Contents

1	Basi	ic 1	
	1.1	.vimrc	
	1.2	Increase Stack Size	
	1.3	digitDP	
	1.4	DP(convex hull optimization)	
	1.5	simulated annealing	
2	Gra	ph 3	
	2.1	HLD	
	2.2	Hungarian	
	2.3	KM	
	2.4	Bi-vertex-connected Subgraph	
	2.5	Bi-edge-connected Subgraph	
	2.6	SCC	
	2.7	Steiner Tree(PECaveros)	
	2.8	Edmond's Matching Algorithm	
	2.9	Tree Decomposition	
	2.10	Tree Longest Path	
	2.10	Tree Dongest Latin	
3	Flov	v 9	
	3.1	Dinic Maxflow	
4	Dat	a Structure 10	
	4.1	Disjoint Set	
	4.2	Sparse Table	
	4.3	Treap	
5	Math 11		
•	5.1	Prime Table	
	5.2	Miller Rabin Prime Test	
	5.3	Extended Euclidean Algorithm	
	5.4	Gauss Elimination	
	5.4 - 5.5	FFT	
	5.6	NNT	
	5.7	Big Number	
	5.7	Dig Number	
6	strin		
	6.1	Palindromic Tree	
	6.2	Suffix Array	
	6.3	Longest Palindromic Substring 17	
7	geometry 18		
	7.1	Point Class	
	7.2	Intersection of Circles/Lines/Segments 18	
	7.3	Convex Hull	
	7.4	Half Plane Intersection	

1 Basic

```
1.1 .vimrc
imap jj <Esc>
sv on
se sw=4 ts=4 sts=4 et nu sc hls cc=69
filet plugin indent on
nm <F5> :!"./%<"<CR>
nm <F6> :!"./%<" < input.txt<CR>
au FileType cpp no <F9> :!g++ % -o
 \ -Wshadow -Wno-unused-result<CR>
no <expr> <silent> <Home> col('.') ==
 \ match(getline('.'),'\S') + 1
 \ ? '0' : '^;
im < silent > < Home > < C-O > < Home >
    Increase Stack Size
//stack resize
asm( "mov %0,%%esp\n" :: "g"(mem+10000000));
//change esp to rsp if 64-bit system
//stack resize (linux)
```

}

 $\operatorname{digitDP}$

}

1.3

 $if(res==0){$

#include <sys/resource.h>
void increase_stack_size() {
 const rlim_t ks = 64*1024*1024;

if(rl.rlim_cur<ks){
 rl.rlim_cur=ks;</pre>

int res=getrlimit(RLIMIT_STACK, &rl);

res=setrlimit(RLIMIT STACK, &rl);

struct rlimit rl;

```
简介
顾名思义,所谓的数位DP就是按照数字的个,十,百,千·····
   位数进行的DP。
数位DP的题目有着非常明显的性质:
  询问[1,r]的区间内, 有多少的数字满足某个性质
做法根据前缀和的思想, 求出[0,1-1]和[0,r]中满足性质的数
   的个数,然后相减即可。
算法核心
LL dfs(int x,int pre,int bo,int limit);
一般需要以上参数 (当然具体情况具体分析)
  x表示当前的数位 (一般都是从高位到低位)
  pre表示前一位的数字
  bo可以表示一些附加条件:是否有前项0,是否当前已经符
      合条件……
  limit 这个很重要! 它表示当前数位是否受到上一位的限
     制, 比较抽象, 举例说明
   如果上限是135,前两位已经是1和3了,现在到了个位,个
     位只能是5以下的数字
注: 如果当前受限,不能够记忆化,也不能返回记忆化的结果
为了避免多次调用时 每次上限不同 而导致的错
//http://acm.csie.org/ntujudge/view_code.php?id=106844
// Multiples
LL x;
int digit [100];
LL ten_pow[ 15 ];
bool ava[15];
LL dp[15][2][1000000];
LL dfs(int len, LL mod, bool bo, bool limit) {
```

```
if ( len < 0 ) return mod == 0;
    if ( !limit && dp[len][bo][mod] != -1 ) return dp[
         len ] [bo] [mod];
    int up = limit? digit[len] : 9;
    LL ret = 0;
    for (int i = 0; i \le up; i++) if (ava[i] || (i=0&&
         bo) ) {
         ret += dfs(len-1, (mod+ten_pow[len]*i)%x, bo
             &&(!i), limit&&(i=up));
    if( !limit ) dp[len][bo][mod] = ret;
    return ret;
LL solve (LL num) {
    int len = 0; digit [0] = 0;
    while( num ) {
         digit[len++] = num\%10;
         num = 10;
    return dfs (len -1, 0,1, 1);
bool check (LL num) {
    while ( num )
         if (!ava[ num%10 ] ) return false;
         num /= 10:
    return true;
int main() {
    LL A, B;
    cin>>x>>A>>B;
    ten_pow[0] = 1;
    \begin{array}{lll} mem(\ dp\,,\ -1)\,; \\ for\,(int\ i\,=\,1;\ i\,<\,15;\ i+\!\!+\!\!\!) \end{array}
         ten_pow[i] = (ten_pow[i-1]*10)\%x;
    string dig; cin>>dig;
    mem(ava, false);
    for (char c : dig) ava [c-'0'] = 1;
    if ( x <= 1000000 ) {
         cout << solve(B) - solve(A-1) << endl;
    }else {
         LL ans = 0;
         LL cur = 0;
         while ( cur < A ) cur += x;
         while ( cur <= B ) {
             if( check(cur) ) ans++;
             cur += x;
         cout << ans << endl;
    }
}
```

1.4 DP(convex hull optimization)

```
// \texttt{http://codeforces.com/contest/311/problem/B}
struct line{
    LL slope,
                 inter;
    LL value(LL x) { return x*slope + inter; }
bool check(line x, line y, line z) {
     return (z.slope - y.slope) * (z.inter - x.inter )
              (z.slope - x.slope) * (z.inter - y.inter);
#define maxn 100005
int n, m, p;
LL \ a\,[\,maxn\,]\;,\;\; d\,[\,maxn\,]\;,\;\; dp\,[\,1\,0\,1\,]\,[\,maxn\,]\;,\;\; s\,[\,maxn\,]\;;
int main() {
     cin>\!\!> n>\!\!> m>\!\!> p;
      for(int i = 2; i \leqslant n; ++i)  {
          d[i] = getint();
          d[i] += d[i-1];
     for(int i = 1; i<=m; ++i) {
    int h; scanf("%d %lld", &h, a+i);
          a[i] -= d[h];
     sort(a+1,a+1+m);
```

```
\label{eq:formalized} \begin{array}{ll} \text{for}\,(\,int\ i \!=\! 1; i \!<\!\!=\! m;\, i \!+\! +\! )\ s\,[\,i\,] \ = \ a\,[\,i\,] \!+\! s\,[\,i\,\!-\! 1\,]\,; \end{array}
//start dp
for (int i=1; i < p; i++) {
      if(i = 1) {
            for (int j=1; j \le m; j++) dp[i][j] = j*a[j] - s
                  [j];
      }else {
           deque<line> dq;
           dq.pb(\{0, 0\});
            for (int j=1; j \le m; j++) {
                  while (dq. size() >= 2 \&\& dq[0]. value(-a[j]) > dq[1]. value(-a[j])) dq.
                        pop_front();
                  dp[i][j] = dq[0].value(-a[j]);
                  line newline{ j, dp[i-1][j]+s[j] };
                  \label{eq:while} \mbox{while} (\mbox{ dq.size}() >= 2 \mbox{ \&\& check}(\mbox{dq}[\mbox{dq}.
                        size()-2], dq.back(), newline)) dq
                        .pop_back();
                  dq.pb( newline );
                  if ( i==1 )
                       dp\,[\,i\,\,]\,[\,j\,\,] \;=\; j\,{}^*a\,[\,j\,\,] \;\; -\;\; s\,[\,j\,\,]\,;
                  }else {
                       LL mn = 0;
                        for (int \ k = 1; \ k < j; \ k+\!\!\!\!+\!\!\!\!+) \ \{
                             mn = min(mn, dp[i-1][k] + s[k]
                                    - a[j]*k);
                        dp[i][j] = mn + a[j]*j-s[j];
                        // apply convex hull optimization
                  dp[i][j] += a[j]*j - s[j];
      }
cout \ll dp[p][m] \ll endl;
```

1.5 simulated annealing

```
//尋找和所有點距離和最小的點
  #include <cstdio>
  #include <cstdlib>
  #include <cmath>
   #define F(n) Fi(i,n)
  #define Fi(i,n) for (int i=0;i< n;i++)
  #define N 1010
    using namespace std;
    int \tilde{X}[N], Y[N], n;
    in line double pow2(double x){
                      return x*x;
   double check(double x,double y){
                      double ans=0;
                      F(n) ans+=sqrt (pow2(x-X[i])+pow2(y-Y[i]));
                      return ans;
    int main(){
                       scanf("d"); while(scanf("F(n)scanf("double
                                          x{=}0, y{=}0, tx, ty, tans, l{=}10000, ans; ans{=}check(x,y); while(l{>}1e{-}check(x,y)) = 0.0000, ans; ans{=}check(x,y); while(x,y); wh
                                          tmp=rand();tx=x+l*cos(tmp);ty=y+l*sin(tmp);tans=check(tx,ty);
                                          l*=0.9;printf("//尋找兩個點使他們跟給定的四個點最小生成樹
                                          最小 include <cstdio>include <cstdlib>include
                                           <cmath>include <algorithm>define F(n) Fi(i,n)define
                                          Fi(i,n) Fl(i,0,n)define Fl(i,l,n) for(int i=l;i < n;i++)define N
                                           10using namespace std;int X[N],Y[N],n,F[N],e;struct Eint
                                          a,b;double c;G[N*2];struct Vdouble x,y;V operator+(double
                                          l)int tmp=rand();return
                                          (V)x+l*cos(tmp),y+l*sin(tmp);v[N];int find(int x)return
                                          x==F[x]?x:F[x]=find(F[x]);inline double pow2(double
                                          x)return x*x;double check(V s1,V s2)double
                                          ans=0;e=0;v[4]=s1,v[5]=s2;F(5)Fl(j,i+1,6)G[e++]=(E)i,j,sqrt(pow2)
                                          v[j].x) + pow2(v[i].y - v[j].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[j].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[j].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[j].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[j].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[j].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[j].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); \\ F(6)F[i] = i; \\ sort(G,G+e,[](E\ a,E)) + pow2(v[i].y - v[i].y)); 
                                          a.c < b.c;; F(e)if(find(G[i].a)! = find(G[i].b))ans + = G[i].c; F[find(G[i].a)]
                                          ans;int main()scanf("while(n-)F(4)scanf("double")
```

```
 \begin{array}{ll} ttans, tans, ans, l1=10000, l2; V \\ s1=(V)0, 0, s2=(V)0, 0, ts1, ts2, tmp; ans=check(s1, s2); while (l1>1e\\ 3) l2=10000; ts1=s1+l1; tans=check(ts1, s2); tmp=s2; while (l2>1e\\ 3) ts2=s2+l2; ttans=check(ts1, ts2); if (ttans< tans) tans=ttans, s2=\\ l2^*=0.9; if (tans< ans) ans=tans, s1=ts1; else\\ l1^*=0.9, s2=tmp; printf(" & in tans) tans=tans, s2=tans, s1=ts1; else & in tans, s2=tans, s2=tans, s3=tans, s3=t
```

