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las4s e pelados

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1 Contest

2 Mathematics

3 Data structures

4 Numerical

5 Number theory

6 Combinatorial

7 Graph

8 Geometry

9 Strings

10 Various

Contest (1)

template.cpp

8 lines

```
#include <bits/stdc++.h>
using namespace std;

#define rep(i, a, b) for(int i = a; i < (b); ++i)
#define all(x) begin(x), end(x)
#define sz(x) (int)(x).size()
using ll = long long;
using pii = pair<int, int>;
```

.bashrc

2 lines

```
alias c='g++ -Wall -Wconversion -Wfatal-errors -g -std=c++17 \
-fsanitize=undefined,address'
```

hash.sh

2 lines

```
# bash hash.sh file.cpp 11 12
sed -n $2'','$3' p' $1 | sed '/^#/d' | cpp -Dd -P -
fpreprocessed | tr -d '[:space:]' | md5sum | cut -c 6
```

troubleshoot.txt

52 lines

Pre-submit:
Write a few simple test cases if sample is not enough.
Are time limits close? If so, generate max cases.
Is the memory usage fine?
Could anything overflow?
Make sure to submit the right file.

Wrong answer:
Print your solution! Print debug output, as well.
Are you clearing all data structures between test cases?
Can your algorithm handle the whole range of input?
Read the full problem statement again.
Do you handle all corner cases correctly?
Have you understood the problem correctly?
Any uninitialized variables?
Any overflows?

1 Confusing N and M, i and j, etc.?
Are you sure your algorithm works?
What special cases have you not thought of?
1 Are you sure the STL functions you use work as you think?
Add some assertions, maybe resubmit.
2 Create some testcases to run your algorithm on.
Go through the algorithm for a simple case.
Go through this list again.
4 Explain your algorithm to a teammate.
Ask the teammate to look at your code.
Go for a small walk, e.g. to the toilet.
6 Is your output format correct? (including whitespace)
Rewrite your solution from the start or let a teammate do it.
7 Runtime error:
Have you tested all corner cases locally?
8 Any uninitialized variables?
Are you reading or writing outside the range of any vector?
Any assertions that might fail?
13 Any possible division by 0? (mod 0 for example)
Any possible infinite recursion?
19 Invalidated pointers or iterators?
Are you using too much memory?
Debug with resubmits (e.g. remapped signals, see Various).
20 Time limit exceeded:
Do you have any possible infinite loops?
What is the complexity of your algorithm?
Are you copying a lot of unnecessary data? (References)
How big is the input and output? (consider scanf)
Avoid vector, map. (use arrays/unordered_map)
What do your teammates think about your algorithm?
Memory limit exceeded:
What is the max amount of memory your algorithm should need?
Are you clearing all data structures between test cases?

Mathematics (2)

2.1 Recurrences

If $a_n = c_1 a_{n-1} + \dots + c_k a_{n-k}$, and r_1, \dots, r_k are distinct roots of $x^k - c_1 x^{k-1} - \dots - c_k$, there are d_1, \dots, d_k s.t.

$$a_n = d_1 r_1^n + \dots + d_k r_k^n.$$

Non-distinct roots r become polynomial factors, e.g.
 $a_n = (d_1 n + d_2) r^n$.

2.2 Trigonometry

$$\begin{aligned}\sin(v+w) &= \sin v \cos w + \cos v \sin w \\ \cos(v+w) &= \cos v \cos w - \sin v \sin w\end{aligned}$$

$$\begin{aligned}\tan(v+w) &= \frac{\tan v + \tan w}{1 - \tan v \tan w} \\ \sin v + \sin w &= 2 \sin \frac{v+w}{2} \cos \frac{v-w}{2} \\ \cos v + \cos w &= 2 \cos \frac{v+w}{2} \cos \frac{v-w}{2}\end{aligned}$$

$$(V+W) \tan(v-w)/2 = (V-W) \tan(v+w)/2$$

where V, W are lengths of sides opposite angles v, w .

$$a \cos x + b \sin x = r \cos(x - \phi)$$

$$a \sin x + b \cos x = r \sin(x + \phi)$$

where $r = \sqrt{a^2 + b^2}, \phi = \text{atan2}(b, a)$.

2.3 Geometry

2.3.1 Triangles

Side lengths: a, b, c

$$\text{Semiperimeter: } p = \frac{a+b+c}{2}$$

$$\text{Area: } A = \sqrt{p(p-a)(p-b)(p-c)}$$

$$\text{Circumradius: } R = \frac{abc}{4A}$$

$$\text{Inradius: } r = \frac{A}{p}$$

Length of median (divides triangle into two equal-area triangles):

$$m_a = \frac{1}{2} \sqrt{2b^2 + 2c^2 - a^2}$$

Length of bisector (divides angles in two):

$$s_a = \sqrt{bc \left[1 - \left(\frac{a}{b+c} \right)^2 \right]}$$

$$\text{Law of sines: } \frac{\sin \alpha}{a} = \frac{\sin \beta}{b} = \frac{\sin \gamma}{c} = \frac{1}{2R}$$

$$\text{Law of cosines: } a^2 = b^2 + c^2 - 2bc \cos \alpha$$

$$\text{Law of tangents: } \frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$$

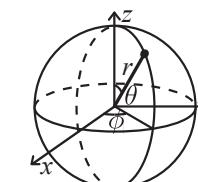
2.3.2 Quadrilaterals

With side lengths a, b, c, d , diagonals e, f , diagonals angle θ , area A and magic flux $F = b^2 + d^2 - a^2 - c^2$:

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2 f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is 180° , $ef = ac + bd$, and $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}$.

2.3.3 Spherical coordinates



$$\begin{aligned}x &= r \sin \theta \cos \phi & r &= \sqrt{x^2 + y^2 + z^2} \\ y &= r \sin \theta \sin \phi & \theta &= \arccos(z / \sqrt{x^2 + y^2 + z^2}) \\ z &= r \cos \theta & \phi &= \text{atan2}(y, x)\end{aligned}$$

2.3.4 Pick's Theorem

The area of a simple polygon whose vertices have integer coordinates is:

$$A = I + \frac{B}{2} - 1$$

where I is the number of interior integer points, and B is the number of integer points in the border of the polygon.

2.3.5 Two Ears Theorem

Every simple polygon with more than 3 vertices has at least two non-overlapping ears (a ear is a vertex whose diagonal induced by its neighbors which lies strictly inside the polygon). Equivalently, every simple polygon can be triangulated.

2.4 Derivatives/Integrals

$$\begin{aligned} \frac{d}{dx} \arcsin x &= \frac{1}{\sqrt{1-x^2}} & \frac{d}{dx} \arccos x &= -\frac{1}{\sqrt{1-x^2}} \\ \frac{d}{dx} \tan x &= 1 + \tan^2 x & \frac{d}{dx} \arctan x &= \frac{1}{1+x^2} \\ \int \tan ax \, dx &= -\frac{\ln |\cos ax|}{a} & \int x \sin ax \, dx &= \frac{\sin ax - ax \cos ax}{a^2} \\ \int e^{-x^2} \, dx &= \frac{\sqrt{\pi}}{2} \operatorname{erf}(x) & \int xe^{ax} \, dx &= \frac{e^{ax}}{a^2} (ax - 1) \end{aligned}$$

Integration by parts:

$$\int_a^b f(x)g(x) \, dx = [F(x)g(x)]_a^b - \int_a^b F(x)g'(x) \, dx$$

2.5 Sums

$$c^a + c^{a+1} + \dots + c^b = \frac{c^{b+1} - c^a}{c - 1}, \quad c \neq 1$$

$$1^2 + 2^2 + \dots + n^2 = \frac{n(2n+1)(n+1)}{6}$$

$$1^3 + 2^3 + \dots + n^3 = \frac{n^2(n+1)^2}{4}$$

$$1^4 + 2^4 + \dots + n^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$$

$$\sum_{i=0}^n ic^i = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c-1)^2}, \quad c \neq 1$$

$$g_k(n) = \sum_{i=1}^n i^k = \frac{1}{k+1} \left(n^{k+1} + \sum_{j=1}^k \binom{k+1}{j+1} (-1)^{j+1} g_{k-j}(n) \right)$$

template .bashrc hash troubleshoot Bit Bit2d

2.6 Series

$$\begin{aligned} e^x &= 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots, \quad (-\infty < x < \infty) \\ \ln(1+x) &= x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots, \quad (-1 < x \leq 1) \\ \sqrt{1+x} &= 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, \quad (-1 \leq x \leq 1) \\ \sin x &= x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, \quad (-\infty < x < \infty) \\ \cos x &= 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, \quad (-\infty < x < \infty) \\ \sum_{i=0}^{\infty} ic^i &= \frac{c}{(1-c)^2}, \quad |c| < 1 \end{aligned}$$

2.7 Probability theory

Let X be a discrete random variable with probability $p_X(x)$ of assuming the value x . It will then have an expected value (mean) $\mu = \mathbb{E}(X) = \sum_x x p_X(x)$ and variance $\sigma^2 = V(X) = \mathbb{E}(X^2) - (\mathbb{E}(X))^2 = \sum_x (x - \mathbb{E}(X))^2 p_X(x)$ where σ is the standard deviation. If X is instead continuous it will have a probability density function $f_X(x)$ and the sums above will instead be integrals with $p_X(x)$ replaced by $f_X(x)$.

Expectation is linear:

$$\mathbb{E}(aX + bY) = a\mathbb{E}(X) + b\mathbb{E}(Y)$$

For independent X and Y ,

$$V(aX + bY) = a^2 V(X) + b^2 V(Y).$$

2.7.1 Discrete distributions

Binomial distribution

The number of successes in n independent yes/no experiments, each which yields success with probability p is

$\text{Bin}(n, p)$, $n = 1, 2, \dots, 0 \leq p \leq 1$.

$$p(k) = \binom{n}{k} p^k (1-p)^{n-k}$$

$$\mu = np, \quad \sigma^2 = np(1-p)$$

$\text{Bin}(n, p)$ is approximately $\text{Po}(np)$ for small p .

First success distribution

The number of trials needed to get the first success in independent yes/no experiments, each which yields success with probability p is $\text{Fs}(p)$, $0 \leq p \leq 1$.

$$p(k) = p(1-p)^{k-1}, \quad k = 1, 2, \dots$$

$$\mu = \frac{1}{p}, \quad \sigma^2 = \frac{1-p}{p^2}$$

Poisson distribution

The number of events occurring in a fixed period of time t if these events occur with a known average rate κ and independently of the time since the last event is $\text{Po}(\lambda)$, $\lambda = t\kappa$.

$$p(k) = e^{-\lambda} \frac{\lambda^k}{k!}, \quad k = 0, 1, 2, \dots$$

$$\mu = \lambda, \quad \sigma^2 = \lambda$$

Data structures (3)

Bit.h

d41d8c, 13 lines

```
d41 struct Bit {
d41     vector<int> bit;
d41     Bit(int n) : bit(n + 1) {}
d41     void update(int i, int v) {
d41         for (i++; i < sz(bit); i += i & -i) bit[i] += v;
d41     }
d41     int query(int i) {
d41         int ret = 0;
d41         for (i++; i > 0; i -= i & -i) ret += bit[i];
d41     }
d41 };
```

Bit2d.h

Description: Points called on the update function NEED to be on the pts vector parameter on build.

Time: $\mathcal{O}((\log N)^2)$

d41d8c, 37 lines

```
d41 struct Bit2d {
d41     vector<vector<int>> ys;
d41     vector<Bit> bit;
d41     vector<int> cmp_x;
d41     Bit2d() {}
d41     void put(int x, int y) {
d41         for (x++; x < sz(ys); x += x & -x) ys[x].push_back(y);
d41     }
d41     int id(const vector<int> &v, int y) {
d41         return (upper_bound(all(v), y) - v.begin()) - 1;
d41     }
d41     void build(vector<pii> pts) {
d41         sort(all(pts));
d41         for (auto p : pts) cmp_x.push_back(p.first);
d41         cmp_x.erase(unique(all(cmp_x)), cmp_x.end());
d41         ys.resize(cmp_x.size() + 1);
d41         for (auto p : pts) put(id(cmp_x, p.first), p.second);
d41         for (auto &v : ys) sort(all(v));
d41         bit.emplace_back(sz(v));
d41     }
d41     void update(int x, int y, int val) {
d41         x = id(cmp_x, x);
d41         for (x++; x < sz(ys); x += x & -x)
d41             bit[x].update(id(ys[x], y), val);
d41     }
d41     int query(int x, int y) {
d41         x = id(cmp_x, x);
d41         int ret = 0;
d41         for (x++; x > 0; x -= x & -x)
d41             ret += bit[x].query(id(ys[x], y));
d41     }
d41     int query(int x1, int y1, int x2, int y2) {
d41         int a = query(x2, y2) - query(x2, y1 - 1);
```

```
d41     return a-query(x1-1, y2)+query(x1-1, y1-1);
d41 }
d41 };
```

LineContainer.h

Description: Container where you can add lines of the form $kx+m$, and query maximum values at points x . Useful for dynamic programming (“convex hull trick”).

Time: $\mathcal{O}(\log N)$

d41d8c, 32 lines

```
d41 struct Line {
d41     mutable ll k, m, p;
d41     bool operator<(const Line& o) const { return k < o.k; }
d41     bool operator<(ll x) const { return p < x; }
d41 };

d41 struct LineContainer : multiset<Line, less<> {
    // (for doubles, use inf = 1/.0, div(a,b) = a/b)
d41     static const ll inf = LLONG_MAX;
d41     ll div(ll a, ll b) { // floored division
d41         return a / b - ((a ^ b) < 0 && a % b); }
d41     bool isect(iterator x, iterator y) {
d41         if (y == end()) return x->p == inf, 0;
d41         if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
d41         else x->p = div(y->m - x->m, x->k - y->k);
d41         return x->p >= y->p;
d41     }
d41     void add(ll k, ll m) {
        auto z = insert({k, m, 0}), y = z++, x = y;
        while (isect(y, z)) z = erase(z);
d41         if (x != begin() && isect(--x, y))
            isect(x, y = erase(y));
        while ((y = x) != begin() && (--x)->p >= y->p)
            isect(x, erase(y));
d41     }
d41     ll query(ll x) {
d41         assert(!empty());
d41         auto l = *lower_bound(x);
d41         return l.k * x + l.m;
d41     }
d41 };
```

Mo.h

Description: For subtree queries, perform an Euler tour and map each node u to the interval $[tin[u], tin[u] + subtree_size[u] - 1]$. A subtree query becomes a range query over this interval.

For path queries between nodes U and V , Let U be the closest to the root. If V lies in U 's subtree, the path corresponds to the interval $[tin[U], tin[V]]$. Otherwise, the path corresponds to the interval $[min(tout[U], tout[V]), max(tin[U], tin[V])]$.

In both cases, nodes on the U - V path appear exactly once in the interval, while all other nodes appear either 0 or 2 times.

Usage: `queries.push(Query(l, r, index of query))`, intervals are $[l, r]$

Time: $\mathcal{O}(N\sqrt(Q))$

d41d8c, 44 lines

```
d41 inline int64_t hilOrd(int x, int y, int pow, int rot) {
d41     if (pow == 0) return 0;
d41     int hpow = 1 << (pow - 1);
d41     int seg = (x < hpow) ? ((y < hpow) ? 0 : 3) : ((y < hpow)
        ) ? 1 : 2;
d41     seg = (seg + rot) & 3;
d41     const int rotDelta[4] = { 3, 0, 0, 1 };
d41     int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
d41     int nrot = (rot + rotDelta[seg]) & 3;
d41     int64_t sub = int64_t(1) << (2 * pow - 2);
d41     int64_t ans = seg * sub;
d41     int64_t add = hilOrd(nx, ny, pow - 1, nrot);
d41     ans += (seg == 1 || seg == 2) ? add : (sub - add - 1);
```

```
d41     return ans;
d41 }

d41 struct Query {
d41     int l, r, idx;
d41     int64_t ord;
d41     Query(int l, int r, int idx) : l(l), r(r), idx(idx) {
d41         ord = hilOrd(l, r, 21, 0);
d41     }
d41     bool operator< (const Query& other) const {
d41         return ord < other.ord;
d41     }
d41 };

d41 vector<Query> queries;
d41 int ans[ms];
d41 void put(int x) {} // F
d41 void remove(int x) {} // F
d41 int getAns() {}

d41 void Mo() {
    int l = 0, r = -1;
    sort(queries.begin(), queries.end());
    for (Query q : queries) {
        while (l > q.l) put(--l);
        while (r < q.r) put(++r);
        while (l < q.l) remove(l++);
        while (r > q.r) remove(r--);
        ans[q.idx] = getAns();
    }
}
```

MoUpdate.h

Description: Block size should be around $(2 * N * N)^{\frac{1}{3}}$

Usage: intervals are $[l, r]$, addQuery(l, r, number of updates happened before this query, index of query), addUpdate(index of updated position, value before update, value after update)

Time: $\mathcal{O}(Q * (2 * N * N)^{\frac{1}{3}} * F)$

d41d8c, 55 lines

```
d41 const int B = 2700;
d41 struct MoUpdate {
d41     struct Query {
d41         int l, r, t, idx;
d41         Query(int l, int r, int t, int idx)
            : l(l), r(r), t(t), idx(idx) {}
d41         bool operator< (const Query& p) const {
d41             if (l / B != p.l / B) return l < p.l;
d41             if (r / B != p.r / B) return r < p.r;
d41             return t < p.t;
d41         }
d41     };
d41     struct Upd {
d41         int i, old, now;
d41         Upd(int i, int old, int now) : i(i), old(old), now(now) {}
d41     };
d41     vector<Query> queries;
d41     vector<Upd> updates;

d41     void addQuery(int l, int r, int t, int idx) {
        queries.push_back(Query(l, r, t, idx));
d41     }
d41     void addUpdate(int i, int old, int now) {
        updates.push_back(Upd(i, old, now));
d41     }

d41     void add(int x) {} // F
d41     void rem(int x) {} // F
d41     int getAns() {}
```

```
d41     void update(int novo, int idx, int l, int r) {
d41         if (l <= idx && idx <= r) rem(idx);
d41         arr[idx] = novo;
d41         if (l <= idx && idx <= r) add(idx);
d41     }

d41     void solve() {
        int l = 0, r = -1, t = 0;
        sort(queries.begin(), queries.end());
        for (Query q : queries) {
            while (l > q.l) add(--l);
            while (r < q.r) add(++r);
            while (l < q.l) rem(l++);
            while (r > q.r) rem(r--);
            while (t < q.t) {
                auto u = updates[t++];
                update(u.now, u.i, l, r);
            }
            while (t > q.t) {
                auto u = updates[--t];
                update(u.old, u.i, l, r);
            }
            ans[q.idx] = getAns();
        }
    }
```

SegmentTree.h

Description: Zero-indexed max-tree. Bounds are inclusive to the left and inclusive to the right. Can be changed by modifying T, f and unit.

Time: $\mathcal{O}(\log N)$

d41d8c, 21 lines

```
d41 struct Tree {
d41     typedef int T;
d41     static constexpr T unit = INT_MIN;
d41     T f(T a, T b) { return max(a, b); } // (any associative fn)
d41     vector<T> s; int n;
d41     Tree(int n = 0, T def = unit) : s(2*n, def), n(n) {}
d41     void update(int pos, T val) {
d41         for (s[pos += n] = val; pos /= 2; )
d41             s[pos] = f(s[pos * 2], s[pos * 2 + 1]);
d41     }
d41     T query(int b, int e) { // query [b, e]
d41         e++;
d41         T ra = unit, rb = unit;
d41         for (b += n, e += n; b < e; b /= 2, e /= 2) {
d41             if (b % 2) ra = f(ra, s[b++]);
d41             if (e % 2) rb = f(s[--e], rb);
d41         }
d41         return f(ra, rb);
d41     };
d41 };
```

OrderStatisticTree.h

Description: A set (not multiset!) with support for finding the n 'th element, and finding the index of an element. To get a map, change null_type.

