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## 1 Basic

.vimrc

```
ap jj <Esc>
  on
  sw=4 ts=4 sts=4 et nu sc hls cc=69
 let plugin indent on
 n <F5> :!"./%<"<CR>
 n <F6> :!"./%<" < input.txt<CR>
  FileType cpp no <F9> :!g++% -o
  \% -std=c++14 -O3 -Wall -Wextra
  -Wshadow -Wno-unused-result<CR>
  <expr> <silent> <Home> col('.') ==
  match(getline('.'),'\S') + 1
  ? '0' : '^'
  <silent><Home><C-O><Home>
    Increase Stack Size
 stack resize
 m("mov %0,%%esp\n"::"g"(mem+10000000));
 change esp to rsp if 64-bit system
 stack resize (linux)
 nclude <sys/resource.h>
 _{
m oid} increase_stack_size() +
 const rlim_t ks = 64*1024*1024;
 struct rlimit rl;
 int res=getrlimit(RLIMIT_STACK, &rl);
  if(res==0){
   if(rl.rlim_cur<ks){</pre>
     rl.rlim_cur=ks;
     res=setrlimit(RLIMIT STACK, &rl);
 }
     digitDP
 \cdot 3
  介
  名思义,所谓的数位DP就是按照数字的个、十、百、千·····
   位数进行的DP。
  位DP的 题 目 有 着 非 常 明 显 的 性 质:
   询问[1,r]的区间内,有多少的数字满足某个性质
  法根据前缀和的思想,求出[0,1-1]和[0,r]中满足性质的数
   的个数,然后相减即可。
 法核心
 dfs(int x, int pre, int bo, int limit);
 般需要以上参数 (当然具体情况具体分析)
   x表示当前的数位 (一般都是从高位到低位)
   pre表示前一位的数字
   bo可以表示一些附加条件:是否有前项0,是否当前已经符
       合条件……
   limit 这个很重要! 它表示当前数位是否受到上一位的限
       制, 比较抽象, 举例说明
   如果上限是135,前两位已经是1和3了,现在到了个位,个
       位只能是5以下的数字
注: 如果当前受限,不能够记忆化,也不能返回记忆化的结果
为了避免多次调用时 每次上限不同 而导致的错
//http://acm.csie.org/ntujudge/view_code.php?id=106844
// Multiples
LL x;
int digit [100];
LL ten_pow[ 15 ];
bool ava[15];
LL dp[15][2][1000000];
```

LL dfs(int len, LL mod, bool bo, bool limit) {

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```
if ( len < 0 ) return mod == 0;
    if (!limit && dp[len][bo][mod]!= -1 ) return dp[
        len ] [bo] [mod];
    int up = limit? digit[len] : 9;
    LL ret = 0;
    for (int i = 0; i \le up; i++) if (ava[i] || (i=0&&
        bo) ) {
        ret += dfs(len-1, (mod+ten_pow[len]*i)%x, bo
            &&(!i), limit&&(i=up));
    if( !limit ) dp[len][bo][mod] = ret;
    return ret;
LL solve (LL num) {
    int len = 0; digit [0] = 0;
    while( num ) {
        digit[len++] = num\%10;
        num = 10;
    return dfs (len -1, 0,1, 1);
bool check (LL num) {
    while ( num )
        if (!ava[ num%10 ] ) return false;
        num /= 10:
    return true;
int main() {
    LL A, B;
    cin>>x>>A>>B;
    ten_pow[0] = 1;
    mem(dp, -1);
    for (int i = 1; i < 15; i++)
        ten_pow[i] = (ten_pow[i-1]*10)\%x;
    string \ dig; \ cin>>dig;
    mem(ava, false);
    for (char c : dig) ava [c-'0'] = 1;
    if ( x <= 1000000 ) {
        cout << solve(B) - solve(A-1) << endl;
    }else {
        LL ans = 0;
        LL cur = 0;
        while ( cur < A ) cur += x;
        while ( cur <= B ) {
            i\dot{f} ( check(cur) ) ans++;
            cur += x;
        cout << ans << endl;
    }
}
```

#### 1.4 DP(convex hull optimization)

```
// \texttt{http://codeforces.com/contest/311/problem/B}
struct line {
     LL slope,
                  inter;
     LL value(LL x) { return x*slope + inter; }
bool check(line x, line y, line z) {
     return (z.slope - y.slope) * (z.inter - x.inter )
              (z.slope - x.slope) * (z.inter - y.inter);
#define maxn 100005
int n, m, p;
 LL \ a \, [\, maxn ] \; , \ d \, [\, maxn ] \; , \ dp \, [\, 1\,0\,1\,] \, [\, maxn \,] \; , \ s \, [\, maxn \,] \; ; 
int main() {
     cin>\!\!> n>\!\!> m>\!\!> p;
     for(int i = 2; i \le n; ++i) {
          d[i] = getint();
          d\,[\,i\,] \; +\!\!= d\,[\,i\,-1\,]\,;
     for(int i = 1; i<=m; ++i) {
    int h; scanf("%d %lld", &h, a+i);
          a[i] -= d[h];
     sort(a+1,a+1+m);
```

```
\label{eq:formalized} \begin{array}{ll} \text{for}\,(\,int\ i\!=\!1; i\!<\!\!=\!\!m;\, i\!+\!\!+\!\!)\ s\,[\,i\,] \ = \ a\,[\,i\,]\!+\!s\,[\,i\,-\,1\,]\,; \end{array}
//start dp
for (int i=1; i < p; i++) {
      if(i = 1) {
            for (int j=1; j \le m; j++) dp[i][j] = j*a[j] - s
                  [j];
      }else {
           deque<line> dq;
           dq.pb(\{0, 0\});
            for (int j=1; j \le m; j++) {
                  while (dq. size() >= 2 \&\& dq[0]. value(-a[j]) > dq[1]. value(-a[j])) dq.
                        pop_front();
                  dp[i][j] = dq[0].value(-a[j]);
                  line newline{ j, dp[i-1][j]+s[j] };
                  \label{eq:while} \mbox{while} (\ \mbox{dq.size}() >= 2 \ \&\& \ \mbox{check}(\mbox{dq}[\mbox{dq}.
                        size()-2], dq.back(), newline)) dq
                        .pop_back();
                  dq.pb( newline );
                  if ( i==1 )
                       dp\,[\,i\,\,]\,[\,j\,\,] \;=\; j\,{}^*a\,[\,j\,\,] \;\; -\;\; s\,[\,j\,\,]\,;
                  }else {
                       LL mn = 0;
                        for (int \ k = 1; \ k < j; \ k+\!\!\!\!+\!\!\!\!+) \ \{
                             mn = min(mn, dp[i-1][k] + s[k]
                                    - a[j]*k);
                        dp[i][j] = mn + a[j]*j-s[j];
                        // apply convex hull optimization
                  dp[i][j] += a[j]*j - s[j];
      }
cout \ll dp[p][m] \ll endl;
```

## 1.5 simulated annealing

```
//http://mikucode.blogspot.tw/2015/03/algorithm.html
//尋找和所有點距離和最小的點
#include <cstdio>
#include <cstdlib>
#include <cmath>
#define F(n) Fi(i,n)
#define Fi(i,n) for (int i=0;i< n;i++)
#define N 1010
using namespace std;
int X[N], Y[N], n;
inline double pow2(double x){
    return x*x;
double check(double x, double y){
    double ans=0;
    F(n) ans += sqrt(pow2(x-X[i]) + pow2(y-Y[i]));
    return ans;
int main(){
     while (~scanf("%d",&n)) {
        F(n) scanf("%d%d",X+i,Y+i);
        double x=0,y=0,tx,ty,tans,l=10000,ans;
        ans=check(x,y);
        while(1>1e-4) {
            int tmp=rand();
             tx=x+l*cos(tmp); ty=y+l*sin(tmp);
            tans=check(tx,ty);
             if(tans < ans) ans=tans, x=tx, y=ty;
            else 1*=0.9;
        printf("%.9f\n",2*ans);
    }
}
//尋找兩個點使他們跟給定的四個點最小生成樹最小
#include <cstdio>
#include <cstdlib>
```

```
#include <cmath>
#include <algorithm>
#define F(n) Fi(i,n)
#define Fi(i,n) Fl(i,0,n)
#define Fl(i,l,n) for (int i=1; i<n; i++)
#define N 10
using namespace std;
int X[N], Y[N], n, F[N], e;
struct E{
    int a,b;
    double c;
G[N*2];
struct V{
    double x, y;
    V operator+(double 1){
         int tmp=rand();
         return (V) \{x+l*cos(tmp), y+l*sin(tmp)\};
}v[N];
int find(int x){
    return x = F[x]?x:F[x] = find(F[x]);
inline double pow2(double x){
    return x*x;
double check (V s1, V s2) {
    double ans=0;
    e=0;v[4]=s1,v[5]=s2;
    F(5)Fl(j,i+1,6)
         G[e++]=(E)\{i,j,sqrt(pow2(v[i].x-v[j].x)+pow2(v[i].x)\}
             i].y-v[j].y))};
    F(6)F[i]=i;
    sort(G,G+e,[](E a,E b)\{return a.c < b.c;\});
    F(e) {
         if (find (G[i].a)!=find (G[i].b)) {
             ans+=G[ i ] . c;
             F[find(G[i].a)]=find(G[i].b);
         }
    return ans;
int main() {
    scanf("%d",&n);
    while (n--) {
         F(4) scanf("%lf%lf",&v[i].x,&v[i].y);
         double ttans, tans, ans, 11=10000, 12;
         V s1=(V) \{0,0\}, s2=(V) \{0,0\}, ts1, ts2, tmp;
         ans=check(s1, s2);
         while (l1>1e-3) {
             12 = 10000;
             ts1=s1+l1;
             tans=check(ts1,s2);
             tmp=s2;
             while (12>1e-3) {
                  ts2=s2+l2;
                  ttans=check(ts1,ts2);
                  if (ttans<tans) tans=ttans, s2=ts2;
                  else 12*=0.9;
             if (tans<ans) ans=tans, s1=ts1;
             else 11 = 0.9, s2 = tmp;
         printf("%f \setminus n", 2*ans);
    }
}
```

# 2 Graph

#### 2.1 HLD

```
//we can reference the problem Greatest graph
///http://acm.csie.org/ntujudge/problemdata/2582.pdf
//this template operate on edges
#define maxn 100005
struct segment_tree{
    #define right(x) x << 1 | 1
    #define left(x) x << 1
    int* arr;
    int m[4*maxn];</pre>
```

```
int tag[4*maxn];
     const int inf = 1e9;
    void init() {
           /\text{memset}(\text{tag}, -1, \text{sizeof}(\text{tag}));
          fill(tag, tag+4*maxn, inf);
     void pull(int ind) {
              m[ind] = min(m[right(ind)], m[left(ind)]);
     void push(int ind) {
         if(tag[ind] != inf) {
              tag[left(ind)] = min(tag[left(ind)], tag[
                   ind]);
              tag[right(ind)] = min(tag[right(ind)], tag[
                   ind]);
              m[left(ind)] = min( m[left(ind)], tag[left(
                   ind)])
              m[right(ind)] = min(m[right(ind)], tag[
                   right (ind)]);
              tag[ind] = inf;
         }
     /// \text{ root} \Rightarrow 1
     void build(int ind, int l, int r) {
          if( r - l == 1) {
              m[ind] = arr[l];
              return;
         \begin{array}{ll} {\bf int} \ \ {\rm mid} \ = \ (\ l{+}r \ ){>}{>}1; \end{array}
         build( left(ind), l, mid );
         build( right(ind), mid, r );
         pull(ind);
     int query_min(int ind, int L, int R, int ql, int qr
         if (L >= qr \mid \mid R <= ql) return 1e9;
         if (R \ll qr \&\& L >= ql) \{
              return m[ind];
         }
         push(ind);
         \quad \text{int} \ \operatorname{mid} = (L\!+\!R) >> 1;
         return min( query_min(left(ind), L, mid, ql, qr
              ), query_min(right(ind), mid, R, ql, qr));
     void modify(int ind, int L, int R, int ql, int qr,
         int x) {
         if (L >= qr \mid \mid R <= ql) return;
         if(R \le qr \&\& L = ql)
              m[ind] = min(m[ind], x);
              tag[ind] = min(tag[ind], x);
              return:
         push(ind);
         int mid = (L+R) >> 1;
         modify(left(ind), L, mid, ql, qr, x);
         modify(\,right\,(ind\,)\,,\,\,mid\,,\,\,R,\,\,ql\,,\,\,qr\,,\,\,x)\,;
         pull(ind);
    }
};
int seg_arr[maxn];
struct Tree{
    segment_tree seg;
    int n;
    struct Edge { int u, v, c; };
     vector<Edge> e;
     \begin{array}{c} void \ addEdge(int \ x, \ int \ y, \ int \ c) \ \{ \\ G[x].pb(\ SZ(e)\ ); \end{array} 
         G[y].pb(SZ(e));
         e.pb(Edge\{x, y, c\});
    int siz [maxn], max_son[maxn], pa[maxn], dep[maxn];
     /*size of subtree index of max_son, parent index >
          depth*/
    \verb|int link_top[maxn]|, \verb|link[maxn]|, \verb|timer|;
     /*chain top index in segtree ime stamp*/
    std::vector<int >G[maxn];
    void init(int N) {
```

