) != to; Rotate(
       x)) if (y->pre != to)
     Rotate(isRight(x) != isRight(y) ? x : y);
  Update(x);
inline tree Access(tree t) {
  tree last = null; for (; t != null; last = t, t = t->
       pre) Splay(t), t->ch[1] = last, Update(t);
  return last:
inline void MakeRoot(tree t) { Access(t); Splay(t);
    MakeRev(t); }
inline tree FindRoot(tree t) { Access(t); Splay(t);
     tree last = null;
   \begin{array}{lll} for & ( & ; & t != null ; & last = t \, , & t = t \hbox{-}\!\!> \hskip -1pt ch [0]) & PushDown(t \\ & & ) \, ; & Splay(last) \, ; & return & last \, ; \end{array} 
inline void Join(tree x, tree y) { MakeRoot(y); y->pre
    = x; 
inline void Cut(tree t) {Access(t); Splay(t); t->ch
     [0]->pre = null; t->ch[0] = null; Update(t);
inline void Cut(tree x, tree y) {
  tree upper = (Access(x), Access(y));
  if (upper == x) { Splay(x); y->pre = null; x->ch[1] =
         null; Update(x); }
  else if (upper == y) { Access(x); Splay(y); x->pre =
  null; y->ch[1] = null; Update(y); }
else assert(0); // 'impossible to happen'
inline int Query(tree a, tree b) { // 'query the cost
    in path a <-> b, lca inclusive
  Access(a); tree c = Access(b); // c is lca
  int v1 = c->ch[1]->maxCost; Access(a);
  \begin{array}{ll} \textbf{int} & v2 \, = \, c\text{--}\!>\! ch\left[1\right]\text{--}\!>\! \max\! Cost\,; \end{array}
  return \max(\max(v1, v2), c->cost);
void Init() {
  null = \&nil; null -> ch[0] = null -> ch[1] = null -> pre =
       null; null->rev = 0;
  Rep(i, 1, N)  { node & m = base[i]; n.rev = 0; n.pre = 0
       n.ch[0] = n.ch[1] = null;
//compressed version
//http://trinklee.blog.163.com/blog/static
     /238158060201521101957375/
const int N=30010;
\begin{array}{l} int \;\; n, fa\left[N\right], son\left[N\right]\left[2\right], val\left[N\right], siz\left[N\right], stmp, rev\left[N\right]; \\ \#define \;\; swap(a,b) \;\; (stmp\!=\!a, a\!=\!b, b\!=\!stmp) \end{array}
void pu(int t) \{ siz[t] = siz[son[t][0]] + siz[son[t][1]] + 1; \}
void pd(int t){rev[t]?rev[t]=0,rev[son[t][0]]^=1,rev[
     son[t][1]]^=1, swap(son[t][0], son[t][1]), 1:1;
```

```
bool \operatorname{nr}(\operatorname{int} t) \{ \operatorname{return} \operatorname{son} [\operatorname{fa}[t]][0] == t || \operatorname{son} [\operatorname{fa}[t]][1] == t || \operatorname{son} [\operatorname{
 void rtt(int t,int f=0,bool p=0){
                          p=son[f=fa[t]][1]==t,
                            fa[t]=fa[f], nr(f)?son[fa[f]][son[fa[f]][1]==f]=t:1,
                           (son[f][p]=son[t][!p])?fa[son[f][p]]=f:1,
                           pu(son[fa[f]=t][!p]=f);
void pv(int t){if(nr(t))pv(fa[t]);pd(t);}
void splay(int t,int f=0){
                            for (pv(t); nr(t); rtt(t)) nr(f=fa[t])?
                           rtt(son[f][1] == t^son[fa[f]][1] == f?t:f), 1:1; pu(t);
 void access(int t, int la=0){for(;t;splay(t),son[t][1]=
                          la , la=t , t=fa [t]);}
 void makeroot(int t){access(t), splay(t), rev[t]^=1;}
 void link(int u, int v)\{makeroot(u), fa[u]=v;\}
void cut(int u,int v){makeroot(u),access(v),splay(v),
                            son[v][0] = fa[u] = 0;
```

4.5 Treap

```
#include <bits/stdc++.h>
using namespace std;
struct Treap{
     Treap *1, *r;
     int pri, key, val;
     Treap(int _val, int _key):
           {\tt val(\_val)}\,,\ {\tt key(\_key)}\,,\ {\tt l(NULL)}\,,\ {\tt r(NULL)}\,,\ {\tt pri(}
                rand()){}
/// We assure that key value in A treap is greater than
       that in treap B
Treap *merge( Treap *a, Treap *b){
       if(a \!\!=\!\!\! NULL \mid\mid b \!\!=\!\!\! NULL) \ return \ (!a) \ ? \ b \ : \ a; 
      if(a->pri > b->pri)
          a \rightarrow r = merge(a \rightarrow r, b);
           return a;
     } else {
          b->l = merge(a, b->l);
           return b;
void split (Treap *t, int k, Treap *&a, Treap *&b) {
     if(!t) a = b = NULL;
     else if (t->key <= k)
          a = t;
          s\, p\, l\, i\, t\, (\, t\, \text{--}{>}\, r\; ,\;\; k\, ,\;\; a\, \text{--}{>}\, r\; ,\;\; b\, )\; ;
     }else{
           s\, p\, l\, i\, t\, (\, t\, \text{--}{>}\, l\; ,\;\; k\, ,\;\; a\, ,\;\; b\text{--}{>}\, l\, )\; ;
Treap* insert( Treap *t, int k, int _val){
     Treap *tl, *tr;
     split(t, k, tl, tr);
     return merge(tl, merge(new Treap(_val, k), tr));
Treap* remove( Treap* t, int k){
     Treap *tl, *tr;
     split(t, k-1, tl, t);
     split(t, k, t, tr);
     return merge(tl, tr);
int main(){
     return 0:
```

5 Math

5.1 Prime Table

```
#include <bits/stdc++.h>
using namespace std;
```

```
struct Prime_table {
    int prime [1000000] = \{2,3,5,7\};
    int sz=4;
     // biggest prime < ub
    int ub=(1<<20);
    int check(int num){
         int k = 0;
         for(k = 0; k < sz \&\& prime[k]*prime[k] <= num;
             if ( num % prime [k]==0) return 0;
         return 1;
    void buildprime(){
        int currentPrime=7;
         int i=4;
         for (sz=4, j=4; currentPrime < ub; sz++, j=6-j)
              currentPrime=currentPrime+j;
              if (check(currentPrime)) {
                 prime[sz] = currentPrime;
              else {
                 SZ - -;
        }
}ptable;
```

5.2 Miller Rabin Prime Test