Time: $\mathcal{O}(\log N)$

d41d8c, 17 lines

```
d41 #include <bits/extc++.h>
d41 using namespace __gnu_pbds;

d41 template<class T>
d41 using Tree = tree<T, null_type, less<T>, rb_tree_tag,
d41     tree_order_statistics_node_update>;

d41 void example() {
d41     Tree<int> t, t2; t.insert(8);
d41     auto it = t.insert(10).first;
```

```
d41 assert(it == t.lower_bound(9));
d41 assert(t.order_of_key(10) == 1);
d41 assert(t.order_of_key(11) == 2);
d41 assert(*t.find_by_order(0) == 8);
d41 t.join(t2); // merge t2 into t
d41 }
```

PersistentSegTree.h**Usage:** SegP(size of the segtree, number of updates)

```
roots = {0}, newRoot = update(roots.back(), ...),
roots.push(newRoot)
```

d41d8c, 42 lines

```
d41 struct SegP {
d41     static constexpr ll neut = 0;
d41     struct Node {
d41         ll v; // start with neutral value
d41         int l, r;
d41         Node(ll v=neut, int l=0, int r=0) : v(v), l(l), r(r){}
d41     };
d41     vector<Node> seg;
d41     int n, CNT;
d41     SegP(int _n, int upd) : seg(20*(upd+_n)), n(_n), CNT(1){}
d41     ll merge(ll a, ll b) { return a + b; }
d41     int update(int root, int pos, int val, int l, int r) {
d41         int p = CNT++;
d41         seg[p] = seg[root];
d41         if (l == r) {
d41             seg[p].v += val;
d41             return p;
d41         }
d41         int mid = (l + r) / 2;
d41         if (pos <= mid) {
d41             seg[p].l = update(seg[p].l, pos, val, l, mid);
d41         } else seg[p].r = update(seg[p].r, pos, val, mid+1, r);

d41         seg[p].v=merge(seg[seg[p].l].v, seg[seg[p].r].v);
d41         return p;
d41     }
d41     int query(int p, int L, int R, int l, int r) {
d41         if (l > R || r < L) return neut;
d41         if (L <= l && r <= R) return seg[p].v;
d41         int mid = (l + r) / 2;
d41         int left = query(seg[p].l, L, R, l, mid);
d41         int right = query(seg[p].r, L, R, mid + 1, r);
d41         return merge(left, right);
d41     }
d41     int update(int root, int pos, int val) {
d41         return update(root, pos, val, 0, n - 1);
d41     }
d41     int query(int root, int L, int R) {
d41         return query(root, L, R, 0, n - 1);
d41     }
};
```

RMQ.h**Usage:** RMQ rmq(values);
rmq.query(inclusive, inclusive);
Time: $\mathcal{O}(|V| \log |V| + Q)$

d41d8c, 17 lines

```
d41 struct RMQ {
d41     vector<vector<int>> dp;
d41     RMQ(const vector<int>& a) : dp(1, a) {
d41         for (int i = 1, pw = 1; pw*2 <= sz(a); pw*=2, i++) {
d41             dp.emplace_back(sz(a) - pw*2 + 1);
d41             for (int j = 0; j < sz(dp[i]); j++) {
d41                 dp[i][j] = min(dp[i-1][j], dp[i-1][j+pw]);
d41             }
d41         }
d41     }
d41 }
```

```
d41     }
d41     int query(int l, int r) {
d41         assert(l <= r);
d41         int k = 31 - __builtin_clz(r - l + 1);
d41         return min(dp[k][l], dp[k][r - (1 << k) + 1]);
d41     }
d41 }
```

Numerical (4)**4.1 Polynomials and recurrences****Polynomial.h**

d41d8c, 19 lines

```
d41     struct Poly {
d41         vector<double> a;
d41         double operator()(double x) const {
d41             double val = 0;
d41             for (int i = sz(a); i--;) (val *= x) += a[i];
d41             return val;
d41         }
d41         void diff() {
d41             rep(i,1,sz(a)) a[i-1] = i*a[i];
d41             a.pop_back();
d41         }
d41         void divroot(double x0) {
d41             double b = a.back(), c; a.back() = 0;
d41             for(int i=sz(a)-1; i--;) c = a[i], a[i] = a[i+1]*x0+b, b=c;
d41             a.pop_back();
d41         }
d41     };
d41 }
```

PolyRoots.h**Description:** Finds the real roots to a polynomial.**Usage:** polyRoots({{2,-3,1}},-1e9,1e9) // solve $x^2-3x+2 = 0$ **Time:** $\mathcal{O}(n^2 \log(1/\epsilon))$

d41d8c, 24 lines

```
"polynomial.h"
d41     vector<double> polyRoots(Poly p, double xmin, double xmax)
d41     {
d41         if (sz(p.a) == 2) { return {-p.a[0]/p.a[1]}; }
d41         vector<double> ret;
d41         Poly der = p;
d41         der.diff();
d41         auto dr = polyRoots(der, xmin, xmax);
d41         dr.push_back(xmin-1);
d41         dr.push_back(xmax+1);
d41         sort(all(dr));
d41         rep(i,0,sz(dr)-1) {
d41             double l = dr[i], h = dr[i+1];
d41             bool sign = p(l) > 0;
d41             if (sign ^ (p(h) > 0)) {
d41                 rep(it,0,60) { // while (h - l > 1e-8)
d41                     double m = (l + h) / 2, f = p(m);
d41                     if ((f <= 0) ^ sign) l = m;
d41                     else h = m;
d41                 }
d41                 ret.push_back((l + h) / 2);
d41             }
d41         }
d41         return ret;
d41     }
d41 }
```

BerlekampMassey.h**Description:** Recovers any n -order linear recurrence relation from the first $2n$ terms of the recurrence. Useful for guessing linear recurrences after brute-forcing the first terms. Should work on any field, but numerical stability for floats is not guaranteed. Output will have size $\leq n$.**Usage:** berlekampMassey({0, 1, 1, 3, 5, 11}) // {1, 2}**Time:** $\mathcal{O}(N^2)$

d41d8c, 21 lines

```
d41     vector<ll> berlekampMassey(vector<ll> s) {
d41         int n = sz(s), L = 0, m = 0;
d41         vector<ll> C(n), B(n), T;
d41         C[0] = B[0] = 1;

d41         ll b = 1;
d41         rep(i,0,n) { ++m;
d41             ll d = s[i] % mod;
d41             rep(j,1,L+1) d = (d + C[j] * s[i - j]) % mod;
d41             if (!d) continue;
d41             T = C; ll coef = d * modpow(b, mod-2) % mod;
d41             rep(j,m,n) C[j] = (C[j] - coef * B[j - m]) % mod;
d41             if (2 * L > i) continue;
d41             L = i + 1 - L; B = T; b = d; m = 0;
d41         }

d41         C.resize(L + 1); C.erase(C.begin());
d41         for (ll& x : C) x = (mod - x) % mod;
d41         return C;
d41     }
```

LinearRecurrence.h**Description:** Generates the k 'th term of an n -order linear recurrence $S[i] = \sum_j S[i-j-1]tr[j]$, given $S[0 \dots \geq n-1]$ and $tr[0 \dots n-1]$. Faster than matrix multiplication. Useful together with Berlekamp–Massey.**Usage:** linearRec({0, 1}, {1, 1}, k) // k 'th Fibonacci number**Time:** $\mathcal{O}(n^2 \log k)$

d41d8c, 27 lines

```
d41     using Poly = vector<ll>;
d41     ll linearRec(Poly S, Poly tr, ll k) {
d41         int n = sz(tr);

d41         auto combine = [&](Poly a, Poly b) {
d41             Poly res(n * 2 + 1);
d41             rep(i,0,n+1) rep(j,0,n+1)
d41                 res[i + j] = (res[i + j] + a[i] * b[j]) % mod;
d41             for (int i = 2 * n; i > n; --i) rep(j,0,n)
d41                 res[i-1-j] = (res[i-1-j] + res[i] * tr[j]) % mod;
d41             res.resize(n + 1);
d41             return res;
d41         };

d41         Poly pol(n + 1), e(pol);
d41         pol[0] = e[1] = 1;

d41         for (++k; k; k /= 2) {
d41             if (k % 2) pol = combine(pol, e);
d41             e = combine(e, e);
d41         }

d41         ll res = 0;
d41         rep(i,0,n) res = (res + pol[i + 1] * S[i]) % mod;
d41         return res;
d41     }
```

4.2 Matrices**SolveLinearBinary.h****Time:** $\mathcal{O}\left(\frac{\min(n,m)nm}{64}\right)$

d41d8c, 33 lines

```
d41     pair<int, bitset<M>> gauss(vector<bitset<M>> eq) {
d41         int n = eq.size(), m = M - 1;
d41         vector<int> where(m, -1);
```

```
d41 for(int col = 0, row = 0; col < m && row < n; col++) {
d41     for (int i = row; i < n; i++) {
d41         if (eq[i][col]) {
d41             swap(eq[i], eq[row]);
d41             break;
d41         }
d41     if (!eq[row][col]) continue;
d41     where[col] = row;
d41
d41     for (int i = 0; i < n; i++) {
d41         if (i != row && eq[i][col]) eq[i] ^= eq[row];
d41     }
d41     ++row;
d41 }
d41
bitset<M> ans;
for (int i = 0; i < m; i++) {
    if (where[i] != -1) ans[i] = eq[where[i]][m];
}
for (int i = 0; i < n; i++) {
    int sum = (ans & eq[i]).count();
    sum %= 2;
    if (sum != eq[i][m]) return pair(0, bitset<M>());
}
for (int i = 0; i < m; i++) {
    if (where[i] == -1) return pair(INF, ans);
}
return pair(1, ans);
d41 }
```

XorGauss.h

d41d8c, 30 lines

```
d41 struct XorGauss {
d41     int N;
d41     vector<ll> basis, who, mask;
d41     XorGauss(int N) : N(N), basis(N), who(N), mask(N) {}
// if(ans & (1ll << j)) who[j] was used to form x
d41     bool belong(ll x) {
d41         ll ans = 0;
d41         for(int i=N-1; i>=0; i--) {
d41             if((x ^ basis[i]) < x) {
d41                 ans ^= mask[i];
d41                 x ^= basis[i];
d41             }
d41         }
d41         return (x == 0);
}
d41     void add(ll v, int idx) {
d41         ll msk = 0;
d41         for (int i = N - 1; i >= 0; i--) {
d41             if (!(v & (1ll << i))) continue;
d41             if (basis[i] == 0) {
d41                 basis[i] = v, who[i] = idx;
d41                 mask[i] = (msk | (1ll << i));
d41                 return;
}
d41             msk ^= mask[i];
d41             v ^= basis[i];
d41         }
d41     }
d41 };
d41 };
```

4.3 Fourier transforms

FastFourierTransform.h

Description: fft(a) computes $\hat{f}(k) = \sum_x a[x] \exp(2\pi i \cdot kx/N)$ for all k . N must be a power of 2. Useful for convolution: $\text{conv}(a, b) = c$, where $c[x] = \sum a[i]b[x-i]$. For convolution of complex numbers or more than two vectors: FFT, multiply pointwise, divide by n , reverse(start+1, end), FFT back. Rounding is safe if $(\sum a_i^2 + \sum b_i^2) \log_2 N < 9 \cdot 10^{14}$ (in practice 10^{16} ; higher for random inputs). Otherwise, use NTT/FFTMod.

Time: $\mathcal{O}(N \log N)$ with $N = |A| + |B|$ ($\sim 1s$ for $N = 2^{22}$)

d41d8c, 44 lines

```
d41     typedef complex<double> C;
d41
d41     void fft(vector<C>& a) {
d41         int n = a.size(), L = 31 - __builtin_clz(n);
d41         static vector<complex<long double>> R(2, 1); // 10% faster if double
d41         static vector<C> rt(2, 1);
d41         for (static int k = 2; k < n; k *= 2) {
d41             R.resize(n);
d41             rt.resize(n);
d41             auto x = polar(1.0L, acos(-1.0L) / k);
d41             rep(i, k, 2*k) rt[i] = R[i] = i&1 ? R[i/2] * x : R[i/2];
d41         }
d41         vector<ll> rev(n);
d41         rep(i, 0, n) rev[i] = (rev[i / 2] | (i & 1) << L) / 2;
d41         rep(i, 0, n) if (i < rev[i]) swap(a[i], a[rev[i]]);
d41
d41         for (int k = 1; k < n; k *= 2) {
d41             for (int i = 0; i < n; i += 2 * k) {
d41                 for (int j = 0; j < k; j++) {
d41                     auto x = (double*)&rt[j + k];
d41                     auto y = (double*)&a[i + j + k];
d41                     C z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]*y[0]);
d41                     a[i + j + k] = a[i + j] - z;
d41                     a[i + j] += z;
d41                 }
d41             }
d41         }
d41     }
d41
vector<ll> conv(const vector<ll>& a, const vector<ll>& b) {
d41     if (a.empty() || b.empty()) return {};
d41     vector<ll> res(sz(a) + sz(b) - 1);
d41     int L = 32 - __builtin_clz(sz(res)), n = 1 << L;
d41     vector<C> in(n), out(n);
d41     copy(all(a), in.begin());
d41     rep(i, 0, sz(b)) in[i].imag(b[i]);
d41     fft(in);
d41     for (C & x : in) x *= x;
d41     rep(i, 0, n) out[i] = in[-i & (n - 1)] - conj(in[i]);
d41     fft(out);
d41     rep(i, 0, sz(res)) res[i] = round(imag(out[i]) / (4 * n));
d41     return res;
}
```

FastFourierTransformMod.h

Description: Higher precision FFT, can be used for convolutions modulo arbitrary integers as long as $N \log_2 N \cdot \text{mod} < 8.6 \cdot 10^{14}$ (in practice 10^{16} or higher). Inputs must be in $[0, \text{mod}]$.

Time: $\mathcal{O}(N \log N)$, where $N = |A| + |B|$ (twice as slow as NTT or FFT)

d41d8c, 23 lines

```
d41     typedef vector<ll> vl;
d41     template<int M> vl convMod(const vl &a, const vl &b) {
d41         if (a.empty() || b.empty()) return {};
d41         vl res(sz(a) + sz(b) - 1);
d41         int B=32-__builtin_clz(sz(res)), n=1<<B,cut=int(sqrt(M));
d41         vector<C> L(n), R(n), outs(n), outl(n);
d41         rep(i,0,sz(a)) L[i]=C((int)a[i]/cut, (int)a[i]%cut);
d41         rep(i,0,sz(b)) R[i]=C((int)b[i]/cut, (int)b[i]%cut);
d41         fft(L), fft(R);
d41         rep(i,0,n) {
```

```
d41             int j = -i & (n - 1);
d41             outl[j] = (L[i] + conj(L[j])) * R[i] / (2.0 * n);
d41             outs[j] = (L[i] - conj(L[j])) * R[i] / (2.0 * n) / i;
d41         }
d41         fft(outl), fft(outs);
d41         rep(i,0,sz(res)) {
d41             ll av = ll(real(outl[i])+.5), cv=ll(imag(outs[i])+.5);
d41             ll bv = ll(imag(outl[i])+.5) + ll(real(outs[i])+.5);
d41             res[i] = ((av % M * cut + bv) % M * cut + cv) % M;
d41         }
d41         return res;
d41     }
```

NumberTheoreticTransform.h

Description: ntt(a) computes $\hat{f}(k) = \sum_x a[x]g^{xk}$ for all k , where $g = \text{root}^{(mod-1)/N}$. N must be a power of 2. Useful for convolution modulo specific nice primes of the form $2^a b + 1$, where the convolution result has size at most 2^a . For arbitrary modulo, see FFTMod. $\text{conv}(a, b) = c$, where $c[x] = \sum a[i]b[x-i]$. For manual convolution: NTT the inputs, multiply pointwise, divide by n , reverse(start+1, end), NTT back. Inputs must be in $[0, \text{mod}]$.

Time: $\mathcal{O}(N \log N)$

d41d8c, 34 lines

```
d41     const int mod = 998244353, root = 62;
d41     typedef vector<ll> vl;
d41     void ntt(vl &a) {
d41         int n = sz(a), L = 31 - __builtin_clz(n);
d41         static vl rt(2, 1);
d41         for (static int k = 2, s = 2; k < n; k *= 2, s++) {
d41             rt.resize(n);
d41             ll z[] = {1, modpow(root, mod >> s)};
d41             rep(i, k, 2*k) rt[i] = rt[i / 2] * z[i & 1] % mod;
d41         }
d41         vector<int> rev(n);
d41         rep(i, 0, n) rev[i] = (rev[i / 2] | (i & 1) << L) / 2;
d41         rep(i, 0, n) if (i < rev[i]) swap(a[i], a[rev[i]]);
d41         for (int k = 1; k < n; k *= 2) {
d41             for (int i = 0; i < n; i += 2 * k) {
d41                 for (int j = 0; j < k; j++) {
d41                     ll z = rt[j+k] * a[i+j+k] % mod, &ai = a[i+j];
d41                     a[i + j + k] = ai - z + (z > ai ? mod : 0);
d41                     ai += (ai + z >= mod ? z - mod : z);
d41                 }
d41             }
d41         }
d41         vl conv(const vl &a, const vl &b) {
d41             if (a.empty() || b.empty()) return {};
d41             int s = sz(a) + sz(b) - 1, B = 32 - __builtin_clz(s),
d41             n = 1 << B;
d41             int inv = modpow(n, mod - 2);
d41             vl L(a), R(b), out(n);
d41             L.resize(n), R.resize(n);
d41             ntt(L), ntt(R);
d41             rep(i, 0, n) out[-i & (n - 1)] = (ll)L[i] * R[i] % mod * inv % mod;
d41             ntt(out);
d41             return {out.begin(), out.begin() + s};
d41         }
d41     }
```

FWHT.h

Description: Transform to a basis with fast convolutions of the form $c[z] = \sum_{z=x \oplus y} a[x] \cdot b[y]$, where \oplus is one of AND, OR, XOR. The size of a must be a power of two.

