Rho-calculus Constructions

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1 Workpaper week 10

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The notation of the \rho-Calculus.[1] P,Q ::= 0 Nil |x(y).P Input |x\langle P\rangle Lift |\neg x \vdash Drop| P|Q Parallel x,y:= \vdash P \vdash Quot
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1.1 Numbers

This example could be a way of expressing numbers in ρ - Calculus. Where the number takes place as x or y in the notation. By quoting the last number you get a higher number, and by droping a number you get a lower number, and at 0 you get the process P.

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\begin{array}{l} 0: \lceil \mathbf{O} \rceil \\ 1: \lceil \lceil \mathbf{O} \rceil \langle \langle \mathbf{O} \rangle \rangle \rceil \\ 2: \lceil \lceil \mathbf{O} \rceil \langle \langle \lceil \mathbf{O} \rceil \langle \langle \mathbf{O} \rangle \rangle \rangle \rceil \\ 3: \lceil \lceil \mathbf{O} \rceil \langle \lceil \mathbf{O} \rceil \langle \lceil \mathbf{O} \rceil \langle \langle \mathbf{O} \rangle \rangle \rangle \rangle \rangle \rceil \\ 4: \lceil \lceil \mathbf{O} \rceil \langle \lceil \mathbf{O} \rceil \langle \lceil \mathbf{O} \rceil \langle \langle \lceil \mathbf{O} \rceil \langle \langle \mathbf{O} \rangle \rangle \rangle \rangle \rangle \rangle \end{array}
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1.2 Structure

We can express some of the programming terms in ρ -Calculus, but more has to be discovered.

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if a == b then P Can be express in \rho-Calculus as: a\langle \mathbf{0} \rangle | b(x).P
We send \mathbf{0} on channel a and if channel a and channel b is equal then channel
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b will get a **0** and execute P.

P+Q Can be express in ρ -Calculus as: $a\langle 0 \rangle | a(x).P | a(x).Q$ We can make a determanistic choice between P and Q by sending $\mathbf{0}$ on channel a and in parallel two channel a tryies both to recive, but only one would succeed and would run its process P or process Q.

1.3 Yet to be solved

We need to express more programming structures, before we can use the ρ -Calculs to model a blockchain structure or even r-chain. These are a few examples of what we need to express, and more should follow.

- if $a \neq b$ then P
- if a<b then P
- if a>b then P
- if a==b then P else Q

References

[1] L. G. Meredith and Matthias Radestock. A reflective higher-order calculus. 141:49–67, 2005. ISSN 1571-0661. doi: 10.1016/j.entcs.2005.05.016.