```
#include <bits/stdc++.h>
using namespace std;
typedef\ long\ long\ LL;
LL mul(LL a, LL b, const LL mod) {
    LL x = 0, y = a \% mod;
    while (b > 0) {
        if (b&1)
            x = (x + y) \% mod;
        y = (y * 2) \% mod;
        b >>= 1;
    return x % mod;
LL mul(LL lhs, LL rhs, const LL mod) {
    return ( lhs * rhs ) \% mod;
LL mypow(LL b, LL e, const LL mod) {
    LL x = 1;
    LL v = b;
    while (e) {
        if (e\&1) x = mul(x, y, mod);
        y = mul(y, y, mod);
        e >>= 1;
    return x;
{\color{red} {\bf const \ int \ testbase}} \; [\;] \; = \; \{2, \; 3, \; 5, \; 7, \; 11, \; 13, \; 17, \; 19, \; 23, \;
     29, 31, 37};
bool isprime(const LL p) {
    if (p < 2) return false;
    if (p != 2 && !(p&1) ) return false;
    LL d = p - 1;
    while ( !(d\&1) ) d >>= 1;
    for( int a : testbase ) {
        LL td = d;
        if ( a >= p-1 ) return true;
        LL st = mypow(a, td, p);
         while (td!=p-1 && st!=1 & st!=p-1)
             st = mul(st, st, p);
             td \ll 1;
         if ( st != p - 1 && !(td&1) ) return false;
```

```
}
return true;
}
int main() {
    int T;
    scanf("%d",&T);
    while(T--) {
        LL q;
        scanf("%lld",&q);
        puts(isprime(q)?"YES":"NO");
    }
    return 0;
}
```

5.3 Extended Euclidean Algorithm

```
/** normal gcd function using recursion **/
int gcd(int a, int b){
   if(b == 0) return a;
   return gcd(b, a%b);
}
// Find solution of ax + by = gcd(a, b)
// ps : x, y may be negative
int extgcd(int a, int b, int& x, int& y){
   int d = a;
   if(b!= 0) {
      d = extgcd(b, a%b, y, x);
      y -= (a/b) * x;
   }else {
      x = 1, y = 0;
   }
   return d;
}
```

5.4 Gauss Elimination

```
// solving linear equations with gauss elimination
#include <iostream>
#include <cmath>
#include <vector>
using namespace std;
void print(vector< vector<double> > A) {
    int n = A. size();
    for (int i=0; i<n; i++) {
        for (int j=0; j< n+1; j++) {
            cout << "|
        cout << "\n";
    cout << endl;</pre>
vector<double> gauss(vector< vector<double> > A) {
    int n = A. size();
    for (int i=0; i< n; i++) {
        // Search for maximum in this column
        double maxEl = abs(A[i][i]);
        int maxRow = i;
        for (int k=i+1; k< n; k++) {
            if (abs(A[k][i]) > maxEl) {
               maxEl = abs(A[k][i]);
               \max Row = k;
            }
        }
        // Swap maximum row with current row (column by
             column)
        for (int k=i; k<n+1;k++) {
            double tmp = A[maxRow][k];
            A[\max Row][k] = A[i][k];
            A[i][k] = tmp;
        }
```

```
// Make all rows below this one 0 in current
              column
         \begin{array}{lll} \text{for (int $k$=$i+1; $k$<$n; $k$++) {} & \\ & \text{double c = -A[k][i]/A[i][i];} \end{array}
               for (int j=i; j< n+1; j++) {
                   if (i==j) {
                        A[k][j] = 0;
                   } else {
                        A[k][j] += c * A[i][j];
              }
         }
     }
     // Solve equation Ax=b for an upper triangular
          matrix A
     vector < double > x(n);
     for (int i=n-1; i>=0; i--) {
         x[i] = A[i][n]/A[i][i];
         for (int k=i-1; k>=0; k--) {
             A[k][n] -= A[k][i] * x[i];
     return x:
int main() {
    int n;
     cin >> n;
     vector < double > line(n+1,0);
     vector< vector<double> > A(n, line);
     // Read input data
     \quad \  \  \text{for (int } \ i\!=\!0; \ i\!<\!n\,; \ i\!+\!+\!) \ \{
          for (int j=0; j< n; j++) {
              cin >> A[i][j];
     for (int i=0; i< n; i++) {
          cin >> A[i][n];
     // Print input
     print(A);
     // Calculate solution
     vector < double > x(n);
     x = gauss(A);
     // Print result
     cout << "Result:\t";</pre>
     for (int i=0; i<n; i++) {
         cout << x[i] << "
     cout << endl;</pre>
5.5 FFT
//pEcaveros
```

```
//pecaveros
// const int MAXN = 262144;
// (must be 2^k)
// before any usage, run pre_fft() first
//
// To implement poly. multiply:
//
// fft( n , a );
// fft( n , b );
// for( int i = 0 ; i < n ; i++)
// c[ i ] = a[ i ] * b[ i ];
// fft( n , c , 1 );
//
// then you have the result in c :: [cplx]
typedef long double ld;
typedef complex<ld> cplx;
const ld PI = acosl(-1);
const cplx I(0, 1);
```

```
cplx omega[MAXN+1];
void pre_fft(){
  for (int i=0; i \le MAXN; i++)
    omega[i] = exp(i * 2 * PI / MAXN * I);
// n must be 2^k
void fft(int n, cplx a[], bool inv=false){
  int basic = MAXN / n;
  int theta = basic:
  \quad \  \  \text{for (int } m=n; \ m>=2; \ m>>=1) \ \{
    int mh = m \gg 1;
    for (int i = 0; i < mh; i++) {
      cplx w = omega[inv ? MAXN (i*theta%MAXN)
                            : i*theta%MAXN];
      \label{eq:formula} \mbox{for (int } j \, = \, i \, ; \ j \, < \, n \, ; \ j \, +\!\! = \, m) \ \{
        int k = j + mh;
         cplx x = a[j] - a[k];
        a[j] += a[k];
        a[k] = w * x;
      }
    theta = (theta * 2) % MAXN;
  int i = 0;
  for (int j = 1; j < n - 1; j++) {
    for (int k = n \gg 1; k > (i \hat{k} > k); k \gg 1);
    if (j < i) swap(a[i], a[j]);
  if (inv)
    for (i = 0; i < n; i++)
      a[i] /= n;
//http://sd-invol.github.io/2016/02/13/FFT-mod-prime/
struct Complex {
    double x , y
    Complex (double _x = 0, double _y = 0) {
        x = \_x , y = \_y;
    Complex operator + (const Complex &r) const {
        Complex operator - (const Complex &r) const {
         return Complex(x - r.x , y - r.y);
    Complex operator * (const Complex &r) const {
         return Complex(x * r.x - y * r.y , x * r.y + y
             * r.x);
    Complex conj () const {
         return Complex(x , -y);
    double operator = (const double a) {
         *this = Complex(a, 0);
         return a;
};
const double pi = acos(-1.0);
//fft with modulo, code referenced from the internet
    fftMod::fftPrepare(len);
    fftMod::convolution(res, le, ri, len, r-l);
namespace fftMod{
    const int N = 1 \ll 18;
    const int Mod = 1e9 + 7;
    // to do, M should be about sqrt (Mod)
    {\color{red}\textbf{const}} \ \ {\color{blue}\textbf{int}} \ \ M = \ 32768;
    Complex w[N];
    int rev[N];
    void fftPrepare(int n) {
         \begin{array}{ll} \textbf{int} & LN = \_\_builtin\_ctz(n); \end{array}
         for (int i = 0; i < n; ++ i) {
    double ang = 2 * pi * i / n;
             w[i] = Complex(cos(ang), sin(ang));
             rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (
                 LN - 1));
         }
```