Time: $\mathcal{O}(N \log N)$

d41d8c, 20 lines

```
d41     void FST(vector<ll>& a, bool inv) {
d41         for (int n = sz(a), step = 1; step < n; step *= 2) {
d41             for (int i = 0; i < n; i += 2 * step) {
d41                 for (int j = i; j < i + step; j++) {
d41                     ll& u = a[j], &v = a[j + step];
d41                 }
d41             }
d41         }
d41     }
```

4.3 Fourier transforms

FastFourierTransform.h

```

d41     tie(u, v) =
d41     inv ? pair(v - u, u) : pair(v, u + v); // AND
d41     inv ? pair(v, u - v) : pair(u + v, u); // OR
d41     pair(u + v, u - v); // XOR
d41   }
d41 }
d41 if(inv) for(ll& x : a) x /= sz(a); // XOR only
d41 }
d41 vector<ll> conv(vector<ll> a, vector<ll> b) {
d41   FST(a, 0); FST(b, 0);
d41   for (int i = 0; i < sz(a); i++) a[i] *= b[i];
d41   FST(a, 1); return a;
d41 }

```

Number theory (5)

5.1 Modular arithmetic

ModInverse.h

Description: Pre-computation of modular inverses. Assumes $\text{LIM} \leq \text{mod}$ and that mod is a prime.

d41d8c, 5 lines

```

d41 const ll mod = 1000000007, LIM = 200000;
d41 inv[1] = 1;
d41 for(int i=2; i<LIM; i++)
d41   inv[i] = mod - (mod / i) * inv[mod % i] % mod;

```

ModMulLL.h

Description: Calculate $a \cdot b \bmod c$ (or $a^b \bmod c$) for $0 \leq a, b \leq c \leq 7.2 \cdot 10^{18}$.
Time: $\mathcal{O}(1)$ for modmul , $\mathcal{O}(\log b)$ for modpow

d41d8c, 12 lines

```

d41 typedef unsigned long long ull;
d41 ull modmul(ull a, ull b, ull M) {
d41   ll ret = a * b - M * ull(1.L / M * a * b);
d41   return ret + M * (ret < 0) - M * (ret >= (11.M));
d41 }
d41 ull modpow(ull b, ull e, ull mod) {
d41   ull ans = 1;
d41   for (; e; b = modmul(b, b, mod), e /= 2)
d41     if (e & 1) ans = modmul(ans, b, mod);
d41   return ans;
d41 }

```

DiscreteLog.h

Description: Returns the smallest x such that $a^x \bmod m = b \bmod m$. If no such x exists, returns -1 .

d41d8c, 32 lines

```

d41 int solve(int a, int b, int m) {
d41   a %= m, b %= m;
d41   if (a == 0) return (b ? -1 : 1);
d41   // caso gcd(a, m) > 1
d41   int k = 1, add = 0, g;
d41   while ((g = gcd(a, m)) > 1) {
d41     if (b == k) return add;
d41     if (b % g) return -1;
d41     b /= g, m /= g, ++add;
d41     k = (k * 111 * a / g) % m;
d41   }

d41   int sq = sqrt(m) + 1;
d41   int big = 1;
d41   for (int i = 0; i < sq; i++) big = (111 * big * a) % m
;

d41   vector<pii> vals;
d41   for (int q = 0, cur = b; q <= sq; q++) {

```

```

d41     vals.push_back({cur, q});
d41     cur = (111 * cur * a) % m;
d41   }
d41   sort(all(vals));
d41   for (int p = 1, cur = k; p <= sq; p++) {
d41     cur = (111 * cur * big) % m;
d41     auto it = lower_bound(all(vals), pair(cur, INF));
d41     if (it != vals.begin() && (--it)->first == cur) {
d41       return sq * p - it->second + add;
d41     }
d41   }
d41   return -1;
d41 }

```

DiscreteRoot.h

Description: Returns x such that $x^k \bmod m = a \bmod m$. If no such x exists, returns -1 .

Time: $\mathcal{O}(\sqrt{m}) * \log(\sqrt{m})$

"primitiveRoot.h", "DiscreteLog.h" d41d8c, 11 lines

// Discrete Root

```

d41 ll discreteRoot(ll k, ll a, ll m) {
d41   ll g = primitiveRoot(m);
d41   ll y = discreteLog(fexp(g, k, m), a, m);
d41   if (y == -1) return y;
d41   return fexp(g, y, m);
d41 }

```

5.2 Primality

MillerRabin.h

Description: Deterministic Miller-Rabin primality test. Guaranteed to work for numbers up to $7 \cdot 10^{18}$; for larger numbers, use Python and extend A randomly.

Time: 7 times the complexity of $a^b \bmod c$.

"ModMulLL.h" d41d8c, 13 lines

```

d41 bool isPrime(ull n) {
d41   if (n < 2 || n % 6 % 4 != 1) return (n | 1) == 3;
d41   ull A[] = {2, 325, 9375, 28178, 450775, 9780504, 17952650
22};
d41   ull s = __builtin_ctzll(n-1), d = n >> s;
d41   for (ull a : A) { // ^ count trailing zeroes
d41     ull p = modpow(a%n, d, n), i = s;
d41     while (p == 1 && p != n - 1 && a % n && i--) {
d41       p = modmul(p, p, n);
d41       if (p != n-1 && i != s) return 0;
d41     }
d41   }
d41   return 1;
d41 }

```

Factor.h

Description: Pollard-rho randomized factorization algorithm. Returns prime factors of a number, in arbitrary order (e.g. 2299 -> {11, 19, 11}).

Time: $\mathcal{O}(n^{1/4})$, less for numbers with small factors.

"ModMulLL.h", "MillerRabin.h" d41d8c, 19 lines

```

d41 ull pollard(ull n) {
d41   ull x = 0, y = 0, t = 30, prd = 2, i = 1, q;
d41   auto f = [&](ull x) { return modmul(x, x, n) + i; };
d41   while (t++ % 40 || gcd(prd, n) == 1) {
d41     if (x == y) x += i, y = f(x);
d41     if ((q = modmul(prd, max(x,y) - min(x,y), n))) prd = q;
d41     x = f(x), y = f(f(y));
d41   }
d41   return gcd(prd, n);
d41 }
d41 vector<ull> factor(ull n) {
d41   if (n == 1) return {};
d41   if (isPrime(n)) return {n};

```

```

d41   ull x = pollard(n);
d41   auto l = factor(x), r = factor(n / x);
d41   l.insert(l.end(), all(r));
d41   return l;
d41 }

```

PrimitiveRoot.h

d41d8c, 15 lines

```

// is n primitive root of p ?
d41 bool test(ll x, ll p) {
d41   ll m = p - 1;
d41   for (ll i = 2; i * i <= m; ++i) if (!(m % i)) {
d41     if (modpow(x, i, p) == 1) return false;
d41     if (modpow(x, m / i, p) == 1) return false;
d41   }
d41   return true;
d41 }

```

// find the smallest primitive root for p

```

d41 ll search(ll p) {
d41   for (ll i = 2; i < p; i++) if (test(i, p)) return i;
d41   return -1;
d41 }

```

5.3 Divisibility

Euclid.h

Description: Find x, y such that $Ax + By = \gcd(A, B)$. If $\gcd(A, B) = 1$, then $x = A^{-1} \pmod{B}$ and $y = B^{-1} \pmod{A}$.

Time: $\mathcal{O}(\log)$

d41d8c, 6 lines

```

d41 tuple<ll, ll, ll> ext_gcd(ll a, ll b) {
d41   if (!a) return {b, 0, 1};
d41   auto [g, x, y] = ext_gcd(b % a, a);
d41   return {g, y - b / a * x, x};
d41 }

```

CRT.h

d41d8c, 32 lines

```

d41 ll modinverse(ll a, ll b, ll s0 = 1, ll s1 = 0) {
d41   return !b ? s0 : modinverse(b, a % b, s1, s0 - s1 * (a / b));
d41 }

d41 ll mul(ll a, ll b, ll m) {
d41   ll q = (long double)a * (long double)b / (long double)m;
d41   ll r = a * b - q * m;
d41   return (r + m) % m;
d41 }

d41 struct Equation {
d41   ll mod, ans;
d41   bool valid;

```

```

d41   Equation(ll a, ll m) { mod = m, ans = a, valid = true; }
d41   Equation() { valid = false; }
d41   Equation(Equation a, Equation b) {
d41     if (!a.valid || !b.valid) {
d41       valid = false;
d41     } else {
d41       ll g = gcd(a.mod, b.mod);
d41       if ((a.ans - b.ans) % g != 0) {
d41         valid = false;
d41       } else {
d41         valid = true;
d41         mod = a.mod * (b.mod / g);
d41         ll x = mul(a.mod, modinverse(a.mod, b.mod), mod);
d41         ans = a.ans + mul(x, (b.ans - a.ans) / g, mod);
d41         ans = (ans % mod + mod) % mod;
d41       }
d41     }
d41   }

```

d41d8c, 15 lines

```
d41 void floor_ranges(int n) {
d41     for (int l = 1, r; l <= n; l = r + 1) {
d41         r = n / (n / l);
d41         // floor(n/y) has the same value for y in [l..r]
d41     }
d41 }
d41 void ceil_ranges(int n) {
d41     for (int l = 1, r; l <= n; l = r + 1) {
d41         int x = (n + l - 1) / l;
d41         if (x == 1) r = n;
d41         else r = (n - 1) / (x - 1);
d41         // ceil(n/y) has the same value for y in [l..r]
d41     }
d41 }
```

5.3.1 Bézout's identity

For $a \neq b \neq 0$, then $d = \gcd(a, b)$ is the smallest positive integer for which there are integer solutions to

$$ax + by = d$$

If (x, y) is one solution, then all solutions are given by

$$\left(x + \frac{kb}{\gcd(a, b)}, y - \frac{ka}{\gcd(a, b)} \right), \quad k \in \mathbb{Z}$$

Description: Euler's ϕ function is defined as $\phi(n) := \#$ of positive integers $\leq n$ that are coprime with n . $\phi(1) = 1$, p prime $\Rightarrow \phi(p^k) = (p-1)p^{k-1}$, m, n coprime $\Rightarrow \phi(mn) = \phi(m)\phi(n)$. If $n = p_1^{k_1}p_2^{k_2}\dots p_r^{k_r}$ then $\phi(n) = (p_1-1)p_1^{k_1-1}\dots(p_r-1)p_r^{k_r-1}$. $\phi(n) = n \cdot \prod_{p|n} (1 - 1/p)$.

$\sum_{d|n} \phi(d) = n$, $\sum_{1 \leq k \leq n, \gcd(k, n)=1} k = n\phi(n)/2$, $n > 1$

Euler's thm: a, n coprime $\Rightarrow a^{\phi(n)} \equiv 1 \pmod{n}$.

Fermat's little thm: p prime $\Rightarrow a^{p-1} \equiv 1 \pmod{p} \forall a$.

d41d8c, 6 lines

```
d41 void calculatePhi() {
d41     for(int i=0; i<LIM; i++) phi[i] = i&1 ? i : i/2;
d41     for (int i = 3; i < LIM; i += 2) if(phi[i] == i)
d41         for (int j = i; j < LIM; j += i) phi[j] -= phi[j] / i;
d41 }
```

5.4 Primes

$p = 962592769$ is such that $2^{21} \mid p - 1$, which may be useful. For hashing use 970592641 (31-bit number), 31443539979727 (45-bit), 3006703054056749 (52-bit). There are 78498 primes less than 1 000 000.

Primitive roots exist modulo any prime power p^a , except for $p = 2, a > 2$, and there are $\phi(\phi(p^a))$ many. For $p = 2, a > 2$, the group $\mathbb{Z}_{2^a}^\times$ is instead isomorphic to $\mathbb{Z}_2 \times \mathbb{Z}_{2^{a-2}}$.

5.5 Estimates

$$\sum_{d|n} d = O(n \log \log n).$$

The number of divisors of n is at most around 100 for $n < 5e4$, 500 for $n < 1e7$, 2000 for $n < 1e10$, 200 000 for $n < 1e19$.

5.6 Möbius Function

$$\mu(n) = \begin{cases} 0 & n \text{ is not square free} \\ 1 & n \text{ has even number of prime factors} \\ -1 & n \text{ has odd number of prime factors} \end{cases}$$

Möbius Inversion:

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu(d)g(n/d)$$

Other useful formulas/forms:

$$\sum_{d|n} \mu(d) = [n = 1] \text{ (very useful)}$$

$$g(n) = \sum_{n|d} f(d) \Leftrightarrow f(n) = \sum_{n|d} \mu(d/n)g(d)$$

$$g(n) = \sum_{1 \leq m \leq n} f(\lfloor \frac{n}{m} \rfloor) \Leftrightarrow f(n) = \sum_{1 \leq m \leq n} \mu(m)g(\lfloor \frac{n}{m} \rfloor)$$

5.7 Theorems

Goldbach's conjecture: Every even integer $n > 2$ can be written as $n = a + b$ with a, b prime.

Legendre's conjecture: There is always at least one prime between n^2 and $(n+1)^2$.

Lagrange's four-square theorem: Every positive integer can be written as

$$n = a^2 + b^2 + c^2 + d^2.$$

Zeckendorf's theorem: Every integer $n \geq 1$ has a unique representation as a sum of non-consecutive Fibonacci numbers:

$$n = F_{i_1} + F_{i_2} + \dots + F_{i_k}, \quad i_j - i_{j+1} \geq 2.$$

Euclid's formula (primitive Pythagorean triples): The Pythagorean triples are uniquely generated by

$$a = k \cdot (m^2 - n^2), \quad b = k \cdot (2mn), \quad c = k \cdot (m^2 + n^2),$$

with $m > n > 0$, $k > 0$, $m \perp n$, and either m or n even.

Wilson's theorem: n is prime iff

$$(n-1)! \equiv -1 \pmod{n}.$$

Chicken McNugget theorem: For coprime n, m , the largest integer not representable as $an + bm$ (with $a, b \geq 0$) is

$$nm - n - m.$$

There are $\frac{(n-1)(m-1)}{2}$ non-representable integers, and for each pair $(k, nm - n - m - k)$ exactly one is representable.

Combinatorial (6)

6.1 Binomial Identities

$${n-1 \choose k} - {n-1 \choose k-1} = \frac{n-2k}{k} {n \choose k} \quad {n \choose h} {n-h \choose k} = {n \choose k} {n-h \choose h}$$

$$\sum_{k=0}^n k {n \choose k} = n2^{n-1} \quad \sum_{k=0}^n k^2 {n \choose k} = (n+n^2)2^{n-2}$$

$$\sum_{j=0}^k {m \choose j} {n-m \choose k-j} = {n \choose k} \quad \sum_{j=0}^m {m \choose j}^2 = {2m \choose m}$$

$$\sum_{m=0}^n {m \choose j} {n-m \choose k-j} = {n+1 \choose k+1} \quad \sum_{m=0}^n {m \choose k} = {n+1 \choose k+1}$$

$$\sum_{r=0}^m {n+r \choose r} = {n+m+1 \choose m} \quad \sum_{k=0}^n {n-k \choose k} = \text{Fib}(n+1)$$

$$\sum_{k=0}^n {r \choose k} {s \choose n-k} = {r+s \choose n}$$

6.2 Permutations

6.2.1 Factorial

n	1	2	3	4	5	6	7	8	9	10
$n!$	1	2	6	24	120	720	5040	40320	362880	3628800
n		11	12	13	14	15	16	17		
$n!$	4.0e7	4.8e8	6.2e9	8.7e10	1.3e12	2.1e13	3.6e14			
n	20	25	30	40	50	100	150	171		
$n!$	2e18	2e25	3e32	8e47	3e64	9e157	6e262	>DBL_MAX		

6.2.2 Cycles

Let $g_S(n)$ be the number of n -permutations whose cycle lengths all belong to the set S . Then

$$\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp \left(\sum_{n \in S} \frac{x^n}{n} \right)$$

6.2.3 Derangements

Permutations of a set such that none of the elements appear in their original position.

$$D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n = \left[\frac{n!}{e} \right]$$

6.2.4 Burnside's lemma

Counts the number of distinct colorings of an object under symmetry.

$$\frac{1}{|G|} \sum_{g \in G} k^{\text{cyc}(g)},$$

where G is the symmetry group, k the number of colors, and $\text{cyc}(g)$ the number of cycles induced by g .

Example: number of ways to color a necklace with n beads using k colors (rotations only):

$$g(n) = \frac{1}{n} \sum_{i=0}^{n-1} k^{(\gcd(n, i))}$$

where rotation i shifts the necklace by i positions.

6.3 Partitions and subsets

6.3.1 Partition function

Number of ways of writing n as a sum of positive integers, disregarding the order of the summands.

$$p(0) = 1, p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k-1)/2)$$

$$p(n) \sim 0.145/n \cdot \exp(2.56\sqrt{n})$$

n	0	1	2	3	4	5	6	7	8	9	20	50	100
$p(n)$	1	1	2	3	5	7	11	15	22	30	627	$\sim 2e5$	$\sim 2e8$

PartitionSolver.h

d41d8c, 61 lines

```
d41 template<const int N>
d41 struct PartitionSolver {
d41     vector<vector<int>> part, to, from;
d41     PartitionSolver() {
d41         vector<int> a;
d41         part.push_back(a);
d41         gen(1, N, a);
d41         sort(all(part));
d41         to.assign(sz(part), vector<int>(N + 1, -1));
d41         from = to;
d41         for (int i = 0; i < sz(part); i++) {
d41             int sum = 0;
d41             auto arr = part[i];
d41             for (auto x : arr) sum += x;
d41             to[i][0] = i;
d41             from[i][0] = i;
d41             for (int j = 1; j + sum <= N; j++) {
d41                 arr = part[i];
d41                 arr.push_back(j);
d41                 sort(all(arr));
d41                 to[i][j] = getIndex(arr);
d41                 from[to[i][j]][j] = i;
d41             }
d41         }
d41         int size() const { return sz(part); }
d41         int getIndex(const vector<int>& arr) const {
d41             return lower_bound(all(part), arr) - part.begin(); }
d41         int add(int id, int num) const { return to[id][num]; }
d41         int rem(int id, int num) const { return from[id][num]; }
d41         vector<int> getPartition(int id) const {
d41             return part[id]; }

d41         void gen(int i, int sum, vector<int>& a) {
d41             if (i > sum) { return; }
d41             a.push_back(i);
d41             part.push_back(a);
d41             gen(i, sum - i, a);
d41             a.pop_back();
d41             gen(i + 1, sum, a);
d41         }
d41     };

// Number of partitions for all integers <= n
d41     vector<ll> partitionNumber(int n) {
d41         vector<ll> ans(n + 1, 0);
d41         ans[0] = 1;
d41         for (int i = 1; i <= n; i++) {
d41             for (int j = 1; j * (3 * j + 1) / 2 <= i; j++) {
d41                 ll here = ans[i - j * (3 * j + 1) / 2];
d41             }
d41         }
d41     }
}
```

PartitionSolver BellmanFord

```
d41         ans[i] = (ans[i] + (j & 1 ? here : -here));
d41     }
d41     for (int j = 1; j * (3 * j - 1) / 2 <= i; j++) {
d41         ll here = ans[i - j * (3 * j - 1) / 2];
d41         ans[i] = (ans[i] + (j & 1 ? here : -here));
d41     }
d41     return ans;
d41 }
```

6.3.2 Lucas' Theorem

Let n, m be non-negative integers and p a prime. Write $n = n_k p^k + \dots + n_1 p + n_0$ and $m = m_k p^k + \dots + m_1 p + m_0$. Then $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \pmod{p}$.

6.4 General purpose numbers

6.4.1 Bernoulli numbers

EGF of Bernoulli numbers is $B(t) = \frac{t}{e^t - 1}$ (FFT-able). $B[0, \dots] = [1, -\frac{1}{2}, \frac{1}{6}, 0, -\frac{1}{30}, 0, \frac{1}{42}, \dots]$

Sums of powers:

$$\sum_{i=1}^n n^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k \cdot (n+1)^{m+1-k}$$

Euler-Maclaurin formula for infinite sums:

$$\begin{aligned} \sum_{i=m}^{\infty} f(i) &= \int_m^{\infty} f(x) dx - \sum_{k=1}^{\infty} \frac{B_k}{k!} f^{(k-1)}(m) \\ &\approx \int_m^{\infty} f(x) dx + \frac{f(m)}{2} - \frac{f'(m)}{12} + \frac{f''(m)}{720} + O(f^{(5)}(m)) \end{aligned}$$

6.4.2 Stirling numbers of the first kind

Number of permutations on n items with k cycles.

$$c(n, k) = c(n-1, k-1) + (n-1)c(n-1, k), c(0, 0) = 1$$

$$\sum_{k=0}^n c(n, k)x^k = x(x+1)\dots(x+n-1)$$

$$c(8, k) = 8, 0, 5040, 13068, 13132, 6769, 1960, 322, 28, 1$$

$$c(n, 2) = 0, 0, 1, 3, 11, 50, 274, 1764, 13068, 109584, \dots$$

6.4.3 Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j :s s.t. $\pi(j) > \pi(j+1)$, $k+1$ j :s s.t. $\pi(j) \geq j$, k j :s s.t. $\pi(j) > j$.

$$E(n, k) = (n-k)E(n-1, k-1) + (k+1)E(n-1, k)$$

$$E(n, 0) = E(n, n-1) = 1$$

$$E(n, k) = \sum_{j=0}^k (-1)^j \binom{n+1}{j} (k+1-j)^n$$

6.4.4 Stirling numbers of the second kind

Partitions of n distinct elements into exactly k groups.

$$S(n, k) = S(n-1, k-1) + kS(n-1, k)$$

$$S(n, 1) = S(n, n) = 1$$

$$S(n, k) = \frac{1}{k!} \sum_{j=0}^k (-1)^{k-j} \binom{k}{j} j^n$$

6.4.5 Bell numbers

Total number of partitions of n distinct elements. $B(n) = 1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, \dots$. For p prime,

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

6.4.6 Labeled unrooted trees

- on n vertices: n^{n-2}
- on k existing trees of size n_i : $n_1 n_2 \dots n_k n^{k-2}$
- with degrees d_i : $(n-2)! / ((d_1-1)! \dots (d_n-1)!)$

6.4.7 Catalan numbers

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, C_{n+1} = \frac{2(2n+1)}{n+2} C_n, C_{n+1} = \sum C_i C_{n-i}$$

$$C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \dots$$

- sub-diagonal monotone paths in an $n \times n$ grid.
- strings with n pairs of parenthesis, correctly nested.
- binary trees with $n+1$ leaves (0 or 2 children).
- ordered trees with $n+1$ vertices.
- ways a convex polygon with $n+2$ sides can be cut into triangles by connecting vertices with straight lines.
- permutations of $[n]$ with no 3-term increasing subseq.

