

Coursework 4

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526- Databases

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1 Question 1

(a) We design the following ER Schema given the specifications:

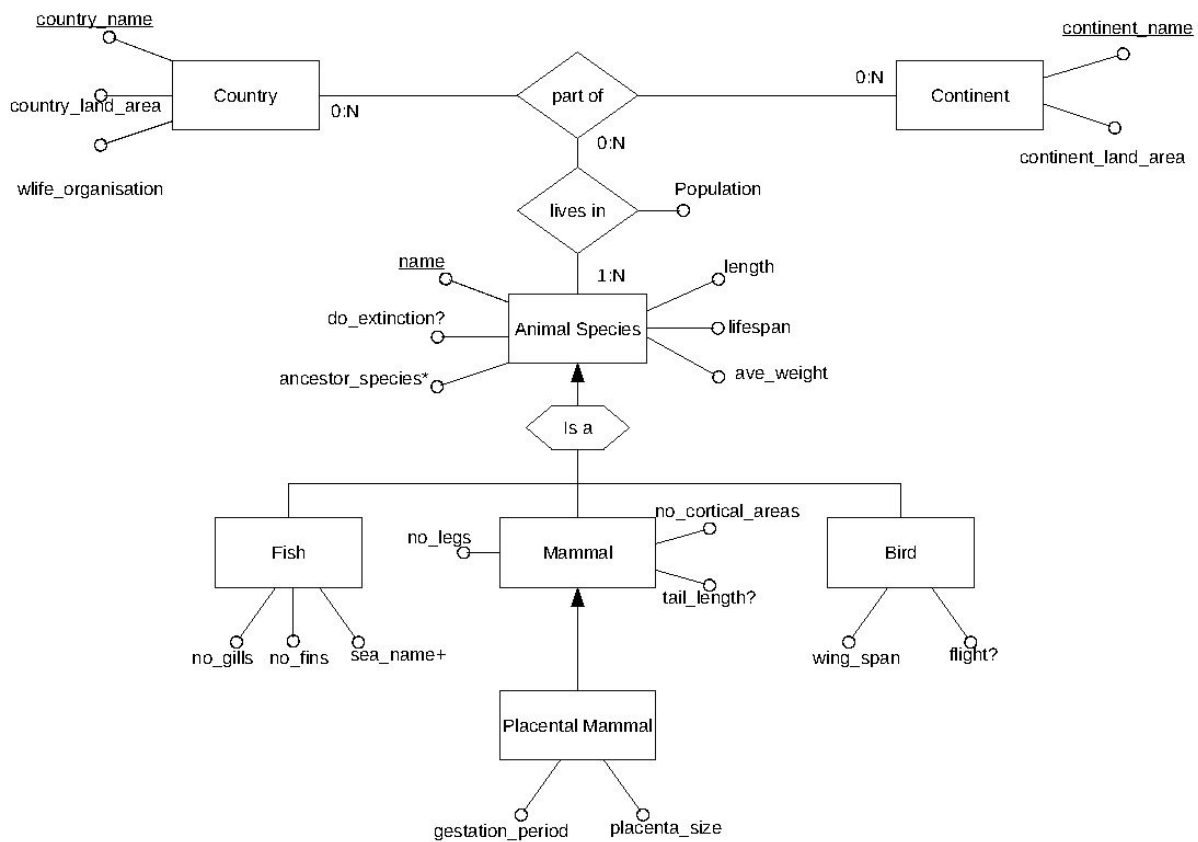


Figure 1: ER Schema to represent the given UoD

(b) We can the translate this ER Schema to an RA Schema. We get:

- animal_species(name, ave_weight,lifespan,do_extinction?)
- country(country_name, country_land_area,wlife_organisation)
- continent(continent_name,continent_land_area)
- fish(name,no_gills,no_fins)
- bird(name,wing_span,flight?)
- mammal(name,no_legs,no_cortical_areas,tail_length)
- placental_mammal(name,gestation_period,placenta_size)
- fish(name) \xRightarrow{fk} animal_species(name)
- bird(name) \xRightarrow{fk} animal_species(name)
- mammal(name) \xRightarrow{fk} animal_species(name)
- placental_mammal(name) \xRightarrow{fk} mammal(name)
- fish_seas(sea_name,name)
- fish_seas(name) \xRightarrow{fk} fish(name)
- ancestor_species(name, ancestor_species)
- ancestor_species(name) \xRightarrow{fk} animal_species(name)
- part_of(country_name,continent_name)
- part_of(country_name) \xRightarrow{fk} country(country_name)
- part_of(contient_name) \xRightarrow{fk} continent(continent_name)
- lives_in(country_name,continent_name,name,population)
- lives_in(country_name,continent_name) \xRightarrow{fk} belongs_to(country_name, continent_name)
- lives_in(name) \xRightarrow{fk} animal_species(name)

2 Question 2

(a) We have :

$S = \{AB \rightarrow HB, E \rightarrow GI, G \rightarrow CDEFHI, EHB \rightarrow CG, H \rightarrow B, F \rightarrow I\}$

Split the RHS of the FD's, we get:

$S' = \{AB \rightarrow H, AB \rightarrow B, E \rightarrow G, E \rightarrow I, G \rightarrow C, G \rightarrow D, G \rightarrow E, G \rightarrow F, G \rightarrow H, G \rightarrow I, EHB \rightarrow C, EHB \rightarrow G, H \rightarrow B, F \rightarrow I\}$

$AB \rightarrow B$ is self evident so remove.