```
void FFT(Complex P[], int n, int oper) {
           for (int i = 0; i < n; i ++) {
                 if (i < rev[i]) {
                      swap(P[i], P[rev[i]]);
           for (int d = 0; (1 \ll d) < n; d++) {
                int m = 1 \ll d, m2 = m * 2, rm = n / m2;
                for (int i = 0; i < n; i += m2) {
                      for (int j = 0; j < m; j++) {
                           Complex &P1 = P[i + j + m], &P2 = P
                                 [i + j];
                           Complex t = w[rm * j] * P1;
                           P1 = P2 - t;
                           P2 = P2 + t;
                      }
                }
           }
      }
      Complex A[N], B[N], C1[N], C2[N];
       \begin{tabular}{ll} \begin{tabular}{ll} void & convolution(vector < int > \& res \ , & vector < int > \& a \ , \\ \end{tabular} 
            vector<int> &b, int len, int K) {
           // a[ 0 .. len ) and b[ 0 .. len ) 's
                 convolution 8 Mod
             / stored in res[ 0 .. K+1 )
           for (int i = 0; i < len; ++ i) {
                A[i] = Complex(a[i] / M , a[i] \% M);
                B[i] = Complex(b[i] / M, b[i] \% M);
           FFT(A, len, 1); FFT(B, len, 1);
           int j = i? len - i : i;
                Complex a1 = (A[i] + A[j].conj()) * Complex
                      (0.5, 0);
                Complex a2 = (A[i] - A[j].conj()) * Complex
                      (0, -0.5);
                Complex b1 = (B[i] + B[j].conj()) * Complex
                      (0.5, 0);
                Complex b2 = (B[i] - B[j].conj()) * Complex
                      (0, -0.5);
                Complex c11 = a1 * b1 , c12 = a1 * b2;
                Complex c21 = a2 * b1 , c22 = a2 * b2;
                C1[j] = c11 + c12 * Complex(0, 1);
                C2[j] = c21 + c22 * Complex(0, 1);
           FFT(C1 , len , -1); FFT(C2 , len , -1);
           for (int i = 0; i \le K; ++ i) {
                \begin{array}{ll} \text{int } x = (LL) \, (C1[\,i\,] \, .x \, / \, \, len \, + \, 0.5) \, \, \% \, \, Mod; \\ \text{int } y1 = (LL) \, (C1[\,i\,] \, .y \, / \, \, len \, + \, 0.5) \, \, \% \, \, Mod; \end{array}
                 \label{eq:int_y2} \begin{array}{lll} \text{int} & y2 \, = \, (LL) \, (C2 [\, i \, ] \, .\, x \, \, / \, \, len \, + \, 0.5) \, \, \% \, \, \text{Mod}; \end{array}
                \begin{array}{ll} \text{int } z = (LL)(C2[i].y \ / \ len \ + \ 0.5) \ \% \ Mod; \\ \text{res}[i] = ((LL)x \ ^*M \ ^*M \ + \ (LL)(y1 \ + \ y2) \ ^*M \end{array}
                       + z) % Mod;
           }
      }
};
```

5.6 NNT

```
//pEcaveros
LL P=2013265921,root=31;
int MAXNUM=4194304;
// Remember coefficient are mod P
   p=a*2^n+1
    2^n
n
                               a.
                                     root
    32
                  97
                               3
                                     5
    64
                  193
                               3
                                     5
7
    128
                  257
                               2
                                     3
    256
                  257
                                     3
    512
                  7681
                                     17
                               15
10
    1024
                  12289
                               12
                                     11
11
    2048
                  12289
                               6
                                     11
                  12289
12
    4096
                               3
                                     11
13
    8192
                  40961
                                     3
14
    16384
                  65537
                               4
                                     3
15
    32768
                  65537
                                     3
```

```
65536
                  65537
16
                              1
17
     131072
                  786433
                              6
                                   10
     262144
                                   10 (605028353, 2308, 3)
                  786433
19
     524288
                  5767169
                             11
                                   3
20
     1048576
                  7340033
                                   3
21
     2097152
                  23068673
                              11
22
     4194304
                  104857601
                                   3
                             25
23
     8388608
                  167772161
                             20
                                   3
24
     16777216
                  167772161
                             10
                                   3
25
     33554432
                  167772161
                             5
                                   3 (1107296257, 33, 10)
                  469762049
26
     67108864
                                   31 */
27
    134217728
                  2013265921 15
LL bigmod(LL a, LL b) {
   if (b==0)return 1;
   LL inv(LL a,LL b){
   if (a==1) return 1;
   return (((LL)(a-inv(b\%a,a))*b+1)/a)\%b;
std::vector < LL > ps(MAXNUM), rev(MAXNUM);
struct poly{
   std::vector <\!\!LL\!\!>\ co\,;
   int n; //polynomial degree = n
   poly(int d) \{n=d; co.resize(n+1,0);\}
   void trans2(int NN){
     int r=0,st,N;
     unsigned int a,b;
     while((1 << r) < (NN > 1)) ++ r;
     for (N=2;N<=NN;N<<=1,--r) {
       for (st=0;st<NN;st+=N)
         int i, ss=st+(N>>1);
         for (i=(N>>1)-1; i>=0;--i)
           a=co[st+i]; b=(ps[i << r]*co[ss+i])%P;
           co[st+i]=a+b; if(co[st+i]>=P)co[st+i]-=P;
           co[ss+i]=a+P-b; if(co[ss+i]>=P)co[ss+i]-=P;
      }
     }
   void trans1(int NN){
     int r=0,st,N;
     unsigned int a,b;
     for (N=NN; N>1; N>>=1,++r) {
       for (st=0;st<NN;st+=N)
         int i, ss=st+(N>>1);
         for (i=(N>>1)-1; i>=0;--i)
           a=co[st+i]; b=co[ss+i];
           co[st+i]=a+b; if(co[st+i]>=P)co[st+i]-=P;
           co[ss+i]=((a+P-b)*ps[i<< r])%P;
       }
     }
   poly operator*(const poly& _b)const{
     poly a=*this, b=_b;
     int k=n+b.n,i,N=1;
     while (N = k)N^* = 2;
     a.co.resize(N,0); b.co.resize(N,0);
     int r=bigmod(root,(P-1)/N),Ni=inv(N,P);
     ps[0]=1;
     for (i=1; i \le N; ++i) ps [i]=(ps[i-1]*r)%P;
     a. trans1(N); b. trans1(N);
     for (i=0; i \le N; ++i) a. co [i] = ((LL) a. co [i] * b. co [i] \%P;
     r=inv(r,P);
     for (i=1; i < N/2; ++i) std :: swap (ps[i], ps[N-i]);
     a.trans2(N);
     for (i=0;i<N;++i)a.co[i]=((LL)a.co[i]*Ni)%P;
     a.n=n+_b.n; return a;
};
constexpr int mod = 1e9+7;
 typedef vector<int> VEC;
 // ntt + Crt, code referenced from the internet
namespace nttCrt {
     constexpr int magic [3] = \{1004535809, 998244353,
         104857601};
     constexpr int \widehat{MOD} = 1000000007:
     constexpr int G = 3;
     int P;
     inline int quick_mod(int x, int k, int MOD) {
```