Graph (7)

7.1 Fundamentals

BellmanFord.h

Description: Calculates shortest paths from s in a graph that might have negative edge weights. Unreachable nodes get dist = inf; nodes reachable through negative-weight cycles get dist = -inf. Assumes $V^2 \max|w_i| < \sim 2^{63}$.

Time: $\mathcal{O}(VE)$

d41d8c, 24 lines

```
d41 const ll inf = LLONG_MAX;
d41 struct Ed { int a, b, w, s() { return a < b ? a : -a; } };
d41 struct Node { ll dist = inf; int prev = -1; };

d41 void bell(vector<Node>& nodes, vector<Ed*>& eds, int s) {
d41     nodes[s].dist = 0;
d41     sort(all(eds), [](Ed a, Ed b) { return a.s() < b.s(); });
d41     int lim = sz(nodes) / 2 + 2; // 3+100 with shuffled
d41     vertices
d41     for(int i=0; i<lim; i++) for (Ed ed : eds) {
d41         Node cur = nodes[ed.a], &dest = nodes[ed.b];
d41         if (abs(cur.dist) == inf) continue;
d41         ll d = cur.dist + ed.w;
```

```
d41  if (d < dest.dist) {
d41    dest.prev = ed.a;
d41    dest.dist = (i < lim-1 ? d : -inf);
d41  }
d41 for(int i=0; i<lim; i++) for (Ed e : eds) {
d41  if (nodes[e.a].dist == -inf)
d41    nodes[e.b].dist = -inf;
d41 }
d41 }
```

FloydWarshall.h

Description: Calculates all-pairs shortest path in a directed graph that might have negative edge weights. Input is an distance matrix m , where $m[i][j] = \text{inf}$ if i and j are not adjacent. As output, $m[i][j]$ is set to the shortest distance between i and j , inf if no path, or $-\text{inf}$ if the path goes through a negative-weight cycle.

Time: $\mathcal{O}(N^3)$

d41d8c, 20 lines

```
d41 void floydWarshall(vector<vector<ll>>& m) {
d41  int n = sz(m);
d41  for(int i=0; i<n; i++) m[i][i] = min(m[i][i], 0LL);
d41  for(int k=0; k<n; k++)
d41    for(int i=0; i<n; i++)
d41      for(int j=0; j<n; j++)
d41        if (m[i][k] != inf && m[k][j] != inf) {
d41          auto newDist = max(m[i][k] + m[k][j], -inf);
d41          m[i][j] = min(m[i][j], newDist);
d41        }

d41  for(int k=0; k<n; k++)
d41    if (m[k][k] < 0)
d41      for(int i=0; i<n; i++)
d41        for(int j=0; j<n; j++)
d41          if (m[i][k] != inf && m[k][j] != inf) {
d41            m[i][j] = -inf;
d41          }
d41 }
```

7.2 Network flow and Matching

Dinic.h

Time: $\mathcal{O}(\min(m \cdot \text{max_flow}, n^2 m))$.

- For graphs with unit capacities: $\mathcal{O}(\min(m\sqrt{m}, mn^{2/3}))$.

- If every vertex has in-degree 1 or out-degree 1: $\mathcal{O}(m\sqrt{n})$.

- With capacity scaling: $\mathcal{O}(nm \log(\text{MAXCAP}))$ with high constant factor.

```
d41 struct Dinic {
d41  const bool scaling = false;
d41  int lim;
d41  struct edge {
d41    int to, rev;
d41    ll cap, flow;
d41    bool res;
d41    edge(int to_, int cap_, int rev_, bool res_)
d41      : to(to_), cap(cap_), rev(rev_), flow(0), res(res_) {}
d41  };

d41  vector<vector<edge>> g;
d41  vector<int> lev, beg;
d41  ll F;
d41  Dinic(int n) : g(n), lev(n), beg(n), F(0) {}

d41  void add(int a, int b, ll c, ll other = 0) {
d41    g[a].emplace_back(b, c, g[b].size(), false);
d41    g[b].emplace_back(a, other, g[a].size()-1, true);
d41  }

d41  bool bfs(int s, int t) {
d41    fill(all(lev), -1);
```

FloydWarshall Dinic MinCost PushRelabel

```
d41    fill(all(beg), 0);
d41    lev[s] = 0;
d41    queue<int> q; q.push(s);
d41    while (q.size()) {
d41      int u = q.front(); q.pop();
d41      for (auto& i : g[u]) {
d41        if (lev[i.to] != -1 || (i.flow == i.cap)) continue;
d41        if (scaling and i.cap - i.flow < lim) continue;
d41        lev[i.to] = lev[u] + 1;
d41        q.push(i.to);
d41      }
d41    }
d41    return lev[t] != -1;
d41  }

ll dfs(int v, int s, ll f = INF) {
d41  if (!f || v == s) return f;
d41  for (int& i = beg[v]; i < g[v].size(); i++) {
d41    auto& e = g[v][i];
d41    if (lev[e.to] != lev[v] + 1) continue;
d41    ll foi = dfs(e.to, s, min(f, e.cap - e.flow));
d41    if (!foi) continue;
d41    e.flow += foi, g[e.to][e.rev].flow -= foi;
d41    return foi;
d41  }
d41  return 0;
}

ll maxFlow(int s, int t) {
d41  for (lim = scaling ? (1<<30) : 1; lim; lim /= 2)
d41    while (bfs(s, t)) while (ll ff = dfs(s, t)) F += ff;
d41  return F;
}

bool inCut(int u){ return lev[u] != -1; }
```

MinCost.h

Description: Min-cost max-flow. If costs can be negative, call setpi before maxflow, but note that negative cost cycles are not supported. To obtain the actual flow, look at positive values only. If graph is a DAG pi can be calculated with DP instead of Bellman ford.

Time: $\mathcal{O}(FE \log(V))$ where F is max flow. $\mathcal{O}(VE)$ for setpi.

```
d41 #include <bits/extc++.h>

d41 const ll INF = numeric_limits<ll>::max() / 4;

d41 struct MCMF {
d41  struct edge {
d41    int from, to, rev;
d41    ll cap, cost, flow;
d41  };
d41  int N;
d41  vector<vector<edge>> ed;
d41  vector<int> seen, vis;
d41  vector<ll> dist, pi;
d41  vector<edge*> par;

d41  MCMF(int N) : N(N), ed(N), seen(N), vis(N),
d41    dist(N), pi(N), par(N) {}

d41  void addEdge(int from, int to, ll cap, ll cost) {
d41    if (from == to || cap == 0) return;
d41    ed[from].push_back(edge{from,to,sz(ed[to]),cap,cost,0});
d41  };
d41  ed[to].push_back(edge{to,from,sz(ed[from])-1,0,-cost,0});
d41  }

d41  void path(int s) {
d41    fill(all(seen), 0);
```

```
d41    fill(all(dist), INF);
d41    dist[s] = 0;
d41    ll di;
d41    __gnu_pbds::priority_queue<pair<ll, int>> q;
d41    vector<decltype(q)::point_iterator> its(N);
d41    q.push({ 0, s });

d41    while (!q.empty()) {
d41      s = q.top().second; q.pop();
d41      seen[s] = 1; di = dist[s] + pi[s];
d41      for (edge& e : ed[s]) {
d41        if (!seen[e.to]) {
d41          ll val = di - pi[e.to] + e.cost;
d41          if(e.cap - e.flow > 0 && val < dist[e.to]){
d41            dist[e.to] = val;
d41            par[e.to] = &e;
d41            if (its[e.to] == q.end()) {
d41              its[e.to] = q.push({ -dist[e.to], e.to });
d41            }
d41            else q.modify(its[e.to], {-dist[e.to], e.to});
d41          }
d41        }
d41      }
d41    }
d41    for (int i = 0; i < N; i++) {
d41      pi[i] = min(pi[i] + dist[i], INF);
d41    }
}

pair<ll, ll> maxflow(int s, int t) {
d41  setpi(s, t);
d41  ll totflow = 0, totcost = 0;
d41  while (path(s), seen[t]) {
d41    ll fl = INF;
d41    for (edge* x = par[t]; x; x = par[x->from]) {
d41      fl = min(fl, x->cap - x->flow);
d41    }
d41    totflow += fl;
d41    for (edge* x = par[t]; x; x = par[x->from]) {
d41      x->flow += fl;
d41      ed[x->to][x->rev].flow -= fl;
d41    }
d41  }
d41  for (int i = 0; i < N; i++) {
d41    for (edge& e : ed[i]) {
d41      totcost += e.cost * e.flow;
d41    }
d41  }
d41  return { totflow, totcost / 2 };
}

// If some costs can be negative, call this before maxflow:
void setpi(int s, int t) {
d41  fill(all(pi), INF);
d41  pi[s] = 0;
d41  int it = N, ch = 1;
d41  ll v;
d41  while (ch-- && it--) {
d41    for (int i = 0; i < N; i++) {
d41      if (pi[i] != INF)
d41        for (edge& e : ed[i]) if (e.cap)
d41          if ((v = pi[i] + e.cost) < pi[e.to])
d41            pi[e.to] = v, ch = 1;
d41    }
d41  }
d41  assert(it >= 0); // negative cost cycle
d41 }
```

Blossom HopcroftKarp WeightedMatching

PushRelabel.h

Description: Push-relabel using the highest label selection rule and the gap heuristic. Quite fast in practice. To obtain the actual flow, look at positive values only.

Time: $\mathcal{O}(V^2 \sqrt{E})$

 [Studydrive](#) [Studydrive](#)

```
d41 struct PushRelabel {
d41     struct Edge {
d41         int dest, back;
d41         ll f, c;
d41     };
d41     vector<vector<Edge>> g;
d41     vector<ll> ec;
d41     vector<Edge*> cur;
d41     vector<vector<int>> hs;
d41     vector<int> H;
d41     PushRelabel(int n) : g(n), ec(n), cur(n), hs(2*n), H(n) {}
```

```
d41 void addEdge(int s, int t, ll cap, ll rcap=0)
d41     if (s == t) return;
d41     g[s].push_back({t, sz(g[t]), 0, cap});
d41     g[t].push_back({s, sz(g[s])-1, 0, rcap});
d41 }
```

```

d41 void addFlow(Edge& e, ll f) {
d41     Edge &back = g[e.dest][e.back];
d41     if (!ec[e.dest] && f) hs[H[e.dest]].push_back(e.dest);
d41     e.f += f; e.c -= f; ec[e.dest] += f;
d41     back.f -= f; back.c += f; ec[back.dest] -= f;
d41 }
d41 ll calc(int s, int t) {
d41     int v = sz(g); H[s] = v; ec[t] = 1;
d41     vector<int> co(2*v); co[0] = v-1;
d41     for(int i=0; i<v; i++) cur[i] = g[i].data();
d41     for (Edge& e : g[s]) addFlow(e, e.c);

```

```

d41     for (int hi = 0;;) {
d41       while (hs[hi].empty()) if (!hi--) return -ec[s];
d41       int u = hs[hi].back(); hs[hi].pop_back();
d41       while (ec[u] > 0) // discharge u
d41         if (cur[u] == g[u].data() + sz(g[u])) {
d41           H[u] = 1e9;
d41           for (Edge& e : g[u]){
d41             if (e.c && H[u] > H[e.dest]+1)
d41               H[u] = H[e.dest]+1, cur[u] = &e;
d41           }
d41           if (++co[H[u]], !--co[hi] && hi < v) {
d41             for(int i=0; i<v; i++){
d41               if (hi < H[i] && H[i] < v)
d41                 --co[H[i]], H[i] = v + 1;
d41             }
d41           }
d41           hi = H[u];
d41         } else if (cur[u]->c && H[u] == H[cur[u]->dest]+1)
d41           addFlow(*cur[u], min(ec[u], cur[u]->c));
d41         }else ++cur[u];
d41       }
d41     }
d41     bool inCut(int a) { return H[a] >= sz(g); }
d41   .

```

Blossom h

Description: Max matching on general Graph. $\text{mate}[i] = \text{match of } i$
Time: $\mathcal{O}(N^3)$

d41d8c, 56 lines

```
d41    vector<int> Blossom(vector<vector<int>>& g) {
```

```
d41      int n = sz(g), timer = -1;
d41      vector<int> mate(n, -1), label(n), par(n), orig(n), aux(n,
```

```

d41 auto lca = [&](int x, int y) {
d41     for (timer++; ; swap(x, y)) {
d41         if (x == -1) continue;
d41         if (aux[x] == timer) return x;
d41         aux[x] = timer;
d41         x=(mate[x] == -1 ? -1 : orig[par[mate[x]]]);
d41     }
d41 }
d41 auto blossom = [&](int v, int w, int a) {
d41     while (orig[v] != a) {
d41         par[v] = w; w = mate[v];
d41         if(label[w] == 1) label[w] = 0, q.push_back(w);
d41         orig[v] = orig[w] = a;
d41         v = par[w];
d41     }
d41 }
d41 auto aug = [&](int v) {
d41     while (v != -1) {
d41         int pv = par[v], nv = mate[pv];
d41         mate[v] = pv; mate[pv] = v; v = nv;
d41     }
d41 };
d41 auto bfs = [&](int root) {
d41     fill(all(label), -1);
d41     iota(all(orig), 0);
d41     q.clear();
d41     label[root] = 0; q.push_back(root);
d41     for (int i = 0; i < sz(q); i++) {
d41         int v = q[i];
d41         for (auto x : g[v]) {
d41             if (label[x] == -1) {
d41                 label[x] = 1; par[x] = v;
d41                 if (mate[x] == -1) return aug(x), 1;
d41                 label[mate[x]] = 0;
d41                 q.push_back(mate[x]);
d41             }
d41             else if(!label[x] && orig[v] != orig[x]){
d41                 int a = lca(orig[v], orig[x]);
d41                 blossom(x, v, a);
d41                 blossom(v, x, a);
d41             }
d41         }
d41     }
d41     return 0;
d41 };
// Time halves if you start with (any) maximal
// matching.
d41 for (int i = 0; i < n; i++) {
d41     if (mate[i] == -1) bfs(i);
d41 }

```

HopcroftKarp.h

Description: *ans* is the size of the max matching.

The match of x is $I[x]$.

Usage: HopcroftKarp(|X|, |Y|, edges(x, y))

Time: $\mathcal{O}(\sqrt{V}E)$

d41d8c, 46 line

```
d41 struct HopcroftKarp {
d41     vector<int> g, l, r;
d41     int ans;
d41     HopcroftKarp(int n, int m, vector<pair<int, int>> e)
d41         : g(e.size()), l(n, -1), r(m, -1), ans(0) {
d41         shuffle(e.begin(), e.end(), rng);
d41         vector<int> deg(n + 1);
d41         for (auto& [x, y] : e) deg[x]++;
d41         for (int i = 1; i <= n; i++) deg[i] += deg[i - 1];
d41     }
d41     void dfs(int v) {
d41         if (l[v] != -1) return;
d41         l[v] = r[ans];
d41         for (int u : g[v]) {
d41             if (r[u] == -1 || (l[u] == -1 && dfs(u))) {
d41                 r[u] = l[v];
d41                 ans++;
d41                 break;
d41             }
d41         }
d41     }
d41     int maxMatching() {
d41         int ans = 0;
d41         while (true) {
d41             vector<int> l(n, -1), r(m, -1);
d41             for (int i = 0; i < n; i++) {
d41                 if (l[i] == -1) dfs(i);
d41             }
d41             if (ans == n) return ans;
d41             int matchCount = 0;
d41             for (int i = 0; i < n; i++) {
d41                 if (l[i] != -1) matchCount++;
d41             }
d41             if (matchCount == n) return ans;
d41             vector<int> deg(n + 1);
d41             for (int i = 0; i < n; i++) {
d41                 if (l[i] == -1) deg[i] = 0;
d41                 else deg[i] = 1;
d41             }
d41             for (int i = 0; i < m; i++) {
d41                 if (r[i] == -1) continue;
d41                 for (int u : g[r[i]]) {
d41                     if (l[u] == -1) deg[u]++;
d41                 }
d41             }
d41             for (int i = 0; i < n; i++) {
d41                 if (l[i] == -1) dfs(i);
d41             }
d41         }
d41     }
d41 };
```

```

d41     for (auto& [x, y] : e) g[--deg[x]] = y;
d41
d41     vector<int> q(n);
d41     while (true) {
d41         vector<int> a(n, -1), p(n, -1);
d41         int t = 0;
d41         for (int i = 0; i < n; i++) {
d41             if (l[i] == -1) {
d41                 q[t++] = a[i] = p[i] = i;
d41             }
d41         }
d41         bool match = false;
d41         for (int i = 0; i < t; i++) {
d41             int x = q[i];
d41             if (~l[a[x]]) continue;
d41             for (int j = deg[x]; j < deg[x + 1]; j++) {
d41
d41                 int y = g[j];
d41                 if (r[y] == -1) {
d41                     while (~y) {
d41                         r[y] = x;
d41                         swap(l[x], y);
d41                         x = p[x];
d41                     }
d41                     match = true, ans++;
d41                     break;
d41                 }
d41                 if (p[r[y]] == -1) {
d41                     q[t++] = y = r[y];
d41                     p[y] = x, a[y] = a[x];
d41                 }
d41             }
d41         }
d41         if (!match) break;
d41     }
d41 }

```

WeightedMatching.h

Description: Given a weighted bipartite graph, matches every node on the left with a node on the right such that no nodes are in two matchings and the sum of the edge weights is minimal. Takes cost[N][M], where $\text{cost}[i][j] = \text{cost}$ for $L[i]$ to be matched with $R[j]$ and returns (min cost, match), where $L[i]$ is matched with $R[\text{match}[i]]$. Negate costs for max cost. Requires $N \leq M$.

Time: $\mathcal{O}(N^2M)$ d41d8c, 41 lines

```

d41 pair<ll, vector<int>> hunga(const vector<vector<ll> & ) {
d41     if (a.empty()) return { 0, {} };
d41     int n = sz(a) + 1, m = sz(a[0]) + 1;
d41     vector<ll> u(n), v(m), p(m);
d41     vector<int> ans(n - 1);
d41     for (int i = 1; i < n; i++) {
d41         p[0] = i;
d41         int j0 = 0;
d41         vector<ll> dist(m, LLONG_MAX), pre(m, -1);
d41         vector<bool> done(m + 1);
d41         do {
d41             done[j0] = true;
d41             ll io = p[j0], jl = -1, delta = LLONG_MAX;
d41             for (int j = 1; j < m; j++) {
d41                 if (!done[j]) {
d41                     ll cur = a[i0-1][j-1] - u[i0] - v[j];
d41                     if (cur < dist[j])
d41                         dist[j] = cur, pre[j] = j0;
d41                     if (dist[j] < delta)
d41                         delta = dist[j], jl = j;
d41                 }
d41             }
d41             for (int j = 0; j < m; j++) {

```

```

d41     if (done[j])
d41         u[p[j]] += delta, v[j] -= delta;
d41     else dist[j] -= delta;
d41 }
d41 assert(j1 != -1);
d41 j0 = j1;
d41 } while (p[j0]);
d41 while (j0) {
d41     int j1 = pre[j0];
d41     p[j0] = p[j1], j0 = j1;
d41 }
d41 for (int j = 1; j < m; j++) {
d41     if (p[j]) ans[p[j] - 1] = j - 1;
d41 }
d41 return { -v[0], ans }; // min cost
d41 }

```

7.2.1 Hall's Theorem

In bipartite graphs, there exists a perfect matching covering the entire side X if and only if for every subset $Y \subseteq X$,

$$|Y| \leq |N(Y)|,$$

where $N(Y)$ denotes the set of neighbors of Y .