2 Graph

2.1 HLD

```
//Greatest graph
///http://acm.csie.org/ntujudge/problemdata/2582.pdf
//this template operate on edges
#define maxn 100005
struct segment_tree{
    \#define right(x) x << 1 | 1
    \#define left(x) x << 1
    int* arr;
    int m[4*maxn];
    int tag[4*maxn];
    const int inf = 1e9;
    void init() {
         //memset(tag, -1, sizeof(tag));
         fill(tag, tag+4*maxn, inf);
    void pull(int ind) {
             m[ind] = min(m[right(ind)], m[left(ind)]);
    void push(int ind) {
         if(tag[ind] != inf) {
              tag[left(ind)] = min(tag[left(ind)], tag[
              {\rm tag}\left[\,{\rm right}\,(\,{\rm ind}\,)\,\right] \,=\, {\rm min}\left(\,{\rm tag}\left[\,{\rm right}\,(\,{\rm ind}\,)\,\right]\,,\ {\rm tag}\left[\,
                   ind]);
              m[left(ind)] = min(m[left(ind)], tag[left(
                   ind)|)
             m[right(ind)] = min( m[right(ind)], tag[
                   right(ind)]);
              tag[ind] = inf;
         }
    \frac{1}{///} root \Rightarrow 1
    void build(int ind, int l, int r) {
         if(r - l = 1) {
             m[ind] = arr[l];
              return;
         int mid = (l+r) >> 1;
         \verb|build(left(ind),l,mid);|\\
         build( right(ind), mid, r );
         pull(ind);
    int query_min(int ind, int L, int R, int ql, int qr
         if ( L >= qr \mid \mid R <= ql ) return 1e9;
         if ( R \le qr \&\& L \ge ql ) {
              return m[ind];
         push(ind);
         int mid = (L+R) >> 1;
         return min( query_min(left(ind), L, mid, ql, qr
              ), query_min(right(ind), mid, R, ql, qr));
    void modify (int ind, int L, int R, int ql, int qr,
         int x) {
         if (L >= qr \mid \mid R <= ql) return;
         if(R \le qr \&\& L \ge ql) {
              m[ind] = min(m[ind], x);
              tag\left[\,ind\,\right] \;=\; min\left(\,tag\left[\,ind\,\right]\,,\;\;x\,\right)\,;
              return;
         }
         int mid = (L+R) >> 1;
         modify(left(ind), L, mid, ql, qr, x);
```

```
modify(right(ind), mid, R, ql, qr, x);
         pull(ind);
    }
int seg_arr[maxn];
struct Tree{
    segment_tree seg;
    int n:
    struct Edge { int u, v, c; };
    vector<Edge> e;
    void addEdge(int x, int y, int c) {
        G[x].pb(SZ(e));
        G[y].pb( SZ(e) );
         e.pb(Edge\{x, y, c\})
    int siz [maxn] ,max_son[maxn] ,pa [maxn] ,dep [maxn];
/*size of subtree index of max_son, parent index index
         depth*/
    \begin{array}{ll} \verb|int| & \verb|link_top[maxn]|, \verb|link[maxn]|, \verb|timer|; \\ \end{array}
    /*chain top index in segtree time stamp*/
    std::vector<int >G[maxn];
    void init(int N) {
        n = N;
         e.clear();
         for (int i = 1; i \le n; i++) G[i]. clear ();
         timer=0;
         pa[1] = 1;
         dep[1] = 0;
    void find_max_son(int x){
         siz[x]=1;
         \max_{} [x] = -1;
         for(int e\_ind : G[x]) {
             int v = e[e\_ind].u == x ? e[e\_ind].v : e[
                  e_ind].u ;
             pa[v] = x; dep[v] = dep[x] + 1;
             find_max_son(v);
              if(max\_son[x] = -1 \mid \mid siz[v] > siz[max\_son]
                  [x])
                  \max_{son}[x] = v;
             siz[x] += siz[v];
    void build_link(int x, int top){
         link[x] = timer++;/*記錄x點的時間戳*/
         link\_top[x] = top;
         if(\max_{x} son[x] != -1)
             build_link( max_son[x], top);/*優先走訪最大
                  孩子*
         for(int e\_ind : G[x]) {
             int v = e[e\_ind].u == x ? e[e\_ind].v : e[
                  e_ind].u ;
             if (v == pa[x]) 
                  \operatorname{seg\_arr}[\operatorname{link}[x]] = \operatorname{e}[\operatorname{e\_ind}].c;
             if (v = \max_{son}[x] \mid | v = pa[x]) continue
             // edge from x \Rightarrow v
             build_link(v, v);
         }
    inline int lca(int a, int b){
         /*求LCA, 可以在過程中對區間進行處理*/
         int ta=link_top[a],tb=link_top[b];
         while (ta != tb) {
             if (dep[ta]<dep[tb]) {
                  std::swap(ta,tb);
                  std::swap(a,b);
             //interval [ link[ta], link[a] ]
             a = pa[ta];
             ta = link\_top[a];
         return dep[a] < dep[b] ? a:b;
    int modify(int a, int b, int c){
```

```
int ta=link_top[a], tb=link_top[b];
         while (ta != tb) {
             if (dep[ta]<dep[tb]) {
                  std::swap(ta,tb);
                  std::swap(a,b);
              //interval [ link[ta], link[a] ]
              //same interval if operate on edges
             seg.modify(1, 1, n, link[ta], link[a]+1, c)
             a = pa[ta];
             ta \, = \, link\_top\,[\,a\,]\,;
         //a, b are on the same chain
         if(a = b); // interval[link[a], link[a]
             ], if operate on edges => no edge
         else {
             if (dep[a]>dep[b])
                  swap(a,b);
              //interval [ link[a], link[b]
             // if operate on edges \Longrightarrow [ link[ max_son[
                  a] ], link[b] ]
              seg.modify(1, 1, n, link[max\_son[a]],
                  link[b]+1, c);
         }
    }
    void modify(int a, int b, int c) {
         if (a=b) return;
         if( link\_top[a] = link\_top[b]) {
             \begin{array}{l} if (\ dep [a] > dep [b] \ ) \ swap (a, \ b) \, ; \\ seg. \, modify (1, \ 1, \ n, \ link [a] + 1, \ link [b] + 1, \ c \end{array}
              assert(link[a]+1 = link[max_son[a]]);
             return;
         if(dep[link\_top[a]] < dep[link\_top[b]])
             swap(a, b);
         // a is the node with deeper link_top
         seg.modify(1, 1, n, link[link\_top[a]], link[a]
              +1, c);
         modify(pa[link\_top[a]], b, c);
    /// Heavy Light Decomposition
    void HLD() {
         // root is indexed 1 here !
         find_max_son(1);
         build_link(1, 1);
}tree;
int main() {
    int T; cin>>T;
    while (T--) {
         int n,m;
         scanf("%d %d",&n, &m);
         int ans = 0;
         tree.init(n);
         for (int i=0; i< n-1; i++) {
             int a, b, c;
             scanf("%d%d%d",&a,&b,&c);
              //a--, b--; be careful here
              tree.addEdge(a, b, c);
             ans += c:
         tree.HLD();
         tree.seg.arr = seg_arr;
         tree.seg.build(1, 1, n);
    return 0;
}
```

2.2 Hungarian

```
// edge and node index starting from 0
// dfs version below
//complexity O ( V*E )
/* to do
#define ___maxNodes
num_left = ?
struct Edge {
   int from;
   int to;
   int weight;
   Edge(int f, int t, int w):from(f), to(t), weight(w)
        {}
vector<int> G[__maxNodes]; /* G[i] 存储顶点 i 出发的边
   的编号 */
vector<Edge> edges;
int num_nodes;
int num_left;
int num_right;
int num_edges;
int matching[__maxNodes]; /* matching result */
int check [___maxNodes];
bool dfs(int u) {
    for (auto i = G[u].begin(); i != G[u].end(); ++i) {
        // 对 u 的每个邻接点
       int v = edges[*i].to;
       if (!check[v]) {
                          // 要求不在交替路中
           check[v] = true; // 放入交替路
           if (matching[v] == -1 || dfs(matching[v]))
               // 如果是未盖点,说明交替路为增广路,则
                   交换路径,并返回成功
               matching[v] = u;
               matching[u] = v;
               return true;
       }
   }
   return false; // 不存在增广路, 返回失败
int hungarian() {
   int ans = 0;
   memset(matching, -1, sizeof(matching));
    if (matching[u] == -1) {
           memset(check, 0, sizeof(check));
           if (dfs(u)) + +ans;
       }
   return ans;
```

2.3 KM

```
最小帶權匹配~km算法
//http://acm.csie.org/ntujudge/contest_view.php?id=836&
      contest id=449
#include <bits/stdc++.h>
using namespace std;
struct bipartite {
      #define maxn 602
      #define INF 0xfffffff
      \begin{array}{ll} \text{int} & \text{sx} \left[ \text{maxn} \right], & \text{sy} \left[ \text{maxn} \right], & \text{mat} \left[ \text{maxn} \right] \left[ \text{maxn} \right]; \end{array}
      int x[maxn], y[maxn], link[maxn];
      int N, M, slack;
      int DFS(int t) {
             int tmp;
             \operatorname{sx}[t] = 1;
             for (int i = 0; i < M; i++) {
                   if (!sy[i]) {
                         tmp \, = \, x \, [\, t \, ] \, + \, y \, [\, i \, ] \, - \, mat \, [\, t \, ] \, [\, i \, ] \, ;
                         if (tmp == 0) {
                               sy[i] = 1;
```

```
if (link[i] = -1 \mid | DFS(link[i]))
                               link[i] = t;
                               return 1;
                     else if (tmp < slack) slack = tmp;
          }
          return 0;
     int KM() {
          for (int i = 0; i < N; i++) {
               x[i] = 0;
               for (int j = 0; j < M; j++) {
                    if (mat[i][j] > x[i]) x[i] = mat[i][j];
          for (int j = 0; j < M; j++) { y[j] = 0; }
          memset(link, -1, sizeof(link));
for (int i = 0; i < N; i++) {
               while (1) {
                    memset(sx\,,\ 0\,,\ \underline{sizeof}(sx))\,;
                    memset(sy, 0, sizeof(sy));
                    slack = INF;
                     if (DFS(i)) break;
                     for (int j = 0; j < N; j++) {
                         if (sx[j]) x[j] = slack;
                     for (int j = 0; j < M; j++) {
                          if (sy[j]) y[j] += slack;
               }
          }
          int ans = 0;
          int cnt = 0;
          int t:
          \  \  \, \text{for} \  \, (\, \text{int} \  \, i \, = \, 0\,; \  \, i \, < M; \  \, i + \! + \!)
               t = link[i];
               \label{eq:if_t} \begin{array}{ll} i\,f & (\,t\,>=\,0\,\,\&\&\,\,\mathrm{mat}\,[\,t\,]\,[\,i\,] \end{array} \stackrel{!}{=} \begin{array}{ll} -\mathrm{INF}) \end{array}
                    cnt ++;
                    ans += mat[t][i];
          // 最大權 : 沒有負號
          return -ans;
     void init(int n, int m) {
          N = n, M = m;
          for (int i = 0; i < N; i++)
               for (int j = 0; j < M; j++)
                    mat[i][j] = -INF;
     void input() {
          for (int i = 0; i < N; i++)
               for (int j = 0; j \triangleleft M; j++) {
                    // fill in mat[i][j]
                    // stands for the weighting , but
                          negative sign !
                    // if 最大權 : 沒有負號
               }
}km;
int main(){
     int n,E;
     while (scanf("%d", &n) != EOF)
          km.init(n, n);
          km.input();
          cout << km.KM() << endl;
     return 0;
}
```