```
n = N:
    e.clear();
    for(int i = 1; i <= n; i++) G[i].clear();
    timer=0;
    pa[1] = 1;
    dep[1] = 0;
void find_max_son(int x){
    siz[x]=1;
    \max_{son}[x] = -1;
    for(int e\_ind : G[x]) {
        int v = e[e\_ind].u == x ? e[e\_ind].v : e[
            e_ind].u
        find_max_son(v);
        if(max\_son[x] = -1 \mid \mid siz[v] > siz[max\_son]
             [x]])
            \max_{son}[x] = v;
        siz[x] += siz[v];
    }
void build_link(int x, int top){
    link[x] = timer++;/*記錄x點的時間戳*/
    link\_top[x] = top;
    if(\max_{x \in \mathbb{R}} |x| != -1)
        build_link( max_son[x], top);/*優先走訪最大
             孩子*/
    for(int e\_ind : G[x]) {
        int v = e[e\_ind].u == x ? e[e\_ind].v : e[
             e\_ind].u;
        if(v = pa[x])
            seg\_arr[link[x]] = e[e\_ind].c;
        if (v = \max_{son}[x] \mid | v = pa[x]) continue
        // edge from x \Rightarrow v
        build_link(v, v);
    }
inline int lca(int a, int b){
    /*求LCA, 可以在過程中對區間進行處理*/
    int ta=link\_top[a],tb=link\_top[b];
    while(ta != tb){
        if (dep[ta]<dep[tb]) {
            std::swap(ta,tb);
            std::swap(a,b);
        //interval [ link[ta], link[a] ]
        a = pa[ta];
        ta = link\_top[a];
    return dep[a] < dep[b] ? a:b;
}
int modify(int a, int b, int c){
    int ta=link_top[a], tb=link_top[b];
    while(ta != tb){
        if (dep[ta]<dep[tb]) {
            std::swap(ta,tb);
            std::swap(a,b);
        //interval [ link[ta], link[a] ]
        //same interval if operate on edges
        seg.modify(1, 1, n, link[ta], link[a]+1, c)
        a = pa[ta];
        ta = link\_top[a];
    //a, b are on the same chain
    if( a = b ); // interval [ link[a], link[a]], if operate on edges \Rightarrow no edge
        if (dep[a]>dep[b])
            swap(a,b)
        //interval [ link[a], link[b] ]
        // if operate on edges \Longrightarrow [ link[ max_son[
            a] ], link[b] ]
```

```
seg.modify(1, 1, n, link[max\_son[a]],
                   link[b]+1, c);
         }
    }
/*
     void modify(int a, int b, int c) {
         if ( a==b ) return;
         if( link\_top[a] = link\_top[b]) {
              if \, (\ dep \, [\, a\, ] \, > \, dep \, [\, b\, ] \ ) \ swap (a \, , \ b) \, ;
              seg.modify(1,\ 1,\ n,\ link\,[a]{+}1,\ link\,[b]{+}1,\ c
              assert(link[a]+1 = link[max_son[a]]);
              return;
         if(dep[link\_top[a]] < dep[link\_top[b]])
              swap(a, b);
         // a is the node with deeper link_top
         seg.modify(1, 1, n, link[link_top[a]], link[a]
               + 1. c):
         modify( pa[link_top[a]], b, c);
    /// Heavy Light Decomposition
    void HLD() {
          // root is indexed 1 here !
         find_max_son(1);
         build_link(1, 1);
}tree;
int main() {
    int T; cin>>T;
     while (T--) {
         int n,m;
         scanf("%d %d",&n, &m);
         int ans = 0;
         tree.init(n);
         for (int i=0; i< n-1; i++) {
              int a, b, c;
              {\rm scanf}\,(\,\text{``'}\!\text{d}\!\text{''}\!\text{d}\!\text{''}\!\text{d}\!\text{''},\!\&a,\!\&b,\!\&c\,)\;;
              //a--, b--; be careful here
              tree.addEdge(a, b, c);
              ans += c;
         tree.HLD();
         tree.seg.arr = seg_arr;
         tree.seg.build(1, 1, n);
    return 0;
2.2 Hungarian
```

```
// edge and node index starting from 0
  dfs version below
//complexity O ( V*E )
/* to do
#define ___maxNodes
num\_left = ?
struct Edge {
    int from:
    int to;
    int weight;
    Edge(int f, int t, int w):from(f), to(t), weight(w)
         {}
};
vector<int> G[__maxNodes]; /* G[i] 存储顶点 i 出发的边
    的编号 */
vector<Edge> edges;
int num_nodes;
int num_left;
int num_right;
int num_edges;
int matching[__maxNodes]; /* matching result */
int check [___maxNodes];
```

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```
bool dfs(int u) {
   for (auto i = G[u].begin(); i != G[u].end(); ++i) {
        // 对 u 的每个邻接点
       int v = edges[*i].to;
       if (!check[v]) { // 要求不在交替路中
           check[v] = true; // 放入交替路
           if (matching[v] == -1 || dfs(matching[v]))
              // 如果是未盖点,说明交替路为增广路,则
                  交换路径,并返回成功
              matching[v] = u;
              matching[u] = v;
              return true;
       }
   }
   return false; // 不存在增广路,返回失败
int hungarian() {
   int ans = 0;
   memset(matching, -1, sizeof(matching));
   for (int u=0; u < num_left; ++u) {</pre>
       if (matching[u] = -1) {
           memset(check, 0, sizeof(check));
           if (dfs(u)) + ans;
       }
   return ans;
```

#### 2.3 KM

```
// 最小帶權匹配~ km算法
//http://acm.csie.org/ntujudge/contest_view.php?id=836&
    contest\_id=449
#include <bits/stdc++.h>
using namespace std;
struct bipartite {
    #define maxn 602
    #define INF 0xfffffff
     \quad \text{int } \operatorname{sx}\left[\operatorname{maxn}\right], \ \operatorname{sy}\left[\operatorname{maxn}\right], \ \operatorname{mat}\left[\operatorname{maxn}\right]\left[\operatorname{maxn}\right];
     int x[maxn], y[maxn], link[maxn];
    int N, M, slack;
    int DFS(int t) {
         int tmp;
         sx[t] = 1;
          for (int i = 0; i < M; i++) {
               if (!sy[i]) {
                    tmp = x[t] + y[i] - mat[t][i];
                    if (tmp == 0) {
                         sy[i] = 1;
                         if (link[i] = -1 || DFS(link[i]))
                              \,l\,i\,n\,k\,\,[\,\,i\,\,]\,\,=\,\,t\,\,;
                              return 1;
                    else if (tmp < slack) slack = tmp;
              }
         }
         return 0;
     int KM() {
          for (int i = 0; i < N; i++) {
              x[i] = 0;
               for (int j = 0; j < M; j++) {
                    if (mat[i][j] > x[i]) x[i] = mat[i][j];
          for (int j = 0; j < M; j++) { y[j] = 0; }
         memset(link, -1, sizeof(link));
          for (int i = 0; i < N; i++) {
               while (1) {
                    memset(sx, 0, sizeof(sx));
                    memset(sy, 0, sizeof(sy));
                    slack = INF;
```

```
if (DFS(i)) break;
                 \quad \  \  \text{for (int } j \, = \, 0; \ j \, < \, N; \ j++) \, \, \{
                      if (sx[j]) x[j] = slack;
                 for (int j = 0; j < M; j++) {
                      if (sy[j]) y[j] += slack;
                 }
             }
        }
         int ans = 0;
         int cnt = 0;
         int t;
         for (int i = 0; i < M; i++)
         {
             t = link[i];
             if (t >= 0 \&\& mat[t][i] != -INF)
                 cnt ++:
                 ans += mat[t][i];
         // 最大權 : 沒有負號
        return -ans;
    void init(int n, int m) {
        N\,=\,n\,,\ M=m;
         for (int i = 0; i < N; i++)
            for (int j = 0; j < M; j++)
                 mat[i][j] = -INF;
    void input() {
         for(int i = 0; i < N; i++)
             for (int j = 0; j < M; j++) {
                 // fill in mat[i][j]
                 // stands for the weighting , but
                      negative sign !
                 // if 最大權 : 沒有負號
             }
}km;
int main(){
    int n,E;
    while (scanf("%d", &n) != EOF)
    {
        km.init(n, n);
        km.input();
         cout<< km.KM() <<endl;
    return 0;
```

## 2.4 Bi-vertex-connected Subgraph

```
#include <bits/stdc++.h>
using namespace std;
#ifdef DEBUG
   #define debug(...) printf(__VA_ARGS__)
   #define debug(...) (void)0
#endif
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair<long long, long long>
#define fi first
#define se second
#define all(x) (x).begin(),(x).end()
#define SZ(x) ((int)(x).size())
const int inf = 0x7ffffffff; //beware overflow
\#define mem(x, y) memset(x, (y), sizeof(x));
#define IOS ios_base::sync_with_stdio(0); cin.tie(0)
template<typename A, typename B>
ostream& operator <<(ostream &s, const pair<A,B> &p) {
    return s<<"("<<p.first<<","<<p.second<<")";
```

```
template<tvpename T>
ostream& operator <<(ostream &s, const vector<T> &c) {
    for (auto it : c) s << it << " ";
    s << "]";
                                                          }
   return s;
template<typename T>
ostream& operator << (ostream &o, const set<T> &st) {
    o << "{";
    for (auto it=st.begin(); it!=st.end(); it++) o << (
       it=st.begin()? "": ", ") << *it;
    return o << "}";
template < typename T1, typename T2>
ostream& operator << (ostream &o, const map<T1, T2> &mp
   ) {
   o << "{";
    for (auto it=mp.begin(); it!=mp.end(); it++) {
       o << (it=mp. begin()?"":", ") << it->fi << ":"
            << it->se;
    o << "}";
    return o;
     regard every vbcc as a set of edges
/** needed for tarjan **/
#define maxn 100005
\#define maxm 100005
int n, m;
struct Edge{int s, t;};
vector<Edge> edge;
int dfn[maxn], low[maxn];
stack<int> st;
bool vis[maxn];
int Time;
bool vis_e [maxm];
int bcnt, vbb[maxm];
vector<int> vb[maxm];
vector<int> G[maxn];
void tarjan(int s){
    dfn[s] = low[s] = ++Time;
    vis[s] = true;
    for (int e_ind : G[s]) {
        if (!vis_e[e_ind]) {
            vis_e[e_ind] = true; st.push(e_ind);
            int to = edge[e\_ind].s + edge[e\_ind].t - s;
            if (! vis[to]) {
                tarjan(to);
                low[s] = min(low[s], low[to]);
                if(low[to] >= dfn[s]){
                    vb[bcnt].clear();
                    while (1) {
                        int t = st.top(); st.pop();
                        vbb[t] = bcnt;
                        vb[bcnt].push_back(t);
                        if(t = e_ind) break;
                    bcnt++;
                }
            }else
                low[s] = min(low[s], dfn[to]);
        }
   }
void init_tarjan() {
   mem(vis, false); mem(vis_e, false);
    Time = bcnt = 0; edge.clear();
    for (int i = 1; i \le n; i++) G[i]. clear ();
int main() {
    cin >> n >> m;
    init_tarjan();
    edge.push\_back(Edge\{a,b\});
```

```
G[a].push_back((int)edge.size()-1);
G[b].push_back((int)edge.size()-1);
}
tarjan(1);
```