7.2.2 König's Theorem

In a bipartite graph, the size of a Minimum Vertex Cover is equal to the size of a Maximum Matching. A Minimum Vertex Cover is a minimum set of vertices such that every edge of the graph has at least one endpoint in the set.

As a consequence,

$$n - \text{Maximum Matching} = \text{Maximum Independent Set},$$

where a Maximum Independent Set is the largest set of vertices with no edges between them.

Recovering the Minimum Vertex Cover Given a maximum matching in a bipartite graph (X, Y) :

- Construct the residual graph by orienting:
 - non-matching edges from X to Y ;
 - matching edges from Y to X .
- Perform a BFS or DFS starting from all free (unmatched) vertices in X .
- Let Z_X be the set of reachable vertices in X , and Z_Y the set of reachable vertices in Y .

The Minimum Vertex Cover is given by:

$$(X \setminus Z_X) \cup Z_Y.$$

Bridges EulerPath SCC

7.2.3 Node-Disjoint Path Cover

A node-disjoint path cover is a set of paths such that each vertex belongs to exactly one path.

In a directed acyclic graph (DAG),

$$\text{Minimum Node-Disjoint Path Cover} = n - \text{Maximum Matching}.$$

The construction is as follows: for each vertex u , create a copy u' . Add an edge $u \rightarrow v'$ if there exists an edge $u \rightarrow v$ in the original graph.

Recovering the Paths

- Vertices that do not appear as destinations in the matching are starting points of paths.
- Each matching edge $u \rightarrow v'$ corresponds to an edge $u \rightarrow v$ in the original DAG.
- Following these edges reconstructs all paths of the path cover.

7.2.4 General Path Cover

A general path cover is a path cover where a vertex may belong to more than one path.

In a DAG, the construction is similar to the node-disjoint case, but an edge $u \rightarrow v'$ exists if there is a path from u to v in the original graph.

Recovering the Cover The vertices can be grouped according to the edges used in the matching to form the path cover.

7.2.5 Dilworth's Theorem

An antichain is a set of vertices such that there is no path between any pair of vertices in the set.

In a directed acyclic graph,

$$\text{Minimum General Path Cover} = \text{Maximum Antichain}.$$

Recovering a Maximum Antichain Given a minimum general path cover, selecting one vertex from each path produces a maximum antichain.

7.3 DFS algorithms

Bridges.h

d41d8c, 24 lines

```

d41 vector<int> g[ms];
d41 int low[ms], tin[ms], vis[ms], t;
d41
d41 void dfs(int u = 0, int p = -1) {
d41     vis[u] = true;
d41     low[u] = tin[u] = t++;
d41     for (auto v : g[u]) {
d41         if (v == p) continue;

```

```

d41         if (vis[v]) {
d41             low[u] = min(low[u], tin[v]);
d41         }
d41     else {
d41         dfs(v, u);
d41         low[u] = min(low[u], low[v]);
d41         // if (low[v] >= tin[u]) && p != -1), U is an
d41         // articulation point
d41         if (low[v] > tin[u]) {
d41             // edge from U to V is a bridge
d41         }
d41         // children++
d41     }
d41 }
d41 // if(children > 1 && p == -1) root is an articulation
d41 point
d41 }

```

EulerPath.h

Description: Receives as input graph(node, edge index), number of edges and source. Returns list of node, index of edge he came from, if path/circuit does not exists returns empty list.

d41d8c, 27 lines

```

d41 vector<pii> eulerPath(const vector<vector<pii>>& g, int
nedges, int src) {
d41     int n = sz(g);
d41     vector<int> deg(n, 0), its(n, 0), used(nedges + 1, 0);
d41     vector<pii> s = { {src, -1} };
d41     //deg[src]++;
d41     vector<pii> ret;
d41     while (!s.empty()) {
d41         int u = s.back().first, &it = its[u];
d41         if (it == sz(g[u])) {
d41             ret.push_back(s.back());
d41             s.pop_back();
d41             continue;
d41         }
d41         auto& [nxt, id] = g[u][it++];
d41         if (!used[id]) {
d41             deg[u]--;
d41             deg[nxt]++;
d41             used[id] = 1;
d41             s.push_back({ nxt, id });
d41         }
d41     }
d41     for (int x : deg) {
d41         if (x < 0 || sz(ret) != (nedges + 1)) return {};
d41     }
d41     reverse(ret.begin(), ret.end());
d41     return ret;
d41 }

```

SCC.h

Description: Kosaraju algorithm for calculating strongly connected components. Components are ordered in topological order.

d41d8c, 36 lines

```

d41 struct SCC {
d41     int n, ncomp;
d41     vector<vector<int>> g, inv;
d41     vector<int> comp, vis, stk;
d41     SCC() {}
d41     SCC(int n)
d41         : n(n), ncomp(0), g(n), inv(n), comp(n, -1), vis(n) {}

d41     void dfs(int u) {
d41         vis[u] = 1;
d41         for (int v : g[u]) if (!vis[v]) dfs(v);
d41         stk.push_back(u);
d41     }
d41     void dfs_inv(int u) {

```

```
d41     comp[u] = ncomp;
d41     for (int v : inv[u]) {
d41         if (comp[v] == -1) dfs_inv(v);
d41     }
d41 }
d41 void solve() {
d41     for (int i = 0; i < n; i++) {
d41         if (!vis[i]) dfs(i);
d41     }
d41     reverse(all(stk));
d41     for (int u : stk) {
d41         if (comp[u] != -1) continue;
d41         dfs_inv(u);
d41         ncomp++;
d41     }
d41 }
d41 void add_edge(int a, int b) {
d41     g[a].push_back(b);
d41     inv[b].push_back(a);
d41 }
d41 }
```

TwoSAT.h

Usage: not A = ~A

```
"scc.h"
d41d8c, 37 lines
d41 struct TwoSat{
d41     int n;
d41     SCC scc;
d41     vector<int> value;
d41     vector<pii> e;
d41     TwoSat(int n) : n(n){}
d41     bool solve(){
d41         value.resize(n);
d41         scc = SCC(2*n);
d41         for(auto &x : e) scc.add_edge(x.first, x.second);
d41         scc.solve();
d41         for(int i=0; i<2*n; i++)
d41             if(scc.comp[i] == scc.comp[i^1]) return false;
d41         for(int i=0; i<n; i++)
d41             value[i] = scc.comp[id(i)] > scc.comp[id(~i)];
d41         return true;
d41     }
d41     void atMostOne(vector<int> &li){
d41         if(sz(li) <= 1) return;
d41         int cur = ~li[0];
d41         for(int i = 2; i < sz(li); i++) {
d41             int next = n++;
d41             addOr(cur, ~li[i]);
d41             addOr(cur, next);
d41             addOr(~li[i], next);
d41             cur = ~next;
d41         }
d41         addOr(cur, ~li[1]);
d41     }
d41     int id(int v) { return v < 0 ? (~v) * 2 ^ 1 : v * 2; }
d41     void add(int a, int b) { e.push_back({id(a), id(b)}); }
d41     void addOr(int a, int b) { add(~a, b); add(~b, a); }
d41     void addImp(int a, int b) { addOr(~a, b); }
d41     void addEqual(int a, int b){ addOr(a, ~b); addOr(~a, b); }
d41 }
d41 void isFalse(int a) { addImp(a, ~a); }
d41 };
```

7.4 Heuristics

MaxClique.h

Description: Quickly finds a maximum clique of a graph (given as symmetric bitset matrix; self-edges not allowed). Can be used to find a maximum independent set by finding a clique of the complement graph.

TwoSAT MaxClique Centroid HLD

Time: Runs in about 1s for n=155 and worst case random graphs (p=.90).
Runs faster for sparse graphs.

d41d8c, 53 lines

```
d41     using vb = vector<bitset<200>>;
d41     struct Maxclique {
d41         double limit=0.025, pk=0;
d41         struct Vertex { int i, d=0; };
d41         using vv = vector<Vertex>;
d41         vb e;
d41         vv V;
d41         vector<vector<int>> C;
d41         vector<int> qmax, q, S, old;
d41         void init(vv& r) {
d41             for (auto& v : r) v.d = 0;
d41             for (auto& v : r) for (auto j : r) v.d += e[v.i][j.i];
d41             sort(all(r), [](auto a, auto b) { return a.d > b.d; });
d41             int mxd = r[0].d;
d41             for(int i=0; i<sz(r); i++) r[i].d = min(i, mxd) + 1;
d41         }
d41         void expand(vv& R, int lev = 1) {
d41             S[lev] += S[lev - 1] - old[lev];
d41             old[lev] = S[lev - 1];
d41             while (sz(R)) {
d41                 if (sz(q) + R.back().d <= sz(qmax)) return;
d41                 q.push_back(R.back().i);
d41                 vv T;
d41                 for(auto v : R)
d41                     if (e[R.back().i][v.i]) T.push_back({v.i});
d41                 if (sz(T)) {
d41                     if (S[lev]++ / ++pk < limit) init(T);
d41                     int j = 0, mxk = 1, mnk = max(sz(qmax)-sz(q)+1, 1);
d41                     C[1].clear(), C[2].clear();
d41                     for (auto v : T) {
d41                         int k = 1;
d41                         auto f = [&](int i) { return e[v.i][i]; };
d41                         while (any_of(all(C[k]), f)) k++;
d41                         if (k > mxk) mxk = k, C[mxk + 1].clear();
d41                         if (k < mnk) T[j++].i = v.i;
d41                         C[k].push_back(v.i);
d41                     }
d41                     if (j > 0) T[j - 1].d = 0;
d41                     for(int k=mnk; k<mxk + 1; k++) {
d41                         for (int i : C[k])
d41                             T[j].i = i, T[j++].d = k;
d41                     }
d41                     expand(T, lev + 1);
d41                 } else if (sz(q) > sz(qmax)) qmax = q;
d41                 q.pop_back(), R.pop_back();
d41             }
d41         }
d41         vector<int> maxClique(){ init(V),expand(V); return qmax;
d41         Maxclique(vb conn) : e(conn),C(sz(e)+1),S(sz(C)),old(S){
d41             for(int i=0; i<sz(e); i++) V.push_back({i});
d41         }
d41     };
d41 }
```

7.5 Trees

Centroid.h

Description: Call decomp(0) to solve, marked array should be initially set to zero.

Time: $\mathcal{O}(N \log N)$

d41d8c, 27 lines

```
d41     int tam[ms], marked[ms];
d41
d41     int calc_tam(int u, int p) {
d41         tam[u] = 1;
d41         for (int v : g[u]) {
d41             if (v != p && !marked[v]) tam[u] += calc_tam(v, u);
d41         }
d41     }
d41 }
```

```
d41     return tam[u];
d41 }
```

```
d41     int get_centroid(int u, int p, int tot) {
d41         for (int v : g[u]) {
d41             if (v != p && !marked[v] && (tam[v] > (tot / 2)))
d41                 return get_centroid(v, u, tot);
d41         }
d41         return u;
d41     }
// Cent is a child of P in the centroid tree
d41     void decomp(int u, int p = -1) {
d41         calc_tam(u, -1);
d41         int cent = get_centroid(u, -1, tam[u]);
d41         marked[cent] = 1;
d41         for (int v : g[cent]) {
d41             if (!marked[v]) decomp(v, cent);
d41         }
d41     }
d41 }
```

HLD.h

Description: If values are stored on edges, set EDGE = true and store each edge's value at the endpoint farther from the root (the deeper node).
rp[i] is the representative (head) of the heavy path containing node i: it is the node in that chain that is closest to the root.

d41d8c, 51 lines

```
d41     template<bool EDGE> struct HLD {
d41         int n, t;
d41         vector<vector<int>> g;
d41         vector<int> pai, rp, tam, pos, val, arr;
d41         Seg seg;
d41         HLD(int n, vector<vector<int>>& g, vector<int>& val)
d41             : n(n), t(0), g(g), pai(n), rp(n), tam(n, 1),
d41               pos(n), val(val), arr(n) {
d41             calc_tam(0, -1);
d41             dfs(0, -1);
d41             seg.build(arr);
d41         }
d41
d41         int calc_tam(int u, int p) {
d41             pai[u] = p;
d41             for (int v : g[u]) {
d41                 if (v == p) continue;
d41                 tam[u] += calc_tam(v, u);
d41                 if(tam[v] > tam[g[u][0]] || g[u][0] == p)
d41                     swap(g[u][0], v);
d41             }
d41             return tam[u];
d41         }
d41
d41         void dfs(int u, int p) {
d41             pos[u] = t++;
d41             arr[pos[u]] = val[u];
d41             for (int v : g[u]) {
d41                 if (v == p) continue;
d41                 rp[v] = (v == g[u][0] ? rp[u] : v);
d41                 dfs(v, u);
d41             }
d41         }
d41
d41         int query(int a, int b) { // query on the path from a
d41             to b
d41             int ans = 0; // neutral value
d41             while (rp[a] != rp[b]) {
d41                 if (pos[a] < pos[b]) swap(a, b);
d41                 ans = max(ans, seg.query(pos[rp[a]], pos[a]));
d41                 a = pai[rp[a]];
d41             }
d41             if(pos[a] > pos[b]) swap(a, b);
d41         }
d41     };
d41 }
```

```
d41     ans = max(ans, seg.query(pos[a] + EDGE, pos[b]));
d41     return ans;
d41   }

d41   void update(int a, int x) {
d41     seg.update(pos[a], x);
d41   }
d41 };
```

LCA.h

Description: LCA algorithm using binary lifting, `is_ancestor(a, b)` returns true if a is an ancestral of b and false otherwise.

Time: $\mathcal{O}(N \log N)$

d41d8c, 26 lines

```
d41 int tin[MAXN], tout[MAXN], timer=0;
d41 int up[MAXN][BITS];
d41 void dfs(int u, int p){
d41   tin[u] = timer++;
d41   up[u][0] = p;
d41   for (int i=1; i<BITS; i++) {
d41     up[u][i] = up[up[u][i-1]][i-1];
d41   }
d41   for (int v: g[u]) if (v != p) dfs(v, u);
d41   tout[u] = timer;
d41 }
```

```
d41 bool is_ancestor(int u, int v){
d41   return (tin[u] <= tin[v] && tout[u] >= tout[v]);
d41 }
```

```
d41 int lca(int u, int v){
d41   if (is_ancestor(u, v)) return u;
d41   if (is_ancestor(v, u)) return v;
d41   for (int i=BITS-1; i>=0; i--) {
d41     if (up[u][i] && !is_ancestor(up[u][i], v)) {
d41       u = up[u][i];
d41     }
d41   }
d41   return up[u][0];
d41 }
```

VirtualTree.h

Description: Given a rooted tree and a subset S of nodes, compute the minimal subtree that contains all the nodes by adding all (at most $|S| - 1$) pairwise LCA's and compressing edges. `virt[u]` is the adjacency list of the virtual tree: it stores pairs (v, dist) , where v is a neighbor of u in the virtual tree and dist is the distance between u and v in the original tree.

Time: $\mathcal{O}(|S| \log |S|)$

LCA.h d41d8c, 24 lines

```
d41 vector<pair<int, int>> virt(ms);
d41

d41 void build_virt(vector<int>& v) {
d41   auto cmp = [&](int i, int j){ return tin[i] < tin[j]; };
d41   sort(all(v), cmp);
d41   for (int i = 0, n = sz(v); i + 1 < n; i++)
d41     v.push_back(lca(v[i], v[i + 1]));
d41   sort(all(v), cmp);
d41   v.erase(unique(all(v)), v.end());
d41   stack<int> st;
d41   for (auto u : v) {
d41     if (st.empty())
d41       st.push(u);
d41     else {
d41       while(sz(st) && !is_ancestor(st.top(), u)) st.pop();
d41       int p = st.top();
d41       virt[p].emplace_back(u, abs(lvl[u] - lvl[p]));
d41       virt[u].emplace_back(p, abs(lvl[u] - lvl[p]));
d41       st.push(u);
d41     }
d41   }
d41 }
```

```
d41   }
d41 }
```

7.6 Math

7.6.1 Number of Spanning Trees

Create an $N \times N$ matrix mat , and for each edge $a \rightarrow b \in G$, do $\text{mat}[a][b]--$, $\text{mat}[b][b]++$ (and $\text{mat}[b][a]--$, $\text{mat}[a][a]++$ if G is undirected). Remove the i th row and column and take the determinant; this yields the number of directed spanning trees rooted at i (if G is undirected, remove any row/column).

7.6.2 Erdős–Gallai theorem

A simple graph with node degrees $d_1 \geq \dots \geq d_n$ exists iff $d_1 + \dots + d_n$ is even and for every $k = 1 \dots n$,

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k).$$

7.7 Planar Graphs

If G has k connected components, then $n - m + f = k + 1$.

Geometry (8)

8.1 Geometric primitives

Point.h

Description: Class to handle points in the plane. T can be e.g. double or long long. (Avoid int.)

d41d8c, 29 lines

```
d41 template <class T> int sgn(T x) { return (x > 0) - (x < 0)
; }
d41 template<class T>
d41 struct Point {
d41   typedef Point P;
d41   T x, y;
d41   explicit Point(T x=0, T y=0) : x(x), y(y) {}
d41   bool operator<(P p) const { return tie(x,y) < tie(p.x,p.y)
; }
d41   bool operator==(P p) const { return tie(x,y)==tie(p.x,p.y)
; }
d41   P operator+(P p) const { return P(x+p.x, y+p.y); }
d41   P operator-(P p) const { return P(x-p.x, y-p.y); }
d41   P operator*(T d) const { return P(x*d, y*d); }
d41   P operator/(T d) const { return P(x/d, y/d); }
d41   T dot(P p) const { return x*p.x + y*p.y; }
d41   T cross(P p) const { return x*p.y - y*p.x; }
d41   T cross(P a, P b) const { return (a-*this).cross(b-*this)
; }
d41   T dist2() const { return x*x + y*y; }
d41   double dist() const { return sqrt((double)dist2()); }
// angle to x-axis in interval [-pi, pi]
d41   double angle() const { return atan2(y, x); }
d41   P unit() const { return *this/dist(); } // makes dist()==1
d41   P perp() const { return P(-y, x); } // rotates +90
degrees
d41   P normal() const { return perp().unit(); }
// returns point rotated 'a' radians ccw around the
origin
d41   P rotate(double a) const {
return P(x*cos(a)-y*sin(a), x*sin(a)+y*cos(a)); }
```

```
d41   friend ostream& operator<<(ostream& os, P p) {
d41     return os << "(" << p.x << ", " << p.y << ")";
d41   };

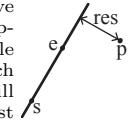
```

lineDistance.h

Description:

Returns the signed distance between point p and the line containing points a and b . Positive value on left side and negative on right as seen from a towards b . $a==b$ gives nan. P is supposed to be `Point<T>` or `Point3D<T>` where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long. Using `Point3D` will always give a non-negative distance. For `Point3D`, call `.dist` on the result of the cross product.

"Point.h" d41d8c, 5 lines



SegmentDistance.h

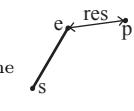
Description:

Returns the shortest distance between point p and the line segment from point s to e .