2.4 Bi-vertex-connected Subgraph

```
#include <bits/stdc++.h>
using namespace std;
#ifdef DEBUG
    #define debug(...) printf(__VA_ARGS__)
   #define debug(...) (void)0
#endif
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair<long long, long long>
#define fi first
#define se second
\#define all(x) (x).begin(),(x).end()
#define SZ(x) ((int)(x).size())
\begin{array}{ll} \textbf{const} & \textbf{inf} & \textbf{inf} = 0\,x\,7\,ffffffff\,; \ //beware \ overflow \end{array}
const LL INF = 0x7ffffffffffffffffff; //beware overflow
\#define mem(x, y) memset(x, (y), sizeof(x));
#define IOS ios_base::sync_with_stdio(0); cin.tie(0)
template<typename A, typename B>
ostream& operator <<(ostream &s, const pair<A,B> &p) {
    return s<<"("<<p.first<<","<<p.second<<")";
template < typename T >
ostream& operator <<(ostream &s, const vector<T> &c) {
    s << "[ ";
    for (auto it : c) s << it << " ";
    s << "]";
    return s;
template<typename T>
ostream& operator << (ostream &o, const set<T> &st) {
    o << "{";
    for (auto it=st.begin(); it!=st.end(); it++) o << (
    it==st.begin() ? "" : ", ") << *it;</pre>
    return o << "}";
template<typename T1, typename T2>
ostream& operator << (ostream &o, const map<T1, T2> &mp
    ) {
    o << "{";
    for (auto it=mp.begin(); it!=mp.end(); it++) {
         o << (it=mp.begin()?"":", ") << it->fi << ":"
             << it->se;
    o << "}";
    return o;
}
      regard every vbcc as a set of edges
/** needed for tarjan **/
#define maxn 100005
#define maxm 100005
int n, m;
struct Edge{int s, t;};
vector<Edge> edge;
int dfn[maxn], low[maxn];
stack<int> st;
bool vis[maxn];
int Time;
bool vis_e [maxm];
int bcnt, vbb [maxm];
vector<int> vb[maxm];
vector < int > G[maxn];
void tarjan(int s){
    dfn[s] = low[s] = ++Time;
    vis[s] = true;
    for (int e_ind : G[s])
         if (!vis_e[e_ind]) {
              vis_e[e_ind] = true; st.push(e_ind);
              int to = edge[e\_ind].s + edge[e\_ind].t - s;
              if (! vis [to]) {
                  tarjan(to);
                  low[s] = min(low[s], low[to]);
                  if(low[to] >= dfn[s])
                       vb[bcnt].clear();
```

```
while (1) {
                         int t = st.top(); st.pop();
                         vbb[t] = bcnt;
                         vb[bcnt].push_back(t);
                         if(t == e_ind) break;
                     bcnt++;
                 }
            } else
                 low[s] = min(low[s], dfn[to]);
        }
    }
void init_tarjan() {
    mem(vis, false); mem(vis_e, false);
    Time = bcnt = 0; edge.clear();
    for (int i = 1; i \le n; i++) G[i]. clear ();
}
int main() {
    cin >> n >> m;
    init_tarjan();
    for (int i = 0; i < m; i++) {
         int a, b; scanf("%d %d", &a, &b);
        edge.push_back(Edge{a,b});
        G[a].push\_back((int)edge.size()-1);
        G[b].push\_back((int)edge.size()-1);
    tarjan(1);
}
```

2.5 Bi-edge-connected Subgraph

```
/** needed for tarjan **/
#define maxn 100005
\#define maxm 100005
int n, m;
int dfn[maxn], low[maxn];
stack<int> st;
int Time;
int bcnt;
vector<int> G[maxn];
bool in_cyc[maxn];
void tarjan(int s, int p){
    dfn[s] = low[s] = ++Time;
    st.push(s);
    for(int to : G[s]) if( to != p ){
        if (!dfn[to]) {
            tarjan(to, s);
            low[s] = min(low[s], low[to]);
            if(low[to] > dfn[s]) {
                 // is cut_edge
                 // pop stack 的過程也可以寫在這
                 // 但最後(after tarjan)還要多判stack
                     not empty的情况
                 if ( low [to] > dfn [s]) {
                 in\_cyc[bcnt] = st.top()!=to;
                 while (1) {
                     int t = st.top(); st.pop();
                     id \, [\, t \, ] \; = \; bcnt \, ;
                     if(t == to) break;
                 bcnt++;
            }
        }else
            low[s] = min(low[s], dfn[to]);
    }
```

```
if(low[s] = dfn[s]){
        in\_cyc[bcnt] = st.top()!=s;
        while (1) {
            int^{\dagger}t^{\dagger} = st.top(); st.pop();
            id[t] = bcnt;
            if (t == s) break;
        bcnt++;
    }
void init_tarjan() {
    Time = bcnt = 0;
int main() {
  cin >> n >> m;
  init_tarjan();
  G[a].pb(b), G[b].pb(a);
 mem( in_cyc , false);
  tarjan(1, 1);
}
```

2.6 SCC

```
#include <bits/stdc++.h>
using namespace std;
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair < long long, long long>
#define fi first
#define se second
const int inf = 1e9;
const LL INF = 1e18;
const int mod = 1e9 + 7;
#define maxn 100050
int n, m;
vector < int > g[maxn];
stack<int> Stack;
int scnt, Time;
int belong[maxn], dfn[maxn], low[maxn], indegree[maxn];
bool instack [maxn];
void input(){
  \texttt{cin} >> \texttt{n} >> \texttt{m};
  for (int i = 0; i < m; i++){
    int a, b; scanf("%d%d", &a, &b);
    g[a].pb(b);
void init() {
  \mathtt{scnt} \, = \, \mathsf{Time} \, = \, 0;
  for (int i = 1; i \le n; i++)
    g[i].clear();
  while (!Stack.empty()) Stack.pop();
  memset(indegree, 0, sizeof(indegree));
  memset(dfn, 0, sizeof(dfn));
memset(instack, false, sizeof(instack));
void dfs(int u)
  dfn[u] = low[u] = ++Time;
  Stack.push(u); instack[u] = true;
  for(int v : g[u]) {
    if (!dfn[v]) {
      dfs(v);
      low[u] = min(low[u], low[v]);
    else if (instack [v])
      low[u] = min(low[u], dfn[v]);
  if(low[u] = dfn[u]) {
    scnt++;
    int tp;
    do{
```

```
tp = Stack.top(); Stack.pop();
      instack[tp] = false;
      belong[tp] = scnt;
    } while (tp != u);
void tarjan() {
  for (int i = 1; i \le n; i++)
    if (!dfn[i])
      dfs(i);
int main(){
  int T; cin >> T;
  while (T--) {
    init();
    input();
    tarjan();
    for (int i = 1; i \le n; i++) {
      for(int v : g[i]) {
  if(belong[v] != belong[i])
           indegree [belong [v]]++;
      }
    LL ans = 0;
    for (int i = 1; i \le scnt; i++)
      if (!indegree[i]) ans++;
    cout << ans << endl;
  }
  return 0;
```

2.7 Steiner Tree (PECaveros)

```
// Minimum Steiner Tree
// O(V 3^T + V^2 2^T)
struct SteinerTree{
#define V 33
#define T 8
#define INF 1023456789
 int n, dst[V][V], dp[1 << T][V], tdst[V]
 void init( int _n ){
   n = \underline{n};
   for ( int i = 0 ; i < n ; i ++ ){
     for (int j = 0 ; j < n ; j ++)
     dst[i][j] = INF;

dst[i][i] = 0;
   }
 }
 void add_edge( int ui , int vi , int wi ){
   void shortest_path(){
   for (int k = 0 ; k < n ; k +++)
     for (int i = 0 ; i < n ; i ++)
      for(int j = 0 ; j < n ; j ++)
        ][ j ] );
 int solve( const vector<int>& ter ){
   int t = (int) ter.size();
   for ( int i = 0 ; i < (1 << t) ; i ++ )
     for(int j = 0 ; j < n ; j ++)
   dp[0][i] = 0;
   for ( int msk = 1 ; msk < (1 << t) ; msk ++ ){
     \begin{array}{ll} if (\ msk == (\ msk \ \& \ (-msk)\ )\ ) \{ \end{array}
       int who = __lg(msk);
       continue;
     for ( int i = 0 ; i < n ; i ++ )
       for ( int submsk = ( msk - 1 ) & msk ; submsk ;
              submsk = (submsk - 1) & msk)
          dp[msk][i] = min(dp[msk][i],
```

2.8 Edmond's Matching Algorithm

2.9 Tree Decomposition

```
//codeforces Digit Tree
 //http://codeforces.com/problemset/problem/715/C
#include <bits/stdc++.h>
using namespace std;
#ifdef DEBUG
         #define debug(...) printf(__VA_ARGS__)
#else
         #define debug(...) (void)0
#endif
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair < long long, long long>
#define fi first
#define se second
\#define all(x) (x).begin(),(x).end()
\#define SZ(x) ((int)(x).size())
const int inf = 0x7ffffffff; //beware overflow
\#define mem(x, y) memset(x, (y), sizeof(x));
#define IOS ios_base::sync_with_stdio(0); cin.tie(0)
template<typename A, typename B>
ostream& operator <<(ostream &s, const pair<A,B> &p) {
    return s<<"("<<p.first<<","<<p.second<<")";
template<typename T>
ostream& operator <<(ostream &s, const vector<T> &c) {
           s << "[ ";
           for (auto it : c) s << it << " ";
           s << `"]";
           return s;
template<typename T>
ostream& operator << (ostream &o, const set<T> &st) {
           o << "{";
            \begin{picture}(10,0) \put(0,0){\line(0,0){150}} \put(0,0){\line(0,0){15
                      it=st.begin() ? "" : ", ") << *it;
           return o << "}";
template<typename T1, typename T2>
ostream& operator << (ostream &o, const map<T1, T2> &mp
           ) {
          o << "{";
           for (auto it=mp.begin(); it!=mp.end(); it++) {
                     o << (it=mp.begin()?"":", ") << it->fi << ":"
                               << it->se;
           o << "}";
          return o;
typedef long long 11;
bool isprime [100005];
vector<LL> primes;
```

```
LL M, PHI;
#define MOD M
ll modpow(ll a, ll b) {
 11 r = 1;
  while(b) {
   if(b\&1) r=(r*a)\%MOD;
    a=(a*a)%MOD;
   b >>= 1;
 }
  return r;
void Sieve(int n) {
  memset(isprime, 1, sizeof(isprime));
  isprime[1] = false;
  for (int i = 2; i \le n; i++) {
    if(isprime[i]) {
      primes.pb(i);
for(int j = 2*i; j <= n; j += i)
        isprime[j] = false;
 }
}
LL phi(LL n) {
  11 \text{ num} = 1; 11 \text{ num} 2 = n;
  for(ll i = 0; primes[i]*primes[i] \le n; i++) {
    if(n\%primes[i]==0) {
      num2/=primes[i];
      num*=(primes[i]-1);
    while (n\%primes[i]==0) {
      n/=primes [ i ];
    }
  if (n>1) {
   num2/=n; num*=(n-1);
 n = 1;
 num \ ^*= \ num2;
  return num;
ll inv(ll a) {
  return modpow(a, PHI-1);
\#define maxn 100005
struct edge{
   int u, v, dig;
    int no(int x) {
        return x = u ? v : u;
vector<edge> e;
vector < int > G[maxn];
LL n, ans;
bool vis[maxn];
int sz[maxn], dep[maxn];
LL tenPow[maxn];
int dfs(int u, int p, int d) {
    sz[u] = 1;
    dep[u] = d;
    for(int eind : G[u])  {
        int v = e[eind].no(u);
        if (v = p \mid | vis[v]) continue;
        sz[u] += dfs(v, u, d+1);
    return sz[u];
int findCenter(int u, int p, int treesize) {
    for (int eind : G[u] )
        int v = e[eind].no(u);
        if (v = p || vis[v]) continue;
if (sz[v]*2 > treesize)
             return findCenter( v, u, treesize);
    return u;
}
LL up[maxn], down[maxn];
int belong[maxn];
map<LL, LL> tot;
vector< map<LL, LL> > vec;
vector<int> pt;
```

```
void calc(int u, int p, int b, int d) {
     pt.pb( u );
     belong[u] = b;
     dep[u] = d;
     int \ id = \ find\_if(\ all(G[u]) \ , [u,p](int \ x) \ \{ \ return
           e[x].no(u) = p; }) - G[u].begin();
     down[u] = (down[p]*10 + e[G[u][id]].dig)M;
     up[u] = (tenPow[d-1]*e[G[u][id]].dig + up[p])
          %M;
     for(int eind : G[u]) {
          int v = e[eind].no(u);
          if (vis[v] || v = p) continue;
          calc(v, u, b, d+1);
     vec[b][ up[u] ]++;
     tot[ up[u] ]++;
LL solve(int cent) {
     //cent is the root now
     vector<int> L;
     \quad \quad \text{for} \, (\, \text{int eind } : \, G[\, \text{cent} \,] \,) \  \, \{ \,
          int v = e[eind].no(cent);
          if (!vis[v]) {
              L.pb( v );
     vec.clear();
     vec.resize(SZ(L), {});
     tot.clear();
     up[cent] = down[cent] = 0;
     dep[cent] = 0;
     pt.clear();
     for (int i = 0; i < SZ(L); i++)
          calc ( L[i], cent, i, 1);
     LL ret = 0;
     for(int u : pt) {
          LL tmp = (-down[u]+M)\%M;
         tmp = (tmp*inv(tenPow[dep[u]]))\%M;
          ret += tot[tmp] - vec[belong[u]][tmp];
     assert( (LL)count_if(all(pt), [] (int x) { return
     \begin{array}{c} \operatorname{up}[x] = 0; \; \} \; ) = \operatorname{tot}[0]); \\ \operatorname{LL} \; \operatorname{tmp} = \operatorname{tot}[0] + (\operatorname{LL})\operatorname{count\_if}(\operatorname{all}(\operatorname{pt}), \; [] \; (\operatorname{int} \; x) \end{array}
          \{ \text{ return down}[x] = 0; \} );
     debug("\%lld \n", tmp);
     return ret+tmp;
void solveAll(int node) {
     dfs(node, -1, 0);
int cent = findCenter(node, -1, sz[node]);
     ans += solve(cent);
     \label{eq:cont_norm} \texttt{debug("\%d\ \%lld\n",\ cent,\ ans);}
     vis[cent] = true;
     for(int eind : G[cent] ) {
          int v = e[eind].no(cent);
          if (vis[v]) continue;
          solveAll(v);
     }
int main() {
     cin>>n>>M;
  Sieve( 100000 );
     PHI = phi(M);
     e.pb( edge{a, b, c});
     //init
     tenPow[0] = 1;
     for (int i = 1; i < maxn; i++) tenPow[i] = (tenPow[i
          -1]*10)%M;
     ans = 0;
    mem( vis, false);
     solveAll(0);
     cout << ans << endl;
```