Since $EHB \rightarrow C, H \rightarrow B \Rightarrow EH \rightarrow C$.

Since $E \rightarrow G, G \rightarrow H$ we get $E \rightarrow H$ and also $EH \rightarrow C \Rightarrow E \rightarrow C$.

Similarly, $EHB \rightarrow G \Rightarrow E \rightarrow G$, and since it is duplicated, remove.

So we get :

$S'' = \{AB \rightarrow H, E \rightarrow G, E \rightarrow I, G \rightarrow C, G \rightarrow D, G \rightarrow E, G \rightarrow F, G \rightarrow H, G \rightarrow I, E \rightarrow C, H \rightarrow B, F \rightarrow I\}$

Now:

$E \rightarrow G, G \rightarrow I$, so remove $E \rightarrow I$

$E \rightarrow G, G \rightarrow C$, so remove $E \rightarrow C$

$G \rightarrow F, F \rightarrow I$, so remove $G \rightarrow I$.

We can remove no other FD and still have the same closure as S.

So the minimal cover of S is:

$$\begin{aligned} Sc &= \{AB \rightarrow H, E \rightarrow G, G \rightarrow C, G \rightarrow D, G \rightarrow E, G \rightarrow F, G \rightarrow H, H \rightarrow B, F \rightarrow I\} \\ &= \{AB \rightarrow H, E \rightarrow G, G \rightarrow CDEFH, H \rightarrow B, F \rightarrow I\} \end{aligned}$$

(b) Start from the minimal cover and the Relation :

$R(A, B, C, D, E, F, G, H, I)$ and $Sc = \{AB \rightarrow H, E \rightarrow G, G \rightarrow CDEFH, H \rightarrow B, F \rightarrow I\}$

Moreover,

$G \rightarrow CDEFH, H \rightarrow B$ and $F \rightarrow I$ so $G \rightarrow BCDEFGHI$. Also $A \rightarrow A$.

So GA is a candidate minimal key of R.

Alternatively,

$E \rightarrow G, G \rightarrow CDEFH$ and $H \rightarrow B$ so $E \rightarrow BCDEFGHI$.

So EA is also a candidate minimal key of R.

From the set S we can also see that there are no other candidate minimal keys for R.

So **EA** and **GA** are the candidate minimal keys of R.

(c) Starting from the minimal cover :

$Sc = \{AB \rightarrow H, E \rightarrow G, G \rightarrow CDEFH, H \rightarrow B, F \rightarrow I\}$

We have: BCDFHI are non prime using question 2)b).

Since $AB \rightarrow H, H \rightarrow B$ violate 3NF, generate $R_1(\underline{A}, B, H)$.

Since $G \rightarrow CDFH$, violates 3NF generate $R_2(\underline{G}, CDFH)$.

Since $F \rightarrow I$ violates 3NF, generate $R_3(\underline{F}, I)$.

We then have the 3NF of R is :

$$R_1(\underline{A}, \underline{B}, H) R_2(\underline{G}, C, D, F, H) R_3(\underline{F}, I) R_4(\underline{A}, \underline{B}, \underline{E}, G) \quad (1)$$

Where:

R_1 contains $AB \rightarrow H$ and $H \rightarrow B$,

R_2 contains $G \rightarrow CDFH$,

R_3 contains $F \rightarrow I$,

R_4 contains $E \rightarrow G$ and $G \rightarrow E$.

(d) Let's decompose the relation 1 into BCNF:

Since $H \rightarrow B$ violates BCNF, generate $R_5(\underline{H}, B)$,

Since $G \rightarrow E$ and $E \rightarrow G$ violates BCNF, generate $R_6(\underline{G}, \underline{E})$.

We then get the BCNF form of R is:

$$R_1(\underline{A}, \underline{H}) R_2(\underline{G}, C, D, F, H) R_3(\underline{F}, I) R_4(\underline{A}, \underline{B}, \underline{E}) R_5(\underline{H}, B) R_6(\underline{G}, \underline{E}) \quad (2)$$

Where:

R_2 contains $G \rightarrow CDFH$, and hence maintains functional dependencies

R_3 contains $F \rightarrow I$, and hence maintains functional dependencies

R_6 contains $E \rightarrow G$ and $G \rightarrow E$, and hence maintains functional dependencies.

R_5 contains $H \rightarrow B$, and hence maintains functional dependencies.

As we can also see decomposition into R_1 and R_4 do not maintain functional dependencies. In particular, R_1 and R_4 do not contain any functional dependencies and $AB \rightarrow H$ cannot be maintained while decomposing R into BCNF.