Usage: `Point<double> a, b(2,2), p(1,1);`

`bool onSegment = segDist(a,b,p) < 1e-10;`

"Point.h" d41d8c, 7 lines



```
d41 typedef Point<double> P;
d41 double segDist(P& s, P& e, P& p) {
d41   if (s==e) return (p-s).dist();
d41   auto d = (e-s).dist2(), t = min(d,max(.0,(p-s).dot(e-s)));
d41   return ((p-s)*d-(e-s)*t).dist()/d;
d41 }
```

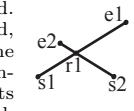
SegmentIntersection.h

Description:

If a unique intersection point between the line segments going from $s1$ to $e1$ and from $s2$ to $e2$ exists then it is returned. If no intersection point exists an empty vector is returned. If infinitely many exist a vector with 2 elements is returned, containing the endpoints of the common line segment. The wrong position will be returned if P is `Point<ll>` and the intersection point does not have integer coordinates. Products of three coordinates are used in intermediate steps so watch out for overflow if using int or long long.

Usage: `vector<P> inter = segInter(s1,e1,s2,e2);`

```
if (sz(inter)==1)
cout << "segments intersect at " << inter[0] << endl;
"Point.h", "OnSegment.h" d41d8c, 14 lines
```

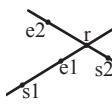


```
d41 template<class P> vector<P> segInter(P a, P b, P c, P d) {
d41   auto oa = c.cross(d, a), ob = c.cross(d, b),
oc = a.cross(b, c), od = a.cross(b, d);
// Checks if intersection is single non-endpoint point.
d41   if (sgn(oa) * sgn(ob) < 0 && sgn(oc) * sgn(od) < 0)
d41     return {(a * ob - b * oa) / (ob - oa)};
d41   set<P> s;
d41   if (onSegment(c, d, a)) s.insert(a);
d41   if (onSegment(c, d, b)) s.insert(b);
d41   if (onSegment(a, b, c)) s.insert(c);
d41   if (onSegment(a, b, d)) s.insert(d);
d41   return {all(s)};
d41 }
```

lineIntersection.h

Description:

If a unique intersection point of the lines going through s_1, e_1 and s_2, e_2 exists $\{1, \text{point}\}$ is returned. If no intersection point exists $\{0, (0,0)\}$ is returned and if infinitely many exists $\{-1, (0,0)\}$ is returned. The wrong position will be returned if P is Point<ll> and the intersection point does not have integer coordinates. Products of three coordinates are used in intermediate steps so watch out for overflow if using int or ll.

**Usage:** auto res = lineInter(s1,e1,s2,e2);

```
if (res.first == 1)
cout << "intersection point at " << res.second << endl;
"Point.h"
```

```
d41 template<class P>
d41 pair<int, P> lineInter(P s1, P e1, P s2, P e2) {
d41 auto d = (e1 - s1).cross(e2 - s2);
d41 if (d == 0) // if parallel
d41 return {-1, (0,0)};
d41 auto p = s2.cross(e1, e2), q = s2.cross(e2, s1);
d41 return {1, (s1 * p + e1 * q) / d};
d41 }
```

sideOf.h

Description: Returns where p is as seen from s towards e. $1/0/-1 \Leftrightarrow$ left/on line/right. If the optional argument *eps* is given 0 is returned if p is within distance *eps* from the line. P is supposed to be Point<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long.

Usage: bool left = sideOf(p1,p2,q)==1;

```
"Point.h"
```

```
d41 template<class P>
d41 int sideOf(P s, P e, P p) { return sgn(s.cross(e, p)); }

d41 template<class P>
d41 int sideOf(const P& s, const P& e, const P& p, double eps) {
d41 auto a = (e-s).cross(p-s);
d41 double l = (e-s).dist()*eps;
d41 return (a > l) - (a < -l);
d41 }
```

OnSegment.h

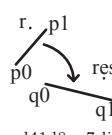
Description: Returns true iff p lies on the line segment from s to e. Use $(\text{segDist}(s, e, p) <= \text{epsilon})$ instead when using Point<double>.

```
"Point.h"
```

```
d41 template<class P> bool onSegment(P s, P e, P p) {
d41 return p.cross(s, e) == 0 && (s - p).dot(e - p) <= 0;
d41 }
```

linearTransformation.h**Description:**

Apply the linear transformation (translation, rotation and scaling) which takes line p_0-p_1 to line q_0-q_1 to point r.



```
"Point.h"
```

```
d41 typedef Point<double> P;
d41 P linearTransformation(const P& p0, const P& p1,
d41 const P& q0, const P& q1, const P& r) {
d41 P dp = p1-p0, dq = q1-q0, num(dp.cross(dq), dp.dot(dq));
d41 return q0 + P((r-p0).cross(num), (r-p0).dot(num))/dp.dist
    2();
d41 }
```

LineProjectionReflection.h

Description: Projects point p onto line ab. Set refl=true to get reflection of point p across line ab instead. The wrong point will be returned if P is an integer point and the desired point doesn't have integer coordinates. Products of three coordinates are used in intermediate steps so watch out for overflow.

```
"Point.h"
```

d41d8c, 6 lines

```
d41 template<class P>
d41 P lineProj(P a, P b, P p, bool refl=false) {
d41 P v = b - a;
d41 return p - v.perp()*(1+refl)*v.cross(p-a)/v.dist2();
d41 }
```

Angle.h

Description: A class for ordering angles (as represented by int points and a number of rotations around the origin). Useful for rotational sweeping. Sometimes also represents points or vectors.

Usage: vector<Angle> v = {w[0], w[0].t360(), ...}; // sorted
int j = 0; rep(i, 0, n) { while (v[j] < v[i].t180()) ++j; }
// sweeps j such that (j-i) represents the number of positively
oriented triangles with vertices at 0 and i

d41d8c, 36 lines

```
d41 struct Angle {
d41 int x, y;
d41 int t;
d41 Angle(int x, int y, int t=0) : x(x), y(y), t(t) {}
d41 Angle operator-(Angle b) const { return {x-b.x, y-b.y, t} }
d41 int half() const {
d41 assert(x || y);
d41 return y < 0 || (y == 0 && x < 0);
d41 }
d41 Angle t90() const { return {-y, x, t + (half() && x >= 0) ? 1 : 0} }
d41 Angle t180() const { return {-x, -y, t + half()} };
d41 Angle t360() const { return {x, y, t + 1}; }
d41 bool operator<(Angle a, Angle b) {
d41 // add a.dist2() and b.dist2() to also compare distances
d41 return make_tuple(a.t, a.half()), a.y * (11)b.x) <
d41 make_tuple(b.t, b.half()), a.x * (11)b.y);
d41 }
```

// Given two points, this calculates the smallest angle
between
// them, i.e., the angle that covers the defined line
segment.

```
d41 pair<Angle, Angle> segmentAngles(Angle a, Angle b) {
d41 if (b < a) swap(a, b);
d41 return (b < a.t180() ?
d41 make_pair(a, b) : make_pair(b, a.t360()));
d41 }
d41 Angle operator+(Angle a, Angle b) { // point a + vector b
d41 Angle r(a.x + b.x, a.y + b.y, a.t);
d41 if (a.t180() < r.t) r.t--;
d41 return r.t180() < a ? r.t360() : r;
d41 }
```

```
d41 Angle angleDiff(Angle a, Angle b) { // angle b - angle a
d41 int tu = b.t - a.t; a.t = b.t;
d41 return (a.x*b.x + a.y*b.y, a.x*b.y - a.y*b.x, tu - (b < a) ? 1 : 0);
d41 }
```

HalfPlane.h

Description: Computes the intersection of a set of half-planes. Half-planes are sorted by angle and processed with a deque, removing redundant or conflicting constraints. Parallel half-planes are handled explicitly. Returns the convex polygon of the intersection, or empty if infeasible.

Time: $\mathcal{O}(n \log n)$

```
"Point.h"
```

```
d41 using ld = long double;
d41 using P = Point<ld>;
d41 struct Hp { // Half plane struct
d41 // 'p' is a passing point of the line and 'pq' is the
d41 // direction vector of the line.
d41 P p, pq;
d41 ld angle;
d41 Hp() {}
d41 Hp(const P& a, const P& b) : p(a), pq(b - a) {
d41 angle = atan2l(pq.y, pq.x);
d41 }
d41 bool out(const P& r) { return pq.cross(r - p) < -eps; }
d41 bool operator < (const Hp& e) const {
d41 return angle < e.angle;
d41 }
d41 friend P inter(const Hp& s, const Hp& t) {
d41 ld alpha = (t.p - s.p).cross(t.pq) / s.pq.cross(t.pq);
d41 return s.p + (s.pq * alpha);
d41 }
d41 };
d41 vector<P> hp_intersect(vector<Hp>& H) {
d41 P box[4] = { P(-inf, -inf), P(-inf, inf),
d41 P(inf, -inf), P(inf, inf) };

d41 for(int i = 0; i < 4; i++) {
d41 Hp aux(box[i], box[(i+1) % 4]);
d41 H.push_back(aux);
d41 }
d41 sort(all(H));
d41 deque<Hp> dq;
d41 int len = 0;
d41 for(int i = 0; i < sz(H); i++) {
d41 while(len > 1 && H[i].out(inter(dq[len-1], dq[len-2]))) {
d41 dq.pop_back();
d41 --len;
d41 }
d41 while (len > 1 && H[i].out(inter(dq[0], dq[1]))) {
d41 dq.pop_front();
d41 --len;
d41 }
d41 if(len && fabsl(H[i].pq.cross(dq[len-1].pq)) < eps) {
d41 if (H[i].pq.dot(dq[len-1].pq) < 0.0)
d41 return vector<P>();
d41 if (H[i].out(dq[len-1].p)) {
d41 dq.pop_back();
d41 --len;
d41 } else continue;
d41 }
d41 dq.push_back(H[i]);
d41 ++len;
d41 }
d41 while(len > 2 && dq[0].out(inter(dq[len-1], dq[len-2]))) {
d41 dq.pop_back();
d41 --len;
d41 }
d41 while (len > 2 && dq[len-1].out(inter(dq[0], dq[1]))) {
d41 dq.pop_front();
d41 --len;
d41 }
d41 if (len < 3) return vector<P>();
vector<P> ret(len);
for(int i = 0; i + 1 < len; i++) {
d41 ret[i] = inter(dq[i], dq[i+1]);
d41 }
```

```
d41     }
d41     ret.back() = inter(dq[len-1], dq[0]);
d41     return ret;
d41 }
```

8.2 Circles

CircleIntersection.h

Description: Computes the pair of points at which two circles intersect. Returns false in case of no intersection.

```
"Point.h"
d41d8c, 12 lines
d41 typedef Point<double> P;
d41 bool circleInter(P a, P b, double r1, double r2, pair<P, P*>
out) {
d41     if (a == b) { assert(r1 != r2); return false; }
d41     P vec = b - a;
d41     double d2 = vec.dist2(), sum = r1+r2, dif = r1-r2;
d41     double p = (d2 + r1*r1 - r2*r2)/(d2*2), h2 = r1*r1 - p*p*
d2;
d41     if (sum*sum < d2 || dif*dif > d2) return false;
d41     P mid = a + vec*p, per = vec.perp() * sqrt(fmax(0, h2) /
d2);
d41     *out = {mid + per, mid - per};
d41     return true;
d41 }
```

CircleTangents.h

Description: Finds the external tangents of two circles, or internal if r2 is negated. Can return 0, 1, or 2 tangents – 0 if one circle contains the other (or overlaps it, in the internal case, or if the circles are the same); 1 if the circles are tangent to each other (in which case .first = .second and the tangent line is perpendicular to the line between the centers). .first and .second give the tangency points at circle 1 and 2 respectively. To find the tangents of a circle with a point set r2 to 0.

```
"Point.h"
d41d8c, 14 lines
d41 template<class P>
d41 vector<pair<P, P>> tangents(P c1, double r1, P c2, double
r2) {
d41     P d = c2 - c1;
d41     double dr = r1 - r2, d2 = d.dist2(), h2 = d2 - dr * dr;
d41     if (d2 == 0 || h2 < 0) return {};
d41     vector<pair<P, P>> out;
d41     for (double sign : {-1, 1}) {
d41         P v = (d * dr + d.perp() * sqrt(h2) * sign) / d2;
d41         out.push_back({c1 + v * r1, c2 + v * r2});
d41     }
d41     if (h2 == 0) out.pop_back();
d41     return out;
d41 }
```

CircleLine.h

Description: Finds the intersection between a circle and a line. Returns a vector of either 0, 1, or 2 intersection points. P is intended to be Point<double>.

```
"Point.h"
d41d8c, 10 lines
d41 template<class P>
d41 vector<P> circleLine(P c, double r, P a, P b) {
d41     P ab = b - a, p = a + ab * (c-a).dot(ab) / ab.dist2();
d41     double s = a.cross(b, c), h2 = r*r - s*s / ab.dist2();
d41     if (h2 < 0) return {};
d41     if (h2 == 0) return {p};
d41     P h = ab.unit() * sqrt(h2);
d41     return {p - h, p + h};
d41 }
```

CirclePolygonIntersection.h
Description: Returns the area of the intersection of a circle with a ccw polygon.
Time: $\mathcal{O}(n)$

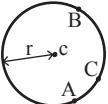
```
"../../../../content/geometry/Point.h"
d41d8c, 20 lines
d41 typedef Point<double> P;
d41 #define arg(p, q) atan2(p.cross(q), p.dot(q))
d41 double circlePoly(P c, double r, vector<P> ps) {
d41     auto tri = [&](P p, P q) {
d41         auto r2 = r * r / 2;
d41         P d = q - p;
d41         auto a = d.dot(p)/d.dist2(), b = (p.dist2()-r*r)/d.dist
2();
d41         auto det = a * a - b;
d41         if (det <= 0) return arg(p, q) * r2;
d41         auto s = max(0., -a-sqrt(det)), t = min(1., -a+sqrt(det)
);
d41         if (t < 0 || 1 <= s) return arg(p, q) * r2;
d41         P u = p + d * s, v = q + d * (t-1);
d41         return arg(p,u) * r2 + u.cross(v)/2 + arg(v,q) * r2;
d41     };
d41     auto sum = 0.0;
d41     rep(i,0,sz(ps))
d41         sum += tri(ps[i] - c, ps[(i + 1) % sz(ps)] - c);
d41     return sum;
d41 }
```

circumcircle.h

Description:

The circumcircle of a triangle is the circle intersecting all three vertices. ccRadius returns the radius of the circle going through points A, B and C and ccCenter returns the center of the same circle.

```
"Point.h"
d41d8c, 10 lines
d41 typedef Point<double> P;
d41 double ccRadius(const P& A, const P& B, const P& C) {
d41     return (B-A).dist()*(C-B).dist()*(A-C).dist()/
abs((B-A).cross(C-A))/2;
d41 }
d41 P ccCenter(const P& A, const P& B, const P& C) {
d41     P b = C-A, c = B-A;
d41     return A + ((b*c.dist2())-c*b.dist2()).perp()/b.cross(c)/2;
d41 }
```



MinimumEnclosingCircle.h

Description: Computes the minimum circle that encloses a set of points.

Time: expected $\mathcal{O}(n)$

```
"circumcircle.h"
d41d8c, 18 lines
d41 pair<P, double> mec(vector<P> ps) {
d41     shuffle(all(ps), mt19937(time(0)));
d41     P o = ps[0];
d41     double r = 0, EPS = 1 + 1e-8;
d41     rep(i,0,sz(ps)) if ((o - ps[i]).dist() > r * EPS) {
d41         o = ps[i], r = 0;
d41         rep(j,0,i) if ((o - ps[j]).dist() > r * EPS) {
d41             o = (ps[i] + ps[j]) / 2;
d41             r = (o - ps[i]).dist();
d41             rep(k,0,j) if ((o - ps[k]).dist() > r * EPS) {
d41                 o = ccCenter(ps[i], ps[j], ps[k]);
d41                 r = (o - ps[i]).dist();
d41             }
d41         }
d41     }
d41     return {o, r};
d41 }
```

8.3 Polygons

InsidePolygon.h

Description: Returns true if p lies within the polygon. If strict is true, it returns false for points on the boundary. The algorithm uses products in intermediate steps so watch out for overflow.

Usage: vector<P> v = {P{4,4}, P{1,2}, P{2,1}};

bool in = inPolygon(v, P{3, 3}, false);

Time: $\mathcal{O}(n)$

```
"Point.h", "OnSegment.h", "SegmentDistance.h"
d41d8c, 12 lines
d41 template<class P>
d41 P inPolygon(vector<P> &p, P a, bool strict = true) {
d41     int cnt = 0, n = sz(p);
d41     rep(i,0,n) {
d41         P q = p[(i + 1) % n];
d41         if (onSegment(p[i], q, a)) return !strict;
d41         //or: if (segDist(p[i], q, a) <= eps) return !strict;
d41         cnt ^= ((a.y<p[i].y) - (a.y<q.y)) * a.cross(p[i], q) >
0;
d41     }
d41     return cnt;
d41 }
```

PolygonArea.h

Description: Returns twice the signed area of a polygon. Clockwise enumeration gives negative area. Watch out for overflow if using int as T!

```
"Point.h"
d41d8c, 7 lines
```

```
d41 template<class T>
d41 T polygonArea2(vector<Point<T>> &v) {
d41     T a = v.back().cross(v[0]);
d41     rep(i,0,sz(v)-1) a += v[i].cross(v[i+1]);
d41     return a;
d41 }
```

PolygonCenter.h

Description: Returns the center of mass for a polygon.

Time: $\mathcal{O}(n)$

```
"Point.h"
d41d8c, 10 lines
d41 typedef Point<double> P;
d41 P polygonCenter(const vector<P>& v) {
d41     P res(0, 0); double A = 0;
d41     for (int i = 0, j = sz(v) - 1; i < sz(v); j = i++) {
d41         res = res + (v[i] + v[j]) * v[j].cross(v[i]);
d41         A += v[j].cross(v[i]);
d41     }
d41     return res / A / 3;
d41 }
```

PolygonCut.h

Description:

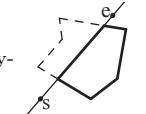
Returns a vector with the vertices of a polygon with everything to the left of the line going from s to e cut away.

Usage: vector<P> p = ...;

p = polygonCut(p, P(0,0), P(1,0));

Time:

```
"Point.h"
d41d8c, 14 lines
d41 typedef Point<double> P;
d41 vector<P> polygonCut(const vector<P>& poly, P s, P e) {
d41     vector<P> res;
d41     rep(i,0,sz(poly)) {
d41         P cur = poly[i], prev = i ? poly[i-1] : poly.back();
d41         auto a = s.cross(e, cur), b = s.cross(e, prev);
d41         if ((a < 0) != (b < 0))
d41             res.push_back(cur + (prev - cur) * (a / (a - b)));
d41         if (a < 0)
d41             res.push_back(cur);
d41     }
d41     return res;
d41 }
```



PolygonUnion.h

Description: Calculates the area of the union of n polygons (not necessarily convex). The points within each polygon must be given in CCW order. (Epsilon checks may optionally be added to sideOf/sgn, but shouldn't be needed.)