## 2.5 Bi-edge-connected Subgraph

```
/** needed for tarjan **/
#define maxn 100005
#define maxm 100005
int n, m;
int dfn[maxn], low[maxn];
stack<int> st;
int Time;
int bcnt;
vector<int> G[maxn];
bool in_cyc[maxn];
void tarjan(int s, int p){
    dfn[s] = low[s] = ++Time;
    st.push(s);
    for (int to : G[s]) if ( to != p ){
        if (!dfn[to]) {
            tarjan(to, s);
            low[s] = min(low[s], low[to]);
            if(low[to] > dfn[s]) {
                // is cut_edge
                 // pop stack 的過程也可以寫在這
                 // 但最後(after tarjan)還要多判stack
                     not empty的情况
                 if ( low [to] > dfn [s]) {
                 in\_cyc[bcnt] = st.top()!=to;
                 while (1) {
                     int t = st.top(); st.pop();
                     id[t] = bcnt;
                     if (t == to) break;
                 bcnt++;
            }
            }
        }else
            low[s] = min(low[s], dfn[to]);
    if(low[s] = dfn[s]){
        in\_cyc[bcnt] = st.top()!=s;
        while (1) {
            int t = st.top(); st.pop();
            id[t] = bcnt;
            if(t = s) break;
        bcnt++;
    }
void init_tarjan() {
    Time = bcnt = 0;
int main() {
  cin >> n >> m;
  init_tarjan();
  for (int i = 0; i < m; i++) {
        int a, b; scanf("%d %d", &a, &b);
        G[a].pb(b), G[b].pb(a);
  mem( in_cyc , false);
  tarjan(1, 1);
}
```

#### 2.6 SCC

```
#include <bits/stdc++.h>
using namespace std;
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair<long long, long long>
#define fi first
#define se second
const int inf = 1e9;
const LL INF = 1e18;
const int mod = 1e9+7;
#define maxn 100050
int n, m;
vector<int> g[maxn];
stack<int> Stack;
int scnt, Time;
int belong[maxn], dfn[maxn], low[maxn], indegree[maxn];
bool instack [maxn];
void input(){
  cin >> n >> m;
  for (int i = 0; i < m; i++){}
    int a, b; scanf("%d%d", &a, &b);
    g[a].pb(b);
void init() {
  scnt = Time = 0;
  for (int i = 1; i \le n; i++)
    g[i].clear();
  while (!Stack.empty()) Stack.pop();
  memset(indegree, 0, sizeof(indegree));
  memset(dfn, 0, sizeof(dfn));
  memset(instack, false, sizeof(instack));
void dfs(int u) {
  dfn\left[u\right] \ = \ low\left[u\right] \ = \ +\!\!\!+\!\! Time;
  Stack.push(u); instack[u] = true;
  for(int v : g[u]) {
    if (!dfn[v]) {
      dfs(v);
      low[u] = min(low[u], low[v]);
    else if (instack [v])
      low[u] = min(low[u], dfn[v]);
  if(low[u] = dfn[u]) {
    scnt++;
    int tp;
      tp = Stack.top(); Stack.pop();
      instack[tp] = false;
      belong[tp] = scnt;
    }  while(tp != u);
  }
void tarjan() {
  for (int i = 1; i \le n; i++)
    if (!dfn[i])
      dfs(i);
int main(){
  int T; cin >> T;
  while (T--) {
    init();
    input();
    tarjan();
    for (int i = 1; i \le n; i++) {
      for(int v : g[i]) {
  if(belong[v] != belong[i])
          indegree [belong [v]]++;
      }
    LL ans = 0;
    for (int i = 1; i \le scnt; i++)
```

if (!indegree[i]) ans++;

```
cout << ans << endl;
}
return 0;
}</pre>
```

# 2.7 Steiner Tree( PECaveros )

```
// Minimum Steiner Tree
 // O(V 3^T + V^2 2^T)
struct SteinerTree{
#define V 33
#define T 8
#define INF 1023456789
   int n , dst [ V ] [ V ] , dp [ 1 << T ] [ V ] , tdst [ V
   void init( int _n ){
    n = \underline{n};
     for (int i = 0 ; i < n ; i ++){
       for (int j = 0 ; j < n ; j ++)
         dst[i][j] = INF;
       dst[i][i] = 0;
   void add_edge( int ui , int vi , int wi ){
  dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
  dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
   void shortest_path(){
     for (int k = 0 ; k < n ; k +++)
       for ( int i = 0 ; i < n ; i +++)
         for (int j = 0 ; j < n ; j ++)
            dst[i][j] = min(dst[i][j]
                                    dst[ i ][ k ] + dst[ k
                                         ][ j ];
   int solve( const vector<int>& ter ){
     int t = (int) ter.size();
     for ( int i = 0 ; i < (1 << t) ; i ++ ) for ( int j = 0 ; j < n ; j ++ )
         dp[i] = INF;
     for (int i = 0; i < n; i ++)
       dp[0][i] = 0;
     for ( int msk = 1 ; msk < (1 << t) ; msk ++ ){
       if (msk = (msk & (-msk))) 
          int who = \underline{\hspace{1cm}} \lg(msk);
          for (int i = 0; i < n; i ++)
           dp[msk][i] = dst[ter[who]][i];
          continue;
       for(int i = 0 ; i < n ; i ++)
          for(int submsk = (msk - 1) \& msk ; submsk ;
                   submsk = ( submsk - 1 ) & msk )
              dp[msk][i] = min(dp[msk][i]
                                       dp[ submsk ][ i ] +
   dp[ msk ^ submsk
       for (int i = 0 ; i < n ; i ++) {
          tdst[i] = INF;
          for (int j = 0 ; j < n ; j ++)
            tdst[i] = min(tdst[i], dp[msk][j] + dst[j][i]
                                     ] );
       for (int i = 0 ; i < n ; i ++)
         dp \, [\ msk\ ] \, [\ i\ ] \, = \, t \, ds \, t \, [\ i\ ] \, ;
     int ans = INF;
     for ( int i = 0 ; i < n ; i +++)
       ans = \min( ans , dp[ ( 1 << t ) - 1 ][ i ] );
} solver;
```

## 2.8 Edmond's Matching Algorithm

```
//http://acm.csie.org/ntujudge/contest_view.php?id=370&contest_id=466
#include <bits/stdc++.h>
```

```
using namespace std:
//带花树,Edmonds's matching algorithm,一般图最大匹配
/// have to be a undirected graph
#define MAXN 505
vector<int>G[MAXN];//用vector存圖
int pa [MAXN], match [MAXN], st [MAXN], S [MAXN], vis [MAXN];
int t,n;
inline int lca(int u,int v){//找花的花托
   for(++t;;swap(u,v)){
        if (u==0)continue;
        if (vis[u]==t)return u;
        vis [u]=t;//這種方法可以不用清空vis陣列
       u=st[pa[match[u]]];
#define qpush(u) q.push(u),S[u]=0
inline void flower(int u, int v, int l, queue<int> &q){
   while(st[u]!=l){
       pa[u]=v; //所有未匹配邊的pa都是雙向的
        if (S[v=match[u]]==1)qpush(v); // 所有奇點變偶點
        st[u]=st[v]=l, u=pa[v];
inline bool bfs(int u){
   for(int i=1;i<=n;++i)st[i]=i;//st[i]表示第i個點的集
   memset(S+1,-1,sizeof(int)*n);//-1:沒走過 0:偶點 1:
        奇點
   queue < int > q; qpush(u);
    while(q.size()){
       u=q.front(),q.pop();
        for(size_t i=0; i < G[u]. size(); ++i){
            int v=G[u][i];
            if(S[v]==-1){
                pa[v]=u, S[v]=1;
                if (!match[v]) {//有增廣路直接擴充
                    for(int lst;u;v=lst,u=pa[v])
                        lst=match[u], match[u]=v, match[v]
                            =u;
                    return 1:
                qpush (match [v]);
            else\ if(!S[v]\&\&st[v]!=st[u]){
                int l=lca(st[v],st[u]);//遇到花,做花的
                flower(v,u,l,q), flower(u,v,l,q);
            }
       }
   }
   return 0;
inline int blossom(){
   memset(pa+1,0,sizeof(int)*n);
   memset(match+1,0,sizeof(int)*n);
    int ans=0;
    for(int i=1;i<=n;++i)
        if (!match[i]&&bfs(i))++ans;
   return ans;
void solve() {
   cin>>n;
    int m; cin>>m;
    while (m--) {
        int a,b;
       scanf("%d %d", &a, &b);
       a++, b++;
       // since node indexed [ 1 .. n ] in this
            template
       #define pb push_back
        //Multiedge and self-cycles are not forbidden
       G[a].pb(b);
       G[b].pb(a);
    cout << blossom() << endl;
   for (int i = 1; i \le n; i++) G[i]. clear();
int main() {
   int t; cin >> t;
    while(t--) solve();
```

# 2.9 Tree Decomposition

```
//codeforces Digit Tree
//http://codeforces.com/problemset/problem/715/C
#include <bits/stdc++.h>
using namespace std;
#ifdef DEBUG
    #define debug(...) printf(__VA_ARGS__)
#else
    #define debug(...) (void)0
#endif
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair<long long, long long>
#define fi first
#define se second
\#define all(x) (x).begin(),(x).end()
#define SZ(x) ((int)(x).size())
const int inf = 0x7ffffffff; //beware overflow
#define mem(x, y) memset(x, (y), sizeof(x));
#define IOS ios_base::sync_with_stdio(0); cin.tie(0)
template<typename A, typename B>
ostream& operator <<(ostream &s, const pair<A,B> &p) {
    return s<<"("<<p.first<<","<<p.second<<")";
template<typename T>
ostream& operator <<(ostream &s , const vector<T> &c) {
    s << "[ ";
    for (auto it : c) s << it << " ";
    s << "]";
    return s;
template<typename T>
ostream& operator << (ostream &o, const set<T> &st) {
    o << "{";
    for (auto it=st.begin(); it!=st.end(); it++) o << ( it=st.begin() ? "" : ", ") << *it;
    return o << "}";
template<typename T1, typename T2>
ostream& operator << (ostream &o, const map<T1, T2> &mp
    ) {
    o << "{";
    for (auto it=mp.begin(); it!=mp.end(); it++) {
        o << (it=mp.begin()?"":", ") << it->fi << ":"
             << it->se;
    o << "}";
    return o;
typedef long long 11;
bool isprime[100005];
vector<LL> primes;
LL M, PHI;
#define MOD M
ll modpow(ll a, ll b) {
  11 r = 1:
  while(b) {
    if(b\&1) r = (r*a)\%MOD;
    a=(a*a)%MOD;
    b >>= 1;
  }
  return r;
void Sieve(int n) {
  memset(isprime, 1, sizeof(isprime));
  isprime[1] = false;
  for (int i = 2; i \le n; i++) {
    if(isprime[i]) {
      primes.pb(i);
      for (int j = 2*i; j \le n; j += i)
        isprime[j] = false;
  }
}
```

```
LL phi(LL n) {
   \begin{array}{lll} \text{ll} & \text{num} = 1; & \text{ll} & \text{num} 2 = n; \\ \end{array} 
  for(ll i = 0; primes[i]*primes[i] \le n; i++) {
    if (n%primes [i]==0) {
      num2/=primes[i];
      num*=(primes[i]-1);
    while (n\%primes [i] == 0) {
      n/=primes [i];
  if(n>1) {
    num2/=n; num*=(n-1);
  }
  n = 1;
  num *= num2;
  return num:
ll inv(ll a) {
  return modpow(a, PHI-1);
#define maxn 100005
struct edge{
    int u, v, dig;
    int no(int x) {
         return x == u ? v : u; 
};
vector<edge> e;
vector < int > G[maxn];
LL n, ans;
bool vis [maxn];
int sz[maxn], dep[maxn];
LL tenPow[maxn];
int dfs(int u, int p, int d) {
    sz[u] = 1;
    dep[u] = d;
    for (int eind : G[u] ) {
         \quad \text{int} \ v = e [eind].no(u);
         if (v = p \mid | vis[v]) continue;
        sz[u] += dfs(v, u, d+1);
    return sz[u];
int findCenter(int u, int p, int treesize) {
    for (int eind : G[u] ) {
int v = e[eind].no(u);
         if (v = p \mid | vis[v]) continue;
         if (sz[v]*2 > treesize)
             return findCenter( v, u, treesize);
    return u;
LL up [maxn], down [maxn];
int belong[maxn];
map<LL, LL> tot;
vector<int> pt;
void calc(int u, int p, int b, int d) {
    pt.pb( u );
    belong[u] = b;
    dep[u] = d;
    int \ id = \ find\_if(\ all(G[u]) \ , [u,p](int \ x) \ \{ \ return
    for (int eind : G[u]) {
         int v = e[eind].no(u);
         if (\ vis [v] \ || \ v == p \ ) \ continue;
         calc(v, u, b, d+1);
    vec[b][up[u]]++;
    tot[ up[u] ]++;
LL solve(int cent) {
```

```
//cent is the root now
     vector<int> L;
     for (int eind : G[cent]) {
         \begin{array}{ll} \hbox{int} & v \, = \, e \, [\, e \, i \, n \, d \, ] \, . \, no \, (\, cent \, ) \, ; \end{array}
         if( !vis[v]) {
              L.pb( v );
    vec.clear();
    {\tt vec.resize}\left(\ SZ(L)\ ,\ \{\}\ \right);
     tot.clear();
    up[cent] = down[cent] = 0;
    dep[cent] = 0;
    pt.clear();
    for(int i = 0; i < SZ(L); i++)
         calc ( L[i], cent, i, 1);
    LL ret = 0;
    for(int u : pt) {
         LL tmp = (-\text{down}[u]+M)\%M;
         tmp = (tmp*inv(tenPow[dep[u]]))%M;
         ret += tot[ tmp ] - vec[ belong[u] ][ tmp ];
    assert ((LL)count_if(all(pt), [] (int x) { return
    \begin{array}{c} up[x] = 0; \; \} \; ) = tot[0]); \\ LL \; tmp = tot[0] + (LL)count\_if(all(pt), \; [] \; (int \; x) \end{array}
         \{ \text{ return down}[x] = 0; \} );
    debug("%lld\n", tmp);
    return ret+tmp;
void solveAll(int node) {
    dfs(node, -1, 0);
    int cent = findCenter(node, -1, sz[node]);
     ans += solve ( cent );
    debug("%d %lld\n", cent, ans);
    vis [cent] = true;
     for (int eind : G[cent] ) {
         int v = e[eind].no(cent);
         if( vis[v] ) continue;
         solveAll(v);
    }
int main() {
    cin >> n >> M:
  Sieve( 100000 );
    PHI = phi(M);
    for (int i = 0; i < n-1; i++) {
         int a, b, c; scanf("%d %d %d", &a, &b, &c);
         G[a].pb(SZ(e)); G[b].pb(SZ(e));
         e.pb(edge{a, b, c});
     //init
    tenPow[0] = 1;
     for(int i = 1; i < maxn; i++) tenPow[i] = (tenPow[i])
         -1]*10)%M;
    ans = 0;
    mem( vis, false);
    solveAll(0);
    cout << ans << endl;
```