|}

2.10 Tree Longest Path

```
/** codeforces 592D - Super M **/
#include <bits/stdc++.h>
using namespace std;
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair < long long, long long>
#define fi first
#define se second
const int inf = 1e9;
const LL INF = 1e18;
const int mod = 1e9+7;
#define maxn 123460
int n, m;
vector<int> g[maxn];
bool is [maxn];
int dep[maxn], R, max_depth, A;
int cnt[maxn], parent[maxn];
bool dfs(int u, int par = 0){
  parent[u] = par;
  dep[u] = dep[par] + 1;
  if(dep[u] > max_depth \&\& is[u])
   \max_{depth} = dep[u], R = u;
  bool ret = is[u];
  for(int v : g[u])
    if(v != par)
      ret = dfs(v, u);
  if(ret) A++;
  return ret;
int find_center(int start) {
 R = start; dep[0] = -1; max_depth = 0;
  dfs(start);
  \max_{\text{depth}} = 0; \text{dep}[R] = -1;
  dfs(R, R);
  int ret = R, d = max_depth/2;
  while(d>0)
    d--:
    ret = parent[ret];
  return ret;
int S, dis, max_length;
bool dfs1(int u, int par = 0) {
  dep[u] = dep[par] + 1;
  if ( is [u])
    if (dep[u] > max_length)
      \max_{\text{length}} = \text{dep}[u], S = u;
    else if (dep[u] = max_length \&\& u < S)
      S = u:
  bool c = false;
  for(int v : g[u])
    if(v!=par)
      dfs1(v, u);
int main(){
  cin >> n >> m;
  for (int i = 0; i < n-1; i++){
    int a, b; scanf("%d%d",&a, &b);
    g[a].pb(b), g[b].pb(a);
  memset(is, false, sizeof(is));
  int tmp;
  for (int i = 0; i < m; i++){
    cin>>tmp; is [tmp] = true;
  int C = find_center(tmp);
  dep[0] = -1;S = inf; dis = (max_depth+1)/2;
```

```
// distance(center, any other node) <= (longestpath + 1) / 2 
 dfs1(C); 
 if ( max_depth & 1) 
  dfs1(parent[C]); 
 cout << S << endl << A-2-max_depth << endl; 
 return 0; 
}
```

3 Flow

3.1 Dinic Maxflow

```
//http://acm.csie.org/ntujudge/problem.php?id=2581
//French Fries Festival
//dinic runs in O( V^2*E )
#include <bits/stdc++.h>
using namespace std;
#ifdef DEBUG
    #define debug(...) printf(__VA_ARGS__)
    #define debug(...) (void)0
#endif
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair<long long, long long>
#define fi first
#define se second
\#define \ all(x) \ (x).begin(),(x).end()
#define SZ(x) ((int)(x).size())
const int inf = 0x7ffffffff; //beware overflow
const LL INF = 0x7ffffffffffffffffff; //beware overflow
\#define mem(x, y) memset(x, (y), sizeof(x));
#define IOS ios_base::sync_with_stdio(0); cin.tie(0)
template<typename A, typename B>
ostream& operator <<(ostream &s, const pair<A,B> &p) {
    return s<<"("<p.first<<","<p.second<")";
template < typename T >
ostream& operator <<(ostream &s, const vector<T> &c) {
    s << "[ ";
    for (auto it : c) s << it << " ";
    s << "]";
    return s;
template < typename T >
ostream& operator << (ostream &o, const set<T> &st) {
    o << "{";
    for (auto it=st.begin(); it!=st.end(); it++) o << (
    it==st.begin() ? "" : ", ") << *it;</pre>
    return o << "}";
template<typename T1, typename T2>
ostream& operator << (ostream &o, const map<T1, T2> &mp
    ) {
    o << "{";
    for (auto it=mp.begin(); it!=mp.end(); it++) {
   o << (it=mp.begin()?"":", ") << it->fi << ":"</pre>
             << it->se;
    o << "}";
    return o;
#define maxn 500
struct Edge{ int to, cap, rev; };
struct Dinic{
    vector < Edge > G[maxn];
    int dis[maxn], iter[maxn];
    void init(int n) {
         //zero based
       for (int i = 0; i < n; i++) G[i]. clear ();
     void addEdge(int from, int to, int cap) {
         vector < Edge > :: iterator it;
         if( ( it=find_if( all(G[from]), [to](Edge& e) {
               return e.to == to; } )) != G[from].end() )
```

```
it -> cap += cap;
              return;
       G[from].pb(Edge\{to, cap, (int)G[to].size()\});
       G[to].pb(Edge\{from, 0, (int)G[from].size()-1\});
         //if undirected 0 will be cap
     bool bfs(int s, int t) {
       memset(dis, -1, sizeof(dis));
       queue<int> que;
       que.push(s); dis[s] = 0;
       while(!que.empty()) {
         int tp = que.front(); que.pop();
         for (Edge &e : G[tp]) {
           if(e.cap > 0 \&\& dis[e.to] = -1)
              dis[e.to] = dis[tp] + 1, que.push(e.to);
       return dis[t] != -1;
     int dfs(int v, int t, int f) {
       if( v == t ) return f;
       for(int \& i = iter[v]; i < G[v].size(); i++) {
         Edge &e = G[v][i];
         if(e.cap > 0 \&\& dis[v] < dis[e.to]) {
           int d = dfs(e.to, t, min(f, e.cap));
            if(d > 0) {
              e \cdot cap -= d;
             G[e.to][e.rev].cap += d;
              f += d;
              return d;
         }
       }
       return 0;
     int maxFlow(int s, int t) {
       int ret = 0;
       while(bfs(s, t)) 
         memset(iter, 0, sizeof(iter));
         while ((f = dfs(s, t, inf)) > 0)
           ret += f;
       return ret;
}dinic, dinic2;
void solve() {
     int n,m,k; cin>>n>>m>>k;
     // flow problem with lower bounds;
     int s = 0, t = n+2, ss = n+3, tt = n+4;
     \label{dinic.init} \mbox{dinic.init} ( \ \ \mbox{n+}5 \ \ ) \ ;
     dinic.addEdge(s, 1, k);
     dinic.addEdge(n+1, t, k);
     int slb = 0;
     while (m--) {
         int l, r, a, b; scanf("%d %d %d %d", &l, &r, &a
              , &b);
         slb += a;
         r++;
         dinic.addEdge(l, r, b-a);
         dinic.addEdge(ss, r, a);
         dinic.addEdge(l, tt, a);
     dinic2 = dinic;
     dinic.addEdge(t, s, k);
     int f1 = dinic.maxFlow(ss, tt);
     if ( \ !all\_of ( \ all ( \ dinic.G[ \ ss \ ] ) \ , \ [ \ ] ( \ Edge \ x ) \ \{ \ return
         x.cap = 0;  ) ) {
         puts("-1"); return;
     }
     \label{eq:dinic2} {\tt dinic2.addEdge(ss\,,\,\,s\,,\,\,1e9)}\,;
     dinic2.addEdge(t, tt, 1e9);
     int f2 = dinic2.maxFlow(ss, tt);
     // maxflow in current graph is f2 - slb
     printf("%d\n", (f2 - slb)*n);
}
```

```
int main() {
    int t; cin>>t;
    while(t--)
        solve();
}
```

4 Data Structure

4.1 Disjoint Set

```
struct Disjoint_set {
     #define MAX_N 500005
     // define MAX_N
     \begin{array}{ll} \verb"int" & pa \left[ \verb"MAX_N" \right] \;, \;\; Rank \left[ \verb"MAX_N" \right] \;; \end{array}
     int sz [MAX_N];
     void init_union_find(int V) {
          for(int i=0; i<\!\!V; i++) {
               pa[i] = i;
               Rank[i] = 0;
               sz[i] = 1;
     int find(int x) {
           return x == pa[x] ? x : pa[x] = find(pa[x]); 
     int unite(int x, int y) {
          x = find(x), y = find(y);
          int S = sz[x]+sz[y];
          if(x != y){
               if(Rank[x] < Rank[y]) {
                    pa\left[\,x\,\right] \;=\; y\,;
                    sz[y]=S;
                    return y;
               }
               else{
                    pa[y] = x;
                    sz[x] = S;
                    if(Rank[x] = Rank[y]) Rank[x] ++;
                    return x;
          }
     bool same(int x, int y) {
          return find(x) = find(y);
}
```

4.2 Djs + Seg

```
/ demo ⇒ undo djs + segtree with offline
// this program doesn't consider the problem of
    overflowing varaible ans
   http://acm.csie.org/ntujudge/view_code.php?id
    =\!108190\&\mathtt{contest\_id}\!=\!472
#include <bits/stdc++.h>
using namespace std;
#ifdef DEBUG
    #define debug(...) printf(__VA_ARGS__)
#else
   #define debug(...) (void)0
#endif
\#define mp make\_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair < long long, long long>
#define fi first
#define se second
\#define all(x) (x).begin(),(x).end()
\#define SZ(x) ((int)(x).size())
const int inf = 0x7ffffffff; //beware overflow
\#define mem(x, y) memset(x, (y), sizeof(x));
#define IOS ios_base::sync_with_stdio(0); cin.tie(0)
template<typename A, typename B>
ostream& operator <<(ostream &s, const pair<A,B> &p) {
```

```
return s<<"("<<p.first<<","<<p.second<<")";</pre>
template<typename T>
ostream& operator <<(ostream &s, const vector<T> &c) {    s << "[";
     for (auto it : c) s << it << " ";
     s << "]";
     return s;
template<typename T>
ostream& operator << (ostream &o, const set<T> &st) {
    o << "{";
     for (auto it=st.begin(); it!=st.end(); it++) o << ( it==st.begin() ? "" : ", ") << *it;
     return o << "}";
template<typename T1, typename T2>
ostream& operator << (ostream &o, const map<T1, T2> &mp
    ) {
    o << "{";
     for (auto it=mp.begin(); it!=mp.end(); it++) {
         o << (it=mp. begin()?"":", ") << it->fi << ":"
               << it->se;
    o << "}";
     return o;
#define maxn 100005
#define maxm 500005
//can be used to solve dynamic connectivity problem
//can be used with segment tree \Longrightarrow offline
struct DisjointSet {
  // save() is like recursive
// undo() is like return
  int n, fa[maxn], sz[maxn];
vector<pair<int*,int>>> h;
  vector<int> sp;
  int ans;
  void init(int tn) {
    ans = 0;
    n=tn:
     for (int i=0; i<n; i++) {
       fa[i]=i;
       sz\;[\;i\;]\!=\!1;
     }
    sp.clear(); h.clear();
  void assign(int *k, int v) {
    h.pb(\{k, *k\});
     *k=v;
  void save() { sp.pb(SZ(h)); }
  void undo() {
     assert(!sp.empty());
     int last=sp.back(); sp.pop\_back();
     while (SZ(h)!=last) {
       \color{red} \textbf{auto} \  \, \textbf{x=} \\ \textbf{h.back()} \  \, ; \  \, \textbf{h.pop\_back()} \  \, ;
        *x.fi=x.se;
    }
  }
  int f(int x) {
     while (fa[x]!=x) = fa[x];
     return x;
  void uni(int x, int y) {
    x=f(x); y=f(y);
     if (x=y) return ;
     if (sz[x] < sz[y]) swap(x, y);
     //nans stands for new answer
     int t = sz[x]+sz[y];
     int nans = ans - (sz[x]*sz[x]-sz[x]) - (sz[y]*sz[y]) - (sz[y]*sz[y]) - (sz[y])*sz[y]
     \operatorname{assign}(\&\operatorname{sz}\left[x\right],\ \operatorname{sz}\left[x\right]\!+\!\operatorname{sz}\left[y\right]);
     assign(\&fa[y], x);
     assign(&ans, nans);
\}djs;
int n, m;
\max < int, int > \max[maxn];
vector<pii> seg[4*maxm];
LL ans [maxm];
```

```
void add(int ql, int qr, int a, int b, int id=1, int l
     =0, int r=m) {
     if (qr \ll l | | ql \gg r) return;
     if(l) = ql \&\& r <= qr)
          seg[id].pb(mp(a, b));
         return ;
     int mid = (l+r) >> 1;
    add( ql, qr, a, b, id*2+1, mid, r);
void dfs(int u=1, int l=0, int r=m) {
     djs.save();
     for(pii v : seg[u] ) djs.uni( v.fi , v.se );
     \begin{array}{cccc} i\,f\,(&r\,\text{-}\,l \ > \ l \ ) & \{\\ &i\,n\,t \ \ mid \ = \ (\,l+r\,) >> 1; \end{array}
          dfs(u*2, 1, mid);
          dfs\left(u^{*}2{+}1,\ mid\,,\ r\,\right);
         // do sth here
         ans[l] = djs.ans;
     djs.undo();
int main() {
     scanf("%d %d", &n, &m);
     \begin{array}{lll} & for(int & i = 0; & i < m; & i++) \; \{ \\ & int & a, & b; & scanf(``\%d \;\%d",\&a, \;\&b); \end{array}
         a--, b--; if (b < a) swap(a, b);
          if( ma[a].count(b) ) {
              add(ma[a][b], i, a, b);
              ma[a].erase(b);
          else ma[a][b] = i;
     for(int i = 0; i < n; i++) if( !ma[i].empty()) 
          for(auto p : ma[i])
              add( p.se, m, i, p.fi);
     djs.init(n);
     dfs();
     for (int i = 0; i < m; i++) printf ("%lld\n", ans [i]);
```