Time: $\mathcal{O}(N^2)$, where N is the total number of points

```
"Point.h", "sideOf.h" d41d8c, 34 lines
d41 typedef Point<double> P;
d41 double rat(P a, P b) { return sgn(b.x) ? a.x/b.x : a.y/b.y;
d41 }
d41 double polyUnion(vector<vector<P>>& poly) {
d41     double ret = 0;
d41     rep(i, 0, sz(poly)) rep(v, 0, sz(poly[i])) {
d41         P A = poly[i][v], B = poly[i][(v + 1) % sz(poly[i])];
d41         vector<pair<double, int>> segs = {{0, 0}, {1, 0}};
d41         rep(j, 0, sz(poly)) if (i != j) {
d41             rep(u, 0, sz(poly[j])) {
d41                 P C = poly[j][u], D = poly[j][(u + 1) % sz(poly[j])];
d41                 int sc = sideOf(A, B, C), sd = sideOf(A, B, D);
d41                 if (sc != sd) {
d41                     double sa = C.cross(D, A), sb = C.cross(D, B);
d41                     if (min(sc, sd) < 0)
d41                         segs.emplace_back(sa / (sa - sb), sgn(sc - sd));
d41                 };
d41             } else if (!sc && !sd && j < i && sgn((B-A).dot(D-C)) > 0) {
d41                 segs.emplace_back(rat(C - A, B - A), 1);
d41                 segs.emplace_back(rat(D - A, B - A), -1);
d41             }
d41         }
d41     }
d41     sort(all(segs));
d41     for (auto& s : segs) s.first = min(max(s.first, 0.0), 1.0);
d41     double sum = 0;
d41     int cnt = segs[0].second;
d41     rep(j, 1, sz(segs)) {
d41         if (!cnt) sum += segs[j].first - segs[j - 1].first;
d41         cnt += segs[j].second;
d41     }
d41     ret += A.cross(B) * sum;
d41 }
d41 return ret / 2;
d41 }
```

ConvexHull.h

Description:

Returns a vector of the points of the convex hull in counter-clockwise order. Points on the edge of the hull between two other points are not considered part of the hull. If you want to keep the collinear points in the convex hull, change the comparison to $h[t-2].cross(h[t-1], p) < 0$ and the size of the vector h to $2 * sz(pts) + 1$.

Time: $\mathcal{O}(n \log n)$

```
"Point.h" d41d8c, 14 lines
d41 typedef Point<11> P;
d41 vector<P> convexHull(vector<P> pts) {
d41     if (sz(pts) <= 1) return pts;
d41     sort(all(pts));
d41     vector<P> h(sz(pts)+1);
d41     int s = 0, t = 0;
d41     for (int it = 2; it--; s = --t, reverse(all(pts)))
d41         for (P p : pts) {
d41             while (t >= s + 2 && h[t-2].cross(h[t-1], p) <= 0) t--;
d41             h[t++] = p;
d41         }
d41 }
```



```
d41     return {h.begin(), h.begin() + t - (t == 2 && h[0] == h[1])};
d41 }
```

HullDiameter.h

Description: Returns the two points with max distance on a convex hull (ccw, no duplicate/collinear points).

Time: $\mathcal{O}(n)$

```
"Point.h" d41d8c, 13 lines
d41 typedef Point<11> P;
d41 array<P, 2> hullDiameter(vector<P> S) {
d41     int n = sz(S), j = n < 2 ? 0 : 1;
d41     pair<11, array<P, 2>> res({0, {S[0], S[0]}});
d41     rep(i, 0, j)
d41         for (; j = (j + 1) % n) {
d41             res = max(res, {(S[i] - S[j]).dist2(), {S[i], S[j]}});
d41         };
d41         if ((S[(j + 1) % n] - S[j]).cross(S[i + 1] - S[i]) >= 0)
d41             break;
d41     }
d41     return res.second;
d41 }
```

PointInsideHull.h

Description: Determine whether a point t lies inside a convex hull (CCW order, with no collinear points). Returns true if point lies within the hull. If strict is true, points on the boundary aren't included.

Time: $\mathcal{O}(\log N)$

```
"Point.h", "sideOf.h", "OnSegment.h" d41d8c, 15 lines
d41 typedef Point<11> P;

d41 bool inHull(const vector<P>& l, P p, bool strict = true) {
d41     int a = 1, b = sz(l) - 1, r = !strict;
d41     if (sz(l) < 3) return r && onSegment(l[0], l.back(), p);
d41     if (sideOf(l[0], l[a], l[b]) > 0) swap(a, b);
d41     if (sideOf(l[0], l[a], p) >= r || sideOf(l[0], l[b], p) <= -r)
d41         return false;
d41     while (abs(a - b) > 1) {
d41         int c = (a + b) / 2;
d41         (sideOf(l[0], l[c], p) > 0 ? b : a) = c;
d41     }
d41     return sgn(l[a].cross(l[b], p)) < r;
d41 }
```

LineHullIntersection.h

Description: Line-convex polygon intersection. The polygon must be ccw and have no collinear points. lineHull(line, poly) returns a pair describing the intersection of a line with the polygon: • $(-1, -1)$ if no collision, • $(i, -1)$ if touching the corner i , • (i, i) if along side $(i, i+1)$, • (i, j) if crossing sides $(i, i+1)$ and $(j, j+1)$. In the last case, if a corner i is crossed, this is treated as happening on side $(i, i+1)$. The points are returned in the same order as the line hits the polygon. extrVertex returns the point of a hull with the max projection onto a line.

Time: $\mathcal{O}(\log n)$

```
"Point.h" d41d8c, 40 lines
d41 #define cmp(i, j) sgn(dir.perp()).cross(poly[(i)%n]-poly[(j)%n])
d41 #define extr(i) cmp(i + 1, i) >= 0 && cmp(i, i - 1 + n) < 0
d41 template <class P> int extrVertex(vector<P>& poly, P dir)
{
    int n = sz(poly), lo = 0, hi = n;
    if (extr(0)) return 0;
    while (lo + 1 < hi) {
        int m = (lo + hi) / 2;
```

```
d41     if (extr(m)) return m;
d41     int ls = cmp(lo + 1, lo), ms = cmp(m + 1, m);
d41     (ls < ms || (ls == ms && ls == cmp(lo, m))) ? hi : lo) = m;
d41 }
d41 return lo;
d41 }
```

```
#define cmpL(i) sgn(a.cross(poly[i], b))
template <class P>
array<int, 2> lineHull(P a, P b, vector<P>& poly) {
    int endA = extrVertex(poly, (a - b).perp());
    int endB = extrVertex(poly, (b - a).perp());
    if (cmpL(endA) < 0 || cmpL(endB) > 0)
        return {-1, -1};
    array<int, 2> res;
    rep(i, 0, 2) {
        int lo = endB, hi = endA, n = sz(poly);
        while ((lo + 1) % n != hi) {
            int m = ((lo + hi + (lo < hi ? 0 : n)) / 2) % n;
            (cmpL(m) == cmpL(endB) ? lo : hi) = m;
        }
        res[i] = (lo + !cmpL(hi)) % n;
        swap(endA, endB);
    }
    if (res[0] == res[1]) return {res[0], -1};
    if (!cmpL(res[0]) && !cmpL(res[1]))
        switch (res[0] - res[1] + sz(poly) + 1) % sz(poly)) {
            case 0: return {res[0], res[0]};
            case 2: return {res[1], res[1]};
        }
    return res;
}
```

Minkowski.h

Description: Computes the Minkowski sum of two convex polygons. Polygons must be convex and given in CCW order. Returns the vertices of the Minkowski sum polygon in CCW order.

Time: $\mathcal{O}(n + m)$

```
"Point.h" d41d8c, 24 lines
d41 using P = Point<11>;
vector<P> minkowski(vector<P> p, vector<P> q) {
    auto fix = [](vector<P>& A) {
        int pos = 0;
        for (int i = 1; i < sz(A); i++) {
            if (A[i].y < A[pos].y || (A[i].y == A[pos].y && A[i].x < A[pos].x))
                pos = i;
        }
        rotate(A.begin(), A.begin() + pos, A.end());
        A.push_back(A[0]), A.push_back(A[1]);
    };
    fix(p), fix(q);
    vector<P> result;
    int i = 0, j = 0;
    while (i < sz(p) - 2 || j < sz(q) - 2) {
        result.push_back(p[i] + q[j]);
        auto cross = (p[i + 1] - p[i]).cross(q[j + 1] - q[j]);
        if (cross >= 0 && i < sz(p) - 2) i++;
        if (cross <= 0 && j < sz(q) - 2) j++;
    }
    return result;
}
```

Extreme.h

Description: Finds an extreme vertex of a convex polygon according to a unimodal comparator. The comparator defines a total order along the polygon (given in CCW order).

Time: $\mathcal{O}(\log n)$

"Point.h" d41d8c, 26 lines

```
d41 using P = Point<ll>;
d41 int extreme(vector<P> &pol, const function<bool>(P, P)&
  cmp) {
d41   int n = pol.size();
d41   auto extr = [&](int i, bool& cur_dir) {
d41     cur_dir = cmp(pol[(i+1)%n], pol[i]);
d41     return !cur_dir and !cmp(pol[(i+n-1)%n], pol[i]);
d41   };
d41   bool last_dir, cur_dir;
d41   if (extr(0, last_dir)) return 0;
d41   int l = 0, r = n;
d41   while (l+1 < r) {
d41     int m = (l+r)/2;
d41     if (extr(m, cur_dir)) return m;
d41     bool rel_dir = cmp(pol[m], pol[1]);
d41     if (!last_dir and cur_dir) or
d41       (last_dir == cur_dir and rel_dir == cur_dir)) {
d41       l = m;
d41       last_dir = cur_dir;
d41     } else r = m;
d41   }
d41   return l;
d41 }
d41 int max_dot(vector<P> &pol, P v) {
d41   return extreme([&](P p, P q) { return p.dot(v) > q.dot(v);
}); });
d41 }
```

Tangents.h

Description: Finds the left and right tangent points from an external point p to a convex polygon given in CCW order. A tangent point is a vertex where the segment $p \rightarrow v$ touches the polygon without intersecting its interior, defining the limits of visibility from p . Returns the indices of the left and right tangent vertices.

Time: $\mathcal{O}(\log n)$

"Point.h", "Extreme.h" d41d8c, 11 lines

```
d41 using P = Point<ll>;
d41
d41 bool ccw(P p, P q, P r) {
d41   return (q-p).cross(r-q) > 0;
d41 }
d41 pair<int, int> tangents(vector<P> &pol, P p) {
d41   auto L = [&](P q, P r) { return ccw(p, r, q); };
d41   auto R = [&](P q, P r) { return ccw(p, q, r); };
d41   return {extreme(pol, L), extreme(pol, R)};
d41 }
```

8.4 Misc. Point Set Problems

ClosestPair.h

Description: Finds the closest pair of points.

Time: $\mathcal{O}(n \log n)$

"Point.h" d41d8c, 18 lines

```
d41 typedef Point<ll> P;
d41 pair<P, P> closest(vector<P> v) {
d41   assert(sz(v) > 1);
d41   set<P> S;
d41   sort(all(v), [](P a, P b) { return a.y < b.y; });
d41   pair<ll, pair<P, P>> ret{LLONG_MAX, {P(), P()}};
d41   int j = 0;
d41   for (P p : v) {
d41     P d{1 + (ll)sqrt(ret.first)}, 0;
```

```
d41       while (v[j].y <= p.y - d.x) S.erase(v[j++]);
d41       auto lo = S.lower_bound(p - d), hi = S.upper_bound(p +
  d);
d41       for (; lo != hi; ++lo)
d41         ret = min(ret, {{*lo - p}.dist2(), {*lo, p}});
d41       S.insert(p);
d41     }
d41   }
d41   return ret.second;
d41 }
```

ManhattanMST.h

Description: Given N points, returns up to 4^*N edges, which are guaranteed to contain a minimum spanning tree for the graph with edge weights $w(p, q) = -p.x - q.x - + -p.y - q.y$. Edges are in the form (distance, src, dst). Use a standard MST algorithm on the result to find the final MST.

Time: $\mathcal{O}(N \log N)$

"Point.h" d41d8c, 24 lines

```
d41 typedef Point<int> P;
d41 vector<array<int, 3>> manhattanMST(vector<P> ps) {
d41   vi id{sz(ps)};
d41   iota(all(id), 0);
d41   vector<array<int, 3>> edges;
d41   rep(k, 0, 4) {
d41     sort(all(id), [&](int i, int j) {
d41       return (ps[i]-ps[j]).x < (ps[j]-ps[i]).y;
}); map<int, int> sweep;
d41     for (int i : id) {
d41       for (auto it = sweep.lower_bound(-ps[i].y);
d41             it != sweep.end(); sweep.erase(it++)) {
d41         int j = it->second;
d41         P d = ps[i] - ps[j];
d41         if (d.y > d.x) break;
d41         edges.push_back({d.y + d.x, i, j});
d41       }
d41       sweep[-ps[i].y] = i;
d41     }
d41     for (P& p : ps) if (k & 1) p.x = -p.x; else swap(p.x, p
.y);
d41   }
d41   return edges;
d41 }
```

kdTree.h

Description: KD-tree (2d, can be extended to 3d)

"Point.h" d41d8c, 64 lines

```
d41 typedef long long T;
d41 typedef Point<T> P;
d41 const T INF = numeric_limits<T>::max();
d41
d41 bool on_x(const P& a, const P& b) { return a.x < b.x; }
d41 bool on_y(const P& a, const P& b) { return a.y < b.y; }

d41 struct Node {
d41   P pt; // if this is a leaf, the single point in it
d41   T x0 = INF, x1 = -INF, y0 = INF, y1 = -INF; // bounds
d41   Node *first = 0, *second = 0;

d41   T distance(const P& p) { // min squared distance to a
  point
d41     T x = (p.x < x0 ? x0 : p.x > x1 ? x1 : p.x);
d41     T y = (p.y < y0 ? y0 : p.y > y1 ? y1 : p.y);
d41     return (P(x,y) - p).dist2();
d41   }

d41   Node(vector<P>&& vp) : pt(vp[0]) {
d41     for (P p : vp) {
d41       x0 = min(x0, p.x); x1 = max(x1, p.x);
d41       y0 = min(y0, p.y); y1 = max(y1, p.y);
```

```
d41     }
d41   }
d41   if (vp.size() > 1) {
// split on x if width >= height (not ideal...)
d41     sort(all(vp), x1 - x0 >= y1 - y0 ? on_x : on_y);
// divide by taking half the array for each child (not
  best performance with many duplicates in the
  middle)
d41     int half = sz(vp)/2;
d41     first = new Node({vp.begin(), vp.begin() + half});
d41     second = new Node({vp.begin() + half, vp.end()});
d41   }
d41 }
```

```
d41 struct KDTree {
d41   Node* root;
d41   KDTree(const vector<P>& vp) : root(new Node(all(vp))) {
d41 }
```

```
d41   pair<T, P> search(Node *node, const P& p) {
d41     if (!node->first) {
// uncomment if we should not find the point itself:
// if (p == node->pt) return {INF, P()};
d41     return make_pair((p - node->pt).dist2(), node->pt);
d41   }

```

```
d41   Node *f = node->first, *s = node->second;
d41   T bfirst = f->distance(p), bsec = s->distance(p);
d41   if (bfirst > bsec) swap(bsec, bfirst), swap(f, s);

```

```
d41   // search closest side first, other side if needed
d41   auto best = search(f, p);
d41   if (bsec < best.first)
d41     best = min(best, search(s, p));
d41   return best;
d41 }
```

```
// find nearest point to a point, and its squared
  distance
// (requires an arbitrary operator< for Point)
```

```
d41   pair<T, P> nearest(const P& p) {
d41     return search(root, p);
d41   }
d41 }
```

FastDelaunay.h

Description: Fast Delaunay triangulation. Each circumcircle contains none of the input points. There must be no duplicate points. If all points are on a line, no triangles will be returned. Should work for doubles as well, though there may be precision issues in 'circ'. Returns triangles in order {t[0][0], t[0][1], t[0][2], t[1][0], ...}, all counter-clockwise.

Time: $\mathcal{O}(n \log n)$

"Point.h" d41d8c, 89 lines

```
d41 typedef Point<ll> P;
d41 typedef struct Quad* Q;
d41 __int128_t l11; // (can be ll if coords are < 2e4)
d41 P arb(LLONG_MAX, LLONG_MAX); // not equal to any other
  point

d41 struct Quad {
d41   Q rot, o; P p = arb; bool mark;
d41   P& F() { return r()->p; }
d41   Q& r() { return rot->rot; }
d41   Q prev() { return rot->o->rot; }
d41   Q next() { return r()->prev(); }
d41 } *H;
```

```

d41 bool circ(P p, P a, P b, P c) { // is p in the
cicumcircle?
d41   ll1 p2 = p.dist2(), A = a.dist2()-p2,
d41   B = b.dist2()-p2, C = c.dist2()-p2;
d41   return p.cross(a,b)*C + p.cross(b,c)*A + p.cross(c,a)*B >
0;
d41 }
d41 Q makeEdge(P orig, P dest) {
d41   Q r = H ? H : new Quad{new Quad{new Quad{new Quad{0}}}};
d41   H = r->o; r->r()->r() = r;
d41   rep(i,0,4) r = r->rot, r->p = arb, r->o = i & 1 ? r : r->
r();
d41   r->p = orig; r->F() = dest;
d41   return r;
d41 }
d41 void splice(Q a, Q b) {
d41   swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
d41 }
Q connect(Q a, Q b) {
d41   Q q = makeEdge(a->F(), b->p);
d41   splice(q, a->next());
d41   splice(q->r(), b);
d41   return q;
d41 }

pair<Q,Q> rec(const vector<P>& s) {
d41   if (sz(s) <= 3) {
d41     Q a = makeEdge(s[0], s[1]), b = makeEdge(s[1], s.back());
d41   }
d41   if (sz(s) == 2) return { a, a->r() };
d41   splice(a->r(), b);
d41   auto side = s[0].cross(s[1], s[2]);
d41   Q c = side ? connect(b, a) : 0;
d41   return {side < 0 ? c->r() : a, side < 0 ? c : b->r()};
d41 }

#define H(e) e->F(), e->p
#define valid(e) (e->F().cross(H(base)) > 0)
d41 Q A, B, ra, rb;
d41 int half = sz(s) / 2;
d41 tie(ra, A) = rec({all(s) - half});
d41 tie(B, rb) = rec({sz(s) - half + all(s)});
d41 while ((B->p.cross(H(A)) < 0 && (A = A->next()) ||
d41         (A->p.cross(H(B)) > 0 && (B = B->r()->o)));
d41 Q base = connect(B->r(), A);
d41 if (A->p == ra->p) ra = base->r();
d41 if (B->p == rb->p) rb = base;

#define DEL(e, init, dir) Q e = init->dir; if (valid(e)) \
d41   while (cire(e->dir->F(), H(base), e->F())) { \
d41     Q t = e->dir; \
d41     splice(e, e->prev()); \
d41     splice(e->r(), e->r()->prev()); \
d41     e->o = H; H = e; e = t; \
d41   }
d41 for (;;) {
d41   DEL(LC, base->r(), o); DEL(RC, base, prev());
d41   if (!valid(LC) && !valid(RC)) break;
d41   if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC)))) \
d41     base = connect(RC, base->r());
d41   else
d41     base = connect(base->r(), LC->r());
d41 }
d41 return {ra, rb};

d41 vector<P> triangulate(vector<P> pts) {
d41   sort(all(pts)); assert(unique(all(pts)) == pts.end());
d41   if (sz(pts) < 2) return {};

```

```

d41   Q e = rec(pts).first;
d41   vector<Q> q = {e};
d41   int qi = 0;
d41   while (e->o->F()).cross(e->F(), e->p) < 0) e = e->o;
d41   #define ADD { Q c = e; do { c->mark = 1; pts.push_back(c->
p); \
d41     q.push_back(c->r()); c = c->next(); } while (c != e); } \
d41   ADD; pts.clear();
d41   while (qi < sz(q)) if (!(e = q[qi++])->mark) ADD;
d41   return pts;
d41 }

```

8.5 3D

PolyhedronVolume.h

Description: Magic formula for the volume of a polyhedron. Faces should point outwards.

```

d41 template<class V, class L>
d41 double signedPolyVolume(const V& p, const L& trilist) {
d41   double v = 0;
d41   for (auto i : trilist) v += p[i.a].cross(p[i.b]).dot(p[i.
c]);
d41   return v / 6;
d41 }

```

Point3D.h

Description: Class to handle points in 3D space. T can be e.g. double or long long.

```

d41 template<class T> struct Point3D {
d41   typedef Point3D P;
d41   typedef const P& R;
d41   T x, y, z;
d41   explicit Point3D(T x=0, T y=0, T z=0) : x(x), y(y), z(z)
{ }
d41   bool operator<(R p) const {
d41     return tie(x, y, z) < tie(p.x, p.y, p.z); }
d41   bool operator==(R p) const {
d41     return tie(x, y, z) == tie(p.x, p.y, p.z); }
d41   P operator+(R p) const { return P(x+p.x, y+p.y, z+p.z); }
d41   P operator-(R p) const { return P(x-p.x, y-p.y, z-p.z); }
d41   P operator*(T d) const { return P(x*d, y*d, z*d); }
d41   P operator/(T d) const { return P(x/d, y/d, z/d); }
d41   T dot(R p) const { return x*p.x + y*p.y + z*p.z; }
d41   P cross(R p) const {
d41     return P(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y - y*p.x);
d41   }
d41   T dist2() const { return x*x + y*y + z*z; }
d41   double dist() const { return sqrt((double)dist2()); }
//Azimuthal angle (longitude) to x-axis in interval [-pi,
pi]
d41   double phi() const { return atan2(y, x); }
//Zenith angle (latitude) to the z-axis in interval [0,
pi]
d41   double theta() const { return atan2(sqrt(x*x+y*y), z); }
d41   P unit() const { return *this/(T)dist(); } //makes dist()
=1
//returns unit vector normal to *this and p
d41   P normal(P p) const { return cross(p).unit(); }
//returns point rotated 'angle' radians ccw around axis
d41   P rotate(double angle, P axis) const {
d41     double s = sin(angle), c = cos(angle); P u = axis.unit
();
d41     return u*dot(u)*(1-c) + (*this)*c - cross(u)*s;
d41   }

```

3dHull.h

Description: Computes all faces of the 3-dimension hull of a point set. *No four points must be coplanar*, or else random results will be returned. All faces will point outwards.