```
int ans = 1:
     while (k) {
          if (k\&1) ans = 1LL * ans * x % MOD;
          x = 1LL * x * x % MOD;
          k >>= 1;
    return ans;
inline void change(int *y, int len) {
    \label{eq:for_int_i} \mbox{for(int } \mbox{$i=1$, $j=len$/ $2$; $i<len - 1$; $i++$) } \{
          if(i < j) swap(y[i], y[j]);</pre>
          //交换互为小标反转的元素, i<j保证交换一次
          //i做正常的+1, j左反转类型的+1,始终保持i和j
               是反转的
          int k = len / 2;
          while(j >= k)  {
               j = k;
               k /= 2;
          if(j < k) j += k;
    }
inline void ntt(int *y, int len, int on) {
    change(y, len);
    \begin{array}{lll} & \text{for}\,(\inf \ h = 2; \ h <= len\,; \ h <<= 1) \ \{ \\ & \text{int} \ wn = quick\_mod(G, \ (P - 1) \ / \ h , \ P)\,; \end{array}
          for (int j = 0; j < len; j += h) {
               int w = 1;
               for (int k = j; k < j + h / 2; k++) {
                    y[k] = (u + t) \% P;
                    y[k + h / 2] = ((u - t) \% P + P) \%
                        р.
                    w = 1LL * w * wn \% P;
               }
          }
     if(on = -1) {
          for (int i = 1; i < len / 2; i++)
              swap(y[i], y[len - i]);
          \begin{array}{lll} & \text{int inv} = \text{quick}\_\text{mod}(\text{len}\,,\,P-2\,,\,P)\,;\\ & \text{for}(\text{int}\,\,i\,=\,0\,;\,\,i\,<\,\text{len}\,;\,\,i+\!\!+\!\!) \end{array}
              y[i] = 1LL * y[i] * inv \% P;
    }
}
int n;
int r[3][3];
inline int CRT(int *a) {
    int sb[3] = \{a[0], a[1], a[2]\};
     for (int i = 0; i < 3; ++i) {
          for (int j = 0; j < i; ++j)
               int t = (sb[i] - sb[j]) \% magic[i];
               if(t < 0) t += magic[i];
               sb[i] = 1LL * t * r[j][i] % magic[i];
     int mul = 1, ans = sb[0] \% MOD;
     ans = (ans + 1LL * sb[i] * mul) % MOD;
    return ans;
int tmp[maxn][3];
int x1[maxn*2], x2[maxn*2];
inline void gao(vector<int>& res, vector<int> &a,
     vector<int> &b ,int len , int kk) {
     for (int ti = 0; ti < 3; ti++) {
         P = magic[ti];
          int k;
          \label{eq:formula} \mbox{for } (\ k = 0; \ k < SZ(a) \ \&\& \ k < len \, ; \ k+\!\!\! +) \ x1[
               \mathbf{k} \,] \,=\, \mathbf{a} \,[\,\mathbf{k}\,] \,;
          for (; k < len; k++) x1[k] = 0;
          \label{eq:continuous} \mbox{for } \mbox{ ( } \mbox{ k} = 0; \mbox{ k} < SZ(b) \&\& k < len; \mbox{ k++) } x2[
               k = b[k];
          for (; k < len; k++) x2[k] = 0;
```

```
ntt(x1, len, 1); ntt(x2, len, 1);
             for (int i = 0; i < len; i++) x1[i] = 1LL *
                  x1[i] * x2[i] % P;
             ntt(x1, len, -1);
             for (int i = 0; i \le kk; i++) tmp[i][ti] =
                 x1[i];
         for (int i = 0; i \le kk; i++) res[i] = CRT(tmp
             [ i ] )
    inline void init() {
         for (int i = 0; i < 3; i++) {
             for (int j = 0; j < 3; j++) {
                 r[i][j] = quick\_mod(magic[i], magic[j])
                     - 2, magic[j]);
        }
    }
};
```

//http://blog.csdn.net/hackbuteer1/article/details

5.7 Big Number

```
/6595881
#include<iostream>
#include<string>
#include<iomanip>
#include<algorithm>
using namespace std;
#define MAXN 9999
#define MAXSIZE 10
#define DLEN 4
class BigNum
private:
 int a[500];
              //可以控制大数的位数
 int len;
              //大数长度
public:
 BigNum() { len = 1; memset(a,0,sizeof(a)); }
                                       //构造函
 BigNum(const int);
                       //将一个int类型的变量转化为
     大数
                       //将一个字符串类型的变量转化
 BigNum(const char*);
     为大数
 BigNum(const BigNum &); //拷贝构造函数
 BigNum & operator=(const BigNum &); //重载赋值运算
     符, 大数之间进行赋值运算
  friend istream& operator>>(istream&, BigNum&);
     重载输入运算符
  friend ostream& operator<<(ostream&, BigNum&);</pre>
     重载输出运算符
 BigNum operator+(const BigNum &) const;
                                      //重载加法
     运算符,两个大数之间的相加运算
 BigNum operator - (const BigNum &) const;
                                      //重载减法
     运算符,两个大数之间的相减运算
 BigNum operator*(const BigNum &) const;
                                      //重载乘法
     运算符, 两个大数之间的相乘运算
 BigNum operator/(const int
                          &) const;
                                      //重载除法
     运算符,大数对一个整数进行相除运算
 BigNum operator (const int &) const;
                                     //大数的n次
     方运算
                                     //大数对一个
       operator%(const int &) const;
     int类型的变量进行取模运算
      operator > (const BigNum & T) const;
                                       //大数和另
     一个大数的大小比较
                                       //大数和一
      operator > (const int & t) const;
     个int类型的变量的大小比较
 void print();
                   //输出大数
```