4.3 Sparse Table

```
//codeforces 689D
#define maxn 200005
template < typename T, typename Cmp = less < T > >
struct RMQ {
      T d[maxn][20];
      Cmp cmp;
      int w[maxn], sz;
      void init (const T *a, int n) {
            int i, j;
             for (w[0] = -1, i = 1; i \le n; ++i) w[i] = (i & 
                     (i - 1)) ? w[i - 1] : w[i - 1] + 1;
              \mbox{for } (sz = n, \ i = 0; \ i < n; \ +\!\!+\! i) \ d[\,i\,][\,0\,] = a[\,i\,]; 
             for (j = 1; (1 << j) <= n; ++j) {
                   for (i = 0; i + (1 \ll j) \ll n; ++i) {
                         \begin{array}{l} d[\,i\,][\,j\,] = cmp(d[\,i\,][\,j\,-\,1]\,,\; d[\,i\,+\,(1<<\\ (\,j\,-\,1)\,)\,][\,j\,-\,1]\,)\;?\; d[\,i\,][\,j\,-\,1]\;:\; d\\ [\,i\,+\,(1<<(\,j\,-\,1)\,)\,][\,j\,-\,1]; \end{array}
                   }
            }
      // index of a [l .. r]
      const T &query(int l, int r) const {
             int x = w[r - l + 1];
             \begin{array}{c} \textbf{return} \ \text{cmp}(d[1][x]\,, \ d[r - (1 <\!\!< x) + 1][x]) \ ? \ d \\ [1][x] \ : \ d[r - (1 <\!\!< x) + 1][x]; \end{array} 
int a[maxn], b[maxn];
```

```
int n:
RMQ < int > s;
RMQ<int, greater<int>> t;
int main() {
      cin>>n;
      \begin{array}{lll} & \text{for} (int & i = 0; i < n; i++) \; scanf(\text{``%d''}, \&a[i]); \\ & \text{for} (int & i = 0; i < n; i++) \; scanf(\text{``%d''}, \&b[i]); \end{array}
      s.init(b, n);
      t.init(a, n);
      int c, d;
     LL ans = 0;
      for (int i=0; i< n; i++) {
            if(a[i] > b[i]) continue;
           int ub = n+1, lb = i;
            while (ub-lb>1) {
                 int mid = (ub+lb)>>1;
                 if(t.query(i, mid-1) - s.query(i, mid-1) >
                        0) ub = mid;
                 else lb = mid;
           int up = ub;
           ub = n+1, lb = i;
             while (ub-lb>1) {
                 int mid = (ub+lb)>>1;
                 \begin{array}{ll} \mbox{if( t.query(i, mid-1) - s.query(i, mid-1)} \\ >= 0) \mbox{ ub = mid;} \end{array}
                 else lb = mid;
           int down = ub;
           ans += up-down;
      cout << ans << endl;
      return 0;
```

4.4 Treap

```
#include <bits/stdc++.h>
using namespace std;
struct Treap{
    Treap *1, *r;
    int pri, key, val;
    {\rm Treap}\left(\begin{smallmatrix} int & \_val \,, & int & \_key \end{smallmatrix}\right):
        val(\_val), key(\_key), l(NULL), r(NULL), pri(
            rand()){}
};
/// We assure that key value in A treap is greater than
     that in treap B
Treap *merge( Treap *a, Treap *b){
    if (a->pri > b->pri){
        a->r = merge(a->r, b);
        return a;
    }else{
        b->l = merge(a, b->l);
        return b;
void split (Treap *t, int k, Treap *&a, Treap *&b) {
    if (!t) a = b = NULL;
    else if (t->key \ll k)
        a = t;
        split(t->r, k, a->r, b);
    }else{
        split(t->l, k, a, b->l);
Treap* insert( Treap *t, int k, int _val){
    Treap *tl, *tr;
    split(t, k, tl, tr);
    return merge(tl, merge(new Treap(_val, k), tr));
```

```
}
Treap* remove( Treap* t, int k){
    Treap *tl, *tr;
    split(t, k-1, tl, t);
    split(t, k, t, tr);
    return merge(tl, tr);
}
int main(){
    return 0;
}
```

5 Math

5.1 Prime Table

```
#include <bits/stdc++.h>
using namespace std;
struct Prime_table {
    int prime [1000000] = \{2,3,5,7\};
    int sz=4:
    // biggest prime < ub
    int ub=(1<<20);
    int check(int num){
        int k = 0;
        for (k = 0; k < sz \&\& prime[k]*prime[k] <= num;
             k++){}
             if ( num \% prime [k]==0) return 0;
        return 1;
    void buildprime(){
        int currentPrime=7;
        int j=4;
        for (sz=4, j=4; currentPrime < ub; sz++, j=6-j){
              currentPrime=currentPrime+j;
              if (check(currentPrime)) {
                 prime[sz] = currentPrime;
              }
              else{
                 SZ - -;
        }
}ptable;
```

5.2 Miller Rabin Prime Test

```
#include <bits/stdc++.h>
using namespace std;
typedef long long LL;
LL mul(LL a, LL b, const LL mod) {
     LL x = 0, y = a \% \text{ mod};
     while (b > 0) {
          if (b&1)
          x = (x + y) \% \text{ mod};

y = (y * 2) \% \text{ mod};
          b >>= 1;
     return x % mod;
LL mul(LL lhs, LL rhs, const LL mod) {
     return (lhs * rhs ) % mod;
LL\ mypow(LL\ b\,,\ LL\ e\,, {\color{red}const}\ LL\ mod)\ \{
     LL x = 1;
     LL y = b;
     while (e) {
          if ( e\&1 ) x = mul(x, y, mod);
          y = \operatorname{mul}(y\,,\ y\,,\ \operatorname{mod})\,;
          e >>= 1;
```

```
return x;
const int testbase [] = \{2, 3, 5, 7, 11, 13, 17, 19, 23,
     29, 31, 37};
bool isprime(const LL p) {
    if (p < 2) return false;
    if (p != 2 && !(p&1) ) return false;
   LL d = p - 1;
    while (!(d\&1)) d >>= 1;
    for( int a : testbase ) {
        LL td = d;
        if ( a \ge p-1 ) return true;
        LL st = mypow(a, td, p);
        while (td!=p-1 && st!=1 & st!=p-1)
            st = mul(st, st, p);
            td \ll 1;
        if ( st != p - 1 && !(td&1) ) return false;
    return true;
int main() {
    int T;
    scanf("%d",&T);
    while (T--) {
        LL q;
        scanf("%lld",&q);
puts(isprime(q)?"YES":"NO");
    return 0;
}
```

5.3 Extended Euclidean Algorithm

```
/** normal gcd function using recursion **/
int gcd(int a, int b){
    if(b == 0) return a;
    return gcd(b, a%b);
}
// Find solution of ax + by = gcd(a, b)
// ps : x, y may be negative
int extgcd(int a, int b, int& x, int& y){
    int d = a;
    if(b!= 0) {
        d = extgcd(b, a%b, y, x);
        y -= (a/b) * x;
    }else {
        x = 1, y = 0;
    }
    return d;
}
```

5.4 Gauss Elimination

```
// solving linear equations with gauss elimination
#include <iostream>
#include <cmath>
#include <vector>

using namespace std;

void print(vector< vector<double> > A) {
   int n = A.size();
   for (int i=0; i<n; i++) {
      for (int j=0; j<n+1; j++) {
        cout << A[i][j] << "\t";
        if (j = n-1) {
        cout << "|";
      }
    }
   cout << "\n";
   }
   cout << endl;
}</pre>
```

```
vector<double> gauss (vector< vector<double> > A) {
     int n = A. size();
     for (int i=0; i<n; i++) {
          // Search for maximum in this column
          double maxEl = abs(A[i][i]);
          int \max Row = i;
          for (int k=i+1; k< n; k++)
              if (abs(A[k][i]) > maxEl) {
                   maxEl = abs(A[k][i]);
                   \max Row = k;
              }
          }
          // Swap maximum row with current row (column by
               column)
          for (int k=i; k<n+1;k++)
              double tmp = A[maxRow][k];
              A[\max Row][k] = A[i][k];
              A[i][k] = tmp;
          // Make all rows below this one 0 in current
              column
          for (int k=i+1; k<n; k++) {
              \begin{array}{ll} \textbf{double} \ c \, = \, -A[\,k\,]\,[\,i\,]\,/A[\,i\,]\,[\,i\,]\,; \end{array}
              for (int j=i; j< n+1; j++) {
                   if (i==j) {
                       A[k][j] = 0;
                   } else {
                       A[k][j] += c * A[i][j];
              }
          }
     // Solve equation Ax=b for an upper triangular
          matrix A
     vector <\!\! double\!\! > x(n)\,;
     for (int i=n-1; i>=0; i--) {
         x[i] = A[i][n]/A[i][i];
          for (int k=i-1; k>=0; k--)
              A[k][n] -= A[k][i] * x[i];
     return x;
 int main() {
     int n;
     cin >> n;
     vector < double > line(n+1,0);
     vector< vector<double> > A(n, line);
     // Read input data
     for (int i=0; i< n; i++) {
          for (int j=0; j<n; j++) {
              cin >> A[i][j];
     }
     for (int i=0; i< n; i++) {
          cin \gg A[i][n];
     // Print input
     print(A);
     // Calculate solution
     vector < double > x(n);
     x = gauss(A);
     // Print result
cout << "Result:\t";</pre>
     for (int i=0; i<n; i++) {
         cout <\!\!< x[\,i\,] <\!\!< "
     cout << endl;
}
```

typedef long double ld;

N must be 2^k and greater than array.size()

5.5 FFT

* FFT(a);