## 2.10 Tree Longest Path

```
/** codeforces 592D - Super M **/
#include <bits/stdc++.h>

using namespace std;

#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair<long long, long long>
#define fi first
#define se second

const int inf = 1e9;
const LL INF = 1e18;
const int mod = 1e9+7;
#define maxn 123460
```

```
int n, m;
vector<int> g[maxn];
bool is [maxn];
int dep[maxn], R, max_depth, A;
int cnt[maxn], parent[maxn];
bool dfs(int u, int par = 0){
  parent[u] = par;
  dep[u] = dep[par] + 1;
  if(dep[u] > max_depth \&\& is[u])
    \max_{\mathbf{depth}} = \operatorname{dep}[\mathbf{u}], \ \mathbf{R} = \mathbf{u};
  bool ret = is[u];
  for(int v : g[u])
    if(v != par)
      ret = dfs(v, u);
  if(ret) A++;
  return ret;
}
int find_center(int start) {
 R = start; dep[0] = -1; max_depth = 0;
  dfs(start);
  \max_{\text{depth}} = 0; \text{dep}[R] = -1;
  dfs(R, R);
  int ret = R, d = max_depth/2;
  while(d>0)
    d--:
    ret = parent[ret];
  return ret;
int S, dis, max_length;
bool dfs1(int u, int par = 0) {
  dep[u] = dep[par] + 1;
  if ( is [u])
    if(dep[u] > max_length)
      max\_length \,=\, dep\left[\,u\,\right]\,,\;\; S \,=\, u\,;
    else if (dep[u] = max_length & u < S)
      S = u;
  bool c = false;
  for(int v : g[u])
    if ( v != par )
       dfs1(v, u);
int main(){
  cin >> n >> m;
  for (int i = 0; i < n-1; i++){
    int a, b; scanf("%d%d", &a, &b);
    g[a].pb(b), g[b].pb(a);
  }
  memset(is, false, sizeof(is));
  int tmp;
  for (int i = 0; i < m; i++){
    cin>>tmp; is [tmp] = true;
  int C = find_center(tmp);
  dep[0] = -1;S = inf; dis = (max_depth+1)/2;
  // distance(center, any other node) <= (longestpath +
        1) / 2
  dfs1(C);
  if ( max_depth & 1)
    dfs1 (parent [C]);
  cout << S << endl << A-2-max_depth << endl;
  return 0;
```

### 3 Flow

#### 3.1 Dinic Maxflow

```
//http://acm.csie.org/ntujudge/problem.php?id=2581
//French Fries Festival
//dinic runs in O( V^2*E )
#include <bits/stdc++.h>
using namespace std;
#ifdef DEBUG
#define debug(...) printf(__VA_ARCS__)
```

```
#else
   #define debug(...) (void)0
#endif
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair < long long, long long>
#define fi first
#define se second
#define all(x) (x).begin(),(x).end()
\#define SZ(x) ((int)(x).size())
const int inf = 0x7ffffffff; //beware overflow
\#define mem(x, y) memset(x, (y), sizeof(x));
\#define IOS ios_base::sync_with_stdio(0); cin.tie(0)
template<typename A, typename B>
ostream& operator <<(ostream &s, const pair<A,B> &p) {
    return s<<"("<<p.first<<","<<p.second<<")";
template<typename T>
ostream& operator <<(ostream &s , const vector<T> &c) {
    s << "[ ";
    for (auto it : c) s << it << " ";
    s << "]";
    return s;
template<typename T>
ostream& operator << (ostream &o, const set<T> &st) {
    o << "{";
    for (auto it=st.begin(); it!=st.end(); it++) o << (
        it=st.begin() ? "" : ", ") << *it;
    return o << "}";
template<typename T1, typename T2>
ostream& operator << (ostream &o, const map<T1, T2> &mp
    o << "{";
    o << (it=mp.begin()?"":", ") << it->fi << ":"
            << it->se;
    o << "}";
    return o;
#define maxn 500
{\tt struct} \ {\tt Edge} \{ \ {\tt int} \ {\tt to} \, , \ {\tt cap} \, , \ {\tt rev} \, ; \ \};
struct Dinic{
    vector < Edge > G[maxn];
    int dis[maxn], iter[maxn];
    void init(int n) {
        //zero based
      for (int i = 0; i < n; i++) G[i].clear();
    void addEdge(int from, int to, int cap) {
        vector < Edge > :: iterator it;
        if( ( it=find_if( all(G[from]), [to](Edge& e) {
             return e.to == to; } )) != G[from].end() )
            i\,t\,\text{-}\!>\!\mathrm{cap}\;+\!\!=\;\mathrm{cap}\,;
            return;
      G[from].pb(Edge\{to, cap, (int)G[to].size()\});
      G[to].pb(Edge\{from, 0, (int)G[from].size()-1\});
        //if undirected 0 will be cap
    bool bfs(int s, int t) {
      memset(dis, -1, sizeof(dis));
      queue<int> que;
      que.push(s); dis[s] = 0;
      while (!que.empty())
        int tp = que.front(); que.pop();
        for (Edge &e : G[tp]) {
           if(e.cap > 0 \&\& dis[e.to] == -1)
             dis[e.to] = dis[tp] + 1, que.push(e.to);
        }
      return dis[t] != -1;
    int dfs(int v, int t, int f) {
      if(v = t) return f;
      for(int \& i = iter[v]; i < G[v].size(); i++) {
```

```
Edge &e = G[v][i];
        if(e.cap > 0 \&\& dis[v] < dis[e.to]) {
          int d = dfs(e.to, t, min(f, e.cap));
          if(d > 0) {
            e.cap -= d;
            G[e.to][e.rev].cap += d;
            f += d;
            return d;
       }
      return 0;
    int maxFlow(int s, int t) {
      int ret = 0;
      while (bfs(s, t))
        memset(iter, 0, sizeof(iter));
        while ((f = dfs(s, t, inf)) > 0)
          ret += f;
      return ret;
}dinic, dinic2;
void solve() {
    int n,m,k; cin>>n>>m>>k;
    // flow problem with lower bounds;
    int s = 0, t = n+2, ss = n+3, tt = n+4;
    dinic.init( n+5);
    dinic.addEdge(s, 1, k);
    dinic.addEdge(n+1, t, k);
    int slb = 0;
    while(m--) {
        int l, r, a, b; scanf("%d %d %d %d", \&l, \&r, \&a
            , &b);
        slb += a;
        r++;
        dinic.addEdge(l, r, b-a);
        dinic.addEdge(ss, r, a);
        dinic.addEdge(l, tt, a);
    dinic2 = dinic;
    dinic.addEdge(t, s, k);
    int f1 = dinic.maxFlow(ss, tt);
    if( !all_of( all(dinic.G[ss]), [](Edge x) { return
        x.cap = 0;  ) ) {
        puts("-1"); return;
    dinic2.addEdge(ss, s, 1e9);
    dinic2.addEdge(t, tt, 1e9);
    int f2 = dinic2.maxFlow(ss, tt);
    // maxflow in current graph is f2 - slb
    printf("%d\n", (f2 - slb)*n);
int main() {
    int t; cin>>t;
    while (t - -)
        solve();
```

### 4 Data Structure

## 4.1 Disjoint Set

```
| struct Disjoint_set {
    #define MAX_N 500005
    // define MAX_N
    int pa [MAX_N], Rank [MAX_N];
    int sz [MAX_N];
    void init_union_find(int V) {
        for(int i=0; i<V; i++) {
            pa[i] = i;
            Rank[i] = 0;
            sz[i] = 1;
```

```
}
int find(int x) {
     return x == pa[x] ? x : pa[x] = find(pa[x]); 
int unite(int x, int y) -
    x = find(x), y = find(y);
    \quad \text{int } S = sz[x] + sz[y];
    if(x != y){
         if(Rank[x] < Rank[y]) {
             pa[x] = y;
             sz[y]=S;
             return y;
         }
         else{

pa[y] = x; \\
sz[x] = S;

             if(Rank[x] = Rank[y]) Rank[x] ++;
             return x;
    }
bool same(int x, int y) {
    return find(x) = find(y);
}
```

# $4.2 \quad \text{Djs} + \text{Seg}$

```
// demo => undo djs + segtree with offline
// this program doesn't consider the problem of
     overflowing varaible ans
// http://acm.csie.org/ntujudge/view_code.php?id
    =\!108190\&\mathtt{contest\_id}\!=\!472
#define maxn 100005
#define maxm 500005
//can be used to solve dynamic connectivity problem
//can be used with segment tree -> offline
struct DisjointSet {
  // save() is like recursive
  // undo() is like return
  int n, fa [maxn], sz [maxn];
  vector<pair<int*,int>>> h;
  vector<int> sp;
  int ans;
  void init(int tn) {
    ans = 0;
    n=tn;
    for (int i=0; i<n; i++) {
      fa[i]=i;
      sz[i]=1;
    sp.clear(); h.clear();
  void assign(int *k, int v) {
    h.pb(\{k, *k\});
    *k=v;
  void save() { sp.pb(SZ(h)); }
  void undo() {
    assert(!sp.empty());
    int last=sp.back(); sp.pop_back();
    while (SZ(h)!=last) {
      auto x=h.back(); h.pop_back();
       *x.fi=x.se;
    }
  int f(int x) {
    while (fa[x]!=x) x=fa[x];
    return x;
  void uni(int x, int y) {
    x=f(x); y=f(y);
    if (x=y) return ;
    \begin{array}{ll} \textbf{if} & (\,sz\,[\,x]\!<\!sz\,[\,y\,]\,) & swap\,(\,x\,,\ y\,)\,; \end{array}
    //nans stands for new answer
    int t = sz[x]+sz[y];
```

```
int nans = ans - (sz[x]*sz[x]-sz[x]) - (sz[y]*sz[y]
         ]-sz[y]) + t*t-t;
    assign(\&sz[x], sz[x]+sz[y]);
    assign(&fa[y], x);
    assign(&ans, nans);
} djs;
int n, m;
map < int, int > ma[maxn];
vector<pii> seg[4*maxm];
LL ans [maxm];
void add(int ql, int qr, int a, int b, int id=1, int l
    =0, int r=m) {
    if (qr \ll l'| ql \gg r) return;
    if(l) = ql \&\& r <= qr)
        seg[id].pb(mp(a, b));
        return ;
    int mid = (l+r) >> 1;
    {\rm add}(\ ql\,,\ qr\,,\ a\,,\ b\,,\ id^{\,*}2\,,\ l\,,\ mid)\,;
    add( ql, qr, a, b, id*2+1, mid, r);
void dfs(int u=1, int l=0, int r=m) {
    djs.save();
    for(pii v : seg[u] ) djs.uni( v.fi , v.se );
    if(r-l > 1) {
        int mid = (l+r) >> 1;
        dfs(u*2, l, mid);
        dfs(u*2+1, mid, r);
    }else {
   // do sth here
        ans[l] = djs.ans;
    djs.undo();
int main() {
    \operatorname{scanf}("\%d \%d", \&n, \&m);
    for(int i = 0; i < m; i++) {
        int a, b; scanf("%d %d",&a, &b);
        a--, b--; if(b < a) swap(a, b);
         if (ma[a].count(b)) {
            add(ma[a][b], i, a, b);
             ma[a].erase(b);
        else ma[a][b] = i;
    for(int i = 0; i < n; i++) if(!ma[i].empty()) 
         for(auto p : ma[i])
            add( p.se, m, i, p.fi);
    djs.init(n);
    dfs();
    for (int i =0; i < m; i++) printf("%lld\n", ans[i]);
```