```
int i:
BigNum::BigNum(const int b)
                                   //将一个int类型的变量转
                                                                  cout << b.a[b.len - 1];
                                                                  for (i = b.len - 2 ; i >= 0 ; i--)
    化为大数
                                                                    cout.width(DLEN);
  int c,d = b;
                                                                    cout. fill('0');
  len = 0;
                                                                    cout << b.a[i];
  memset(a,0,sizeof(a));
  while(d > MAXN)
                                                                  return out;
                                                               }
    c = d - (d / (MAXN + 1)) * (MAXN + 1);
    d = d / (MAXN + 1);
                                                               BigNum BigNum::operator+(const BigNum & T) const
    a[len++] = c;
                                                                     个大数之间的相加运算
  a[len++] = d;
                                                                  \operatorname{BigNum}\ t\left(*this\right);
BigNum::BigNum(const char*s)
                                    //将一个字符串类型的变
                                                                  int i, big;
                                                                                   //位数
                                                                  big = T.len > len ? T.len : len;
    量转化为大数
                                                                  for(i = 0 ; i < big ; i++)
  int t,k,index,l,i;
                                                                    t \, . \, a \, [ \, i \, ] \, +\!\!\!=\!\! T . \, a \, [ \, i \, ] \, ;
  memset(a, 0, sizeof(a));
                                                                    if(t.a[i] > MAXN)
  l=strlen(s);
  len=l/DLEN;
                                                                      t.a[i + 1]++;
  if (1%DLEN)
                                                                      t.a[i] -=MAXN+1;
    len++;
  index=0;
  for (i=l-1; i>=0; i-=DLEN)
                                                                  if (t.a[big] != 0)
                                                                    t.len = big + 1;
    k=i -DLEN+1;
                                                                    t.len = big;
    if(k<0)
                                                                  return t;
      k=0;
    for (int j=k; j \le i; j++)
      t=t*10+s[j]-'0';
                                                               BigNum BigNum::operator-(const BigNum & T) const
    a [index++]=t;
                                                                     个大数之间的相减运算
                                                                  int i,j,big;
                                                                  bool flag;
BigNum::BigNum(const BigNum & T) : len(T.len) //拷贝构
                                                                  BigNum t1, t2;
    造函数
                                                                  if (*this>T)
{
  int i;
                                                                  {
                                                                    t1 = *this;
  memset(a,0,sizeof(a));
                                                                    t2=T;
  for(i = 0 ; i < len ; i++)
    a[i] = T.a[i];
                                                                    flag = 0;
                                                                  else
BigNum & BigNum::operator=(const BigNum & n)
                                                   //重载赋
    值运算符, 大数之间进行赋值运算
                                                                    t1=T:
                                                                    t2 = *this;
  int i;
                                                                    flag=1;
  len = n.len;
  memset(a, 0, sizeof(a));
                                                                  big=t1.len;
  for(i = 0 ; i < len ; i++)
                                                                  for(i = 0 ; i < big ; i++)
   a[i] = n.a[i];
  return *this;
                                                                     if(t1.a[i] < t2.a[i])</pre>
istream& operator>>(istream & in, BigNum & b)
                                                                      j = i + 1;
    输入运算符
                                                                      \frac{\mathbf{while}(\mathbf{t}1.\mathbf{a}[\mathbf{j}] = 0)}{\mathbf{mhile}(\mathbf{t}1.\mathbf{a}[\mathbf{j}] = 0)}
                                                                        j++;
  char ch[MAXSIZE*4];
                                                                      t1.a[j--]--;
  int i = -1;
                                                                      while (j > i)
  in >> ch:
                                                                        t1.a[j--] += MAXN;
  int l=strlen(ch);
                                                                      t1.a[i] += MAXN + 1 - t2.a[i];
  int count=0,sum=0;
  for (i=l-1; i>=0;)
                                                                    else
                                                                      t1.a[i] -= t2.a[i];
    sum = 0;
    int t=1;
                                                                  t1.len = big;
    for(int j=0; j<4&&i>=0; j++,i--,t*=10)
                                                                  while (t1.a[len - 1] = 0 \&\& t1.len > 1)
      sum + = (ch [i] - '0') *t;
                                                                    t1.len--;
                                                                    big - -;
    b.a[count]=sum;
    count++;
                                                                  if (flag)
                                                                    t1.a[big-1]=0-t1.a[big-1];
  b.len =count++;
                                                                  return t1;
  return in;
                                                               }
                                                               BigNum BigNum::operator*(const BigNum & T) const
                                                     //重载
ostream& operator << (ostream& out, BigNum& b)
                                                                     个大数之间的相乘运算
    输出运算符
                                                                  BigNum ret;
```

```
\begin{array}{ll} \textbf{int} & i \;, j \;, up \,; \end{array}
  int temp, temp1;
  for (i = 0 ; i < len ; i++)
    up = 0;
    for(j = 0 ; j < T.len ; j++)
       temp \, = \, a \, [\, i \, ] \ * \ T.\, a \, [\, j \, ] \ + \ ret.\, a \, [\, i \ + \ j \, ] \ + \ up \, ;
       if(temp > MAXN)
         temp1 = temp - temp / (MAXN + 1) * (MAXN + 1);
         up = temp / (MAXN + 1);
         ret.a[i + j] = temp1;
       }
       else
       {
         up = 0;
         ret.a[i + j] = temp;
    if(up != 0)
       ret.a[\,i \,+\, j\,] \,=\, up\,;
  ret.len = i + j;
  while (ret.a[ret.len - 1] = 0 \&\& ret.len > 1)
    \operatorname{ret.len}-;
                                                                     {
  return ret;
BigNum BigNum::operator/(const int & b) const
                                                       //大数
    对一个整数进行相除运算
  {\rm BigNum\ ret}\;;
  int i, down = 0;
  for(i = len - 1 ; i >= 0 ; i--)
    ret.a[i] = (a[i] + down * (MAXN + 1)) / b;
    down \, = \, a \, [\, i \, ] \, + \, down \, * \, (MAXN \, + \, 1) \, - \, ret \, . \, a \, [\, i \, ] \, * \, b \, ;
                                                                     {
  ret.len = len;
  while(ret.a[ret.len - 1] == 0 \&\& ret.len > 1)
    ret.len--:
  return ret;
                                                        //大数对
int BigNum::operator %(const int & b) const
     一个int类型的变量进行取模运算
  int i, d=0;
                                                                    }
  for (i = len -1; i>=0; i--)
    d = ((d * (MAXN+1))\% b + a[i])\% b;
  }
  return d:
BigNum BigNum::operator^(const int & n) const
                                                          //大数
    的n次方运算
 BigNum t, ret(1);
  int i;
  if (n<0)
    exit(-1);
  if(n==0)
    return 1;
  if(n==1)
    return *this;
  int m=n;
  while (m>1)
    t=*this:
    for (i=1; i << 1 <= m; i << =1)
       t=t*t;
    }
    m=i;
    ret=ret*t;
    if (m==1)
       ret=ret*(*this);
  }
  return ret;
bool BigNum::operator>(const BigNum & T) const
                                                          //大数
     和另一个大数的大小比较
```