* FFT(b);

```
* for (int i = 0; i < N; ++i) c[i] = conj(a[i] * b[i]);
 * FFT( c );
 * for (int i = 0; i < N; ++i) c[i] = conj(c[i]);
 * for (int i = 0; i < N; ++i) c[i] /= N;
void FFT(vector< complex<ld>>& v) {
     int \dot{N} = v.size();
                                                                                            [i + j];
     P1 = P2 - t;
         if(i>j) swap(v[i],v[j]);
                                                                                       P2 = P2 + t;
                                                                                   }
     for (int k = 2; k \le N; k < =1) {
                                                                              }
         1d w = -2.0* pi/k;
                                                                          }
         complex < ld > deg(cos(w), sin(w));
         for (int j = 0; j < N; j + = k) {
              complex < ld > theta(1,0);
              for (int i = j; i < j+k/2; ++i) {
                  complex < ld > a = v[i];
                  complex < ld > b = v[i+k/2]*theta;
                  v\,[\;i\;] \;=\; a\!\!+\!\!b\,;
                                                                              convolution % Mod
                  v[i+k/2] = (a-b);
                                                                          // stored in res[ 0 .. K+1 )
                  theta *= deg;
         }
    }
}
//http://sd-invol.github.io/2016/02/13/FFT-mod-prime/
                                                                              int j = i ? len - i : i;
struct Complex {
     double x , y;
     Complex \ (\begin{array}{c} \hbox{\tt double} \ \underline{\hspace{0.1cm}} x = 0 \end{array} , \ \begin{array}{c} \hbox{\tt double} \ \underline{\hspace{0.1cm}} y = 0) \end{array} \{
                                                                                  (0.5, 0);
         x = _x , y = _y;
                                                                                   (0, -0.5);
     Complex operator + (const Complex &r) const {
         return Complex(x + r.x , y + r.y);
                                                                                   (0.5, 0);
     Complex operator - (const Complex &r) const {
                                                                                   (0, -0.5);
         return Complex(x - r.x , y - r.y);
     Complex operator * (const Complex &r) const {
         return Complex(x * r.x - y * r.y , x * r.y + y
              * r.x);
     Complex conj () const {
         return Complex(x , -y);
     double operator = (const double a) {
         *this = Complex(a, 0);
         return a;
                                                                                    + z) \% Mod;
const double pi = acos(-1.0);
                                                                          }
                                                                     }
//fft with modulo, code referenced from the internet
                                                                };
     fftMod::fftPrepare(len);
     fftMod::convolution(res, le, ri, len, r-1);
                                                                 5.6 NNT
namespace fftMod{
     const int N = 1 \ll 18;
     const int Mod = 1e9 + 7;
                                                                 NTT( a );
     // to do, M should be about sqrt(Mod)
                                                                 NTT( b );
     const int M = 32768;
     for (int i = 0; i < N: ++i)
                                                                 NTT( c, true );
                                                                 for (int i = 0; i < N; ++i)
    Complex w[N];
                                                                     c\,[\,i\,]\ =\ (786433LL\text{--}12)\ \ *\ \ c\,[\,i\,]\ \ \%\ \ mod\,;
     int rev[N];
     int LN = \__builtin\_ctz(n);
                                                                 constexpr int mod = 786433;
         for (int i = 0; i < n; ++ i) {
    double ang = 2 * pi * i / n;
                                                                 constexpr int N = 65536;
              w[i] = Complex(cos(ang), sin(ang));
              rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (
                  LN - 1));
                                                                     for (int i = 1, j = 0; i < N; ++i)
```

```
}
void FFT(Complex P[], int n, int oper) {
    for (int i = 0; i < n; i ++) {
    if (i < rev[i]) {
              swap(P[i], P[rev[i]]);
     for (int d = 0; (1 \ll d) < n; d++) {
          for (int i = 0; i < n; i += m2) {
              for (int j = 0; j < m; j++) {
                   Complex &P1 = P[i + j + m], &P2 = P
                   Complex \ t \ = \ w[rm \ * \ j \ ] \ * \ P1;
\begin{array}{lll} Complex~A[N]~,~B[N]~,~C1[N]~,~C2[N]\,;\\ void~convolution \big(vector{<}int{>}~\&res~,~vector{<}int{>}~\&a~, \end{array}
     vector<int> &b, int len, int K) {
    // a[ 0 .. len ) and b[ 0 .. len )'s
    for (int i = 0; i < len; +++i) {
    A[i] = Complex(a[i] / M, a[i] % M);
    B[i] = Complex(b[i] / M, b[i] % M);
    FFT(A, len, 1); FFT(B, len, 1);
     for (int i = 0; i < len ; ++ i) {
         Complex a1 = (A[i] + A[j].conj()) * Complex
         Complex a2 = (A[i] - A[j].conj()) * Complex
         Complex b1 = (B[i] + B[j].conj()) * Complex
         Complex b2 = (B[i] - B[j].conj()) * Complex
         Complex c11 = a1 * b1 , c12 = a1 * b2;
         Complex c21 = a2 * b1 , c22 = a2 * b2;
         C1[j] = c11 + c12 * Complex(0, 1);
         C2[j] = c21 + c22 * Complex(0, 1);
    \overline{FFT}(C1 \ , \ len \ , \ -1) \, ; \ \overline{FFT}(C2 \ , \ len \ , \ -1) \, ;
     int y2 = (LL)(C2[i].x / len + 0.5) \% Mod;

int z = (LL)(C2[i].y / len + 0.5) \% Mod;

res[i] = ((LL)x * M * M + (LL)(y1 + y2) * M
```

```
c\,[\,i\,] \,=\, (\,long\ long\,)\ \dot{a}\,[\,i\,]\ *\ b\,[\,i\,]\ \%\ mod\,;
void NTT(vector< int >& v, bool flag = false)
```

```
for (int k = N > 1; !((j^=k)&k); k > = 1);
         if(i>j) swap(v[i],v[j]);
    for (int k = 2; k \le N; k < = 1)
         \label{eq:int_deg} \begin{array}{lll} \mbox{int} & \mbox{deg} \, = \, \mbox{mypow}(\, \mbox{flag} \, \ ? \, \, 524289 \, : \, \, 3 \, , \, \, \mbox{N} \, \, / \, \, \, k) \, ; \end{array}
         for(int j = 0; j < N; j + = k)
              int theta = 1;
              for (int i = j; i < j+k/2; ++i)
              {
                   \quad \quad \text{int} \ a \, = \, v \, [\, i \, ] \, ; \quad
                   int b = (long long) v[i+k/2]*theta%mod;

v[i] = (a+b) \% mod;
                   v[i+k/2] = (a-b+mod)\%mod;
                   theta = (long long) theta * deg % mod;
         }
    }
constexpr int mod = 1e9+7;
typedef vector<int> VEC;
// ntt + Crt, code referenced from the internet
namespace nttCrt {
    constexpr int magic[3] = \{1004535809, 998244353,
         104857601};
    constexpr int MOD = 1000000007;
    constexpr int G = 3;
    int P:
    inline int quick_mod(int x, int k, int MOD) {
         int ans = 1;
         while (k) {
              if (k\&1) ans = 1LL * ans * x % MOD;
              x = 1LL * x * x % MOD;
              k >>= 1:
         return ans;
     inline void change(int *y, int len) {
         for(int i = 1, j = len / 2; i < len - 1; i++) {
              i\,f\,(\,i\,<\,j\,)\ swap\,(\,y\,[\,i\,]\,,\ y\,[\,j\,]\,)\;;
              //交换互为小标反转的元素, i<j保证交换一次
              //i做正常的+1, j左反转类型的+1,始终保持i和j
                   是反转的
              int k = len / 2;
              while(j >= k)  {
                   j -= k;
                   k /= 2;
              if(j < k) j += k;
         }
     inline void ntt(int *y, int len, int on) {
         change(y, len);
         for (int h = 2; h \le len; h \le 1)
              int wn = quick\mod(G, (P - 1) / h, P);
              for(int j = 0; j < len; j += h) {
                   int w = 1;
                   y[k] = (u + t) \% P;
                        y[k + h / 2] = ((u - t) \% P + P) \%
                        w = 1LL * w * wn \% P;
                   }
              }
         if(on = -1) {
              for (int i = 1; i < len / 2; i++)
                  swap(y[i], y[len - i]);
              int inv = quick_mod(len, P - 2, P);
for(int i = 0; i < len; i++)
    y[i] = 1LL * y[i] * inv % P;</pre>
         }
    }
    int n;
    int r[3][3];
```

```
inline int CRT(int *a) {
    int sb[3] = \{a[0], a[1], a[2]\};
    for (int i = 0; i < 3; ++i) {
         for (int j = 0; j < i; ++j) {
              int t = (sb[i] - sb[j]) % magic[i];
              }
    int mul = 1, ans = sb[0] \% MOD;
    for (int i = 1; i < 3; ++i) {
    mul = 1LL * mul * magic[i - 1] % MOD;
         ans = (ans + 1LL * sb[i] * mul) % MOD;
    return ans;
int tmp[maxn][3];
int x1[maxn*2], x2[maxn*2];
inline void gao(vector<int>& res , vector<int> &a ,
    vector<int> &b ,int len , int kk) {
    for (int ti = 0; ti < 3; ti++) {
         P = magic [ti];
         int k;
         for (k = 0; k < SZ(a) & k < len; k++) x1[
              k] = a[k];
         \quad \text{for } (\ ; k < \ len\, ; \ k+\!\!\!\! +) \ x1[\,k\,] \ = \ 0\, ;
         for (k = 0; k < SZ(b) & k < len; k++) x2[
              k] = b[k];
         for (; k < len; k++) x2[k] = 0;
         ntt(x1, len, 1); ntt(x2, len, 1);
         for (int i = 0; i < len; i++) x1[i] = 1LL *
               x1\,[\,i\,]\ *\ x2\,[\,i\,]\ \%\ P;
         ntt(x1, len, -1);
         \  \  \, \hbox{for}\  \, (\,\hbox{int}\  \, i\,=\,0\,;\  \, i\,<=\,kk\,;\  \, i+\!+\!)\,\,tmp\,[\,i\,\,]\,[\,\,t\,i\,\,]\,=\,
              x1[i];
    for (int i = 0; i \le kk; i++) res[i] = CRT(tmp
         [ i ])
inline void init() {
    for (int i = 0; i < 3; i++) {
         for (int j = 0; j < 3; j++) {
              r[i][j] = quick\_mod(magic[i], magic[j])
                   - 2, magic[j]);
    }
}
```

5.7 Big Number

};

```
//http://blog.csdn.net/hackbuteer1/article/details
    /6595881
#include<iostream>
#include<string>
#include<iomanip>
#include < algorithm >
using namespace std;
#define MAXN 9999
#define MAXSIZE 10
#define DLEN 4
class BigNum
private:
 int a[500];
                //可以控制大数的位数
 int len;
                //大数长度
public:
 BigNum() \{ len = 1; memset(a,0,sizeof(a)); \}
                                             //构造函
  BigNum(const int);
                          //将一个int类型的变量转化为
      大数
```

```
BigNum(const char*);
                        //将一个字符串类型的变量转化
     为大数
 BigNum(const BigNum &); //拷贝构造函数
 BigNum & operator=(const BigNum &); //重载赋值运算
     符,大数之间进行赋值运算
 friend istream& operator>>(istream&, BigNum&);
     重载输入运算符
 friend ostream& operator<<(ostream&, BigNum&);</pre>
     重载输出运算符
 BigNum operator+(const BigNum &) const;
                                       //重载加法
     运算符,两个大数之间的相加运算
 BigNum operator-(const BigNum &) const;
                                       //重载减法
     运算符,两个大数之间的相减运算
 BigNum operator*(const BigNum &) const;
                                       //重载乘法
     运算符,两个大数之间的相乘运算
 BigNum operator/(const int
                                       //重载除法
                           &) const;
     运算符,大数对一个整数进行相除运算
 BigNum operator (const int &) const;
                                      //大数的n次
     方运算
       operator%(const int &) const;
                                      //大数对一个
     int类型的变量进行取模运算
      operator > (const BigNum & T) const;
                                        //大数和另
     一个大数的大小比较
                                        //大数和一
 bool operator > (const int & t) const;
     个int类型的变量的大小比较
                   //输出大数
 void print();
BigNum::BigNum(const int b)
                          //将一个int类型的变量转
   化为大数
 int c, d = b;
 len = 0:
 memset(a, 0, sizeof(a));
 while (d > MAXN)
   c = d - (d / (MAXN + 1)) * (MAXN + 1);
   d = d / (MAXN + 1);
   a[len++] = c;
 a[len++] = d;
                             //将一个字符串类型的变
BigNum::BigNum(const char*s)
   量转化为大数
 int t,k,index,l,i;
 memset(a, 0, sizeof(a));
 l=strlen(s);
 len=l/DLEN;
 if (1%DLEN)
   len++;
 index=0:
 for (i=l-1; i>=0; i=DLEN)
 {
   t=0:
   k=i -DLEN+1;
   if(k<0)
     k=0;
   for (int j=k; j<=i; j++)
     t=t*10+s[j]-'0';
   a[index++]=t;
BigNum::BigNum(const BigNum & T): len(T.len) //拷贝构
   造函数
 int i:
 memset(a, 0, sizeof(a));
 for(i = 0 ; i < len ; i++)
   a[i] = T.a[i];
BigNum & BigNum::operator=(const BigNum & n)
                                         //重载赋
   值运算符, 大数之间进行赋值运算
 int i;
```

```
memset(a,0,sizeof(a));
  for(i = 0 ; i < len ; i++)
    a[i] = n.a[i];
  return *this;
istream& operator>>(istream & in, BigNum & b)
    输入运算符
  char ch [MAXSIZE*4];
  int i = -1;
  in>>ch;
  int l=strlen(ch);
  int count=0,sum=0;
  for (i=l-1; i>=0;)
    sum = 0;
    int t=1;
    for (int j=0; j<4\&\&i>=0; j++,i--,t*=10)
      sum + = (ch[i] - '0') * t;
    b.a[count]=sum;
    {\rm count}{++};
  b.len =count++;
  return in;
ostream& operator << (ostream& out, BigNum& b)
                                                 //重载
    输出运算符
  int i;
  cout \ll b.a[b.len - 1];
  for(i = b.len - 2 ; i >= 0 ; i--)
    cout.width(DLEN);
    cout. fill('0');
    cout << b.a[i];
  return out;
}
BigNum BigNum::operator+(const BigNum & T) const
    个大数之间的相加运算
  BigNum t(*this);
                  //位数
  int i, big;
  big = T.len > len ? T.len : len;
  for(i = 0 ; i < big ; i++)
    t.a[i] +=T.a[i];
    if(t.a[i] > MAXN)
      t.a[i + 1]++;
      t . a [i] = MAXN+1;
  if(t.a[big] != 0)
    t.len = big + 1;
  else
    t.len = big;
  return t;
BigNum BigNum::operator-(const BigNum & T) const
    个大数之间的相减运算
  int i,j,big;
  bool flag;
  BigNum t1, t2;
  if (*this>T)
    t1=*this;
    t2=T;
    f \log = 0;
  else
    t1=T;
    t2 = *this;
```