### 4.3 Sparse Table

```
//codeforces 689D
#define maxn 200005
template < typename T, typename Cmp = less < T > >
struct RMQ {
   T d[maxn][20];
   Cmp cmp;
    int w[maxn], sz;
    void init (const T *a, int n) {
        int i, j;
        for (w[0] = -1, i = 1; i \le n; ++i) w[i] = (i & 
             (i - 1)) ? w[i - 1] : w[i - 1] + 1;
        for (sz = n, i = 0; i < n; +i) d[i][0] = a[i];
        for (j = 1; (1 << j) <= n; ++j) {
            for (i = 0; i + (1 << j) <= n; ++i)
                d[i][j] = cmp(d[i][j - 1], d[i + (1 <<
                     (j - 1))][j - 1]) ? d[i][j - 1] : d
```

```
[i + (1 << (j - 1))][j - 1];
          }
     // index of a [l .. r]
     const T &query(int 1, int r) const {
           \begin{array}{l} \textbf{return} \ cmp(d[1][x]\,, \ d[r - (1 <\!\!< x) + 1][x]) \ ? \ d \\ [1][x] \ : \ d[r - (1 <\!\!< x) + 1][x]; \end{array} 
    }
};
int a[maxn], b[maxn];
int n;
RMQ<int>s;
RMQ\!\!<\!\!int\;,\;\;greater<\!\!int>>\;t\;;
int main() {
     cin >> n:
     for (int i = 0; i < n; i++) scanf ("%d", &a[i]);
     for (int i = 0; i < n; i++) scanf ("%d", &b[i]);
     s.init(b, n);
     t.\,i\,n\,i\,t\,\left(\,a\,,\ n\,\right)\,;
     int c, d;
    LL ans = 0;
     for (int i=0; i< n; i++) {
          if(a[i] > b[i]) continue;
          int ub = n+1, lb = i;
          while (ub-lb>1) {
              int mid = (ub+lb)>>1;
               if(t.query(i, mid-1) - s.query(i, mid-1) >
                    0) ub = mid;
               else lb = mid;
          int up = ub;
         ub = n+1, lb = i;
           while (ub-lb>1) {
              int mid = (ub+lb)>>1;
               if(t.query(i, mid-1) - s.query(i, mid-1)
                   >= 0) ub = mid;
               else lb = mid;
          int down = ub;
         ans += up-down;
     cout << ans << endl;
     return 0;
```

# 4.4 Link Cut Tree

```
//\,https://\,github.com/yzgysjr/ACM-ICPC-Templates/blob/
                master/Data\%20Structure/Link\%20Cut\%20Tree.cpp
 nil, *null;
 typedef node *tree;
#define isRoot(x) (x->pre->ch[0] != x && x->pre->ch[1]
               != x)
\hat{}= 1; \operatorname{swap}(t->\operatorname{ch}[0], t->\operatorname{ch}[1]); \}
 inline void PushDown(tree t) { if (t->rev) { MakeRev(t
                ->ch [0]); MakeRev(t->ch [1]); t->rev = 0; }
 inline void Rotate(tree x) {
        tree y = x->pre; PushDown(y); PushDown(x);
         int d = isRight(x);
          if \ (!isRoot(y)) \ y\text{-}pre\text{-}ch[isRight(y)] = x; \ x\text{-}pre = (!isRoot(y)) \ y\text{-}pre\text{-}ch[isRight(y)] = x; \ x\text{-}pre = (!isRoot(y)) \ y\text{-}pre\text{-}ch[isRoot(y)] = x; \ x\text{-}pre = (!isRoot(y)) \ y\text{-}pre = (!isRoot(y)) \ y\text{-
                        y->pre;
         if ((y-ch[d] = x-ch[!d]) != null) y-ch[d]-pre = y
        x->ch[!d] = y; y->pre = x; Update(y);
 inline void Splay(tree x) {
       PushDown(x); for (tree y; !isRoot(x); Rotate(x)) {
               y = x-pre; if (!isRoot(y)) Rotate(isRight(x))!=
                               isRight(y) ? x : y);
```

```
} Update(x);
inline void Splay(tree x, tree to) {
   PushDown(x); for (tree y; (y = x->pre) != to; Rotate(
            x)) if (y-pre != to)
        Rotate(isRight(x) != isRight(y) ? x : y);
   Update(x);
inline tree Access(tree t) {
    tree last = null; for (; t \stackrel{!}{=} null; last = t, t = t->
            pre) Splay(t), t->ch[1] = last, Update(t);
    return last;
inline void MakeRoot(tree t) { Access(t); Splay(t);
       MakeRev(t); }
inline tree FindRoot(tree t) { Access(t); Splay(t);
       tree last = null;
    for ( ; t != null; last = t, t = t->ch[0]) PushDown(t
            ); Splay(last); return last;
inline void Join(tree x, tree y) { MakeRoot(y); y->pre
       = x; 
inline void Cut(tree t) {Access(t); Splay(t); t->ch
        [0]->pre = null; t->ch[0] = null; Update(t);
inline void Cut(tree x, tree y) {
    tree upper = (Access(x), Access(y));
    if (upper == x) { Splay(x); y->pre = null; x->ch[1] =
              null; Update(x); }
    else if (upper = y) \{ Access(x); Splay(y); x->pre = access(x) \}
   null; y->ch[1] = null; Update(y); }
else assert(0); // 'impossible to happen
inline int Query(tree a, tree b) { // 'query the cost
       in path a <-> b, lca inclusive
    Access(a); tree c = Access(b); // c is lca
    int v1 = c->ch[1]->maxCost; Access(a);
    int v2 = c->ch[1]->maxCost;
   return \max(\max(v1, v2), c > \cot);
void Init() {
   null = \&nil; null -> ch[0] = null -> ch[1] = null -> pre =
            null; null -> rev = 0;
    Rep(i, 1, N) { node &n = base[i]; n.rev = 0; n.pre =
           n.ch[0] = n.ch[1] = null;
//compressed version
//http://trinklee.blog.163.com/blog/static
        /238158060201521101957375/
const int N=30010;
int n, fa[N], son[N][2], val[N], siz[N], stmp, rev[N];
\#define swap(a,b) (stmp=a,a=b,b=stmp)
void pu(int t) \{ siz[t] = siz[son[t][0]] + siz[son[t][1]] + 1; \}
void pd(int t) {rev[t]? rev[t]=0, rev[son[t][0]]^=1, rev[
       son[t][1]]^=1, swap(son[t][0], son[t][1]), 1:1;
bool nr(int t) \{return son[fa[t]][0] == t | | son[fa[t]][1] == t | | son[fa[t]][1] == t | s
void rtt(int t, int f=0, bool p=0){
       p=son[f=fa[t]][1]==t,
        fa[t] = fa[f], nr(f)? son[fa[f]][son[fa[f]][1] == f] = t:1,
       (son[f][p]=son[t][!p])?fa[son[f][p]]=f:1,
       pu(son [fa [f]=t][!p]=f);
void pv(int t){if(nr(t))pv(fa[t]);pd(t);}
void splay(int t,int f=0){
       for (pv(t); nr(t); rtt(t)) nr(f=fa[t])?
        {\rm rtt}\,({\rm son}\,[\,f][1]\!=\!=\!t\,\hat{}\,{\rm son}\,[\,fa\,[\,f]][1]\!=\!=\!f?t:f)\;,1\!:\!1;pu(\,t\,)\;;
void access(int t, int la=0)\{for(;t;splay(t),son[t][1]=
        la , la=t , t=fa [t]);}
void makeroot(int t){access(t),splay(t),rev[t]^=1;}
void link(int u, int v)\{makeroot(u), fa[u]=v;\}
void cut(int u, int v){makeroot(u), access(v), splay(v),
        son[v][0] = fa[u] = 0;
```

## 4.5 Treap

```
#include <bits/stdc++.h>
using namespace std;
struct Treap{
```

```
Treap *1, *r;
    int pri, key, val;
    Treap(int _val, int _key):
        val(val), key(key), l(NULL), r(NULL), pri(
            rand()){}
};
/// We assure that key value in A treap is greater than
     that in treap B
Treap *merge( Treap *a, Treap *b){
    if(a=NULL || b=NULL) return (!a) ? b : a;
    if (a->pri > b->pri) {
        a->r = merge(a->r, b);
        return a;
    else{
        b>1 = merge(a, b>1);
        return b:
void split (Treap *t , int k , Treap *&a , Treap *&b) {
    if(!t) a = b = NULL;
    else if (t->key <= k){
        a = t:
        split(t->r, k, a->r, b);
    }else{
        b = t:
        split(t->l, k, a, b->l);
Treap* insert( Treap *t, int k, int _val){
    Treap *tl, *tr;
    split(t, k, tl, tr);
    return merge(tl, merge(new Treap(_val, k), tr));
Treap* remove( Treap* t, int k){
    Treap *tl, *tr;
    split(t, k-1, tl, t);
    split(t, k, t, tr);
    return merge(tl, tr);
int main(){
    return 0;
```

#### 5 Math

#### 5.1 Prime Table

```
#include <bits/stdc++.h>
using namespace std;
struct Prime_table {
    int prime [1000000] = \{2,3,5,7\};
    int sz=4;
    // biggest prime < ub
    int ub=(1 << 20);
    int check(int num){
        int k = 0;
        for(k = 0; k < sz \&\& prime[k]*prime[k] <= num;
            k++){}
             if ( num % prime [k]==0) return 0;
        return 1;
    void buildprime(){
        int currentPrime=7;
        int j=4;
        for(sz=4,j=4; currentPrime < ub; sz++, j=6-j){
             currentPrime=currentPrime+j;
              if (check(currentPrime)) {
                 prime[sz] = currentPrime;
             else{
                 SZ - -;
        }
```

| } ptable;

5.2

Miller Rabin Prime Test

```
#include <bits/stdc++.h>
using namespace std;
typedef long long LL;
LL mul(LL a, LL b, const LL mod) {
    LL x = 0, y = a \% \text{ mod};
     while (b > 0) {
         if (b&1)
         x = (x + y) \% \text{ mod};

y = (y * 2) \% \text{ mod};
         b >>= 1:
     return x % mod;
LL mul(LL lhs, LL rhs, const LL mod) {
     return ( lhs * rhs ) % mod;
LL mypow(LL b, LL e, const LL mod) {
    LL x = 1;
    LL y = b;
     while ( e ) {
   if ( e&1 ) x = mul(x, y, mod);
         y = \operatorname{mul}(y, y, \operatorname{mod});
         e >>= 1:
     }
     return x;
const int testbase [] = \{2, 3, 5, 7, 11, 13, 17, 19, 23,
      29, 31, 37};
bool isprime (const LL p) {
     if (p < 2) return false;
     if (p != 2 && !(p&1) ) return false;
    LL\ d\,=\,p\ \textbf{-}\ 1;
     while (!(d\&1)) d >>= 1;
     for( int a : testbase ) {
         LL td = d:
         if ( a \ge p-1 ) return true;
         LL st = mypow(a, td, p);
          while (td!=p-1 && st!=1 && st!=p-1)
              st = mul(st, st, p);
              td \ll 1;
          if ( st != p - 1 && !(td&1) ) return false;
     return true;
int main() {
     int T;
     scanf("%d",&T);
     while (T--) {
         LL q;
         scanf("%lld",&q);
         puts(isprime(q)?"YES":"NO");
     return 0;
}
```