```
int ln;
  if (len > T.len)
    return true;
  else if (len = T.len)
    ln = len - 1;
    while (a[ln] = T.a[ln] \&\& ln >= 0)
     ln - -;
    if(ln >= 0 \&\& a[ln] > T.a[ln])
      return true;
    else
      return false;
  }
  else
    return false;
bool BigNum::operator >(const int & t) const
                                                  //大数
    和一个int类型的变量的大小比较
 BigNum b(t);
  return *this>b;
void BigNum::print()
                        //输出大数
  int i;
  cout \ll a[len - 1];
  for (i = len - 2 ; i >= 0 ; i--)
    cout.width(DLEN);
    cout. fill('0');
    cout << a[i];
  cout << endl;
int main(void)
  int i,n;
  BigNum x[101];
                       //定义大数的对象数组
  x[0]=1;
  for (i=1; i<101; i++)
  x[i]=x[i-1]*(4*i-2)/(i+1);
while (scanf("%d",&n)==1 && n!=-1) {
   x[n].print();
5.8 Simplex
```

```
//reference from bcw's codebook
const int maxn = 111;
const int maxm = 111:
const double eps = 1E-10;
double a [maxn] [maxm], b [maxn], c [maxm], d [maxn] [maxm];
double x [maxm];
int ix [maxn + maxm]; // !!! array all indexed from 0
// \max\{cx\} \text{ subject to } \{Ax \le b, x > = 0\}
// n: constraints, m: vars !!!
// x[] is the optimal solution vector
// usage :
// value = simplex(a, b, c, N, M);
double simplex (double a [maxn] [maxm], double b [maxn],
    double c [maxm], int n, int m) {
    ++m;
    \begin{array}{l} \text{int } r=n, \ s=m-1; \\ \text{memset}(d,\ 0,\ \text{sizeof}(d)); \end{array}
    for (int i = 0; i < n + m; ++i) ix[i] = i;
    for (int i = 0; i < n; ++i) {
         for (int j = 0; j < m - 1; ++j) d[i][j] = -a[i]
              ][j];
         d[i][m - 1] = 1;
         d[i][m] = b[i];
         if (d[r][m] > d[i][m]) r = i;
    for (int j = 0; j < m - 1; ++j) d[n][j] = c[j];
    d[n + 1][m - 1] = -1;
    for (double dd;; ) {
```

```
if (r < n) {
           int t = ix[s]; ix[s] = ix[r+m]; ix[r+m]
           \begin{array}{l} d[\,r\,][\,s\,] \ = \ 1.0 \ / \ d[\,r\,][\,s\,]; \\ for \ (int \ j = 0; \ j <= m; \ +\!\!\!+\!\!\! j) \ if \ (j \ !\!\! = s) \ d[ \end{array}
                r][j] *= -d[r][s];
           for (int i = 0; i \le n + 1; ++i) if (i != r
                for (int j = 0; j <= m; ++j) if (j != s
) d[i][j] += d[r][j] * d[i][s];
                d[i][s] *= d[r][s];
           }
     }
     r = -1; s = -1;
     \label{eq:formula} \mbox{for (int } j = 0; \ j < m; \ +\!\!\!+\!\!\! j) \ \mbox{if (} s < 0 \ || \ \mbox{ix} [s]
           > ix[j]) {
           \begin{array}{l} if \ (d[n+1][j] > eps \ || \ (d[n+1][j] > - \\ eps \ \&\& \ d[n][j] > eps)) \ s = j; \end{array}
     if (s < 0) break;
     for (int i = 0; i < n; ++i) if (d[i][s] < -eps)
           if (r < 0 | | (dd = d[r][m] / d[r][s] - d[i]
                 [m] / d[i][s] < -eps || (dd < eps &&
                ix[r + m] > ix[i + m]) r = i;
     if (r < 0) return -1; // not bounded
if (d[n + 1][m] < -eps) return -1; // not
     executable
double ans = 0;
for (int i=0; i < m; i++) x[i] = 0;
for (int i = m; i < n + m; ++i) { // the missing
     enumerated x[i] = 0
     if (ix[i] < m - 1)
     {
           ans += d[i - m][m] * c[ix[i]];
           x[ix[i]] = d[i-m][m];
return ans;
```

6 string

6.1 Palindromic Tree

```
回文自動機包含以下元素:
  狀態St, 所有節點的集合, 一開始兩個節點, 0表示偶數長
     度串的根和1表示奇數長度串的根
  last 新增加一個字符後所形成的最長回文串的節點編號
  s 當前的字符串(一開始設s[0]=-1(可以是任意一個在串S
     中不會出現的字符))
  n 表示添加的字符個數
每個節點代表一個不同的回文子字串,我們在每個節點會儲存
  一些數值:
  len 表示所代表的回文子字串長度
  next[c] 表示所代表的回文子字串在頭尾各增加一個字符c
     後的回文字串其節點編號
  sufflink 表示所代表的回文子字串不包括本身的最長後綴
     回文子串的節點編號
  cnt(非必要) 表示所代表的回文子字串在整體字串出現的
     次數(在建構完成後呼叫count()才能計算)
  //num(非必要) 表示所代表的回文子字串其後綴為回文字
     串的個數 <== not included
struct palindromic_tree{
  struct node{
     int next [26], sufflink, len; /*這些是必要的元素*/int l, r; // this node is s[ l .. r ]
     int cnt, num;
                     /*這些是額外維護的元素*/
```