len = n.len:

```
flag=1:
                                                                           d = ((d * (MAXN+1))\% b + a[i])\% b;
                                                                        }
  big=t1.len;
  for(i = 0 ; i < big ; i++)
                                                                        return d:
    if(t1.a[i] < t2.a[i])</pre>
                                                                      BigNum BigNum::operator^(const int & n) const
                                                                                                                                 //大数
                                                                           的n次方运算
       j = i + 1;
       \frac{\text{while}(t1.a[j]}{} = 0)
                                                                        BigNum t, ret(1);
                                                                        int i:
       t1.a[j--]--;
                                                                         if (n<0)
       while (j > i)
                                                                           exit(-1);
        t1.a[j--] += MAXN;
                                                                         if(n==0)
       t1.a[i] += MAXN + 1 - t2.a[i];
                                                                           return 1;
    }
                                                                         if(n==1)
    else
                                                                           return *this;
       t1.a[i] -= t2.a[i];
                                                                         int m=n;
                                                                        while (m>1)
  t1.len = big;
  while (t1.a[len - 1] = 0 \&\& t1.len > 1)
                                                                           t=*this;
                                                                           for (i=1; i << 1 <= m; i << =1)
    t1.len - -;
    big - -;
                                                                             t=t*t;
                                                                           }
  if (flag)
                                                                          m-i;
    t1.a[big-1]=0-t1.a[big-1];
                                                                           ret=ret*t;
                                                                           i f (m==1)
                                                                             ret=ret*(*this);
                                                             //两
BigNum BigNum::operator*(const BigNum & T) const
     个大数之间的相乘运算
                                                                      bool BigNum::operator>(const BigNum & T) const
  BigNum ret;
                                                                           和另一个大数的大小比较
  {\color{red}int} \quad i\ , j\ , up\,;
  int temp, temp1;
                                                                        int ln;
  for (i = 0 ; i < len ; i++)
                                                                        if (len > T.len)
                                                                          return true;
    up = 0;
                                                                         else if (len = T.len)
    for (j = 0 ; j < T.len ; j++)
                                                                           ln = len - 1:
       temp = a[i] * T.a[j] + ret.a[i + j] + up;
                                                                           \mathbf{while}(\mathbf{a}[\ln] = \mathbf{T}.\mathbf{a}[\ln] \&\& \ln >= 0)
       if(temp > MAXN)
                                                                             ln - -;
                                                                           if(ln >= 0 \&\& a[ln] > T.a[ln])
         \begin{array}{l} temp1 \, = \, temp \, \, - \, \, temp \, \, / \, \, \left( MAXN \, + \, \, 1 \right) \, \, * \, \, \left( MAXN \, + \, \, 1 \right) \, ; \\ up \, = \, \, temp \, \, / \, \, \left( MAXN \, + \, \, 1 \right) \, ; \end{array}
                                                                             return true;
                                                                           else
         ret.a[i + j] = temp1;
                                                                             return false;
       }
                                                                        }
       else
                                                                        else
       {
                                                                           return false;
         up = 0;
                                                                      }
         ret.a[i + j] = temp;
                                                                      bool BigNum::operator >(const int & t) const
                                                                                                                                //大数
                                                                           和一个int类型的变量的大小比较
    if(up != 0)
                                                                        BigNum b(t);
       ret.a[i + j] = up;
                                                                        return *this>b;
  ret.len = i + j;
  while (ret.a[ret.len - 1] == 0 && ret.len > 1)
                                                                                                  //输出大数
                                                                      void BigNum::print()
    ret.len--;
  return ret;
                                                                        cout << a[len - 1];
BigNum BigNum::operator/(const int & b) const
                                                         //大数
                                                                         for (i = len - 2 ; i >= 0 ; i--)
    对一个整数进行相除运算
                                                                           cout.width(DLEN);
 BigNum ret;
                                                                           cout. fill('0');
  int i, down = 0;
                                                                           cout <\!\!< a[i];
  for(i = len - 1 ; i >= 0 ; i--)
                                                                        cout << endl;
    ret.a[i] = (a[i] + down * (MAXN + 1)) / b;
    down \, = \, a \, [\, i \, ] \, + \, down \, * \, (MAXN \, + \, 1) \, - \, ret \, . \, a \, [\, i \, ] \, * \, b \, ;
                                                                      int main(void)
  ret.len = len:
                                                                        int i,n;
  while (\text{ret.a} [\text{ret.len} - 1] = 0 \&\& \text{ret.len} > 1)
                                                                                                //定义大数的对象数组
                                                                        BigNum x[101];
    ret.len - -;
                                                                        x[0]=1;
  return ret;
                                                                        for (i=1; i<101; i++)
                                                                        x[i]=x[i-1]*(4*i-2)/(i+1);
while (scanf("%d",&n)==1 && n!=-1) {
                                                        //大数对
int BigNum::operator %(const int & b) const
      一个int类型的变量进行取模运算
                                                                          x[n].print();
                                                                        }
  int i.d=0:
                                                                     }
  for (i = len -1; i >= 0; i --)
```

6 string

6.1 Palindromic Tree

```
回文自動機包含以下元素:
    狀態St, 所有節點的集合, 一開始兩個節點, 0表示偶數長
         度串的根和1表示奇數長度串的根
    last 新增加一個字符後所形成的最長回文串的節點編號
    s 當前的字符串(一開始設s[0]=-1(可以是任意一個在串S
         中不會出現的字符))
    n 表示添加的字符個數
每個節點代表一個不同的回文子字串,我們在每個節點會儲存
    一些數值:
    len 表示所代表的回文子字串長度
    next[c] 表示所代表的回文子字串在頭尾各增加一個字符c
         後的回文字串其節點編號
    sufflink 表示所代表的回文子字串不包括本身的最長後綴
         回文子串的節點編號
    cnt(非必要) 表示所代表的回文子字串在整體字串出現的
         次數(在建構完成後呼叫count()才能計算)
    //num(非必要) 表示所代表的回文子字串其後綴為回文字
         串的個數 <── not included
struct palindromic_tree{
    struct node{
        int next[26], sufflink, len; /*這些是必要的元素*/int l, r; // this node is s[\ l\ ..\ r\ ]
                                   /*這些是額外維護的元素*/
        int cnt, num;
        \mathsf{node}(\hspace{.05cm} \mathsf{int} \hspace{.1cm} \hspace{.1cm} l\hspace{-.05cm}=\hspace{-.05cm} 0)\hspace{.05cm} :\hspace{.05cm} \mathsf{sufflink}\hspace{.05cm} (\hspace{.05cm} 0)\hspace{.1cm} \hspace{.1cm} ,\hspace{-.05cm} \mathsf{len}\hspace{.05cm} (\hspace{.05cm} 1)\hspace{.1cm} \hspace{.1cm} ,\hspace{-.05cm} \mathsf{cnt}\hspace{.05cm} (\hspace{.05cm} 0)\hspace{.1cm} \hspace{.1cm} ,\hspace{-.05cm} \mathsf{num}\hspace{.05cm} (\hspace{.05cm} 0)\hspace{.1cm} \}
             for (int i=0; i<26;++i) next [i]=0;
    };
    std::vector<node> St;
    std::string s; //current string [ 1 .. n ]
    int last ,n;
    palindromic\_tree():St(2), last(1), n(0)
        St[0].sufflink=1;
        St[1].len=-1;
        s.push\_back(-1);
    inline void clear(){
        St.clear();
        s.clear();
        last=1;
        n=0:
        St.push_back(0);
        St.push back(-1);
        St[0].sufflink=1;
        s.push_back(-1);
    inline int get_sufflink(int x){
        while (s[n-St[x].len-1] != s[n]) x=St[x].
             sufflink;
        return x;
    inline void add(int c){
        s.push_back(c-= 'a',);
         int cur=get_sufflink(last);
         if (!St[cur].next[c]) {
             int now=St.size();
             St.push_back(St[cur].len+2);
             St[now].sufflink=St[get_sufflink(St[cur].
                  sufflink)].next[c];
              /*不用擔心會找到空節點,由證明的過程可知*/
             St[cur].next[c]=now;
             St[now].num=St[St[now].sufflink].num+1;
             St[now] \cdot l = n - St[now] \cdot len + 1, St[now] \cdot r
                 = n;
         last=St[cur].next[c];
        ++St[last].cnt;
    }
```

```
inline void count(){/*cnt必須要在構造完後呼叫count
()去計算*/
std::vector<node>::reverse_iterator i=St.rbegin
();
for(;i!=St.rend();++i) {
    St[i->sufflink].cnt+=i->cnt;
    }
}
inline int size(){/*傳回其不同的回文子串個數*/
    return St.size()-2;
}
```

