

Lab 6: Distributed Arithmetic

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1. Tables

a) Resources

i. Slice Logic

Site Type	Used	Fixed	Available	Util%
Slice LUTs	39	0	20800	0.19
LUT as Logic	39	0	20800	0.19
LUT as Memory	0	0	9600	0.00
Slice Registers	18	0	41600	0.04
Register as Flip Flop	18	0	41600	0.04
Register as Latch	0	0	41600	0.00
F7 Muxes	0	0	16300	0.00
F8 Muxes	0	0	8150	0.00

ii. IO and GT Specific

Site Type	Used	Fixed	Available	Util%
Bonded IOB	31	0	106	29.25
IOB Master Pads	15			
IOB Slave Pads	15			
Bonded IPADs	0	0	10	0.00
Bonded OPADs	0	0	4	0.00
PHY_CONTROL	0	0	5	0.00
PHASER_REF	0	0	5	0.00
OUT_FIFO	0	0	20	0.00
IN_FIFO	0	0	20	0.00
IDELAYCTRL	0	0	5	0.00
IBUFDS	0	0	104	0.00
GTPE2_CHANNEL	0	0	2	0.00
PHASER_OUT/PHASER_OUT_PHY	0	0	20	0.00
PHASER_IN/PHASER_IN_PHY	0	0	20	0.00
IDELAYE2/IDELAYE2_FINEDELAY	0	0	250	0.00
IBUFDS_GTE2	0	0	2	0.00
ILOGIC	0	0	106	0.00
OLOGIC	0	0	106	0.00

iii. Primitives

Ref Name	Used	Functional Category
IBUF	19	I/O
FDRE	18	Flop & Latch
LUT5	16	LUT
OBUF	12	I/O
LUT4	10	LUT
LUT6	8	LUT
LUT3	8	LUT
CARRY4	4	CarryLogic
LUT2	2	LUT
LUT1	2	LUT
BUFG	1	Clock

b) Power

Total On-Chip Power (W)	0.080
Dynamic (W)	0.010
Device Static (W)	0.070
Effective TJA (C/W)	5.0
Max Ambient (C)	84.6
Junction Temperature (C)	25.4
Confidence Level	Low
Setting File	---
Simulation Activity File	---
Design Nets Matched	NA

Dynamic: 0.010W Static: 0.070W

c) Worst Negative Slack

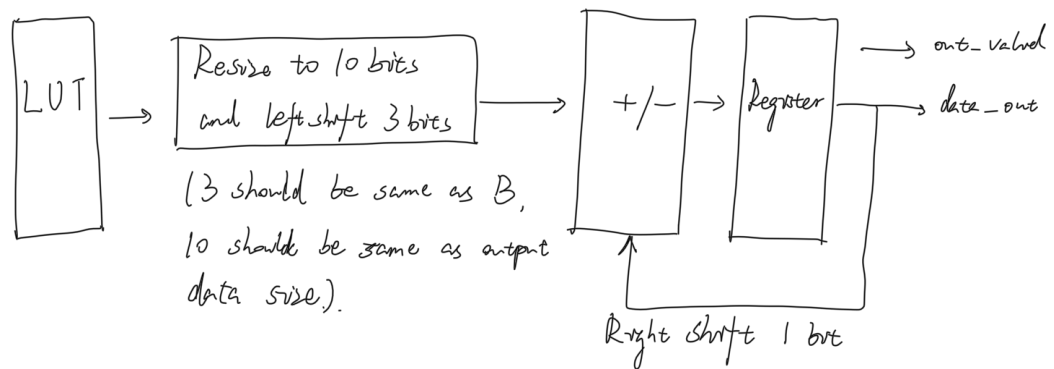
Design Timing Summary					
WNS(ns)	TNS(ns)	TNS Failing Endpoints	TNS Total Endpoints	WHS(ns)	THS(ns)
5.844	0.000	0	20	0.157	0.000
THS Failing Endpoints	THS Total Endpoints	WPWS(ns)	TPWS(ns)	TPWS Failing Endpoints	TPWS Total Endpoints
0	20	4.500	0.000	0	19

2. Questions and Answers

Is there an area-efficient way to do the left-shift which is after the output of LUT?

Yes.

- a) If yes, draw the updated block diagram starting from the LUT on the left to the output data_out on the right.



- b) Write down the mathematical recurrence relation induced from Eq. 3 that proves your thought above.

$$y = -2^B \left(\sum_{n=0}^{N-1} c[n] \times x_B[n] \right) + \sum_{b=0}^{B-1} 2^b \left(\sum_{n=0}^{N-1} c[n] \times x_b[n] \right)$$

$$= -2^B \left(\sum_{n=0}^{N-1} c[n] \times x_B[n] \right) + \sum_{b=0}^{B-1} 2^{b-B} \left(\left(\sum_{n=0}^{N-1} c[n] \times x_b[n] \right) \times 2^B \right)$$

If we denote $\sum_{n=0}^{N-1} c[n] \times x_b[n] = r_b$, we set $B=3$ as an example.

$$y = -2^B \cdot r_3 + \sum_{b=0}^2 2^{b-3} r_b \times 2^3$$

$$= -2^3 \cdot r_3 + 2^{-3} \cdot r_0 \times 2^3 + 2^{-2} \cdot r_1 \times 2^3 + 2^{-1} \cdot r_2 \times 2^3$$

$$= (r_0 \times 2^3 \times 2^{-1} + r_1 \times 2^3) \times 2^{-1} + r_2 \times 2^3 + r_3 \times 2^3$$

And that's how my code accomplish the calculation.

- c) Why the original left-shift is not a good idea and why the updated one is area-efficient? (Answer should not exceed 4 sentences)

Because the original left-shift requires different shift bits, which requires different shift registers with more area to do it. As for the updated one, it uses the same bits shift, so it is area-efficient.