```
node(int l=0): sufflink(0), len(l), cnt(0), num(0)
            for (int i=0; i<26;++i) next [i]=0;
    };
    std::vector<node> St;
    std::string s; //current string [ 1 .. n ]
    int last ,n;
    palindromic\_tree():St(2), last(1), n(0){
        St[0].sufflink=1;
        St[1].len=-1;
        s.push_back(-1);
    inline void clear(){
        St.clear();
        s.clear();
        last=1;
        n=0;
        St.push\_back(0);
        St.push_back(-1);
        St[0].sufflink=1;
        s.push_back(-1);
    inline int get_sufflink(int x){
        while (s[n-St[x].len-1] != s[n]) x=St[x].
            sufflink;
        return x;
    inline void add(int c){
        s.push_back(c-= 'a');
        int cur=get_sufflink(last);
        if (!St[cur].next[c]) {
            int now=St.size();
            St.push_back(St[cur].len+2);
            St [now]. sufflink=St [get_sufflink(St [cur].
                 sufflink)].next[c];
             /*不用擔心會找到空節點,由證明的過程可知*/
            St [cur].next[c]=now;
            St[now].num=St[St[now].sufflink].num+1;
            St[now].l = n - St[now].len + 1, St[now].r
        last=St [cur].next[c];
        ++St[last].cnt;
    inline void count(){/*cnt必須要在構造完後呼叫count
        ()去計算*/
        std::vector<node>::reverse_iterator i=St.rbegin
            ();
        for (; i!=St.rend();++i) {
            St[i->sufflink].cnt+=i->cnt;
    inline int size(){/*傳回其不同的回文子串個數*/
        return St.size()-2;
}ptree;
```

6.2 Suffix Array

6.3 Longest Palindromic Substring

```
//ntu judge Earse
#define maxn 200001
char t [maxn];
char s [maxn * 2];
int z [maxn * 2];
int N;
int longest_palindromic_substring() {
    // t穿插特殊字元, 存放到s。
    int n = strlen(t);
    N = n * 2 + 1;
    memset(s, '.', N);
    for (int i=0; i<n; ++i) s[i*2+1] = t[i];
    s[N] = '\0';
    z[0] = 1;    // if無須使用, then無須計算。

int L = 0, R = 0;
    for (int i=1; i<N; ++i) // 從z[1]開始
```

```
{
       z[i] = (R > i) ? min(z[2*L-i], R-i) : 1;
       while (i - z[i]) = 0 \&\& i + z[i] < N \&\&
              s[i-z[i]] = s[i+z[i]]) z[i]++;
       if (i+z[i] > R) L = i, R = i+z[i];
   }
    // 尋找最長迴文子字串的長度
   n = 0;
   int p = 0;
    for (int i=1; i≪N; ++i) // 從z[1]開始
       if (z[i] > n)
           n = z[p = i];
    // longest 從中心到外端的長度 ⇒ (n-2)/2
                                                       //PECaveros
    //cout << "最長迴文子字串的長度是" << (2*n-1) / 2;
                                                            r2 ){
    // 印出最長迴文子字串, 記得別印特殊字元。
                                                         double d = sqrt(d2);
       for (int i=p-z[p]+1; i<=p+z[p]-1; ++i)
                                                         if(d > r1 + r2) return {};
           if (i & 1) {
               cout \ll s[i];
                                                            +d));
   return (2*n-1)/2;
                                                         return {u+v, u-v};
int nxt[maxn * 2];
int main() {
                                                      int T; cin>>T;
    while (T--) {
       scanf("%s", t);
       #ifdef DEBUG
                                                         double f = (f1 + f2);
           cout << longest_palindromic_substring() <<</pre>
       #else
           longest_palindromic_substring();
       #endif
       memset(nxt, -1, sizeof(nxt));
       for (int i = 0; i < N; i++) {
           nxt[i-z[i]+1] = i+1;
       int leftmost = 0;
       for (int i = 0; i < N; i++) {
           leftmost = max(leftmost, nxt[i]);
           nxt[i] = max(leftmost, nxt[i]);
       int ans = 0;
       for (int cur = 0; cur < N-1;) {
           cur = nxt[cur];
           ans++:
       cout << ans << endl;
    return 0;
}
```

geometry

Point Class 7.1

```
const double eps = 1e-10;
#define N 100
struct P {
   {\color{red} \textbf{double}} \ x\,,\ y\,;
   P(double \_x=0, double \_y=0) : x(\_x), y(\_y) {};
   void read() {
       scanf("%lf%lf",&x,&y);
   void print() {
       printf("%f %f\n", x, y);
} p[N];
bool operator <( Pa, Pb) { return tie(a.x,a.y)<tie(b
   .x,b.y);
```

```
P operator *( P b, double a ) { return P{a*b.x,a*b.y};
P operator /( P a, double b ) { return P{a.x/b,a.y/b};
P& operator /=( P &a, double b ) { return a=a/b; }
double operator *( Pa, Pb) { return a.x*b.x+a.y*b.y;
double operator ^( Pa, Pb) { return a.x*b.y-a.y*b.x;
double x( P o, P a, P b ) { return (a-o)^(b-o); }
double dot(Po, Pa, Pb) { return (a-o)*(b-o); }
```

7.2 Intersection of Circles/Lines/Segments

```
vector<P> interCircle( P o1 , double r1 , P o2 , double
    double d2 = (01 - 02) * (01 - 02);
   P\ u = (o1 + o2)*0.5 + (o1 - o2)*((r2*r2 - r1*r1)/(2*d2));
    double A = sqrt((r1+r2+d)*(r1-r2+d)*(r1+r2-d)*(-r1+r2-d)
   P \ v = P(\ o1.y-o2.y \ , \ -o1.x + o2.x \ ) * A / (2*d2);
                                            ( q1 - p1 );
    double f^2 = (p^2 - p^2) - (q^2 - p^2);
     \begin{array}{l} if (\ fabs (\ f\ ) < eps\ ) \ return\ Pt (\ nan ("")\ ,\ nan ("")\ )\ ; \\ return\ q1\ *\ (\ f2\ /\ f\ )\ +\ q2\ *\ (\ f1\ /\ f\ )\ ; \\ \end{array} 
return ret / max( 111 , abs( ret ) );
 // p1 == p2 || q1 == q2 need to be handled
bool banana (const PLL& p1 , const PLL& p2
                      \begin{array}{l} \mbox{if} (\ (\ p2\ -\ p1\ )\ ^- (\ q2\ -\ q1\ )\ ) == 0\ )\{\ //\ \ parallel \\ \mbox{if} (\ ori(\ p1\ ,\ p2\ ,\ q1\ )\ )\ return\ false; \\ \mbox{return} \ (\ (\ p1\ -\ q1\ )\ ^* \ (\ p2\ -\ q1\ )\ ) <= 0\ || \\ \mbox{(} \ (\ p1\ -\ q2\ )\ ^* \ (\ p2\ -\ q2\ )\ ) <= 0\ || \\ \mbox{(} \ (\ q1\ -\ p1\ )\ ^* \ (\ q2\ -\ p1\ )\ ) <= 0\ || \\ \mbox{(} \ (\ q1\ -\ p2\ )\ ^* \ (\ q2\ -\ p2\ )\ ) <= 0; \\ \end{array} 
    return (ori( p1, p2, q1 ) * ori( p1, p2, q2 )<=0) &&
               (ori(q1, q2, p1) * ori(q1, q2, p2)<=0);
```

7.3 Convex Hull