6.2 Suffix Array

6.3 Longest Palindromic Substring

```
//ntu judge Earse
#include <bits/stdc++.h>
using namespace std;
//#define DEBUG
#ifdef DEBUG
    #define debug(...) printf(__VA_ARGS__)
    #define debug(...) (void)0
#endif
#define mp make_pair
#define pb push_back
#define LL long long
#define pii pair<int,int>
#define PII pair < long long, long long>
#define fi first
#define se second
\#define all(x) (x).begin(),(x).end()
#define SZ(x) ((int)(x).size())
const int inf = 0x7ffffffff; //beware overflow
const LL mod = 1e9+7;
#define IOS ios_base::sync_with_stdio(0); cin.tie(0)
template<typename A, typename B>
ostream& operator <<(ostream &s, const pair<A,B> &p) {
    return s<<"("<<p.first<<","<<p.second<<")";
template < typename T >
ostream& operator <<(ostream &s, const vector<T> &c) {
    s << "[ ";
    for (auto it : c) s << it << " ";
    s << `"]";
    return s;
template<typename T>
ostream& operator << (ostream &o, const set<T> &st) {
    o << "{";
    return o << "}";</pre>
template<typename T1, typename T2>
ostream& operator << (ostream &o, const map<T1, T2> &mp
    o << "{";
    for (auto it=mp.begin(); it!=mp.end(); it++) {
       o << (it=mp.begin()?"":", ") << it->fi << ":"
            << it->se;
    o << "}";
    return o;
#define maxn 200001
char t[maxn];
char s [maxn * 2];
int z [maxn * 2];
int N;
int longest_palindromic_substring() {
    // t穿插特殊字元, 存放到s。
    int n = strlen(t);
```

```
\begin{array}{lll} N = n & * & 2 & + & 1 \, ; \\ memset (\, s \, , & \cdot \, , & \cdot \, , & N) \, ; \end{array}
    for (int i=0; i< n; ++i) s[i*2+1] = t[i];
    s[N] = ' \setminus 0';
    z[0] = 1; // if 無須使用, then 無須計算。
    int L = 0, R = 0;
    for (int i=1; i<N; ++i) // 從z[1] 開始
         z[i] = (R > i) ? min(z[2*L-i], R-i) : 1;
         while (i - z[i]) = 0 \&\& i + z[i] < N \&\&
                 s[i-z[i]] = s[i+z[i]]) z[i]++;
         if (i+z[i] > R) L = i, R = i+z[i];
    }
    // 尋找最長迴文子字串的長度
    n = 0;
    int p = 0;
    for (int i=1; i≪N; ++i) // 從z[1]開始
         if (z[i] > n)
             n = z[p = i];
    // longest 從中心到外端的長度 => (n-2)/2
    //cout << "最長迴文子字串的長度是" << (2*n-1) / 2;
    // 印出最長迴文子字串, 記得別印特殊字元。
         for (int i=p-z[p]+1; i \le p+z[p]-1; ++i)
              if (i & 1) {
                  cout \ll s[i];
    return (2*n-1)/2;
int nxt[maxn * 2];
int main() {
    int T; cin>>T;
    while (T--) {
         scanf("%s", t);
         #ifdef DEBUG
              {\tt cout} <\!\!< {\tt longest\_palindromic\_substring}\,() <\!\!<
         #else
             longest_palindromic_substring();
         \begin{array}{lll} \text{memset(nxt, -1, sizeof(nxt));} \\ \text{for(int } i = 0; i < N; i++) \end{array} 
              nxt[i-z[i]+1] = i+1;
         int leftmost = 0;
         for (int i = 0; i < N; i++) {
              leftmost = max(leftmost, nxt[i]);
              nxt[i] = max(leftmost, nxt[i]);
         int ans = 0;
         for(int cur = 0; cur<N-1;) {
              cur = nxt[cur];
         cout << ans << endl;
    return 0;
```

geometry

7.1 Point Class

```
const double eps = 1e-10;
#define N 100
struct P {
         \quad \quad \text{double } x\,,\ y\,;
        P(\  \, \text{double } \  \, \underline{\hspace{0.1cm}} x{=}0, \  \, \text{double } \  \, \underline{\hspace{0.1cm}} y{=}0) \  \, :x(\underline{\hspace{0.1cm}} x) \, , \  \, y(\underline{\hspace{0.1cm}} y) \  \, \{\};
         void read() {
                  scanf("%lf%lf",&x,&y);
         void print() {
```

```
\texttt{printf("\%f \%f \n", x, y);}
} p[N];
bool operator <( Pa, Pb) { return tie(a.x,a.y)<tie(b
    .x,b.y);
 P \text{ operator } + (Pa, Pb) \{ \text{ return } P\{a.x+b.x,a.y+b.y\}; \} 
P operator -( Pa, Pb) { return P{a.x-b.x,a.y-b.y}; }
P \ operator \ *( \ P \ b, \ double \ a \ ) \ \{ \ return \ P\{a*b.x,a*b.y\};
P operator /( P a, double b ) { return P{a.x/b,a.y/b};
P& operator /=( P &a, double b ) { return a=a/b; }
double operator *( Pa, Pb) { return a.x*b.x+a.y*b.y;
double operator ^( Pa, Pb) { return a.x*b.y-a.y*b.x;
double x( P o, P a, P b ) { return (a-o)^(b-o); }
double dot( P o, P a, P b ) { return (a-o)*(b-o); }
```

Intersection of Circles/Lines/Segments

```
//PECaveros
vector <\!\!P\!\!> interCircle (\ P\ o1\ ,\ double\ r1\ ,\ P\ o2\ ,\ double
        r2 ){
    double d2 = (01 - 02) * (01 - 02);
   double d = sqrt(d2);
    if (d > r1 + r2) return \{\};
   P u = (o1+o2)*0.5 + (o1-o2)*((r2*r2-r1*r1)/(2*d2));
   \begin{array}{ll} \textbf{double} \ A = \ sqrt \, ((\,r1 + r2 + d)\,*(\,r1 - r2 + d)\,*(\,r1 + r2 - d)\,*(\,-\,r1 + r2 \\ \end{array}
   P \ v = P(\ o1.y-o2.y \ , \ -o1.x + o2.x \ ) * A / (2*d2);
   return \{u+v, u-v\};
| P interPnt( P p1, P p2, P q1, P q2){
| double f1 = ( p2 - p1 ) ^ ( q1 - p
                                      (q1 - p1);
    double f2 = (p2 - p1) ^ (p1 - q2);
   double f = (f1 + f2);
if(fabs(f) < eps) return Pt(nan(""), nan(""));</pre>
```

return q1 * (f2 / f) + q2 * (f1 / f);

```
int ori( const PLL& o , const PLL& a , const PLL& b ){ LL ret = ( a - o ) ^ ( b - o );
              return ret / max( 111 , abs( ret ) );
    // p1 == p2 || q1 == q2 need to be handled
 bool banana<br/>( {\tt const} PLL& p1 , {\tt const} PLL& p2
                                                                                       const PLL& q1 , const PLL& q2 ){
                          \begin{array}{l} f(\ (\ p2\ -\ p1\ )\ (\ q2\ -\ q1\ )\ ) == 0\ )\{\ //\ parallel\ if(\ ori(\ p1\ ,\ p2\ ,\ q1\ )\ )\ return\ false; \\ return\ (\ (\ p1\ -\ q1\ )\ *\ (\ p2\ -\ q1\ )\ ) <= 0\ ||\ (\ p1\ -\ q2\ )\ *\ (\ p2\ -\ q2\ )\ ) <= 0\ ||\ (\ p1\ -\ q1\ )\ *\ (\ p1\ -\ q2\ )\ *\ (\ p2\ -\ q2\ )\ ) <= 0\ ||\ (\ p1\ -\ q1\ )\ *\ (\ p2\ -\ q2\ )\ ) <= 0\ ||\ (\ p1\ -\ q1\ )\ *\ (\ p1\ -\ q1\ )
                if( ( p2 - p1 ) ^
                                                                            ((q1 - p1) * (q2 - p1)) <= 0
                                                                            ((q1 - p2) * (q2 - p2)) <= 0;
```

return (ori(p1, p2, q1) * ori(p1, p2, q2)<=0) && (ori(q1, q2, p1) * ori(q1, q2, p2)<=0);

7.3 Convex Hull

```
#define REP(i,n) for ( int i=0; i< int(n); i++)
int n;
void input() {
   scanf("%d",&n);
   REP(i,n) p[i].read();
}
P findCenter() {
   p[n]=p[0];
   P center=P\{0,0\};
   REP(i,n) {
       double v=p[i]*p[i+1];
       center.y += (p[i].y+p[i+1].y)*v;
```

```
double area=0;
    REP(i,n) area+=p[i]*p[i+1];
    area \neq 2;
    center /= 6*area;
    return center;
P q1[N], q2[N], q[N];
void convex() {
    sort(p,p+n);
    int m1=0,m2=0;
    REP(\,i\,\,,n\,)\  \  \{
        while (m1>=2 \&\& X(q1[m1-2],q1[m1-1],p[i]) >= 0
              ) m1--;
         while (m2)=2 \&\& X(q2[m2-2],q2[m2-1],p[i]) <= 0
              ) m2--;
        q1 [m1++]=q2 [m2++]=p[i];
    int m=0;
    REP(i, m1) q[m++]=q1[i];
    for ( int i=m2-2; i>=1; i-- ) q[m++]=q2[i];
    q[m]=q[0];
void solve() {
    convex();
    // continue ...
```

7.4 Half Plane Intersection

```
//http://acm.csie.org/ntujudge/problemdata/2575.pdf
//http://www.csie.ntnu.edu.tw/~u91029/Half-
    planeIntersection.html
預先使用四個半平面, 設定一個極大的正方形邊界, 讓半平面
    交集擁有邊界。
 1、逐一加入每個半平面,求出當下的半平面交集(凸多邊
online 演算法,隨時維護一個半平面交集。每次更新需時 O(N
    ),總時間複雜度為 O(N<sup>2</sup>), N 是半平面數目。
#include <bits/stdc++.h>
using namespace std;
#define mp make_pair
typedef complex<double> Point;
typedef vector<Point> Polygon;
typedef pair<Point, Point> Line;
#define x real()
#define y imag()
// 兩向量叉積
double cross (Point& a, Point& b) {
    return a.x * b.y - a.y * b.x;
// 向量oa與向量ob進行叉積
double cross(Point& o, Point& a, Point& b) {
    return (a.x-o.x) * (b.y-o.y) - (a.y-o.y) * (b.x-o.x)
        );
}
// 多邊形面積
double area (Polygon& p) {
    double a = 0;
    int n = p.size();
    for (int i=0; i< n; ++i)
        a \; +\!\!= \; c \, ross \, (\, p \, [\, i \, ] \, , \; \, p \, [\, (\, i \, +\! 1)\!\%\! n \, ] \, ) \; ;
    return fabs(a) / 2;
}
// 兩線交點
Point intersection (Point& a1, Point& a2, Point& b1,
    Point& b2) {
    Point a = a2 - a1, b = b2 - b1, s = b1 - a1;
    return a1 + a * cross(b, s) / cross(b, a);
```

```
/ 一個凸多邊形與一個半平面的交集
Polygon halfplane_intersection(Polygon& p, Line& line)
          Polygon q;
          Point p1 = line.first, p2 = line.second;
          // 依序窮舉凸多邊形所有點,判斷是否在半平面上。
          // 如果凸多邊形與半平面分界線有相交, 就求交點。
          _{\hbox{\scriptsize int}}\ n=p.\,\hbox{\scriptsize size}\,(\,)\,;
          for (int i=0; i<n; ++i)
                    double c = cross(p1, p2, p[i]);
                    \begin{array}{lll} \textbf{double} \ d = \, cross \, (p1 \, , \ p2 \, , \ p \, [\, (\, i \, + 1) \% n \, ] \, ) \, ; \end{array}
                    if (c >= 0) q.push_back(p[i]);
                    if (c * d < 0) q.push_back(intersection(p1, p2,
                               p[i], p[(i+1)\%n]);
          return q;
#define maxn 550
 //Line line[maxn];
Point v[maxn];
double ans[maxn];
int main() {
         int T; cin>>T;
          while (T--) {
                    int n;
                    double
                                     w, h;
                    scanf("%d %lf %lf", &n, &w, &h);
                    // 預先設定一個極大的正方形邊界
                    Polygon p, org;
                    /** initialize
                     p.push_back(Point(-1e9,-1e9));
                    p.push_back(Point(-1e9,+1e9));
                   p.push_back(Point(+1e9,-1e9));
                    p.push_back(Point(+1e9,+1e9));
                   p.push\_back(Point(0,0));
                   p.push\_back(Point(0,h));
                   p.push_back(Point(w,h));
                   p.push\_back(Point(w,0));
                    org = p;
                    for (int i =0; i < n; i ++) {
                             double a, b;
                             scanf("%lf %lf", &a, &b);
                             v[i] = Point(a, b);
                    // 每一個半平面都與目前的半平面交集求交集
                    for (int i=0; i<n; ++i)
                             p = org;
                             for (int j = 0; j < n; j++) {
                                       if(i==j) continue;
                                       Line line;
                                        // find perpendicular line to line i_j
                                       Point a( (v[i].x+v[j].x)/2, (v[i].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j].y+v[j]
                                                 [.y)/2;
                                       Point b(a.x+(v[i].y-v[j].y), a.y-(v[i].
                                                 x-v[j].x));
                                       line = cross(a, b, v[i]) >= 0 ? mp(a, b)
                                                 : mp(b,a);
                                       p = halfplane_intersection(p, line);
                                       if (area(p) = 0) break;
                                                                                                      // 退化或者
                                                  空集合
                             }
                             ans[i] = area(p);
                    for (int i = 0; i < n; i ++) printf ("%.9f\n", ans [
                              i]);
         }
}
/*
10
3 4 4
1 1 2 2 3 3
```