# 5.3 Extended Euclidean Algorithm

```
/** normal gcd function using recursion **/
int gcd(int a, int b){
   if(b = 0) return a;
   return gcd(b, a%b);
}
// Find solution of ax + by = gcd(a, b)
// ps : x, y may be negative
int extgcd(int a, int b, int& x, int& y){
   int d = a;
   if(b != 0) {
```

```
d = extgcd(b, a%b, y, x);
    y -= (a/b) * x;
}else {
    x = 1, y = 0;
}
return d;
}
```

```
5.4 Gauss Elimination
// solving linear equations with gauss elimination
#include <iostream>
#include <cmath>
#include <vector>
using namespace std;
void print(vector< vector<double> > A) {
    int n = A. size();
    for (int i=0; i<n; i++) {
         for (int j=0; j< n+1; j++) {
             if (j = n-1) { cout << "|";
         }
         cout << "\n";
    cout << endl;
}
vector<double> gauss(vector< vector<double> > A) {
    int n = A. size();
    for (int i=0; i< n; i++) {
         // Search for maximum in this column
         \frac{double \ maxEl = abs(A[i][i]);}{double \ maxEl = abs(A[i][i]);}
         int \max Row = i;
         for (int k=i+1; k<n; k++) {
             if \ (abs(A[k][i]) > maxEl) \ \{\\
                 maxEl = abs(A[k][i]);
                 \max Row = k;
             }
         }
        // Swap maximum row with current row (column by
              column)
         for (int k=i; k<n+1;k++)
             double tmp = A[maxRow][k];
             A[maxRow][k] = A[i][k];
             A[i][k] = tmp;
        // Make all rows below this one 0 in current
             column
         for (int k=i+1; k< n; k++)
             double c = -A[k][i]/A[i][i];
             \quad \  \   \text{for (int } j{=}i\;;\; j{<}n{+}1;\; j{+}{+}) \ \{
                 if (i==j)
                     A[k][j] = 0;
                 } else {
                     A[k][j] += c * A[i][j];
             }
        }
    // Solve equation Ax=b for an upper triangular
         matrix A
    vector < double > x(n);
    for (int i=n-1; i>=0; i--) {
        x[i] = A[i][n]/A[i][i];
         for (int k=i-1; k>=0; k--) {
             A[k][n] -= A[k][i] * x[i];
    return x;
```

int main() {

\* r.x);

```
int n:
    cin >> n;
                                                                      Complex conj () const {
                                                                          return Complex(x , -y);
    vector < double > line(n+1,0);
    vector < vector < double > > A(n, line);
                                                                      double operator = (const double a) {
                                                                          *this = Complex(a, 0);
     // Read input data
                                                                          return a;
    for (int i=0; i< n; i++) {
         for (int j=0; j<n; j++) {
                                                                 };
             cin >> A[i][j];
                                                                 const double pi = acos(-1.0);
                                                                 //fft with modulo, code referenced from the internet
    }
                                                                      fftMod::fftPrepare(len);
    for (int i=0; i<n; i++) {
                                                                      fftMod::convolution(res, le, ri, len, r-1);
         cin >> A[i][n];
                                                                 namespace fftMod{
                                                                     const int N = 1 \ll 18;
    // Print input
                                                                      const int Mod = 1e9 + 7;
                                                                      // to do, M should be about sqrt(Mod)
    print(A);
                                                                      const int M = 32768;
                                                                     // Calculate solution
    vector < double > x(n);
    x = gauss(A);
                                                                      Complex w[N];
    // Print result
                                                                      int rev[N];
    cout << "Result:\t";
    for (int i=0; i<n; i++) {
                                                                      void fftPrepare(int n) {
         cout <\!\!< x[\,i\,] <\!\!< "\ ";
                                                                          int LN = __builtin_ctz(n);
                                                                          for (int i = 0; i < n; ++ i) {
    double ang = 2 * pi * i / n;
    w[i] = Complex(cos(ang) , sin(ang));
    rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (
    cout << endl;
}
                                                                                   LN - 1));
                                                                          }
5.5 FFT
                                                                      void FFT(Complex P[], int n, int oper) {
   for (int i = 0 ; i < n ; i ++) {</pre>
typedef long double ld;
  N must be 2<sup>k</sup> and greater than array.size()
                                                                               if (i < rev[i]) {
 * FFT( a );
                                                                                   swap(P[i], P[rev[i]]);
 * FFT( b );
 * for(int i = 0; i < N; ++i) c[i] = conj(a[i] * b[i]);
 * FFT( c );
                                                                          \quad \text{for (int } d = 0; \ (1 <\!\!< d) < n; \ d+\!\!+\!\!) \ \{
 * for(int'i = 0; i < N; ++i) c[i] = conj(c[i]);
                                                                               int m = 1 \ll d, m2 = m * 2, rm = n / m2;
 * for (int i = 0; i < N; ++i) c[i] /= N;
                                                                               for (int i = 0; i < n; i += m2) {
                                                                                   for (int j = 0; j < m; j++) {
void FFT(vector< complex<ld>>& v) {
                                                                                       Complex &P1 = P[i + j + m], &P2 = P
    int N = v.size();
                                                                                            [i + j];
    Complex t = w[rm * j] * P1;
                                                                                        P1 = P2 - t;
         if(i>j) swap(v[i],v[j]);
                                                                                       P2 = P2 + t;
                                                                                   }
    for (int k = 2; k \le N; k < = 1) {
                                                                              }
         ld w = -2.0* pi/k;
                                                                          }
         complex < ld > deg(cos(w), sin(w));
         for(int j = 0; j < N; j + = k) 
             complex< ld > theta(1,0);
                                                                      Complex \ A[N] \ \ , \ B[N] \ \ , \ C1[N] \ \ , \ C2[N] \, ;
             for (int i = j; i < j+k/2; +++i) {
                                                                      void convolution (vector < int > &res, vector < int > &a,
                  complex < ld > a = v[i];
                                                                          vector < int > \&b, int len, int K) {
                  complex<ld> b = v[i+k/2]*theta;
                                                                          // a[ 0 .. len ) and b[ 0 .. len ) 's
                  v[i] = a+b;
                                                                              convolution % Mod
                  v[i+k/2] = (a-b);
                                                                          // stored in res[ 0 .. K+1 )
                  theta *= deg;
                                                                          for (int i = 0; i < len; +++i) {
A[i] = Complex(a[i] / M, a[i] \% M);
         }
                                                                              B[i] = Complex(b[i] / M, b[i] \% M);
    }
                                                                          \widetilde{FFT}(A, len, 1); FFT(B, len, 1);
// \texttt{http://sd-invol.github.io} / 2016/02/13/FFT-mod-prime/
                                                                          \quad \text{for (int } i = 0 \ ; \ i < len \ ; \ +\!\!\!+ i) \ \{
                                                                               int j = i ? len - i : i;
struct Complex {
                                                                               Complex a1 = (A[i] + A[j].conj()) * Complex
    double x , y:
    Complex (double \underline{x} = 0, double \underline{y} = 0) {
                                                                                   (0.5, 0);
                                                                               Complex a2 = (A[i] - A[j].conj()) * Complex
         x = _x , y = _y;
                                                                                   (0, -0.5);
                                                                              Complex b1 = (B[i] + B[j].conj()) * Complex
    Complex operator + (const Complex &r) const {
         return Complex(x + r.x , y + r.y);
                                                                               Complex b2 = (B[i] - B[j].conj()) * Complex
                                                                                   (0, -0.5);
    Complex operator - (const Complex &r) const {
                                                                               Complex c11 = a1 * b1 , c12 = a1 * b2;
         return Complex(x - r.x , y - r.y);
                                                                               Complex c21 = a2 * b1, c22 = a2 * b2;
                                                                               C1[j] = c11 + c12 * Complex(0, 1);
    Complex operator * (const Complex &r) const {
         return Complex(x * r.x - y * r.y , x * r.y + y
                                                                               C2[j] = c21 + c22 * Complex(0, 1);
```

}

```
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          FFT(C1 , len , -1); FFT(C2 , len , -1);
                                                                                               i -= k:
                                                                                              k /= 2;
           for (int i = 0; i \le K; ++ i) {
                                                                                         if(j < k) j += k;
                int x = (LL)(C1[i].x / len + 0.5) \% Mod;
                }
                \begin{array}{l} \text{int } z = (LL)(C2[\,i\,].\,y\,\,/\,\,len\,\,+\,\,0.5)\,\,\%\,\,Mod; \\ \text{res}\,[\,i\,] = ((LL)x\,\,*\,\,M\,\,*\,\,M\,+\,\,(LL)(y1\,\,+\,\,y2)\,\,*\,\,M \end{array}
                                                                               inline void ntt(int *y, int len, int on) {
                                                                                    change(y, len);
                                                                                    for (int h = 2; h \le len; h \le 1) {
                      + z) % Mod;
                                                                                          \begin{array}{lll} \textbf{int} & wn = quick\_mod(G, \ (P \ \textbf{-} \ 1) \ / \ h, \ P); \end{array}
          }
                                                                                          for (int j = 0; j < len; j += h) {
     }
};
                                                                                              int w = 1;
5.6 NNT
NTT( a );
                                                                                              }
NTT( b );
                                                                                         }
for (int i = 0; i < N: ++i)
     c[i] = (long long) a[i] * b[i] % mod;
                                                                                    if (on == -1) {
NTT( c, true );
for (int i = 0; i < N; ++i)
     c[i] = (786433LL-12) * c[i] \% mod;
constexpr int mod = 786433;
                                                                                    }
constexpr int N = 65536;
                                                                               }
 void NTT(vector< int >& v, bool flag = false)
                                                                               int n;
                                                                               int r[3][3];
      for (int i = 1, j = 0; i < N; ++i)
                                                                               inline int CRT(int *a) {
           for (int k = N > 1; !((j^=k)&k); k > = 1);
           if(i>j) swap(v[i],v[j]);
      for (int k = 2; k \le N; k < = 1)
           for(int j = 0; j < N; j + = k)
                                                                                         }
                int theta = 1;
                for(int i = j; i < j+k/2; ++i)
                {
                     int a = v[i];
                    \begin{array}{ll} \text{int } b = (\underset{i=1}{\text{long long}}) \ v[\,i+k/2\,]^* \\ \text{theta\%mod}; \\ v[\,i\,] = (\underset{i=1}{\text{a+b}}) \% \ \text{mod}; \end{array}
                                                                                    return ans;
                     v[i+k/2] = (a-b+mod)\%mod;
                     theta = (long long) theta * deg % mod;
                                                                               int tmp[maxn][3];
          }
     }
}
constexpr int mod = 1e9+7;
typedef vector<int> VEC;
                                                                                         P\,=\,\mathrm{magic}\,[\;\mathrm{t}\,\mathrm{i}\;]\,;
// ntt + Crt, code referenced from the internet
                                                                                          int k;
namespace nttCrt {
     constexpr int magic [3] = \{1004535809, 998244353,
```

104857601};

constexpr int G = 3;

int ans = 1;

while (k) {

return ans;

k >>= 1;

int P:

constexpr int  $\widehat{MOD} = 1000000007$ ;

 ${\tt inline\ int\ quick\_mod(int\ x,\ int\ k,\ int\ MOD)\ \{}$ 

inline void change(int \*y, int len) {

是反转的

int k = len / 2;while(j >= k) {

if (k&1) ans = 1LL \* ans \* x % MOD; x = 1LL \* x \* x % MOD;

for (int i = 1, j = len / 2; i < len - 1; i++) {
 if (i < j) swap(y[i], y[j]);

//交换互为小标反转的元素, i<j保证交换一次

//i做正常的+1, j左反转类型的+1,始终保持i和j

```
for (int k = j; k < j + h / 2; k++) {
                    int u = y[k] \% P;
int t = 1LL * w * y[k + h / 2] \% P;
                    y[k] = (u + t) \% P;
                    y[k + h / 2] = ((u - t) \% P + P) \%
                    w = 1LL * w * wn \% P;
          for (int i = 1; i < len / 2; i++)
              swap(y[i], y[len - i]);
          int inv = quick\_mod(len, P - 2, P);
          for (int i = 0; i < len; i++)

y[i] = 1LL * y[i] * inv % P;
     int sb[3] = \{a[0], a[1], a[2]\};
for (int i = 0; i < 3; ++i) {
          for (int j = 0; j < i; ++j) {
               int t = (sb[i] - sb[j]) % magic[i];
               if(t < 0) t += magic[i];
               sb[i] = 1LL * t * r[j][i] % magic[i];
     int mul = 1, ans = sb[0] \% MOD;
     int x1[maxn*2], x2[maxn*2];
inline void gao(vector<int>& res, vector<int> &a,
     vector<int> &b ,int len , int kk) {
     for (int ti = 0; ti < 3; ti++) {
          for (k = 0; k < SZ(a) & k < len; k++) x1[
               \mathbf{k}] = \mathbf{a}[\mathbf{k}];
          \begin{array}{lll} & \text{for } (\ ; k < len \, ; \ k++) \ x1 \, [k] \, = \, 0 \, ; \\ & \text{for } (\ k = \, 0 \, ; \ k < \, SZ(b) \, \&\& \, k < \, len \, ; \ k++) \, \, x2 \, [ \end{array}
               k = b[k];
          for (; k < len; k++) x2[k] = 0;
          ntt(x1, len, 1); ntt(x2, len, 1);
          for (int i = 0; i < len; i++) x1[i] = 1LL *
                x1[i] * x2[i] % P;
          ntt(x1, len, -1);
          for (int i = 0; i \le kk; i++) tmp[i][ti] =
               x1[i];
     for (int i = 0; i \le kk; i++) res[i] = CRT(tmp
          [ i ])
inline void init() {
     for (int i = 0; i < 3; i++) {
          for (int j = 0; j < 3; j++) {
               r\left[\:i\:\right]\left[\:j\:\right] \: = \: quick\_mod\left(\:magic\left[\:i\:\right]\:,\:\: magic\left[\:j\:\right]\:
                     - 2, magic[j]);
```

```
} };
```