```
#define REP(i,n) for ( int i=0; i<int(n); i++)
int n;
void input() {
     scanf("%d",&n);
    REP(i,n) p[i].read();
P findCenter() {
    p[n]=p[0];
    P center=P\{0,0\};
    REP(i,n) {
         double v=p[i]*p[i+1];
         \text{center.x} \; +\!\!=\; (p\,[\,i\,\,]\,.\,x\!+\!\!p\,[\,i\,+\!1].\,x)\,^*v\,;
         center.y += (p[i].y+p[i+1].y)*v;
     double area=0;
    REP(i,n) area+=p[i]*p[i+1];
     area \neq 2;
     center /= 6*area;
     return center;
```

```
P q1 [N], q2 [N], q [N];
void convex() {
    sort(p,p+n);
    int m1=0,m2=0;
   REP(i,n) {
        while (m1>=2 && X(q1[m1-2],q1[m1-1],p[i]) >= 0
             ) m1--;
        while (m2>=2 \&\& X(q2[m2-2],q2[m2-1],p[i]) <= 0
             ) m2--
        q1 [m1++]=q2 [m2++]=p[i];
    int m=0;
   REP(i, m1) q[m++]=q1[i];
    for ( int i=m2-2; i>=1; i-- ) q[m++]=q2[i];
   q[m]=q[0];
void solve() {
    convex();
    // continue ...
```

7.4 Half Plane Intersection

```
// \texttt{http://acm.csie.org/ntujudge/problemdata/2575.pdf}
//http://www.csie.ntnu.edu.tw/~u91029/Half-
   planeIntersection.html
預先使用四個半平面, 設定一個極大的正方形邊界, 讓半平面
    交集擁有邊界。
 二、逐一加入每個半平面,求出當下的半平面交集(凸多邊
online 演算法,隨時維護一個半平面交集。每次更新需時 O(N
   ),總時間複雜度為 O(N<sup>2</sup>), N 是半平面數目。
#include <bits/stdc++.h>
using namespace std;
#define mp make_pair
typedef complex<double> Point;
typedef vector<Point> Polygon;
typedef pair<Point, Point> Line;
#define x real()
#define y imag()
// 兩向量叉積
double cross (Point& a, Point& b) {
   return a.x * b.y - a.y * b.x;
// 向量oa與向量ob進行叉積
double cross(Point& o, Point& a, Point& b) {
    return (a.x-o.x) * (b.y-o.y) - (a.y-o.y) * (b.x-o.x)
       );
}
// 多邊形面積
double area (Polygon& p) {
   double a = 0;
   int n = p.size();
   for (int i=0; i< n; ++i)
       a += cross(p[i], p[(i+1)\%n]);
   return fabs(a) / 2;
}
 / 兩線交點
Point intersection (Point& a1, Point& a2, Point& b1,
   Point& b2) {
   Point a = a2 - a1, b = b2 - b1, s = b1 - a1;
   return a1 + a * cross(b, s) / cross(b, a);
 / 一個凸多邊形與一個半平面的交集
Polygon halfplane_intersection(Polygon& p, Line& line)
   Polygon q;
```

Point p1 = line.first, p2 = line.second;

```
// 依序窮舉凸多邊形所有點,判斷是否在半平面上。
     // 如果凸多邊形與半平面分界線有相交, 就求交點。
     int n = p.size();
     for (int i=0; i<n; ++i)
         double c = cross(p1, p2, p[i]);
         double d = cross(p1, p2, p[(i+1)\%n]);
         if (c >= 0) q.push_back(p[i]);
          \begin{tabular}{ll} \textbf{if} & (c * d < 0) \\ \end{tabular} \ q. push\_back (intersection (p1, p2, p2, p3)) \\ \end{tabular} 
               p[i], p[(i+1)\%n]);
     return q;
#define maxn 550
//Line line[maxn];
Point v[maxn];
double ans[maxn];
int main() {
    int T; cin>>T;
     while (T--) {
         int n;
         double
                  w, h;
         scanf("%d %lf %lf", &n, &w, &h);
          // 預先設定一個極大的正方形邊界
         Polygon p, org;
         /** initialize
          p.push_back(Point(-1e9,-1e9));
         p.push_back(Point(-1e9,+1e9));
         p.push_back(Point(+1e9,-1e9));
         p.push\_back(Point(+1e9,+1e9));
         p.push\_back(Point(0,0));
         p.push\_back(Point(0,h));
         p.push_back(Point(w,h));
         p.push_back(Point(w,0));
         org = p;
         for(int i = 0; i < n; i ++) {
              double a, b;
              scanf("%lf %lf", &a, &b);
              v\left[\,i\,\right] \;=\; Point\left(\,a\,,\;\;b\,\right)\,;
         // 每一個半平面都與目前的半平面交集求交集
         for (int i=0; i<n; ++i)
              p = org;
              for (int j = 0; j < n; j++) {
                   if(i==j) continue;
                   Line line;
                   // find perpendicular line to line i_j
                   Point a( (v[i].x+v[j].x)/2, (v[i].y+v[j].x)
                        [.y)/2;
                   Point b(a.x+(v[i].y-v[j].y), a.y-(v[i].
                        x-v[j].x));
                   \label{eq:line} \mbox{line} \, = \, \mbox{cross} \, (\mbox{a} \, , \, \, \mbox{b} \, , \, \, \mbox{v} \, [\, \mbox{i} \, ] \, ) \, > \!\! = \!\! 0 \, \; ? \, \, \mbox{mp} (\mbox{a} \, , \, \, \mbox{b})
                        : mp(b,a);
                   p = halfplane_intersection(p, line);
                   if (area(p) == 0) break; // 退化或者
                        空集合
              }
              ans[i] = area(p);
         for (int i = 0; i < n; i ++) printf ("%.9f\n", ans [
              i]);
    }
}
/*
10
3 4 4
1 1 2 2 3 3
```