Time: $\mathcal{O}(n^2)$

"Point3D.h" d41d8c, 50 lines

```

d41   typedef Point3D<double> P3;
d41   struct PR {
d41     void ins(int x) { (a == -1 ? a : b) = x; }
d41     void rem(int x) { (a == x ? a : b) = -1; }
d41     int cnt() { return (a != -1) + (b != -1); }
d41     int a, b;
d41   };
d41   struct F { P3 q; int a, b, c; };
d41   vector<F> hull3d(const vector<P3>& A) {
d41   assert(sz(A) >= 4);
d41   vector<vector<PR>> E(sz(A)), vector<PR>(sz(A), {-1, -1});
d41   #define E(x,y) E[f.x][f.y]
d41   vector<F> FS;
d41   auto mf = [&](int i, int j, int k, int l) {
d41     P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
d41     if (q.dot(A[l]) > q.dot(A[i])) {
d41       q = q * -1;
d41     f(q, i, j, k);
d41     E(a,b).ins(k); E(a,c).ins(j); E(b,c).ins(i);
d41     FS.push_back(f);
d41   };
d41   rep(i,0,4) rep(j,i+1,4) rep(k,j+1,4)
d41   mf(i, j, k, 6 - i - j - k);

d41   rep(i,4,sz(A)) {
d41     rep(j,0,sz(FS)) {
d41       F f = FS[j];
d41       if (f.q.dot(A[i]) > f.q.dot(A[f.a])) {
d41         E(a,b).rem(f.c);
d41         E(a,c).rem(f.b);
d41         E(b,c).rem(f.a);
d41         swap(FS[j--], FS.back());
d41         FS.pop_back();
d41       }
d41     }
d41     int nw = sz(FS);
d41     rep(j,0,nw) {
d41       F f = FS[j];
d41     }
d41   }
#define C(a, b, c) if (E(a,b).cnt() != 2) mf(f.a, f.b, i,
f.c);
d41   C(a, b, c); C(a, c, b); C(b, c, a);
d41   }
d41   for (F& it : FS) if ((A[it.b] - A[it.a]).cross(
d41     A[it.c] - A[it.a]).dot(it.q) <= 0) swap(it.c, it.b);
d41   return FS;
d41 };

```

sphericalDistance.h

Description: Returns the shortest distance on the sphere with radius radius between the points with azimuthal angles (longitude) $f_1(\phi_1)$ and $f_2(\phi_2)$ from x axis and zenith angles (latitude) $t_1(\theta_1)$ and $t_2(\theta_2)$ from z axis (0 = north pole). All angles measured in radians. The algorithm starts by converting the spherical coordinates to cartesian coordinates so that is what you have you can use only the two last rows. $dx \cdot radius$ is then the difference between the two points in the x direction and $d \cdot radius$ is the total distance between the points.

"sphericalDistance.h" d41d8c, 9 lines

```

d41   double sphericalDistance(double f1, double t1,
d41   double f2, double t2, double radius) {

```

```
d41 double dx = sin(t2)*cos(f2) - sin(t1)*cos(f1);
d41 double dy = sin(t2)*sin(f2) - sin(t1)*sin(f1);
d41 double dz = cos(t2) - cos(t1);
d41 double d = sqrt(dx*dx + dy*dy + dz*dz);
d41 return radius*2*asin(d/2);
d41 }
```

Strings (9)

AhoCorasick.h

d41d8c, 46 lines

```
d41 int trie[ms][sigma], fail[ms], terminal[ms], superfail[ms];
d41 bool present[ms];
d41 int z = 1;

d41 int val(char c) { return c - 'a'; }

d41 void add(string& p) {
d41     int cur = 0;
d41     for (int i = 0; i < (int)p.size(); i++) {
d41         int& nxt = trie[cur][val(p[i])];
d41         if (nxt == 0) nxt = z++;
d41         cur = nxt;
d41     }
d41     present[cur] = true;
d41     terminal[cur]++;
d41 }

d41 void build() {
d41     queue<int> q;
d41     for (q.push(0); !q.empty(); q.pop()) {
d41         int on = q.front();
d41         for (int i = 0; i < sigma; i++) {
d41             int to = trie[on][i];
d41             int f = (on == 0 ? 0 : trie[fail[on]][i]);
d41             int sf = (present[f] ? f : superfail[f]);
d41             if (!to) {
d41                 to = f;
d41             } else {
d41                 fail[to] = f;
d41                 superfail[to] = sf;
d41                 // merge infos (ex: terminal[to] += terminal[f])
d41                 q.push(to);
d41             }
d41         }
d41     }
d41 }

d41 void search(string& s) {
d41     int cur = 0;
d41     for (char c : s) {
d41         cur = trie[cur][val(c)];
d41         // process infos on current node (ex: occurrences
d41         // += terminal[cur])
d41     }
d41 }
```

Hash.h

Description: C can also be random, operator is $[l, r]$

d41d8c, 28 lines

```
d41 using ull = uint64_t;
d41 struct H {
d41     ull x; H(ull x = 0) : x(x) {}
d41     H operator+(H o) { return x + o.x + (x + o.x < x); }
d41     H operator-(H o) { return *this - ~o.x; }
```

```
d41     H operator*(H o) {
d41         auto m = (_uint128_t)x * o.x;
d41         return H((ull)m + (ull)(m >> 64));
d41     }
d41     ull get() const { return x + !~x; }
d41     bool operator==(H o) const { return get() == o.get(); }
d41     bool operator<(H o) const { return get() < o.get(); }
d41 };
d41 static const H C = (1111e11 + 3;
d41 struct Hash {
d41     vector<H> h, pw;
d41     Hash(string& str) : h(str.size()), pw(str.size()) {
d41         pw[0] = 1, h[0] = str[0];
d41         for (int i = 1; i < str.size(); i++) {
d41             h[i] = h[i - 1] * C + str[i];
d41             pw[i] = pw[i - 1] * C;
d41         }
d41     }
d41     H operator()(int l, int r) {
d41         return h[r] - (l ? h[l - 1] * pw[r - l + 1] : 0);
d41     }
d41 };
```

Kmp.h

Description: $\pi_i[x]$ computes the length of the longest prefix of s that ends at x , other than $s[0..x]$ itself (abacaba -> 0010123).

d41d8c, 10 lines

```
d41 vector<int> pi(const string& s) {
d41     vector<int> p(sz(s));
d41     for(int i = 1; i < sz(s); i++) {
d41         int g = p[i-1];
d41         while (g && s[i] != s[g]) g = p[g-1];
d41         p[i] = g + (s[i] == s[g]);
d41     }
d41     return p;
d41 }
```

KmpAutomaton.h

Description: $go[i][j]$ = length of the longest prefix of s that is a suffix of $s[0..i]$ followed by the letter j (i.e., the next matched prefix length if, at state i , we read letter j).

d41d8c, 17 lines

```
d41 int go[ms][sigma];
d41 int val(char c) { return c - 'a'; }
d41 void automaton(string& s) {
d41     for (int i = 0; i < sigma; i++)
d41         go[0][i] = (i == val(s[0]));

d41     for (int i = 1, bdr = 0; i <= (int)s.size(); i++) {
d41         for (int j = 0; j < sigma; j++) {
d41             go[i][j] = go[bdr][j];
d41         }
d41         if (i < (int)s.size()) {
d41             go[i][val(s[i])] = i + 1;
d41             bdr = go[bdr][val(s[i])];
d41         }
d41     }
d41 }
```

Manacher.h

Description: $p[0][i+1]$ is the length of matches of even length palindrome, starting from $[i, i+1]$.

$p[1][i]$ is the length of matches of odd length palindrome, starting from $[i, i]$.
 $(abaxx \rightarrow p[0] = 00001)$
 $(abaxx \rightarrow p[1] = 01000)$

d41d8c, 17 lines

```
d41 array<vector<int>, 2> manacher(const string& s) {
d41     int n = sz(s);
```

```
d41     array<vector<int>, 2> p={vector<int>(n+1), vector<int>(n
d41     );
d41     for (int z = 0; z < 2; z++) {
d41         for (int i = 0, l = 0, r = 0; i < n; i++) {
d41             int t = r - i + !z;
d41             if (i < r) p[z][i] = min(t, p[z][l + t]);
d41             int L = i - p[z][i], R = i + p[z][i] - !z;
d41             while(L >= 1 && R+1 < n && s[L-1] == s[R+1]){
d41                 p[z][i]++;
d41                 L--;
d41                 R++;
d41             }
d41             if (R > r) l = L, r = R;
d41         }
d41     }
d41     return p;
d41 }
```

MinRotation.h

Description: Finds the lexicographically smallest rotation of a string.

Usage: $\text{rotate}(s.\text{begin}(), s.\text{begin}() + \text{minRotation}(s), s.\text{end}())$

Time: $\mathcal{O}(N)$

d41d8c, 14 lines

```
d41     int minRotation(string s) {
d41         int a = 0, N = s.size(); s += s;
d41         for (int b = 0; b < N; b++) {
d41             for (int k = 0; k < N; k++) {
d41                 if (a+k == b || s[a+k] < s[b+k]) {
d41                     b += max(0, k-1);
d41                     break;
d41                 }
d41                 if (s[a+k] > s[b+k]) { a = b; break; }
d41             }
d41         }
d41         return a;
d41 }
```

SuffixArray.h

Description: $lcp[i]$ is the length of the longest common prefix between the suffixes $s[sa[i]..n-1]$ and $s[sa[i-1]..n-1]$.

If we concatenate multiple strings using separator characters, the separator that appears furthest to the right must be the smallest character in the alphabet.

d41d8c, 31 lines

```
d41     struct SuffixArray {
d41         vector<int> sa, lcp;
d41         SuffixArray(string s, int lim=256) {
d41             s.push_back('$');
d41             int n = sz(s), k = 0, a, b;
d41             vector<int> x(all(s)), y(n), ws(max(n, lim));
d41             sa = lcp = y, iota(all(sa), 0);
d41             for(int j = 0, p = 0; p < n; j = max(1, j*2), lim = p) {
d41                 p = j, iota(all(y), n - j);
d41                 for(int i=0; i<n; i++) {
d41                     if (sa[i] >= j) y[p++] = sa[i] - j;
d41                 }
d41                 fill(all(ws), 0);
d41                 for(int i=0; i<n; i++) ws[x[i]]++;
d41                 for(int i=1; i<lim; i++) ws[i] += ws[i - 1];
d41                 for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
d41                 swap(x, y), p = 1, x[sa[0]] = 0;
d41                 for(int i=1; i<n; i++) {
d41                     a = sa[i - 1], b = sa[i];
d41                     x[b] = p-1;
d41                     if(y[a] != y[b] || y[a+j] != y[b+j]) x[b] = p++;
d41                 }
d41             }
d41             for (int i = 0, j; i < n - 1; lcp[x[i++]] = k) {
d41                 for (k && k--, j = sa[x[i] - 1];
d41                     s[i + k] == s[j + k]; k++);
d41             sa = vector<int>(sa.begin() + 1, sa.end());
d41 }
```

```
d41     lcp = vector<int>(lcp.begin() + 1, lcp.end());
d41 }
d41 };
```

Zfunc.h

Description: $z[i]$ computes the length of the longest common prefix of $s[i:]$ and s , except $z[0] = 0$. (abacaba -> 0010301)

d41d8c, 13 lines

```
d41 vector<int> ZFunc(const string& s) {
d41     int n = sz(s), a = 0, b = 0;
d41     vector<int> z(n, 0);
d41     if (!z.empty()) z[0] = 0;
d41     for (int i = 1; i < n; i++) {
d41         int end = i;
d41         if (i < b) end = min(i + z[i - a], b);
d41         while (end < n && s[end] == s[end - i]) ++end;
d41         z[i] = end - i; if (end > b) a = i, b = end;
d41     }
d41     return z;
d41 }
```

Various (10)

10.1 Misc. algorithms

Dates.h

Description: dateToInt converts Gregorian date to integer (Julian day number). intToDate converts integer (Julian day number) to Gregorian date: month/day/year. intToDay converts Julian day number to day of the week

d41d8c, 23 lines

```
d41 string day[] = { "Mon", "Tue", "Wed", "Thu", "Fri", "Sat",
d41     "Sun" };
d41 int dateToInt(int m, int d, int y) {
d41     return
d41         1461 * (y + 4800 + (m - 14) / 12) / 4 +
d41         367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
d41         3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
d41         d - 32075;
d41 }
d41 void intToDate(int jd, int& m, int& d, int& y) {
d41     int x, n, i, j;
d41     x = jd + 68569;
d41     n = 4 * x / 146097;
d41     x -= (146097 * n + 3) / 4;
d41     i = (4000 * (x + 1)) / 1461001;
d41     x -= 1461 * i / 4 - 31;
d41     j = 80 * x / 2447;
d41     d = x - 2447 * j / 80;
d41     x = j / 11;
d41     m = j + 2 - 12 * x;
d41     y = 100 * (n - 49) + i + x;
d41 }
d41 string intToDay(int jd) { return day[jd % 7]; }
```

MultisetHash.h

d41d8c, 8 lines

```
d41 ull hashify(ull sum) {
d41     sum += FIXED_RANDOM;
d41     sum += 0x9e3779b97f4a7c15;
d41     sum = (sum ^ (sum >> 30)) * 0xbff58476d1ce4e5b9;
d41     sum = (sum ^ (sum >> 27)) * 0x94d049bb133111eb;
d41     return sum ^ (sum >> 31);
d41 }
```

Rand.h

d41d8c, 8 lines

```
d41 mt19937 rng(chrono::steady_clock::now().time_since_epoch()
.dcount());
```

```
// -64
d41 int uniform(int l, int r) { // [l, r]
d41     uniform_int_distribution<int> uid(l, r);
d41     return uid(rng);
d41 }
```

10.2 Dynamic programming

KnuthDP.h

Description: When doing DP on intervals: $dp[i][j] = \min_{i < k < j} (dp[i][k] + dp[k][j]) + f(i, j)$, where the (minimal) optimal k increases with both i and j . This is known as Knuth DP. Sufficient criteria for this are if $f(b, c) \leq f(a, d)$ and $f(a, c) + f(b, d) \leq f(a, d) + f(b, c)$ for all $a \leq b \leq c \leq d$. Another sufficient criteria is: $opt[i][j-1] \leq opt[i][j] \leq opt[i+1][j]$

Time: $\mathcal{O}(N^2)$

```
d41 for (int k = optL; k < opt; k++) add(k);
d41 }

d41 int solve(int N, int M) { // 1-based
d41     for (int i = 0; i <= M; i++) {
d41         for (int j = 0; j <= N; j++) {
d41             dp[i][j] = inf; // base case
d41         }
d41     }
d41     cost = 0; // neutral value
d41     for (int i = 1; i <= N; i++) add(i);
d41     for (int i = 1; i <= M; i++) {
d41         deC(i, 1, N, 1, N);
d41     }
d41     return dp[M][N];
d41 }
```

d41d8c, 22 lines

```
d41 ll knuth() {
d41     memset(opt, -1, sizeof opt);
d41     for (int i = n - 1; i >= 0; i--) {
d41         dp[i][i] = 0; // base case
d41         opt[i][i] = i;
d41         for (int j = i + 1; j < n; j++) {
d41             int optL = (!j ? 0 : opt[i][j - 1]);
d41             int optR = (~opt[i + 1][j] ? opt[i + 1][j] : n - 1);
d41             ll cst = cost(i, j);
d41             dp[i][j] = INF;
d41             optL = max(i, optL), optR = min(j - 1, optR);
d41             for (int k = optL; k <= optR; k++) {
d41                 ll now = dp[i][k] + dp[k + 1][j] + cst;
d41                 if (now <= dp[i][j]) {
d41                     dp[i][j] = now;
d41                     opt[i][j] = k;
d41                 }
d41             }
d41         }
d41     }
d41 }
```

DivideAndConquerDP.h

Description: Divide and Conquer DP maintaining cost, can be used when $opt[i][j] \leq opt[i][j+1]$. In this code everything is 1-based. Memory can be optimized by keeping only the last row

Time: $\mathcal{O}(MN \log N)$

d41d8c, 42 lines

```
d41 void add(int idx) {}
d41 void rem(int idx) {}

d41 void deC(int i, int l, int r, int optL, int optR) {
d41     if (l > r) return;
d41     int j = (l + r) / 2;
d41     for (int k = r; k > j; k--) rem(k);
d41     int opt = optL;
d41     for (int k = optL; k <= min(optR, j); k++) {
d41         // cost = cost[k, j]
d41         int val = dp[i - 1][k - 1] + cost;
d41         if (val < dp[i][j]) {
d41             dp[i][j] = val;
d41             opt = k;
d41         }
d41     }
d41     rem(k);
d41     for (int k = min(optR, j); k >= optL; k--) add(k);
d41     rem(j);
d41     deC(i, l, j - 1, optL, opt);

d41     for (int k = j; k <= r; k++) add(k);
d41     for (int k = optL; k < opt; k++) rem(k);
d41     deC(i, j + 1, r, opt, optR);
d41 }
```

10.3 Optimization tricks

10.3.1 Bit hacks

- `for (int x = m; x; x = (x - 1)&m) { ... }` loops over all subset masks of m (except 0).
- $c = x \& -x$, $r = x + c$; $((r^x) >> 2)/c$ | r is the next number after x with the same number of bits set.
- `rep(b, 0, K) rep(i, 0, (1 << K)) if (i & 1 << b) D[i] += D[i^(1 << b)];` computes all sums of subsets.

10.3.2 Pragmas

- `#pragma GCC optimize ("Ofast")` will make GCC auto-vectorize loops and optimizes floating points better.
- `#pragma GCC target ("avx2")` can double performance of vectorized code, but causes crashes on old machines.
- `#pragma GCC target ("bmi,bmi2,popcnt,lzcnt")` improve bit operations.
- `#pragma GCC optimize("unroll-loops")` self explanatory.