## 5.7 Big Number

```
//http://blog.csdn.net/hackbuteer1/article/details
    /6595881
#include<iostream>
#include<string>
#include<iomanip>
#include<algorithm>
using namespace std;
#define MAXN 9999
#define MAXSIZE 10
#define DLEN 4
class BigNum
private:
 int a[500];
               //可以控制大数的位数
 int len;
               //大数长度
public:
 BigNum() { len = 1; memset(a,0,sizeof(a)); }
                                          //构造函
 BigNum(const int);
                        //将一个int类型的变量转化为
     大数
                        //将一个字符串类型的变量转化
 BigNum(const char*);
     为大数
 BigNum(const BigNum &); //拷贝构造函数
 BigNum & operator = (const BigNum &);
                                  //重载赋值运算
     符,大数之间进行赋值运算
  friend \ istream \& \ operator >> (istream \&, \ BigNum \&); \\
                                               //
     重载输入运算符
 {\tt friend} \ \ ostream \& \ \ operator << (ostream \&, \ \ BigNum \&);
     重载输出运算符
 BigNum operator+(const BigNum &) const;
                                       //重载加法
     运算符,两个大数之间的相加运算
 BigNum operator - (const BigNum &) const;
                                        //重载减法
     运算符,两个大数之间的相减运算
 BigNum operator*(const BigNum &) const;
                                        //重载乘法
     运算符,两个大数之间的相乘运算
 BigNum operator/(const int
                          &) const;
                                        //重载除法
     运算符,大数对一个整数进行相除运算
                                       //大数的n次
 BigNum operator (const int &) const;
     方运算
                                       //大数对一个
 int
       operator%(const int &) const;
     int类型的变量进行取模运算
      operator > (const BigNum & T) const;
                                        //大数和另
     一个大数的大小比较
       operator > (const int & t) const;
                                        //大数和一
     个int类型的变量的大小比较
 void print();
                   //输出大数
                            //将一个int类型的变量转
BigNum::BigNum(const int b)
   化为大数
 int c, d = b;
 len = 0;
 memset(a,0,sizeof(a));
 while (d > MAXN)
   c = d - (d / (MAXN + 1)) * (MAXN + 1);
   d = d / (MAXN + 1);
   a[len++] = c;
 a[len++] = d;
BigNum::BigNum(const char*s)
                             //将一个字符串类型的变
   量转化为大数
```

```
int t,k,index,l,i;
  memset(a, 0, sizeof(a));
  l=strlen(s);
  len=l/DLEN;
  if (1%DLEN)
    len++;
  index=0;
  for (i=l-1; i>=0; i-=DLEN)
    t = 0:
    k=i -DLEN+1;
    if(k<0)
      k=0;
    for(int j=k; j<=i; j++)
      t=t*10+s[j]-'0';
    a[index++]=t;
BigNum::BigNum(const BigNum & T) : len(T.len) //拷贝构
    造函数
{
  int i;
  memset(a, 0, sizeof(a));
  for(i = 0 ; i < len ; i++)
    a\,[\;i\;]\;=T.\,a\,[\;i\;]\,;
                                                   //重载赋
BigNum & BigNum::operator=(const BigNum & n)
    值运算符, 大数之间进行赋值运算
{
  int i;
  len = n.len;
  memset(a, 0, sizeof(a));
  for(i = 0 ; i < len ; i++)
  a[i] = n.a[i];
return *this;
istream& operator>>(istream & in, BigNum & b)
    输入运算符
  char ch[MAXSIZE*4];
  int i = -1;
  in >> ch:
  int l=strlen(ch);
  int count=0,sum=0;
  for (i=l-1; i>=0;)
  {
    sum = 0;
    int t=1;
    for (int j=0; j<4&&i>=0; j++, i--, t*=10)
      sum + = (ch[i] - '0') *t;
    b.a[count]=sum;
    count++;
  b.len =count++;
  return in;
ostream& operator << (ostream& out, BigNum& b)
                                                    //重载
     输出运算符
{
  int i;
  cout << b.a[b.len - 1];
  for (i = b.len - 2 ; i >= 0 ; i--)
    \operatorname{cout}.\operatorname{width}\left(\operatorname{DLEN}\right);
    cout. fill('0');
    cout << b.a[i];
  return out;
}
BigNum BigNum::operator+(const BigNum & T) const //两
     个大数之间的相加运算
  BigNum t(*this);
                   //位数
  int i, big;
  big = T.len > len ? T.len : len;
```

```
for(i = 0 ; i < big ; i++)
                                                                                     ret.a[i + j] = temp;
                                                                                  }
     t.a[i] +=T.a[i];
     if(t.a[i] > MAXN)
                                                                                if(up != 0)
                                                                                  ret.a[i + j] = up;
       t.a[i + 1]++;
       t.a[i] -=MAXN+1;
                                                                             ret.len = i + j;
                                                                             while(ret.a[ret.len - 1] == 0 \&\& ret.len > 1)
                                                                               ret.len--:
  if(t.a[big] != 0)
                                                                             return ret;
     t.len = big + 1;
                                                                          BigNum BigNum::operator/(const int & b) const //大数
     t.len = big;
                                                                                对一个整数进行相除运算
  return t;
                                                                             BigNum ret;
                                                                 //两
BigNum BigNum::operator-(const BigNum & T) const
                                                                             \begin{array}{ll} \textbf{int} & i \ , down \ = \ 0 \, ; \end{array}
     个大数之间的相减运算
                                                                             for (i = len - 1 ; i >= 0 ; i--)
                                                                                ret.a[\,i\,] \;=\; (\,a\,[\,i\,] \;+\; down \;\;*\;\; (M\!A\!X\!N + \,1)\,) \;\;/\;\; b\,;
  int i,j,big;
                                                                               down \, = \, a \, [\, i \, ] \, + \, down \, * \, (MAXN \, + \, 1) \, - \, ret \, . \, a \, [\, i \, ] \, * \, b \, ;
  bool flag;
  BigNum t1, t2;
  if(*this>T)
                                                                             while (ret.a[ret.len - 1] == 0 && ret.len > 1)
  {
     t1=*this:
                                                                               ret.len - -;
     t2=T;
                                                                             return ret;
     flag=0;
                                                                          int BigNum::operator %(const int & b) const
                                                                                                                                       //大数对
  else
                                                                                一个int类型的变量进行取模运算
  {
     t1=T;
                                                                             int i, d=0;
     t2 = *this;
                                                                             for (i = len -1; i>=0; i--)
     flag=1;
                                                                               d = ((d * (MAXN+1))\% b + a[i])\% b;
  big=t1.len;
  for(i = 0 ; i < big ; i++)
                                                                             return d:
     i\,f\,(\,t\,1\,.\,a\,[\,i\,]\,<\,t\,2\,.\,a\,[\,i\,]\,)
                                                                          BigNum BigNum::operator^(const int & n) const
                                                                                                                                          //大数
     {
                                                                                的n次方运算
       j = i + 1;
       \mathbf{while}\,(\,\mathrm{t1.a}\,[\,\mathrm{j}\,]\,=\!\!\!=\,0)
                                                                             BigNum t, ret(1);
         j++;
                                                                             int i:
       t1.a[j--]--:
                                                                             if(n<0)
       while(j > i)
                                                                               exit(-1):
         t1.a[j--] += MAXN;
                                                                             if(n==0)
       t1.a[i] += MAXN + 1 - t2.a[i];
                                                                               return 1;
                                                                             if(n==1)
     else
                                                                                return *this;
       t1.a[i] -= t2.a[i];
                                                                             int m=n;
                                                                             while (m>1)
  t1.len = big;
                                                                             {
  while (t1.a[len - 1] = 0 \&\& t1.len > 1)
                                                                                t=*this;
                                                                               for (i=1; i << 1 <= m; i << =1)
     t1.len - -:
                                                                                1
     big --;
                                                                                  t=t*t:
  if (flag)
                                                                               m=i;
     t1.a[big-1]=0-t1.a[big-1];
                                                                                ret=ret*t;
  return t1;
                                                                                if (m=1)
}
                                                                                  ret=ret*(*this);
BigNum BigNum::operator*(const BigNum & T) const
                                                                             return ret;
     个大数之间的相乘运算
                                                                          bool BigNum::operator>(const BigNum & T) const
                                                                                                                                         //大数
  BigNum ret;
                                                                                和另一个大数的大小比较
  int i,j,up;
  int temp, temp1;
                                                                             int ln;
  \quad \  \  \, \text{for} \, (\, i \, = \, 0 \ ; \ i \, < \, len \ ; \ i + \! + \! )
                                                                             if (len > T.len)
  {
                                                                               return true;
     up = 0;
                                                                             else if (len = T.len)
     for(j = 0 ; j < T.len ; j++)
                                                                             {
                                                                               ln = len - 1;
       temp \, = \, a \, [\, i \, ] \ * \ T.\, a \, [\, j \, ] \ + \ ret.\, a \, [\, i \ + \ j \, ] \ + \ up \, ;
                                                                                while (a[ln] = T.a[ln] & ln >= 0)
       if (temp > MAXN)
                                                                                  ln - -:
                                                                                if(ln >= 0 \&\& a[ln] > T.a[ln])
         temp1 \, = \, temp \, \left. \begin{array}{cccc} - & temp \end{array} \right. / \left. \left( \begin{array}{cccc} MAXN \, + \, 1 \right) \right. \, * \left. \left( MAXN \, + \, 1 \right) ; \end{array} \label{eq:temp1}
                                                                                  return true;
          up = temp / (MAXN + 1);
                                                                                else
          ret.a[i + j] = temp1;
                                                                                  return false;
                                                                             }
       else
                                                                             else
                                                                                return false;
         up \, = \, 0\,;
                                                                          }
```

```
//大数
bool BigNum::operator >(const int & t) const
    和一个int类型的变量的大小比较
  BigNum b(t);
  return *this>b;
                         //输出大数
void BigNum::print()
{
  int i;
  cout \ll a[len - 1];
  for (i = len - 2; i >= 0; i--)
    cout.width(DLEN);
    cout.fill('0');
    cout << a[i];
  cout << endl;
int main(void)
{
  int i,n;
  BigNum x[101];
                       //定义大数的对象数组
  x[0]=1;
  for (i=1; i<101; i++)
  x[i]=x[i-1]*(4*i-2)/(i+1);
while (scanf("%d",&n)==1 && n!=-1) {
    x[n].print();
}
```

# 6 string

### 6.1 Palindromic Tree

```
回文自動機包含以下元素:
   狀態St, 所有節點的集合, 一開始兩個節點, 0表示偶數長
       度串的根和1表示奇數長度串的根
   last 新增加一個字符後所形成的最長回文串的節點編號
    s 當前的字符串(一開始設s[0]=-1(可以是任意一個在串S
       中不會出現的字符))
   n 表示添加的字符個數
每個節點代表一個不同的回文子字串,我們在每個節點會儲存
    一些數值:
   len 表示所代表的回文子字串長度
   next[c]表示所代表的回文子字串在頭尾各增加一個字符c
       後的回文字串其節點編號
   sufflink 表示所代表的回文子字串不包括本身的最長後綴
       回文子串的節點編號
   cnt(非必要) 表示所代表的回文子字串在整體字串出現的
       次數(在建構完成後呼叫count()才能計算)
   //num(非必要) 表示所代表的回文子字串其後綴為回文字
       串的個數 <== not included
{\color{red} \textbf{struct}} \hspace{0.1cm} \textbf{palindromic\_tree} \{
   struct node{
       int next[26], sufflink, len; /*這些是必要的元素*/int l, r; // this node is s[l...r]
                             /*這些是額外維護的元素*/
       int cnt, num;
       \mathsf{node}(\mathsf{int}\ l{=}0) \colon \mathsf{sufflink}\left(0\right), \mathsf{len}\left(1\right), \mathsf{cnt}\left(0\right), \mathsf{num}\left(0\right) \{
           for (int i=0; i<26;++i) next [i]=0;
   std::vector<node> St;
   std::string s; //current string [ 1 .. n ]
   int last, n;
   palindromic\_tree(): St(2), last(1), n(0)
       St[0].sufflink=1;
       St[1].len=-1;
       s.push\_back(-1);
   }
```

```
inline void clear(){
        St.clear();
        s.clear();
        last=1:
        n=0;
        St.push_back(0);
        St.push\_back(-1);
        St[0].sufflink=1;
        s.push\_back(-1);
    inline int get_sufflink(int x){
        sufflink;
        return x;
    inline void add(int c){
        s.push_back(c-= 'a');
        ++n;
        int cur=get_sufflink(last);
        if (!St[cur].next[c]) {
            int now=St.size();
            St.push_back(St[cur].len+2);
            St [now]. sufflink=St [get_sufflink(St [cur].
                 sufflink)].next[c];
             /*不用擔心會找到空節點,由證明的過程可知*/
            St[cur].next[c]=now;
            St \left[ now \right].num\!\!=\!\!St \left[ \, St \left[ now \right].\,sufflink \, \right].num\!+\!1;
            St[now] \cdot l = n - St[now] \cdot len + 1, St[now] \cdot r
                 = n:
        last=St[cur].next[c];
        ++St[last].cnt;
    inline void count(){/*cnt必須要在構造完後呼叫count
        ()去計算*/
        std::vector<node>::reverse_iterator i=St.rbegin
            ();
        for (; i!=St.rend();++i) {
            St[i->sufflink].cnt+=i->cnt;
    inline int size(){/*傳回其不同的回文子串個數*/
        return St. size()-2;
}ptree;
```

#### 6.2 Suffix Array

#### 6.3 Longest Palindromic Substring

```
//ntu judge Earse
#define maxn 200001
char t [maxn];
char s [maxn * 2];
int z [maxn * 2];
int N:
int longest_palindromic_substring() {
   // t穿插特殊字元, 存放到s。
   int n = strlen(t);
   N = n * 2 + 1;
memset(s, '.', N);
   for (int i=0; i< n; ++i) s[i*2+1] = t[i];
   s[N] = ' \setminus 0';
              // if無須使用, then無須計算。
   z[0] = 1;
   int L = 0, R = 0;
   for (int i=1; i<N; ++i) // 從z[1] 開始
       z[i] = (R > i) ? min(z[2*L-i], R-i) : 1;
       s[i-z[i]] = s[i+z[i]]) z[i]++;
        if (i+z[i] > R) L = i, R = i+z[i];
   }
   // 尋找最長迴文子字串的長度
   n = 0;
   int p = 0;
   for (int i=1; i<N; ++i) // 從z[1] 開始
```

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```
if (z[i] > n)
            n = z[p = i];
   // longest 從中心到外端的長度 ⇒ (n-2)/2
   //cout << "最長迴文子字串的長度是" << (2*n-1) / 2;
   // 印出最長迴文子字串, 記得別印特殊字元。
        for (int i=p-z[p]+1; i<=p+z[p]-1; ++i)
            if (i & 1) {
               cout \ll s[i];
   return (2*n-1)/2;
int nxt[maxn * 2];
int main() {
   int T; cin>>T;
   while (T--) {
    scanf("%s", t);
       #ifdef DEBUG
           cout << longest_palindromic_substring() <<</pre>
                endl;
       #else
           longest_palindromic_substring();
       #endif
       memset(nxt, -1, sizeof(nxt));
       for (int i = 0; i < N; i++) {
           nxt[i-z[i]+1] = i+1;
       int leftmost = 0;
        for(int i = 0; i < N; i++) {
            leftmost = max(leftmost, nxt[i]);
            nxt[i] = max(leftmost, nxt[i]);
        int ans = 0;
        for (int cur = 0; cur < N-1;) {
            cur = nxt[cur];
            ans++;
       cout << ans << endl;
   return 0;
```

# 7 geometry

#### 7.1 Point Class

```
const double eps = 1e-10;
#define N 100
struct P {
    \quad \quad double \ x, \ y;
    P({\color{red} \textbf{double}} \ \_x{=}0, \ {\color{red} \textbf{double}} \ \_y{=}0) \ : x(\_x) \ , \ y(\_y) \ \{\};
    void read() {
         scanf("%lf%lf",&x,&y);
    void print() {
         printf("%f %f\n", x, y);
} p[N];
bool operator <( P a, P b ) { return tie(a.x,a.y)<tie(b
.x,b.y); }
P operator +( P a, P b ) { return P{a.x+b.x,a.y+b.y}; }
P operator -( Pa, Pb) { return P{a.x-b.x,a.y-b.y}; }
P operator *( P b, double a ) { return P\{a*b.x,a*b.y\};
P operator /(Pa, double b) { return P\{a.x/b, a.y/b\};
P& operator /=( P &a, double b ) { return a=a/b; }
double operator *( Pa, Pb) { return a.x*b.x+a.y*b.y;
double operator ^( Pa, Pb) { return a.x*b.y-a.y*b.x;
double x( P o, P a, P b ) { return (a-o)^(b-o); }
double dot(Po, Pa, Pb) { return (a-o)*(b-o); }
```

## 7.2 Intersection of Circles/Lines/Segments

```
//PECaveros
 vector<P> interCircle( P o1 , double r1 , P o2 , double
      r2 ){
   double d2 = (01 - 02) * (01 - 02);
   \frac{double}{double} d = sqrt(d2);
   if (d > r1 + r2) return \{\};
   P u = (o1+o2)*0.5 + (o1-o2)*((r2*r2-r1*r1)/(2*d2));
   double A = sqrt((r1+r2+d)*(r1-r2+d)*(r1+r2-d)*(-r1+r2
       +d));
   P\ v \,=\, P(\ o1\,.\,y\,\text{-}\,o2\,.\,y\ ,\ \text{-}\,o1\,.\,x\,\,+\,\,o2\,.\,x\ )\ *\ A\ /\ (2*d2)\,;
   return \{u+v, u-v\};
| P interPnt( P p1, P p2, P q1, P q2){
   double f1 = (p2 - p1)
   double f = (f1 + f2);
   if( fabs( f ) < eps ) return Pt( nan(""), nan("") );
return q1 * ( f2 / f ) + q2 * ( f1 / f );</pre>
int ori( const PLI& o , const PLI& a , const PLI& b ){ LL ret = ( a - o ) ^ ( b - o );
   return ret / max('lll', abs('ret'));
 // p1 == p2 || q1 == q2 need to be handled
bool banana (const PLL& p1 , const PLL& p2
                const PLL& q1 , const PLL& q2 ){
   if( ( ( p2 - p1 ) ^ ( q2 - q1 ) ) == 0 ){ // parallel
     if (ori(p1, p2, q1)) return false;
return ((p1 - q1) * (p2 - q1)) <= 0 |
             ((p1 - q2) * (p2 - q2)) <= 0
              ((q1 - p1) * (q2 - p1)) <= 0
             ((q1 - p2) * (q2 - p2)) <= 0;
   return (ori( p1, p2, q1 ) * ori( p1, p2, q2 )<=0) &&
           (\,{\rm ori}\,(\ q1\,,\ q2\,,\ p1\ )\ *\ {\rm ori}\,(\ q1\,,\ q2\,,\ p2\ )<\!\!=\!\!0);
```

## 7.3 Convex Hull

```
#define REP(i,n) for ( int i=0; i<int(n); i++)
void input() {
    scanf("%d",&n);
    REP(i,n) p[i].read();
}
P findCenter() {
    p[n]=p[0];
    P center=P\{0,0\};
    REP(i,n) {
         double v=p[i]*p[i+1];
         center.x += (p[i].x+p[i+1].x)*v;
         center.y += (p[i].y+p[i+1].y)*v;
    double area=0;
    REP(i,n) area+=p[i]*p[i+1];
    area \neq 2;
    center /= 6*area;
    return center;
P q1[N], q2[N], q[N];
void convex() {
    sort(p,p+n);
    int m1=0, m2=0;
    REP(\:i\:\:,n\:) \ \ \{
         while (m1>=2 \&\& X(q1[m1-2],q1[m1-1],p[i]) >= 0
              ) m1--;
         while (m2>=2 \&\& X(q2[m2-2],q2[m2-1],p[i]) <= 0
              ) m2--
         q1 [m1++]=q2 [m2++]=p[i];
    int m=0;
    REP(i, m1) q[m++]=q1[i];
```

```
q[m]=q[0];
void solve() {
  convex();
  // continue ...
```

#### 7.4 Half Plane Intersection

```
//http://acm.csie.org/ntujudge/problemdata/2575.pdf
// http://www.\,csie.ntnu.edu.tw/{\sim}u91029/Half-
    planeIntersection.html
預先使用四個半平面, 設定一個極大的正方形邊界, 讓半平面
    交集擁有邊界。
二、逐一加入每個半平面,求出當下的半平面交集(凸多邊
    形)。
online 演算法, 隨時維護一個半平面交集。每次更新需時 O(N
   ) , 總時間複雜度為 O(N<sup>2</sup>) , N 是半平面數目。
#include <bits/stdc++.h>
using namespace std;
#define mp make_pair
typedef complex<double> Point;
typedef vector<Point> Polygon;
typedef pair<Point, Point> Line;
#define x real()
\#define y imag()
// 兩向量叉積
double cross (Point& a, Point& b) {
   return a.x * b.y - a.y * b.x;
// 向量oa與向量ob進行叉積
double cross(Point& o, Point& a, Point& b) {
    return (a.x-o.x) * (b.y-o.y) - (a.y-o.y) * (b.x-o.x)
       );
// 多邊形面積
double area (Polygon& p) {
   double a = 0;
   int n = p.size();
   for (int i=0; i< n; ++i)
       a += cross(p[i], p[(i+1)\%n]);
    return fabs(a) / 2;
}
// 兩線交點
Point intersection (Point& a1, Point& a2, Point& b1,
    Point& b2) {
    Point a = a2 - a1, b = b2 - b1, s = b1 - a1;
   return a1 + a * cross(b, s) / cross(b, a);
// 一個凸多邊形與一個半平面的交集
Polygon halfplane_intersection(Polygon& p, Line& line)
   Polygon q;
    Point p1 = line.first, p2 = line.second;
    // 依序窮舉凸多邊形所有點, 判斷是否在半平面上。
    // 如果凸多邊形與半平面分界線有相交, 就求交點。
   int n = p.size();
   for (int i=0; i< n; ++i)
       double c = cross(p1, p2, p[i]);
       double d = cross(p1, p2, p[(i+1)\%n]);
        if (c >= 0) q.push_back(p[i]);
        if \ (c\ *\ d\ <\ 0)\ q.push\_back(intersection(p1,\ p2,
            p[i], p[(i+1)\%n]);
    return q;
```

```
#define maxn 550
//Line line [maxn];
Point v[maxn];
double ans [maxn];
int main() {
    int T; cin>>T;
    while (T--) {
        int n:
        double w, h;
        scanf("%d %lf %lf", &n, &w, &h);
        // 預先設定一個極大的正方形邊界
        Polygon p, org;
        /** initialize
        p.push_back(Point(-1e9,-1e9));
        p.push\_back(Point(-1e9,+1e9));
        p.push_back(Point(+1e9,-1e9));
        p.push_back(Point(+1e9,+1e9));
        p.push\_back(Point(0,0));
        p.push_back(Point(0,h));
        p.push\_back(Point(w,h));
        p.push\_back(Point(w,0));
        org = p;
        for (int i =0; i < n; i ++) {
            double a, b;
            scanf("%lf %lf", &a, &b);
v[i] = Point(a, b);
        // 每一個半平面都與目前的半平面交集求交集
        for (int i=0; i<n; ++i)
            p = org;
            for (int j = 0; j < n; j++) {
                if (i=j) continue;
                Line line;
                // find perpendicular line to line i_j
                Point a( (v[i].x+v[j].x)/2, (v[i].y+v[j].x)
                    [1,y]/2;
                Point b(a.x+(v[i].y-v[j].y), a.y-(v[i].
                    x-v[j].x));
                line = cross(a, b, v[i]) >= 0 ? mp(a, b)
                     : mp(b,a);
                p = halfplane_intersection(p, line);
                if (area(p) == 0) break; // 退化或者
                     空集合
            ans[i] = area(p);
        for (int i = 0; i < n; i ++) printf ("%.9f\n", ans [
            i]);
    }
}
10
3 4 4
1 1 2 2 3 3
```